

FOCS HW4

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Problem 7.4(c)

$A_0=1; A_1=2; A_n=2A_{n-1}-A_{n-2}+2$ for $n \geq 2$

$$A_2=2A_1-A_0+2=5$$

$$A_3=2A_2-A_1+2=10$$

Guess $P(n)$: $A_n=n^2+1$, for $n \geq 0$

$P(0)$: $A_0=0^2+1=1$, which is True

Let's Assume $P(n)$ and prove $P(n+1)$:

$$A_{n+1} = 2A_n - A_{n-1} + 2$$

$$= 2(n^2+1) - ((n-1)^2+1) + 2$$

$$= 2n^2+2-n^2+2n-1-1+2$$

$$= n^2+2n+2$$

$$= (n+1)^2+1$$

$P(n+1)$ is True

By induction, Our guess is True

Problem 7.56

(a) $M(0,k) = 0$, don't need drop the egg from the first floor

$M(n,1) = n$, drop one time per floor

$M(n,k) = \log_2(n)$, binary relationship (try $n/2$, then $n/4$ or $3n/4$ depend on if it breaks)

(b) $M(n,k) \rightarrow (x) \rightarrow$:

If it breaks, $M(n,k) = M(x,k-1) + 1$

If it doesn't break, $M(n,k) = M(n-x,k) + 1$

Basis: $M(0,k) = 0$, $M(n,1) = n$

(worst cases)

Since we have infinite eggs, a binary search will be a good solution. So the number of egg drop will be logarithm.

When the egg breaks, the second egg drop will be in range of 0 to x

We need $\log_2 x$ of eggs to drop if this egg break

When the egg survives, we need to move up to a range of x to n

So we need $\log_2(n-x)$ of eggs if the egg survive

(c) $M(7,3) = 3$, $M(8,3) = 3$, $M(9,3) = 4$

Problem 8.6

Recursive Definition:

1. $1 \in A$ (base case)
2. $x \in A \rightarrow 2 \cdot x \in A$ (constructor)

(a) Every element of your set is a non-negative power of 2

$P(n)$: n is a non-negative power of 2

$P(1)$: $1=2^0$, which is True

Let's assume $P(n)$ and prove $P(n+1)$:

$a_n = 2^k$ for $k \geq 0$,

child of n is $2x = 2^{k+1}$ for $k+1 \geq 0$

The child is also a non-negative power of two

By structural induction, every element in set A is a non-negative power of 2

(b) Every non-negative power of 2 is in your set

Let's assume $x = 2^k$ and $x \notin A$

Because of $x/2 = 2^{k-1}$ that x can be constructed by $x/2$

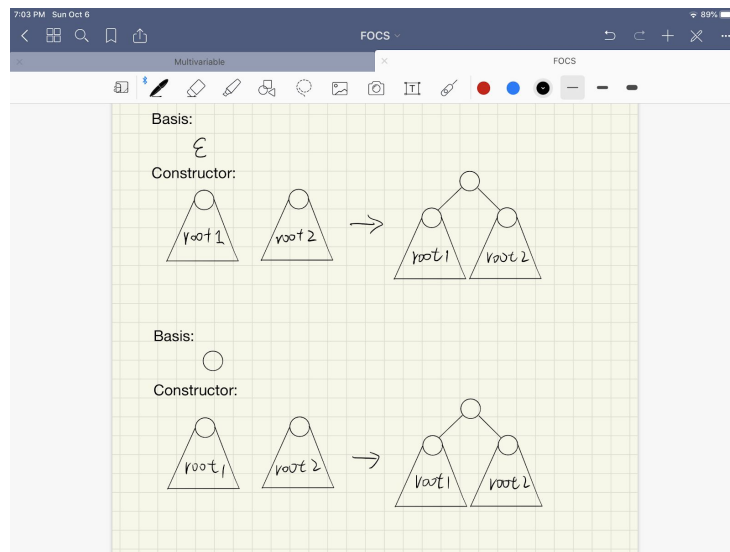
So, $x/2 \notin A$

When it goes on: $x/4 = 2^{k-2} \notin A$, $x/8 \notin A \dots x/2^k = 2^{k-k} = 1 \notin A$

However, 1 is the base case of A

By contradiction, Every non-negative power of 2 is in set A

Problem 8.18



(a) Give examples, with derivations, of RBTs and RFBTs with 5,6 and 7 vertices

	RBT	RFBT
5 vertices		
6 vertices		Not possible for even number
7 vertices		

(b) Prove by structural induction that every RFBT has an odd number of vertices

$P(n)$: when RFBT tree has the height of n , the number of vertices is $2^{n+1}-1$ which is odd, for $n \geq 0$.

