

# Linked Lists

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# Outline

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- ▶ A linked list is just as slow as an array (actually slower).
- ▶ But a skip list is MUCH faster.



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  - ▶ No hope of a fast addOrChange method for large  $n$ .



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  - ▶ References to the next and previous entries in the list.



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- ▶ The **slides** show how to use this structure to implement a phone directory.



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  - ▶ How can we tell?
  - ▶ Two possibilities.



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  - ▶ Similar to entries in array with index bigger than size.



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- ▶ Draw the diagram of what should happen.
- ▶ Write the line that makes that change happen.



speed



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  - ▶ add is now  $O(1)$
  - ▶ but `put (addOrChangeEntry)` must call find
  - ▶ so it is still  $O(n)$ .
  - ▶ One step forward, two steps back!



# Summary



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  - ▶ Draw the diagram of each change.
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    - ▶ with only two variables.
  - ▶ Keep each step simple!



# This week's application



## This week's application

- ▶ We need a nice application for our Map.



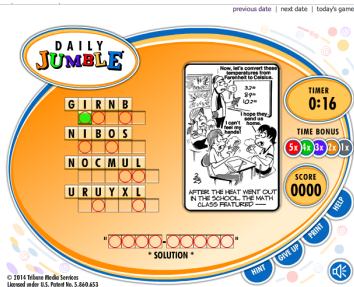
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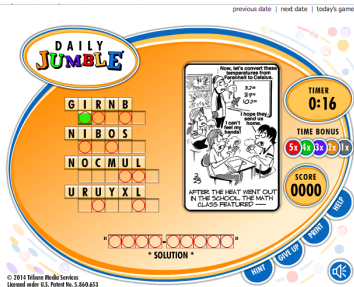
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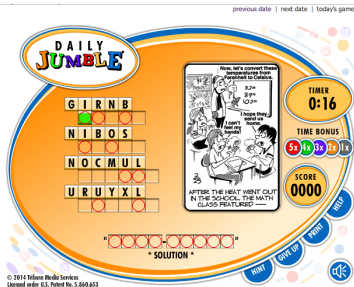
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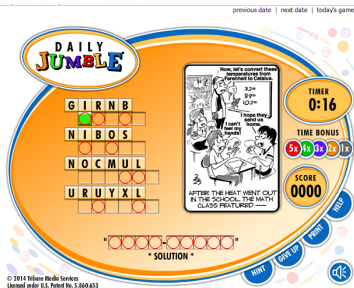
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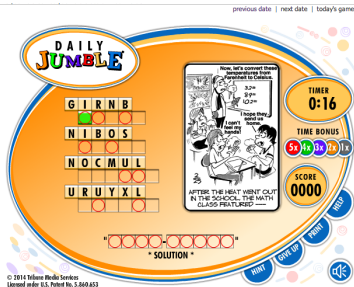


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- ▶ Daily Jumble
- ▶ Need to unscramble words.
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  - ▶ Unscrambled is “computer”.
  - ▶ How can a Map help us to do that?

# Slow Way

## Slow Way

- ▶ We have a dictionary file.



# Slow Way

- ▶ We have a dictionary file.
  - ▶ Read it in.



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- ▶ We have a dictionary file.
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- ▶ What is the running time?
  - ▶ Lookup might be  $O(\log n)$  time, good.
  - ▶ But the number of orderings is  $8! = 40,320$ , bad!.



# Using a Map



# Using a Map

- ▶ Let's use a Map.



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  - ▶ That is "cemoprut".
- ▶ To get ready:





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  - ▶ Read each word from the dictionary file,
  - ▶ Put it into the Map.



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- ▶ To solve a scramble "rtpmceuo":
  - ▶ Alphabetize it to "cemoprtu".



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- ▶ Does anyone see a problem?
  - ▶ The words "dare", "dear", and "read"



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  - ▶ The words "dare", "dear", and "read"
  - ▶ will all be stored under the key "ader".



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  - ▶ The words "dare", "dear", and "read"
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  - ▶ So the value will be "read" because it is last.



