

Functors, Applicative Functors and Monoids

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2017-02-14

Outline

- 1 Functors, Applicative Functors and Monoids
- 2 Functors redux
- 3 Functor Laws
- 4 Applicative functors
- 5 The newtype keyword
- 6 Monoids
- 7 Lists are monoids

Polymorphism on a higher level

- Types are not part of a hierarchy
- We can think about how they should act
 - then connect them with typeclasses

Functors defined

Typeclasses define functions

- Eq *defines* concrete types that are equatable
 - functions ('=') and ('/')
- Ord *defines* concrete types that 'orderable'
 - implements the 'compare' function
- Enum *defines* concrete types that enumerable
 - defines '..' a range

Example (List Functor Examples)

- `map:: (a -> b) -> [a] -> [b]`

```
instance Functor [] where  
  fmap = map
```

Functor code in the repl

List Functor in the repl

```
:t map
fmap (*2) [1..3]
map (*2) [1..3]
i==i

map :: (a -> b) -> [a] -> [b]
[2,4,6]
[2,4,6]
```

Maybe Functor examples

Example (Maybe Functor Examples)

- type `Maybe a = Nothing | Just a`

```
instance Functor Maybe where
  fmap f (Just x) = Just (f x)
  fmap f Nothing = Nothing
```


Maybe Functor code in the repl

Maybe Functor in the repl

```
:t fmap
fmap (++) " HEY GUYS IM INSIDE THE JUST") (Just "Something serious")
fmap (++) " HEY GUYS IM INSIDE THE JUST") Nothing
fmap (*2) (Just 200)
fmap (*2) Nothing
i==i

fmap :: Functor f => (a -> b) -> f a -> f b
Just "Something serious. HEY GUYS IM INSIDE THE JUST"
Nothing
Just 400
Nothing
```

Either Functor examples

Example (Either Functor Examples)

- `data Either e a = Left e | Right a`

```
instance Functor (Either a) where
    fmap f (Right x) = Right (f x)
    fmap f (Left x)  = Left x
```

IO is a Functor

IO is a Functor

```
instance Functor IO where
    fmap f action = do
        result <- action
        return (f result)
```

Play with the IO

Reversing a string

```
main = do line <- getLine
  let line' = reverse line
  putStrLn $ "You said " ++ line' ++ " backwards!"
  putStrLn $ "Yes, you really said" ++ line' ++ " backwards"
```

Now we can do the same with fmap

```
main = do line <- fmap reverse getLine
  putStrLn $ "You said " ++ line ++ " backwards!"
  putStrLn $ "Yes, you really said" ++ line ++ " backwards!"
```

A general Functor example

A general Functor example

- Let's say we want to reverse a string and upcase it
- And we want to interleave '-'

```
import Data.Char
import Data.List

main = do line <- fmap (intersperse '-'
    . reverse
    . map toUpper) getLine
    putStrLn line
```

Functor in the repl

Let's look at the type of

- a partially applied fmap

`fmap (replicate 3) (Functor f) => f a -> f [a]`

```
fmap (replicate 3) [1,2,3,4]
fmap (replicate 3) (Just 4)
fmap (replicate 3) (Right "blah")
fmap (replicate 3) (Left "foo")
1==1
```

```
[[1,1,1],[2,2,2],[3,3,3],[4,4,4]]
Just [4,4,4]
Right ["blah","blah","blah"]
Left "foo"
```

If functors mean that something can be mapped over...

- then calling 'fmap' on a functor should
 - map a function over the functor

Functor Law intuition

If functors mean that something can be mapped over...

- then calling 'fmap' on a functor should
 - map a function over the functor
- Nothing else

The First Functor Laws

Definition (The First Functor Law)

states that if we map the identity (`id`) function over a functor, we get the functor

- $\text{fmap id} = \text{id}$

Identity in the Repl

Identity functions in the repl

```
fmap id (Just 3)
```

```
id (Just 3)
```

```
fmap id [1..5]
```

```
id [1..5]
```

```
fmap id []
```

```
fmap id Nothing
```

```
1==1
```

```
Just 3
```

```
Just 3
```

```
[1,2,3,4,5]
```

```
[1,2,3,4,5]
```

```
[]
```

```
Nothing
```

The Second Functor Law

Definition (The Second Functor Law says)

The Second Functor Law says that composing two functions and then mapping the composed function over a functor is the same as first mapping one function over the functor and then mapping the other one.

- $\text{fmap } (f.g) = \text{fmap } f . \text{fmap } g$
- $\text{fmap } (f.g) F = \text{fmap } f (\text{fmap } g F)$

Composition in the Repl

Composition functions in the repl

```
fmap ((+1).(*2)) (Just 3)
fmap (+1) (fmap (*2) (Just 3))
fmap ((+1).(*2)) [1..5]
fmap (+1) (fmap (*2) [1..5])
1==1
```

```
Just 7
```

```
Just 7
```

```
[3,5,7,9,11]
```

```
[3,5,7,9,11]
```

What if we map a multi-parameter function over a functor?