$P(0)$: number of vertices when the height is 0: $2-1=1$. Which is odd. True

Let's assume $P(n)$ and prove $P(n+1)$:

$$2^{n+1+1}-1=2*2^{n+1}-2+1=2(2^{n+1}-1)+1$$

$2(2^{n+1}-1)$ is even, so $2(2^{n+1}-1)+1$ is odd

$P(n+1)$ is True, By induction, every RFBT has an odd number of vertices

Problem 9.3

(b)

$$\begin{aligned} & \sum_{i=1}^n \sum_{j=1}^n (i+j) \\ &= \sum_{i=1}^n \left(\sum_{j=1}^n i + \sum_{j=1}^n j \right) \\ &= \sum_{i=1}^n \left(i^2 + \frac{1}{2}i(i+1) \right) \\ &= \sum_{i=1}^n \left(\frac{3}{2}i^2 + \frac{1}{2}i \right) \\ &= \sum_{i=1}^n \frac{3}{2}i^2 + \sum_{i=1}^n \frac{1}{2}i \\ &= \frac{3}{2} \sum_{i=1}^n i^2 + \frac{1}{2} \sum_{i=1}^n i \\ &= \frac{3}{2} \cdot \frac{1}{6}n(n+1)(2n+1) + \frac{1}{2} \cdot \frac{1}{2}n(n+1) \\ &= \frac{1}{4}n(n+1)(2n+2) \\ &= \frac{1}{2}n(n+1)^2 \end{aligned}$$

(e)

$$\begin{aligned} & \sum_{i=0}^n \sum_{j=0}^m 2^{i+j} \\ &= \sum_{i=0}^n \left(\sum_{j=0}^m 2^i * 2^j \right) \\ &= \sum_{i=0}^n \left(2^i \sum_{j=0}^m 2^j \right) \\ &= \sum_{i=0}^n \left(2^i * (2^{m+1} - 1) \right) \\ &= (2^{m+1} - 1) \sum_{i=0}^n 2^i \\ &= (2^{m+1} - 1) (2^{n+1} - 1) \end{aligned}$$

Problem 9.37

	a	b	c	d	e	f
i	$i \in O(a)$	$b \in O(i)$	$c \in O(i)$	$i \in O(d)$	$i \in O(e)$	$i \in O(f)$
ii	both	$ii \in O(b)$	$ii \in O(c)$	$ii \in O(d)$	both	$ii \in O(f)$
iii	$a \in O(iii)$	$b \in O(iii)$	$c \in O(iii)$	both	$e \in O(iii)$	neither
iv	both	$b \in O(iv)$	$c \in O(iv)$	$iv \in O(d)$	$e \in O(iv)$	$iv \in O(f)$
v	$v \in O(a)$	both	$c \in O(v)$	both	both	$f \in O(v)$

Problem 9.44(a)

Give upper and lower bounds and the asymptotic (big-Theta) behavior for

$$\sum_{i=1}^n \frac{i^2}{i^3+1}$$

It is decreasing and approaching to 0

$$\int_{m-1}^n f(x) dx \geq \sum_{i=m}^n f(i) \geq \int_m^{n+1} f(x) dx$$

Let $u=i^3+1$, $du=3i^2 dx$

$$f(x) dx = \frac{1}{3u}$$

$$\int_0^n \frac{1}{3u} du \geq \sum_{i=1}^n \frac{i^2}{i^3+1} \geq \int_1^{n+1} \frac{1}{3u} du$$

$$\int_0^n \frac{1}{3} \ln|i^3+1| \geq \sum_{i=1}^n \frac{i^2}{i^3+1} \geq \int_1^{n+1} \frac{1}{3} \ln|i^3+1|$$

$$\frac{1}{3} \ln|n^3+1| \geq \sum_{i=1}^n \frac{i^2}{i^3+1} \geq \frac{1}{3} \ln|(n+1)^3+1| - \frac{1}{3} \ln 2$$