

CS2500 Homework 4

Evan Wilcox

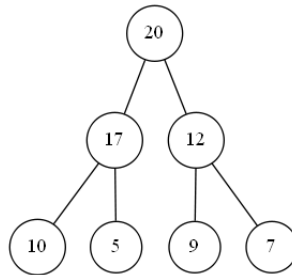
Due March 12, 2019

1. 6.1-6

Index	1	2	3	4	5	6	7	8	9	10
Value	23	17	14	6	13	10	1	5	7	12

No this array is not a max heap because the right child at the 9th index of the node at the 4th index has a value greater than the 4th node.

2. 6.1-7



Index	1	2	3	4	5	6	7
Node	20	17	12	10	5	9	7

This is a 7 element tree and there are 4 leafs.

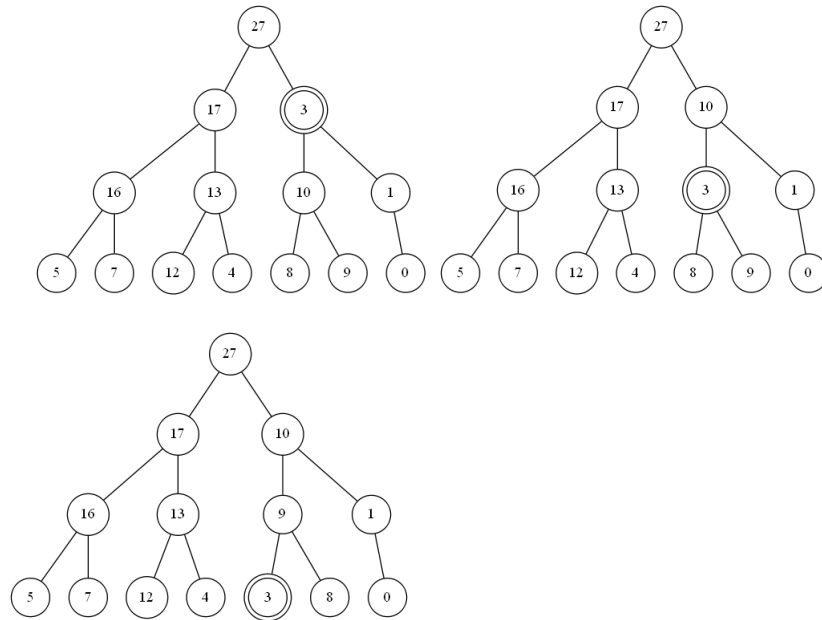
The first leaf is at index $\lfloor n/2 \rfloor + 1 = 4$

