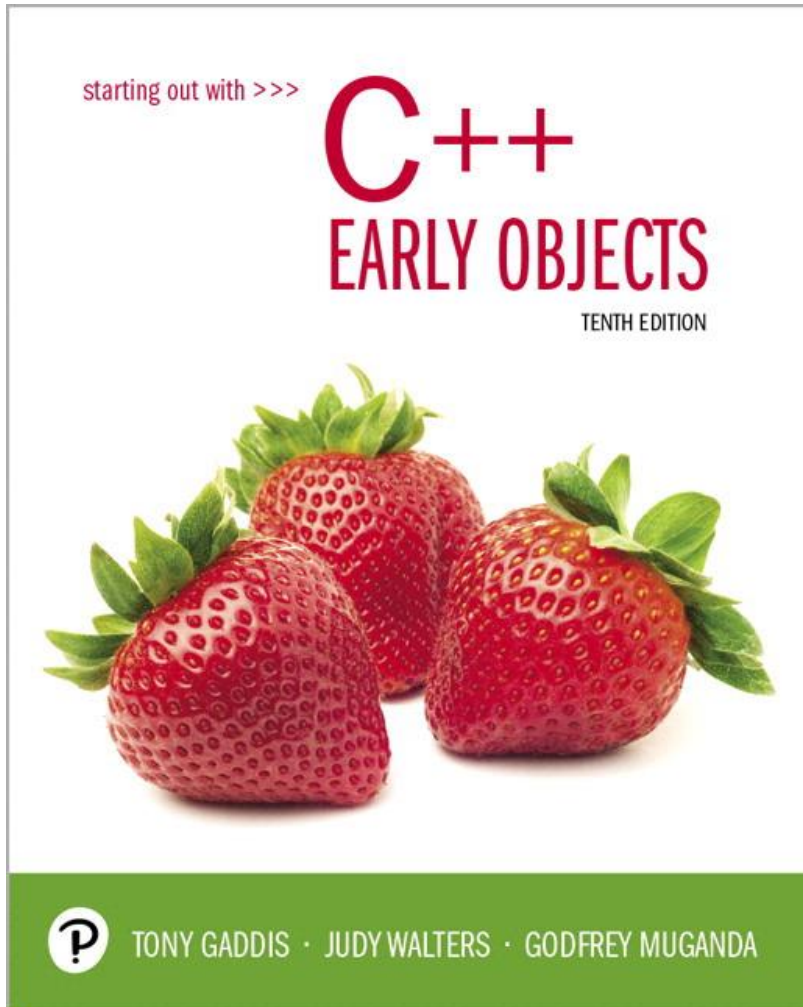


Starting Out with C++ Early Objects

Tenth Edition



Chapter 14

Recursion

Topics

- 14.1 Introduction to Recursion
- 14.2 The Recursive Factorial Function
- 14.3 The Recursive gcd Function
- 14.4 Solving Recursively Defined Problems
- 14.5 A Recursive Binary Search Function
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- 14.7 The Towers of Hanoi
- 14.8 Exhaustive and Enumeration Algorithms
- 14.9 Recursion Versus Iteration

14.1 Introduction to Recursion

- A **recursive** function is a function that calls itself.
- Recursive functions can be useful in solving problems that can be broken down into smaller or simpler subproblems of the same type. A **base case** should eventually be reached, at which time the breaking down (recursion) will stop.

Recursive Functions 1 of 2

Consider a function for solving the count-down problem from some number `num` down to `0`:

- The base case is when `num` is already `0`: the problem is solved and we “blast off!”
- If `num` is greater than `0`, we count off `num` and then recursively count down from `num-1`

Recursive Functions 2 of 2

A recursive function for counting down to 0:

