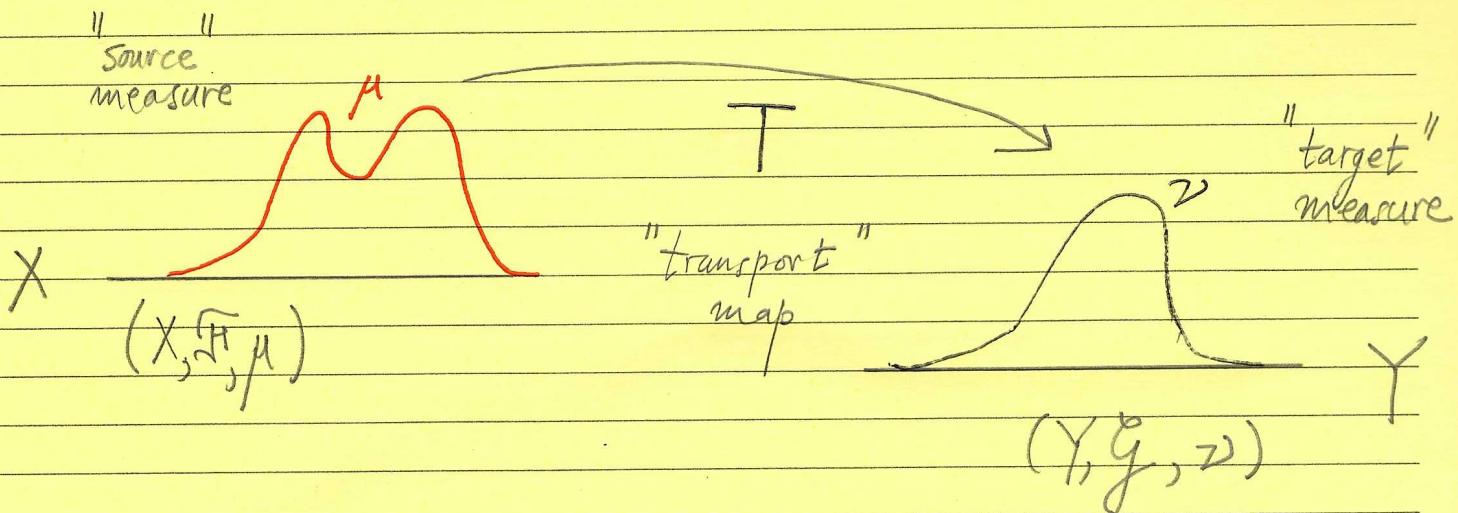


NOTIONS OF OPTIMAL TRANSPORT



Suppose there is a pile of soil in some location X and that, for one reason or another, it needs to be transported to a different location Y . (Construction managers and civil engineers face such questions a lot.)

You know how the pile is "distributed" over X , and how you want it to be distributed over Y , once you have managed to transport it.

But transportation is costly: you need hands to load the soil on trucks in X , drivers and gas to drive to Y , then hands to upload in Y the way you want. So suppose it costs $c(x, y) \geq 0$ to transport a unit of soil from $x \in X$ to $y \in Y$.

Denote the "source" distribution of soil by a probability measure μ on the source space (X, f_i) (the total amount of soil is finite, so you assume $\mu(\Omega) = 1$), and the "target" distribution by a probability

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measure ν on (Y, \mathcal{G}) ($\nu(Y) = 1$, because you assume all soil is transported, without loss or gain of mass). And you devise a "transport map" $T: X \rightarrow Y$ which is measurable, i.e., $T^{-1}(B) \in \mathcal{F}$, $\forall B \in \mathcal{G}$, and has the property

(1)

meaning that

$$\nu(B) = \mu(T^{-1}(B)), \forall B \in \mathcal{G},$$

(2)

$\nu = \mu \circ T^{-1} = T \# \mu$ is the "push-forward" (or "induced") measure on \mathcal{G} , through the action of T on μ .

