CS2200 Homework 4

Evan Wilcox

Due March 19, 2019

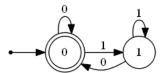
1. Write a Python NDFSA+ ϵ simulator.

```
# MDFSA+e Simulator
class Simulator():
    def __init__(self, file):
        f = open(file, "r")
              self.states = []
self.delta = []
outputs = []
l = []
                       1.append(line)
                      if line[0] == 'A':
    self.alphabet = line[2:]
                      if line[0] == 'S':
    self.states.append([line[2:len(line)-4], int(line[len(line)-2:len(line)-1])])
                       if line[0] == 'B':
    self.beginState = line[2:-1]
                      if line[0] == 'D':
    s = line[2:line.index(',')]
    c = line[line.index(',')+2:line.index(',', line.index(',')+1)]
    e = line[line.index(',', line.index(',')+1)+2:-1]
    self.delta.append([s, c, e])
                       if line[0] == 'T':
    t = line[2:-1]
    o = self.run(t)
    outputs.append(o)
              w = open(file, 'w')
for line in 1:
    if line[0] == '0':
        line = line[:2] + outputs[0] + line[-1:]
        outputs = outputs[1:]
       elif d[0] == state and d[1] == c:
| state = d[2]
               if [state, 1] in self.states:
    return "Accepted"
```

2. Write a program that can generate a Graphviz file from either a .fsa or .ndfsa file.

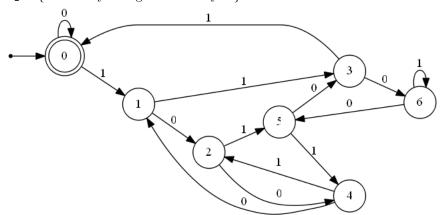
```
class Generator():
    def __init__(self, file):
    f = open(file, "r")
         self.states = []
         self.delta = []
              if line[0] == 'A':
                  self.alphabet = line[2:]
              if line[0] == 'S':
                  self.states.append([line[2:len(line)-4], int(line[len(line)-2:len(line)-1])])
              if line[0] == 'B':
                  self.beginState = line[2:-1]
              if line[0] == 'D':
                  s = line[2:line.index(',')]
                  c = line[line.index(',')+2:line.index(',', line.index(',')+1)]
e = line[line.index(',', line.index(',')+1)+2:-1]
self.delta.append([s, c, e])
         f.close()
         output = file[:file.find('.')] + ".dot"
         w = open(output, 'w')
         w.write("digraph finite_state_machine {\n")
         w.write(tab + 'rankdir=LR;\n')
w.write(tab + '_ize="8,5"\n\n')
         w.write(tab + 'node [shape = point] x\n')
          for state in self.states:
               if state[1]:
                   w.write("doublecircle] " + state[0] + "\n")
                   w.write("circle] " + state[0] + "\n")
          w.write("\n" + tab + "x \rightarrow " + self.beginState + "\n")
          for delta in self.delta:
              w.write(tab + delta[0] + " -> " + delta[2] + ' [label = "')
               for d in self.delta:
                   if delta[0] == d[0] and delta[2] == d[2] and delta != d:
                        delta[1] += (", " + d[1])
self.delta.remove(d)
               w.write(delta[1] + '"]\n')
          w.write("}")
          w.close()
```

- 4. For each of the following, find a FSA automaton that recognizes the language or prove that there is no FSA that recognizes the language.
 - (a) $L_1 = \{$ all binary strings divisible by 2 $\}$



- A 01
- S 0, 1
- S 1, 0
- В 0
- D 0, 0, 0
- D 0, 1, 1
- D 1, 0, 0
- D 1, 1, 1
- Т
- O Accepted
- T 1
- ${\tt O}$ Rejected
- T 0
- O Accepted
- T 11
- O Rejected
- T 00
- O Accepted
- T 10
- O Accepted
- T 01
- O Rejected
- T 100101010100
- O Accepted
- T 0011001101001101
- O Rejected
- T 11111110000000110
- O Accepted

(b) $L_2 = \{$ all binary strings divisible by 7 $\}$



- A 01
- S 0, 1
- S 1, 0
- S 2, 0
- S 3, 0
- S 4, 0
- S 5, 0
- S 6, 0
- D 0,
- ВО
- D 0, 0, 0
- D 0, 1, 1
- D 1, 0, 2
- D 1, 1, 3
- D 2, 0, 4
- D 2, 1, 5
- D 3, 0, 6
- D 3, 1, 0
- D 4, 0, 1
- D 4, 1, 2
- D 5, 0, 3
- D 5, 1, 4
- D 6, 0, 5
- D 6, 1, 6
- T
- O Accepted
- T 1
- O Rejected
- T 0
- O Accepted
- T 111
- O Accepted
- T 110
- O Rejected
- T 1110

- O Accepted
- T 11101
- O Rejected
- T 11100
- O Accepted
- T 0011001101001100
- O Accepted
- T 111111100000011111
- O Accepted
- (c) $L_3 = \{$ all unary strings that represent prime numbers $\}$

Let $w = 1^p$ where p is a prime number. $w \in L_3$. L_3 can not be represented by a FSA because FSA can only represent regular languages and L_3 does not produce a regular language. We know this because using the Pumping Lemma, 1^p can be pumped.

(d) $L_4 = \{$ all unary strings that represent composite numbers $\}$

Let $w = 1^p$ where p is a composite number. $w \in L_3$. L_3 can not be represented by a FSA because FSA can only represent regular languages and L_3 does not produce a regular language. We know this because using the Pumping Lemma, 1^p can be pumped.

(e) $L_5 = \{ w \in \{a, b, c\}^* \text{ such that } w \text{ is a palindrome } \}$

Using the pumping lemma, \mathbf{a}^n can be pumped to create an infinite length palindrome.