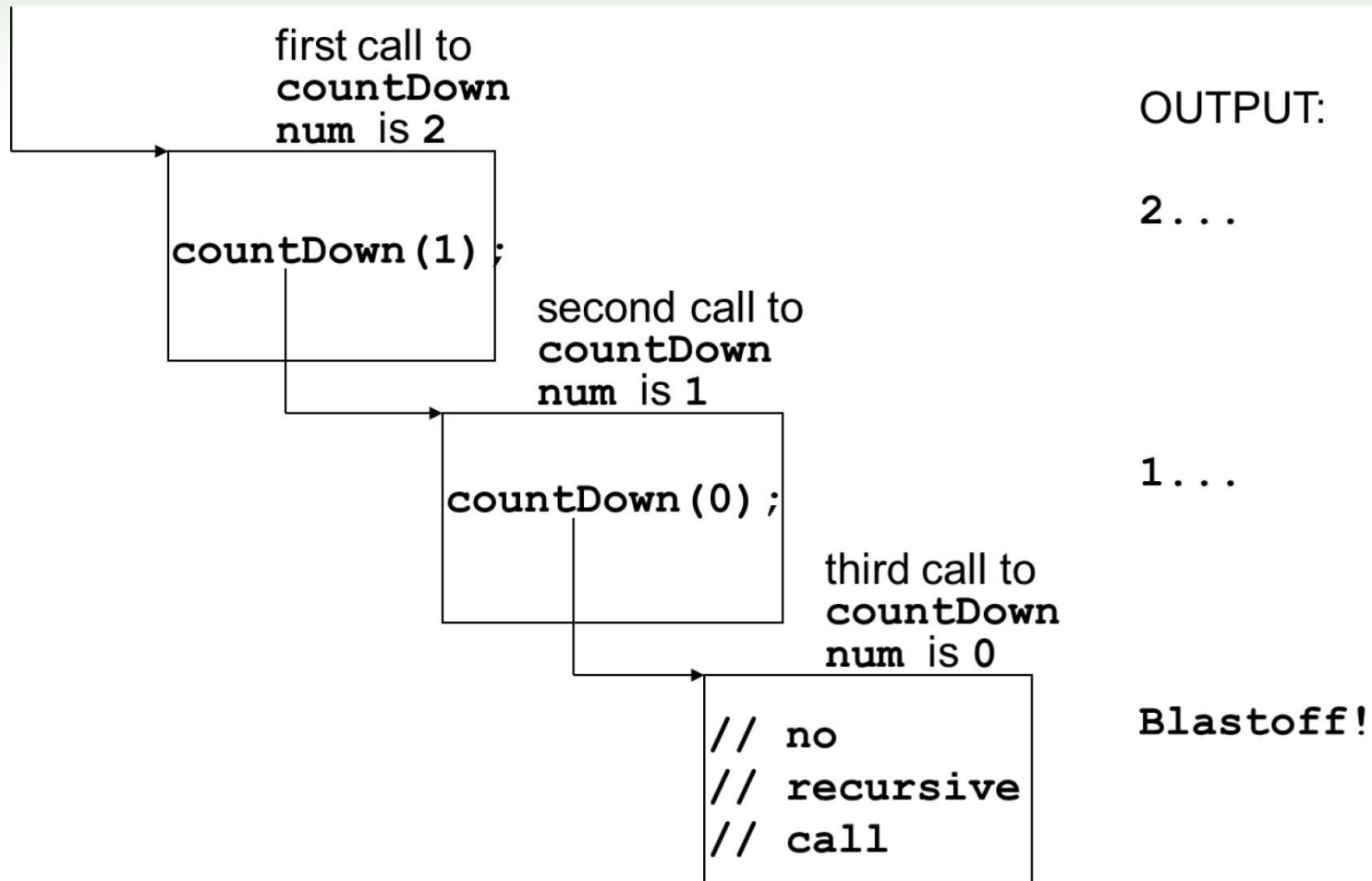
```
void countDown(int num)
{
    if (num == 0)
        cout << "Blast off!";
    else
    {
        cout << num << ". . .";
        countDown(num-1); // recursive
                           // call
    }
}
```