- Look at the type signature

```
let a = fmap (*) [1..4]
```

```
:t a
```

```
fmap (\f -> f 9) a
```

```
1==1
```

```
a :: [Integer -> Integer]
```

```
[9,18,27,36]
```

What if we want to take a function out of a Just

Let's take a Just (3 *) and map

and map it over Just 5

```
class (Functor f) => Applicative f where
  pure :: a -> f a
  (<*>) :: f (a -> b) -> f a -> f b
```

Maybe Applicative

Let's look at the Applicative for Maybe

```
instance Applicative Maybe where
  pure = Just
  Nothing <*> _ = Nothing
  (Just f) <*> something = fmap f something
```

Maybe Applicative inside the repl

Using the Maybe Applicative

```
:m Control.Applicative
Just (+3) <*> Just 9
pure (*2) <*> Just 10
pure (+3) <*> Just 9
Just (++"!!") <*> Just "Go now"
Nothing <*> Just "woot"
1==1
```

```
Just 12
Just 20
Just 12
Just "Go now!!"
Nothing
```


Fmap as an infix operator

Control.Applicative exports a function called `<$>`

which is `fmap` as an infix operator

```
(<$>) :: (Functor f) => (a->b) -> f a -> f b
```

```
f <$> x = fmap f x
```

Compare Applicatives in the repl

Infix fmap in the repl

```
(++) <$> Just "John " <*> Just "Travolta"
```

```
(++) "John " "Travolta"
```

```
1==1
```

```
Just "John Travolta"
```

```
John Travolta
```

Lists are Applicative Functors

Definition (Definition of the Applicative for a list)

- Literally a Cartesian product of functions and list values

```
instance Applicative [] where
  pure x = [x]
  fs <*> xs = [f x | f <- fs, x<- xs]
```

Applicative Functors of lists in the repl

Applicative Functors of lists in the repl

```
[(*0),(+100),(^2)] <*> [1..4]
```

```
[(+),(*)] <*> [1,2]<*> [3,4]
```

```
(++) <$> ["ha","heh","hmm"] <*> ["?","!","."]
```

```
1==1
```

```
[0,0,0,0,101,102,103,104,1,4,9,16]
```

```
[4,5,5,6,3,4,6,8]
```

```
["ha?","ha!","ha.","heh?","heh!","heh.","hmm?","hmm!","hmm."]
```

IO is an Applicative

Let's see how the IO Applicative is implemented:

```
instance Applicative IO where
    pure = return
    a <*> b = do
        f <- a
        x <- b
        return (f x)
```

Concatenating IO strings

Two ways to concatenate two lines of user input string

- Imperative code

```
myAction :: IO String
myAction = do
  a <- getLine
  b <- getLine
  return $ a ++ b
```

Applicative way to concatenate two lines of user input string

- Applicative code

```
myAction :: IO String
myAction = (++)
  <$> getLine
  <*> getLine
```

The first Applicative Functor Law

Theorem (The first Applicative Functor Law)

$$\text{pure } f \langle * \rangle x = \text{fmap } f \ x$$

Some lessons we've skipped

Defining types

- *data* will define a new algebraic type
- *type* creates a type synonym
- *newtype* creates new types from old types

Applicative Functor in two ways

function left, each argument right

```
:m Control.Applicative
[(+1),(*100),(*5)] <*> [1..3]
1==1

[2,3,4,100,200,300,5,10,15]
```

function left, every argument right

```
getZipList $ ZipList [(+1),(*100),(*5)]
-- getZipList $
-- ZipList [(+1),(*100),(*5)]
-- <*> ZipList [1,2,3]
1==1

[2,200,15]
```

The newtype keyword

'newtype' takes one type and wrap it

- to present it as another type

`ZipList [a] }`

- data can have multiple value constructors

'data' to make new types

- Here are additive and multiplicative types with multiple constructors

```
data Profession = Fighter | Archer | Wizard
data Species = Human | Elf | Orc | Goblin
data PlayerCharacter = PlayerCharacter Species Profession
```

Using newtype to drive typeclass properties

newtype

```
newtype CharList = CharList {getCharList :: [Char]} deriving(Eq)
CharList "this will be shown!"
CharList "benny" == CharList "benny"
CharList "benny" == CharList "oisters"
1==1

CharList {getCharList = "this will be shown!"}
True
False
```

Monoid Definition

Definition (Monoid definition)

A data type, category or set is a **monoid** if it has a binary operation \bullet which is associative and has an identity.

- $\forall a, b, c \in S, (a \bullet b) \bullet c = a \bullet (b \bullet c)$
- $e \bullet a = a \bullet e = a$

```
class Monoid m where
  mempty  :: m
  mappend :: m -> m -> m
  mconcat :: [m] -> m
  mconcat = foldr mappend mempty
```

Monoid functions defined

Defining the monoid functions

- 'mempty' is just the identity function
- mappend is the binary function
 - it doesn't just append
- mconcat reduces a list of monoid values and reduces them to one by applying mappend

Theorem (The Monoid Laws are just the definition in Haskell)

- $\text{mappend mempty } x = x$
- $\text{mappend } x \text{ mempty} = x$
- $\text{mappend } (\text{mappend } x \ y) \ z = \text{mappend } x \ (\text{mappend } y \ z)$

Example (List is a monoid)

- `[]` with `(++)` is a monoid
 - `id = ""`
- Natural numbers with `(*)` is a monoid
 - `id = 1`
- Natural numbers with `(+)` is a monoid
 - `id = 0`