MONGE (1781) TRANSPORTATION PROBLEM :

Given a cost function $c: X \times Y \rightarrow [0, \infty)$, and two probability measures μ (on (X, \mathcal{F}) , "source") and ν (on (Y, \mathcal{G}) , "target"), find a transport map $T: X \rightarrow Y$ that minimizes the total cost

$$(3) \quad G(T) \triangleq \int_X c(x, T(x)) \mu(dx) \quad \text{over all transport maps.}$$

For reasons we shall try to explain, this is a difficult problem; two centuries had to pass before the first major, significant result about it was proved by Y. BRENIER (1987).

Let us present then, a different, "relaxed" formulation of this problem which DOES admit general solution.

Linear Programming: A motivational example

Suppose X consists of N bakeries at locations x_1, \dots, x_N , and Y of M coffee shops at locations y_1, \dots, y_M . The i^{th} bakery produces an amount $\alpha_i \geq 0$ of "bread", and the j^{th} coffee shop needs an amount $\beta_j \geq 0$; assume that "supply equals demand", and normalize: $\sum_{i=1}^N \alpha_i = 1 = \sum_{j=1}^M \beta_j$.

(Unlike in the MONGE

formulation (where all the supply of a given bakery needs to be sent to only one coffee shop), we allow now each bakery to supply several coffee shops; and each coffee shop to buy from multiple bakeries.

Assuming a cost $c(x_i, y_j)$ for moving one unit from x_i to y_j , let us look for a matrix $\Gamma = (\gamma_{ij})_{1 \leq i \leq N, 1 \leq j \leq M}$ such that

- the amount of bread from x_i to y_j is $\gamma_{ij} \geq 0$,
- $\alpha_i = \sum_{j=1}^M \gamma_{ij}, \forall i=1, \dots, N$ (all supply is consumed)
- $\beta_j = \sum_{i=1}^N \gamma_{ij}, \forall j=1, \dots, M$ (all demand is met)
- Γ minimizes the total transportation cost

$$\sum_{i=1}^N \sum_{j=1}^M \gamma_{ij} c(x_i, y_j).$$

This leads to the

KANTOROVICH (1942, 1948) relaxation of the MONGE problem.

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COUPLINGS: With the probability spaces (X, \mathcal{F}, μ) and (Y, \mathcal{G}, ν) as before, we introduce now the product measurable space $(X \times Y, \mathcal{F} \otimes \mathcal{G})$ and on it the collection of all probability measures π with the properties

$$(4) \quad \pi(A \times Y) = \mu(A), \forall A \in \mathcal{F}; \quad \pi(X \times B) = \nu(B), \forall B \in \mathcal{G}.$$

We call such π a coupling of μ, ν .

An equivalent, and often more convenient, formulation of property (4) is

$$(5) \quad \iint_{X \times Y} [\varphi(x) + \psi(y)] \pi(dx, dy) = \int_X \varphi(x) \mu(dx) + \int_Y \psi(y) \nu(dy)$$

for every $\varphi \in L^1(\mu)$, $\psi \in L^1(\nu)$ (or equivalently
 $\varphi \in L^\infty(\mu)$, $\psi \in L^\infty(\nu)$)

and very often $\varphi \in C_b(X)$, $\psi \in C_b(Y)$, depending on context).

The collection of coupling probability measures

$$(6) \quad \mathcal{T}(\mu, \nu) \triangleq \{\pi \in \mathcal{P}(X \times Y) : (4) \text{ holds}\}$$

is non-empty: it contains (at least) the product measure $\mu \otimes \nu$. Every supplier should ship uniformly over all consumers; every consumer should select uniformly over all suppliers.

