

SE465/ECE453/ECE653/CS447/CS647 Midterm Sample Solutions

Question 1 Solutions

(a) *Fault*: A static defect in software (incorrect lines of code). It is often referred to as bug. For example, the expected behavior is to calculate the sum of **a** and **b**. The line **return a-b;** in the following code snippet is a fault.

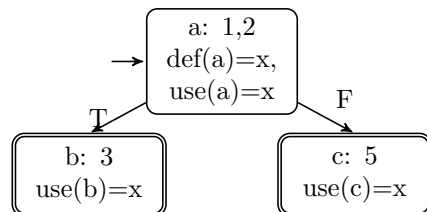
```
1 // should return the sum of a and b
2 int calSum(int a, int b) {
3     return a-b;
4 }
```

Failure: Incorrect behaviour that can be observed. For example, in the code snippet above, when $a = 2$ and $b = 1$, the expected output is 3, while the actual output is 1, which is a failure.

Note that a test case must include expected output. If not, deduct 0.5 mark. However, do not deduct marks for the same mistake twice.

(b)

```
1 int foo(int x) {
2     if (x>0)
3         return 100-x;
4     else
5         return 100+x;
6 }
```



With respect to x : $TR_{ADC} = \{ [a, b] \}$, and $TR_{AUC} = \{ [a, b], [a, c] \}$.

The test set that contains a test case $x = 1$ with expected output 99 satisfies ADC, but not AUC because the path $[a, c]$ is not covered.

It must go by the definition of subsumption. For example, it must use a test case to show that ADC doesn't subsume AUC. If no test case is given, deduct 1 mark. In addition, a test case must contain input and expected output.

(c) Valgrind is a dynamic tool, while all others are static tools.

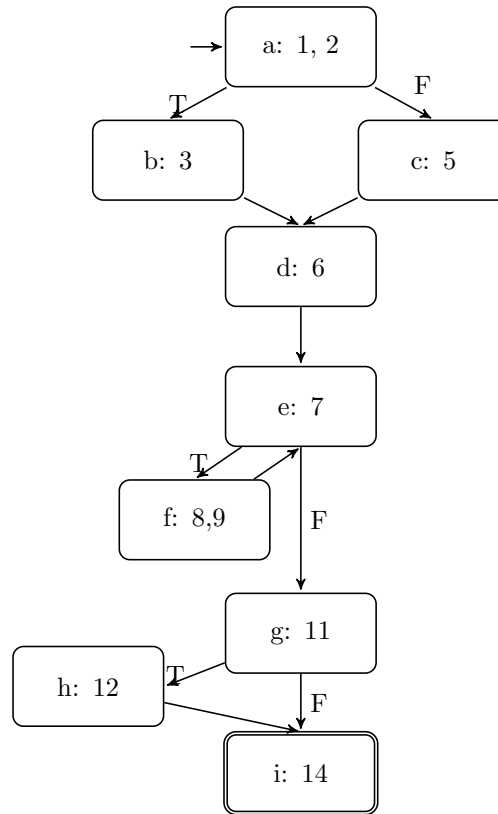
We should accept other reasonable answers.

(d) It can lead to deadlock because the lock is not released before executing **return -EBUSY;**. This leaves the lock unreleased after the method returns, making it impossible for other threads to acquire the lock.

A fix is to add **unlock_kernel();** before **return -EBUSY;**.

Question 2 Solutions

(a)



(b) $TR_{EC} = \{ [a,b],[a,c],[b,d],[c,d],[d,e],[e,f],[e,g],[f,e],[g,h],[g,i],[h,i] \}$ There are 11 edges.

Test set $T = \{t1, t2\}$ is a test set that satisfies EC.

Test case $t1$'s mapping test path is $path(t1) = [a,c,d,e,f,e,g,i]$. $t1$'s input is $(2, 1, _)$ and its expected output is 2, where $_$ denotes any valid integer.

Test case $t2$'s mapping test path is $path(t2) = [a,b,d,e,f,e,g,h,i]$. $t2$'s input is $(2, -1, _)$ and its expected output is 0.5.

(c) No, T doesn't trigger the division by zero bug.

Test case $t3$'s input is $(0, -1, _)$ and its expected output is infinity.

$t3$ triggers a division by zero bug.

Depending on the answer in b), the answer of c) can be yes. If $t3$ is part of test set given b), then the answer for c) is yes.

Question 3 Solutions

Similar to the assignment question. No sample solutions provided.

Marking Scheme: 1 mark for mutation operator, 3 marks for each non-equivalent mutation, 2 marks for test case for each mutation, and 2 marks for mutant and original output (per mutation).

Question 4 Solutions

(a)

```
1 private void visitRegularCall(Node node,int type,Node child, boolean firstArgDone) {
2     while(firstArgDone && (child !=null)) {
3         child = child.getNext();
4     }
5     if(firstArgDone && (child !=null)) {
6         child = child.getNext();
7     }
8     boolean isSpecialCall = node.getProp(Node.SPECIALCALL_PROP) != null;
9     String simpleCallName = null;
10    simpleCallName = getSimpleCallName(node);
11    if (simpleCallName != null && !isSpecialCall) {
12        child = child.getNext().getNext();
13    }
14 }
```

The line numbers in the original code snippet is: 1876, 1878 (optional), 1897, 1899, 1900 (optional), 1906, 1907, 1911 (optional), 1918, 1920, 1922, 1923, 1927, 1931 (optional), and 1953 (optional).

Please pay attention to the lines that are simply {, or }. No need to be strict on them. These lines are optional.