What Happens When Called? 1 of 2

If a program contains a line like **countDown (2) ;**

1. **countDown (2)** generates the output **2 . . .** , then it calls **countDown (1)**
2. **countDown (1)** generates the output **1 . . .** , then it calls **countDown (0)**
3. **countDown (0)** generates the output **Blast off!** , then returns to **countDown (1)**
4. **countDown (1)** returns to **countDown (2)**
5. **countDown (2)** returns to the calling function

What Happens When Called? 2 of 2



Stopping the Recursion 1 of 4

- A recursive function should include a test for the base cases
- In the sample program, the test is:

```
if (num == 0)
```


Stopping the Recursion 2 of 4

```
void countdown(int num)
{
    if (num == 0) // test
        cout << "Blast off!";
    else
    {
        cout << num << "... \n";
        countdown(num-1); // recursive
    } // call
}
```

Stopping the Recursion 3 of 4

- With each recursive call, the parameter controlling the recursion should move closer to the base case
- Eventually, the parameter reaches the base case and the chain of recursive calls terminates

Stopping the Recursion 4 of 4

```
void countDown(int num)
{
    if (num == 0)           // base case
        cout << "Blast off!";
    else
    {
        cout << num << "... \n";
        countDown(num-1); // decrement
        // parameter to approach base case
    }
}
```

What Happens When Called?

- Each time a recursive function is called, a new copy of the function runs, with new instances of parameters and local variables being created
- As each copy finishes executing, it returns to the copy of the function that called it
- When the initial copy finishes executing, it returns to the part of the program that made the initial call to the function

Types of Recursion

- **Direct recursion**
 - a function calls itself
- **Indirect recursion**
 - function A calls function B, and function B calls function A. Or,
 - function A calls function B, which calls ..., which then calls function A

14.2 The Recursive Factorial Function 1 of 2

- The factorial of a nonnegative integer n is the product of all positive integers less than or equal to n
- The factorial of n is denoted by $n!$
- The factorial of 0 is 1

$$0! = 1$$

$$n! = n \times (n-1) \times \dots \times 2 \times 1 \text{ if } n > 0$$

Recursive Factorial Function 2 of 2

- The factorial of n can be expressed in terms of the factorial of $n-1$

$$0! = 1$$

$$n! = n \times (n-1)!$$

- Recursive function:

```
int factorial(int n)
{ if (n == 0)
    return 1; // base
  else
    return n *factorial(n-1);
}
```

14.3 The Recursive gcd Function 1 of 2

- Greatest common divisor (gcd) of two integers x and y is the largest number that divides both x and y with no remainder.
- The Greek mathematician Euclid discovered that
 - If y divides x , then $\text{gcd}(x, y)$ is just y
 - Otherwise, the $\text{gcd}(x, y)$ is the gcd of y and the remainder of dividing x by y

The Recursive gcd Function 2 of 2

```
int gcd(int x, int y)
{
    if (x % y == 0) //base case
        return y;
    else
        return gcd(y, x % y);
}
```

14.4 Solving Recursively Defined Problems

- The natural definition of some problems leads to a recursive solution
- Example: Fibonacci numbers:
 $0, 1, 1, 2, 3, 5, 8, 13, 21, \dots$
- After the initial 0 and 1, each term is the sum of the two preceding terms
- Recursive calculation of the n th Fibonacci number:
$$\text{fib}(n) = \text{fib}(n - 1) + \text{fib}(n - 2);$$
- Base cases: $n == 0, n == 1$

Recursive Fibonacci Function

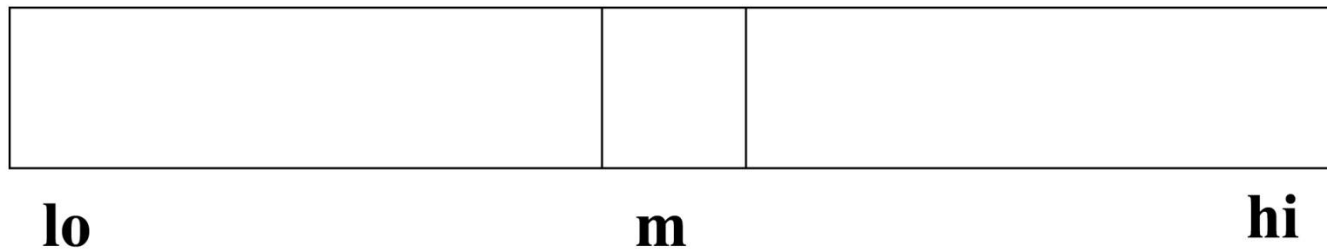
```
int fib(int n)
{
    if (n <= 0)                // base case
        return 0;
    else if (n == 1)           // base case
        return 1;
    else
        return fib(n - 1) + fib(n - 2);
}
```