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  - ▶ Solution is to use **List<String>** as the value type.



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- ▶ Does anyone see a problem?
  - ▶ The words "dare", "dear", and "read"
  - ▶ will all be stored under the key "ader".
  - ▶ So the value will be "read" because it is last.
  - ▶ Solution is to use **List<String>** as the value type.
  - ▶ But we won't do that this time.



We're going to need a bigger...dictionary.



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- ▶ words.txt doesn't have the solution to "zagboe"



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- ▶ words.txt doesn't have the solution to "zagboe"
- ▶ Let's try dict.txt





We're going to need a bigger...dictionary.

- ▶ words.txt doesn't have the solution to "zagboe"
- ▶ Let's try dict.txt with 483423 words!!



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- ▶ Run Jumble using LinkedHashMap.



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- ▶ Run Jumble using LinkedMap. Seems to be taking a while...



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- ▶  $n$  times  $O(n)$ ?



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- ▶  $n \cdot (c \cdot n) = c \cdot n^2$





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- ▶  $n \cdot (c \cdot n) = c \cdot n^2$
- ▶  $O(n^2)$ !



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- ▶  $O(n^2)$ !
- ▶  $n^2 = 233,697,796,929$ .



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- ▶  $n^2 = 233,697,796,929$ .
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- ▶  $O(n^2)$ !
- ▶  $n^2 = 233,697,796,929$ .
- ▶ A half million squared is a quarter trillion.
- ▶ A computer that can do a billion operations in a second

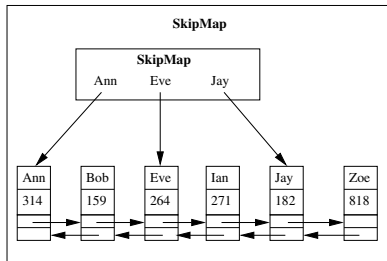


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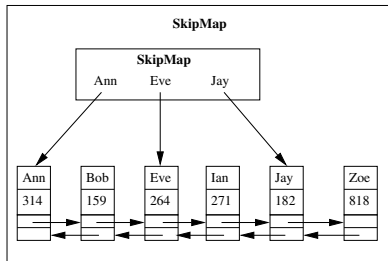
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- ▶ put (addOrChangeEntry) has to call find. Which is  $O(n)$ .
- ▶ What is the  $O()$  to read all  $n$  words?
- ▶  $n$  times  $O(n)$ ?
- ▶  $n \cdot (c \cdot n) = c \cdot n^2$
- ▶  $O(n^2)$ !
- ▶  $n^2 = 233,697,796,929$ .
- ▶ A half million squared is a quarter trillion.
- ▶ A computer that can do a billion operations in a second
- ▶ will take 233 seconds times the number of operations per find.



# SkipMap

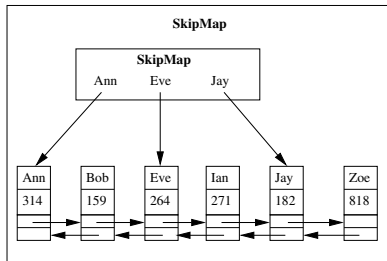


# SkipMap



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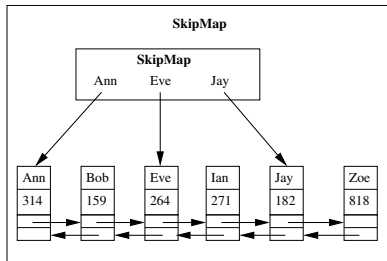
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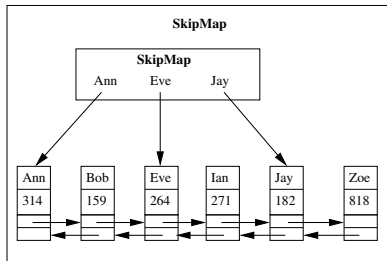