The second leaf is at index $\lfloor n/2 \rfloor + 2 = 5$

The third leaf is at index $\lfloor n/2 \rfloor + 3 = 6$

The fourth leaf is at index $\lfloor n/2 \rfloor + 4 = 7$

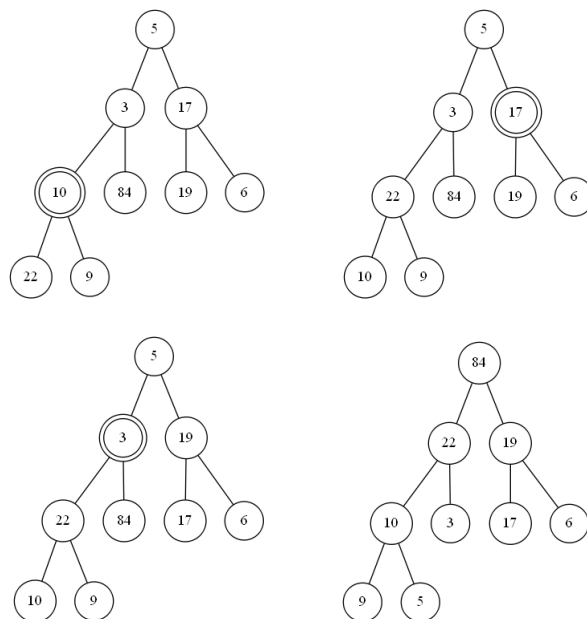
3. 6.2-1



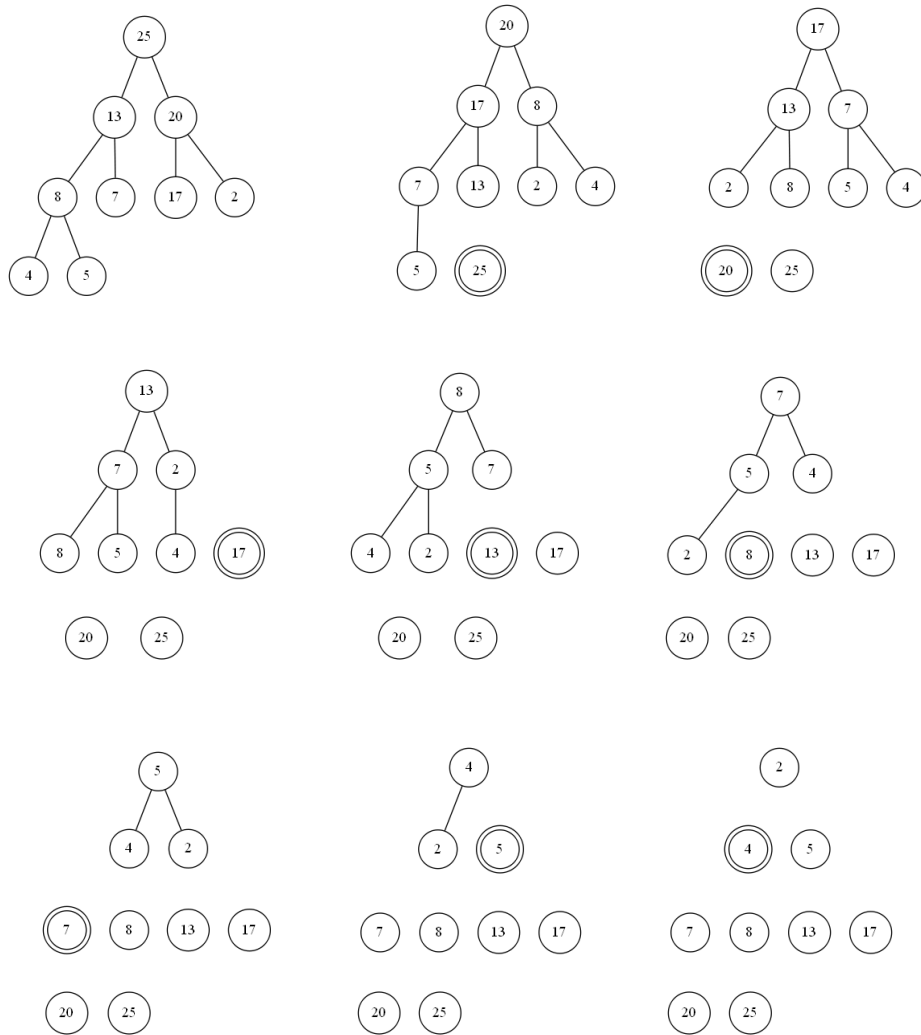
4. 6.2-3

Calling $\text{MAX-HEAPIFY}(A, i)$ when the element $A[i]$ is larger than its children has no effect because that node is already a max heap.

5. 6.3-1



6. 6.4-1



A

2	4	5	7	8	13	17	20	25
---	---	---	---	---	----	----	----	----

7. **6.4-2** – show this invariant implemented as a c assert statement

```
for(int j = 1; j <= i; j++)
{
    if(2 * j < heap-size)
    {
        # left child
        assert(A[j] >= A[2 * j]);
    }

    if(2 * j + 1 < heap-size)
    {
        # right child
        assert(A[j] >= A[2 * j + 1]);
    }

    for(int k = j+1; k < A.length; k++)
    {
        # elements at end of array
        assert(A[j] < A[k])
    }
}
```

8. **6.4-3**

The running time of HEAPSORT on an array A of length n that is already sorted in increasing order is $n \lg n$ because the array must still be converted to a max heap. The running time on an array that is sorted in decreasing order is already a max heap but MAX-HEAPIFY still cost $n \lg n$ so the runtime is $n \lg n$.