14.5 A Recursive Binary Search Function

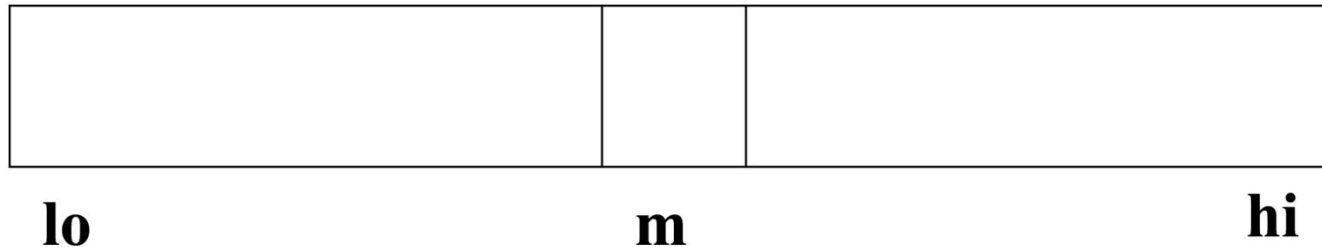
- Assume an array **a** that is sorted in ascending order, and an item **x** to search for
- We want to write a function that searches for **x** within the array **a**, returning the index of **x** if it is found, and returning **-1** if **x** is not in the array

Recursive Binary Search 1 of 3

A recursive strategy for searching a portion of the array from index **lo** to index **hi** is to
set **m** to the index of the middle element of the
array:



Recursive Binary Search 2 of 3



If $a[m] == x$, we found x , so return m

If $a[m] > x$, recursively search $a[lo..m-1]$

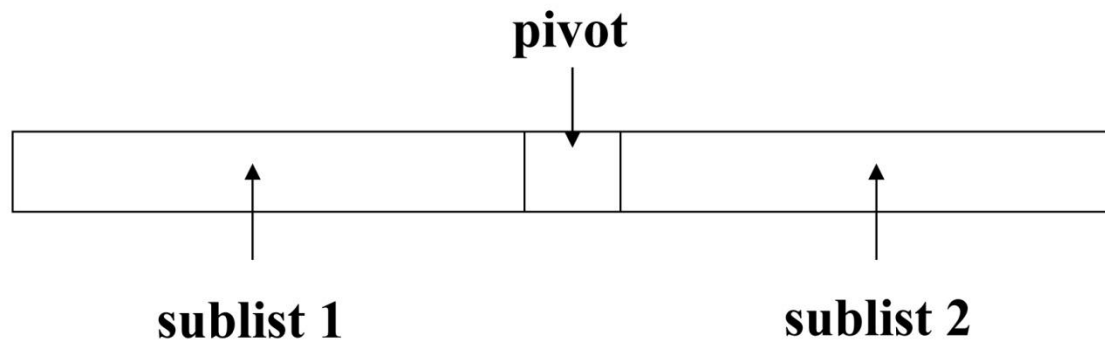
If $a[m] < x$, recursively search $a[m+1..hi]$

