CS 2200 Spring 2019 HW 05

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Due Tuesday, April 2, 2019, 11:59PM

Submit your HW as a SINGLE PDF file that includes your answers, your programs, and their output. Also include any data files if you used them. Please make sure that your PDF is well-organized so we can quickly find all the pieces of each problem that belong together.

PROBLEMS

- 1. (10 points) Find a deterministic finite-state automaton that recognizes the language, L, consisting of all strings in {a, b}* that contain an odd number of b's such that there is at least one a between every two b's in the string. Verify your solution by running some non-trivial cases on your simulator. Include your graphviz drawing, your .gv file, and your .ndfsa file.
- 2. (5 points) Construct a regular expression that generates the language, L, defined in Problem 1.
- 3. (10 points) Let $A = \{c,d\}$. Let $L_2 \subseteq A^*$ be the language consisting of all strings not in the language, L_3 , which is generated by the regular expression $c^*d^*c^*$. Find a deterministic finite-state automaton that recognizes L_2 . Verify your solution by running some non-trivial cases on your simulator. Include your graphviz drawing, your .gv file, and your .ndfsa file.
- 4. (5 points) Find a regular expression that generates the language, L_2 , defined in Problem 3
- 5. (10 points) Let $L_3 \subseteq \{f, g\}^*$ be the language of all strings generated by the regular expression (fg + ffg + fgf)*. Construct a non-deterministic finite-state machine that recognizes L_3 . Verify your solution by running some non-trivial cases on your simulator. Include your graphviz drawing, your .gv file, and your .ndfsa file.
- 6. (10 points) Using the subset construction described in class, construct a deterministic finite-state machine based on the non-deterministic finite-state machine that

you constructed in Problem 5. Verify your solution by running some non-trivial cases on your simulator. Include your graphviz drawing, your .gv file, and your .fsa file.

- 7. (10 points) Let G be the context-free grammar { {Q,V}, {c,d},{ Q \rightarrow QQ|V, V \rightarrow cVd|cd}, Q}.
 - (a) (5 points) Describe the language $L_6 = L(G)$. Prove that G is ambiguous.
 - (b) (5 points) Give an unambiguous grammar H, such that L(H) = L(G). Give some justification for the claim that H is unambiguous.
- 8. (20 points) Consider the CFG G = { {S}, {a, b}, {S → aSbS|bSaS|ε}, S }. Prove by induction that L(G) = { w ∈ { a, b}* | the number of a's in w = the number of b's in w }. Note that there are two things to prove here. First, you have to prove that if a string of terminals is generated by G, then the number of a's in it is equal to the number of b's. Second, you have to prove that if a string of a's and b's has an equal number of a's and b's, then it can be generated by G.
- 9. (20 points) Let G = (Vars, Alph, Rules, Start) be a CFG as described below. Using the algorithms discussed in class, convert G into an equivalent grammar that is in Chomsky Normal Form.
 - (a) $Vars = \{ P, Q, R, S, T, U. V. W. X, Y, Z \}.$
 - (b) $Alph = \{ a, b, c \}.$
 - (c) Start = S.
 - (d) Rules =
 - i. $S \rightarrow aXT|YbT|UbZ|UWc|PQT$.
 - ii. $P \rightarrow aT$.
 - iii. $Q \to QT|aQ$.
 - iv. $T \to cT | \epsilon$.
 - v. $U \to aU | \epsilon$.
 - vi. $X \to aX|R|\epsilon$.
 - vii. $R \to aRb|\epsilon$.
 - viii. $Y \to Yb|R|\epsilon$.
 - ix. $V \to bVc|\epsilon$.
 - x. $W \to Wc|V|\epsilon$.
 - xi. $Z \to bZ|V|\epsilon$.