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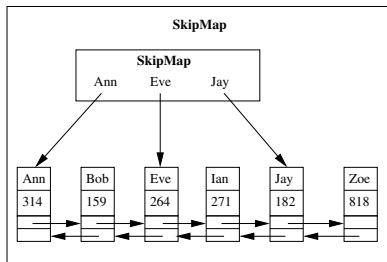
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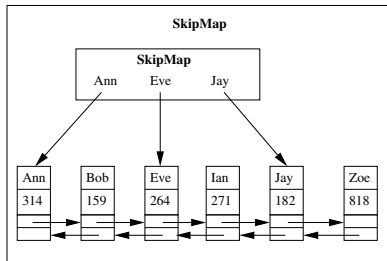
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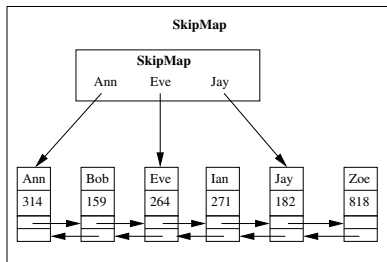
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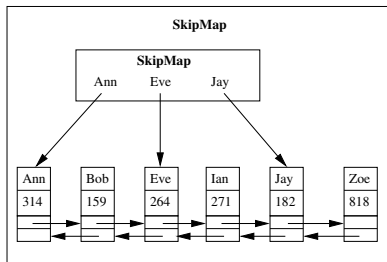
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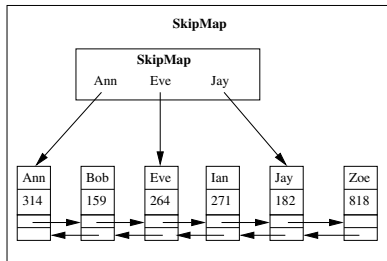
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- ▶ How many people do you expect between Eve and Jay?

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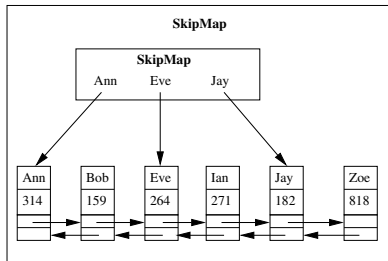
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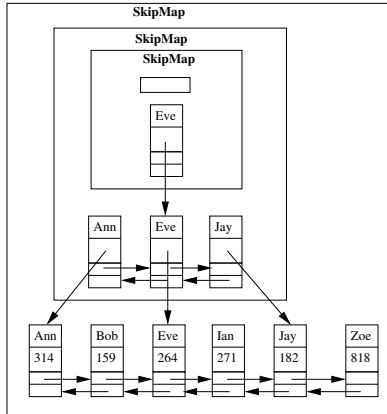
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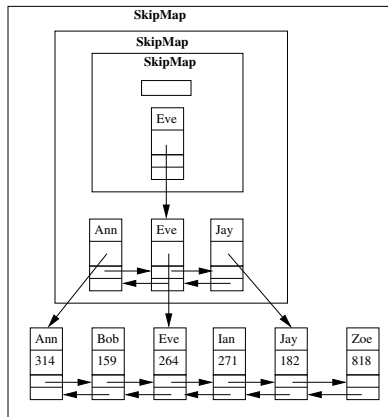
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- ▶ How many people do you expect between Eve and Jay? Just one.
- ▶ So lookup in a SkipMap requires lookup in a SkipMap half the size plus one extra step.
- ▶ Gold coin idea! Find and get are  $O(\log n)$ .