Recursive Binary Search 3 of 3

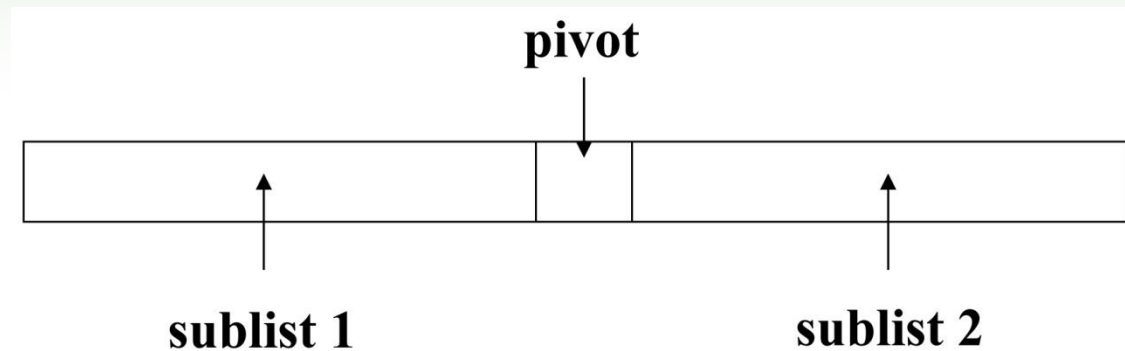
```
int bSearch(const int a[],int lo,
            int hi,int X)
{
    int m = (lo + hi) /2;
    if(lo > hi) return -1;        // base
    if(a[m] == X) return m;      // base
    if(a[m] > X)
        return bsearch(a,lo,m-1,X) ;
    else
        return bsearch(a,m+1,hi,X) ;
}
```

14.6 The QuickSort Algorithm 1 of 2

- Recursive algorithm that can sort an array
- First, determine an element to use as **pivot value**:



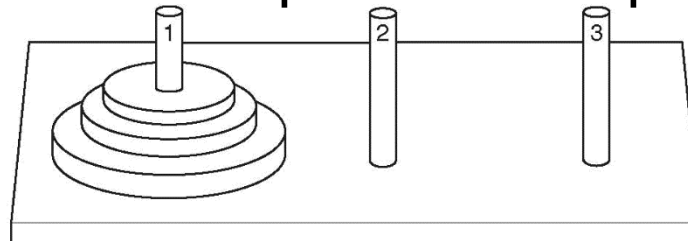
The QuickSort Algorithm 2 of 2



- Then, values are shifted so that elements in sublist1 are $<$ pivot and elements in sublist2 are \geq pivot
- Algorithm then recursively sorts sublist1 and sublist2
- Base case: a sublist has size ≤ 1

14.7 The Towers of Hanoi 1 of 2

- Setup: 3 pegs, one has n disks on it, the other two pegs empty. The disks are arranged in increasing diameter, top \rightarrow bottom
- Objective: move the disks from peg 1 to peg 3, observing these rules:
 - only one disk moves at a time
 - all remain on pegs except the one being moved
 - a larger disk cannot be placed on top of a smaller disk at any time



The Towers of Hanoi 2 of 2

How it works:

| Number of Disks | Sequence of Moves |
|-----------------|---|
| n=1 | Move disk from peg 1 to peg 3. Done. |
| n=2 | Move top disk from peg 1 to peg 2. Move remaining disk from peg 1 to peg 3. Move disk from peg 2 to peg 3. Done. |

Outline of the Recursive Algorithm

If $n==0$, do nothing (base case)

If $n>0$, then

- a. Move the topmost $n-1$ disks from peg1 to peg2
- b. Move the n^{th} disk from peg1 to peg3
- c. Move the $n-1$ disks from peg2 to peg3

end if

14.8 Exhaustive and Enumeration Algorithms

- **Enumeration algorithm**: generate all possible combinations

Example: all possible ways to make change for a certain amount of money

- **Exhaustive algorithm**: search a set of combinations to find an optimal one

Example: change for a certain amount of money that uses the fewest coins

14.9 Recursion vs. Iteration

- Benefits (+), disadvantages(-) for recursion:
 - + Natural formulation of solution to certain problems
 - May not execute very efficiently
- Benefits (+), disadvantages(-) for iteration:
 - + Executes more efficiently than recursion
 - May not be as natural a method of solution as recursion for some problems

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