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KANTOROVICH (1942, 1948) TRANSPORT PROBLEM

Minimize over $\pi \in \Pi(\mu, \nu)$ the linear transportation cost

$$(7) \quad I(\pi) := \iint_{X \times Y} c(x, y) \pi(dx, dy).$$

Ambition: Use the optimal transport cost

$$(8) \quad \inf_{\pi \in \Pi(\mu, \nu)} I(\pi) =: \mathcal{T}_c(\mu, \nu) \leq \mathcal{M}_c(\mu, \nu) := \inf_{\substack{T: X \rightarrow Y \\ T \# \mu = \nu}} C(T)$$

as a kind of "distance"

between the probability measures μ and ν .

MONGE

$T: X \rightarrow Y$
 $T \# \mu = \nu$



Can easily be $\equiv \infty$

- It is easy to construct examples of probability measures μ, ν on $\mathcal{B}(\mathbb{R}^n)$ which charge "small" sets (e.g., line segments, if $n=2$), such that the optimal transfer plans in the KANTOROVICH problem have to split mass; and thus the inequality in (8) is strict, even with both quantities finite.

Remark: We have the probabilistic interpretation

$$(9) \quad \mathcal{T}_c(\mu, \nu) \equiv \inf_{(U, V)} \mathbb{E}^P(c(U, V)).$$

$U \sim \mu, V \sim \nu$

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The special case $X=Y=\mathbb{R}^n$, $c_r(x,y)=|x-y|^r$ for $1 \leq r < \infty$ has special significance, in that

$$(10) \quad W_n(\mu, \nu) \triangleq \left(\mathcal{G}_{c_n}(\mu, \nu) \right)^{1/r}, \quad (\mu, \nu) \in \mathcal{P}_n(\mathbb{R}^n)$$

defines a distance on $\mathcal{P}_n(\mathbb{R}^n)$, the so-called WASSERSTEIN distance of order n . Very important special case is $n=2$, which corresponds to the quadratic WASSERSTEIN distance.

Remark: We can replace here \mathbb{R}^n by a locally compact, separable metric space

BIG BENEFIT of the KANTOROVICH Formulation:

The "relaxed" problem of (7)-(9) admits a solution under very general conditions.

NOBEL Prize for KANTOROVICH.

Methodologies of convex duality.

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THEOREM : KANTOROVICH DUALITY

- X, Y Polish spaces (think of \mathbb{R}^n);
- $c: X \times Y \rightarrow [0, \infty]$ a l.s.c. cost function is given;
- $\mu \in \mathcal{P}(X)$ and $\nu \in \mathcal{P}(Y)$ are given. Consider probability measures $\pi \in \mathcal{P}(X \times Y)$ and functions $\varphi \in L^1(\mu)$, $\psi \in L^1(\nu)$ **NOT** required to be nonnegative;
- $\Pi(\mu, \nu)$ is the collection of probability measures π in $\mathcal{P}(X \times Y)$ with marginals μ, ν as in (6);

$$\Phi_c \triangleq \{(\varphi, \psi) \in L^1(\mu) \times L^1(\nu) : \varphi(x) + \psi(y) \leq c(x, y) \text{ for } \mu\text{-a.e. } x \in X, \nu\text{-a.e. } y \in Y\}$$

PRIMAL COST: $I(\pi) \triangleq \iint_{X \times Y} c(x, y) \pi(dx, dy) \quad \text{as in (7)}$

DUAL PAYOFF:
$$\begin{aligned} J(\varphi, \psi) &\triangleq \int_X \varphi(x) \mu(dx) + \int_Y \psi(y) \nu(dy) \\ &= \iint_{X \times Y} [\varphi(x) + \psi(y)] \pi(dx, dy) \end{aligned}$$
 from (5).

• PRIMAL PROBLEM : $\underline{V} \triangleq \inf_{\pi \in \Pi(\mu, \nu)} I(\pi)$

• DUAL PROBLEM : $\bar{V} \triangleq \sup_{(\varphi, \psi) \in \Phi_c} J(\varphi, \psi)$

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Then the primal and dual value functions are the same, and the primal problem admits a solution:

there exists $\pi^* \in \mathcal{P}(\mu, \omega)$ with

$$I(\pi^*) = \underline{V} = \bar{V}.$$