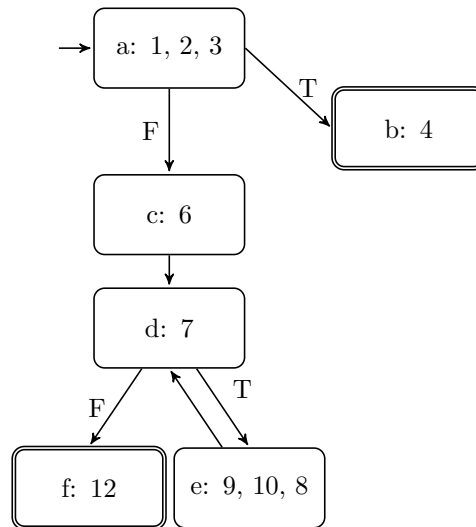
(b) The answers are not unique.

Line 3 in the slice is a definition node for `child` (`child = child.getNext()`). Line 2 is a use node of `child` (`while([...] (child !=null))`). The path from the node of line 3 to the node of line 2 is a du-path with respect to `child` because of the while loop. This path is a du-path because it is def-clear with respect to `child` and it is a simple path.

Another correct answer: The def of `child` in the first `child=child.getNext()` (Line 3) statement reaches the use of `child` of the second `if (child !=NULL)` (Line 5) statement. This path is a du-path because it is def-clear with respect to `child` and it is a simple path.

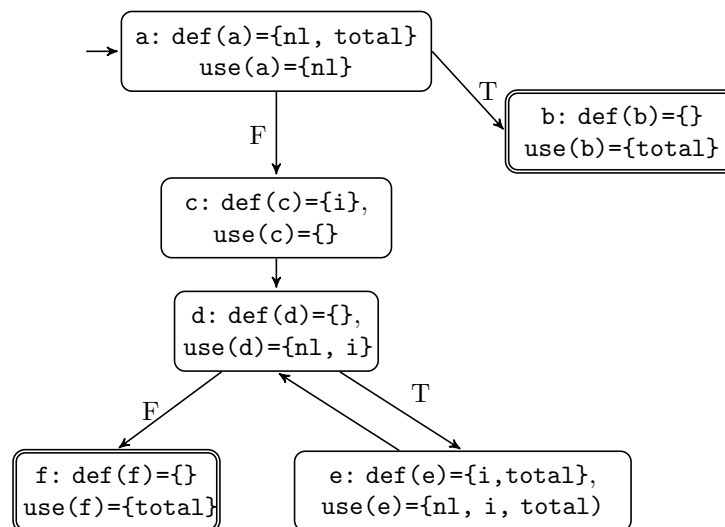
Question 5 Solutions

(a)



Marking Scheme: 1 point for each correct node including outgoing edges.

(b)



Note that **n** is omitted per question (b)'s requirement. Also, a detailed & ordered def-sets & use-sets for node **e** is as follows:

```

// Line 9: use(e) : {i, nl}
           def(e) : {n}
// Line 10: use(e) : {n, total}
           def(e) : {total}
// Line 8: use(e) : {i}
           def(e) : {i}
  
```

It is not required to have this latter presentation of the def and use sets for node **e**, however the order shown here is the correct set for determining AUC requirements in part d).

Marking Scheme: 1 point for each correct node, assuming graph in part a) is correct.

(c) $TR_{PPC} = \{ [a,b], [a,c,d,e], [a,c,d,f], [d,e,d], [e,d,e], [e,d,f] \}$

Marking Scheme: 1 point for each correct test requirement, + 1 additional point if all 6 TRs are correct as a set.

(d)

- $du(nl, a) = \{ [a,c,d], [a,c,d,e] \}$ Note that it is unnecessary to list $[a]$ since the def and use are in the same node, but do not deduct points if students list $[a]$. So $[a]$ is optional.
- $du(i, c) = \{ [c,d], [c,d,e] \}$ Note that TR $[c,d,e]$ ensures the initial def of $i = 0$ reaches the first execution of line 9 and line 8.
- $du(i, e) = \{ [e,d], [e,d,e] \}$ Note that TR $[e,d,e]$ ensures the value from $i++$ reaches the next occurrence of line 9 and $i++$ via looping.

Therefore, $TR_{AUC} = du(nl, a) \cup du(i, c) \cup du(i, e) = \{ [a,c,d], [a,c,d,e], [c,d], [c,d,e], [e,d], [e,d,e] \}$

Marking Scheme: 1 point for each correct test requirement, + 1 additional point if all 6 TRs are correctly enumerated.

(e) Notation: Last Def -> First Use

Listing any of the following 3 excluding the one that was provided in the original question (marked below) receive full marks.

Return from CountLeaves

$(CountLeaves, total, 6) \rightarrow (CountLeavesInNodeList, total, 10)$ for return values (This one was given in the question.)

$(CountLeaves, total, 3) \rightarrow (CountLeavesInNodeList, total, 10)$ for return values

CountLeavesInNodeList calling CountLeaves

$(CountLeavesInNodeList, n, 9) \rightarrow (CountLeaves, m, 2)$ for parameter passing

Return from CountLeavesInNodeList

$(CountLeavesInNodeList, total, 2) \rightarrow (CountLeaves, total, 8)$ for return values (Note that the actual use is at line 8 not line 6).

$(CountLeavesInNodeList, total, 10) \rightarrow (CountLeaves, total, 8)$ for return values (Note that the actual use is at line 8 not line 6).

CountLeaves calling CountLeavesInNodeList

$(CountLeaves, children, 5) \rightarrow (CountLeavesInNodeList, nl, 3)$ for parameter passing

Marking Scheme: 0.5 points for each component of an inter-procedural DU pair. Each pair must be a last-def & first-use pair related to a variable passed as a parameter, or a variable used in a return; has 6 components, and is worth 3 points. If more than 3 pairs were listed, a maximum of 9 points was awarded.