# SkipMap – Gold Coin?

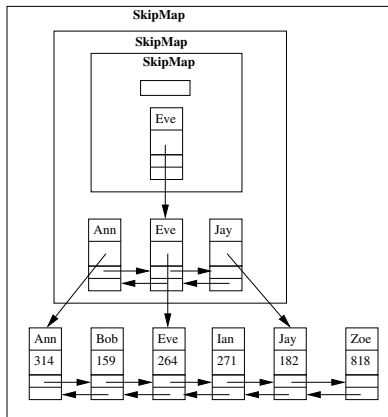


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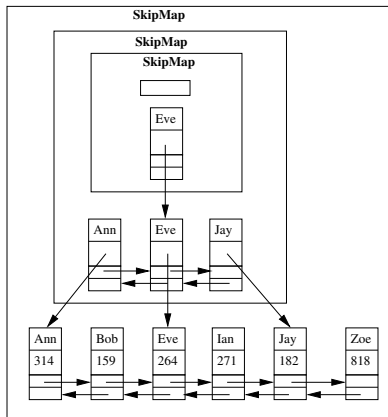
- Suppose the LinkedHashMap in the (outer) SkipMap has 1000 entries.

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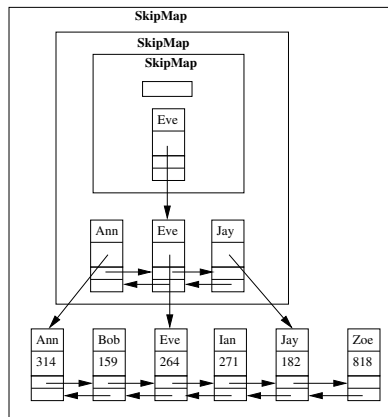
- ▶ Suppose the LinkedMap in the (outer) SkipMap has 1000 entries.
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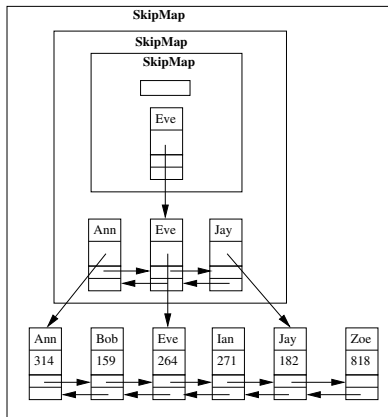
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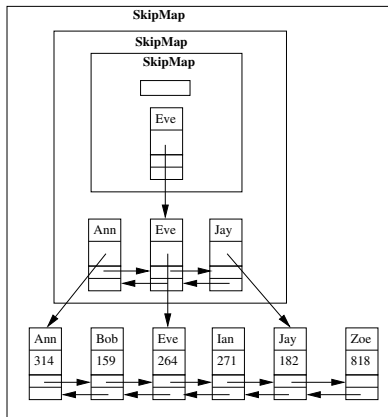
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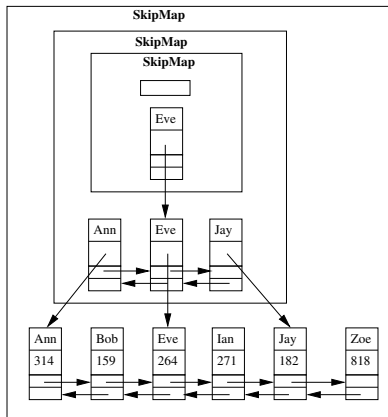
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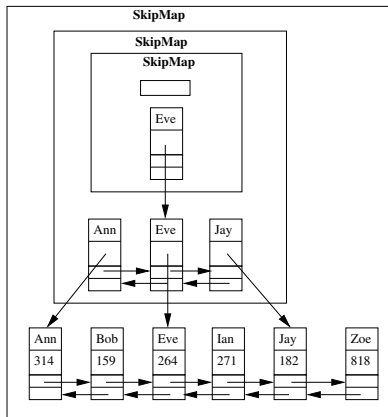


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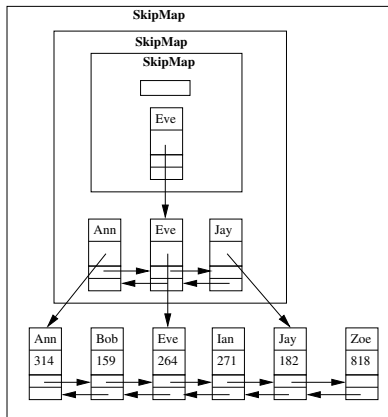
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- ▶ Finding the gold coin in 1000 required
- ▶ 1 weighing

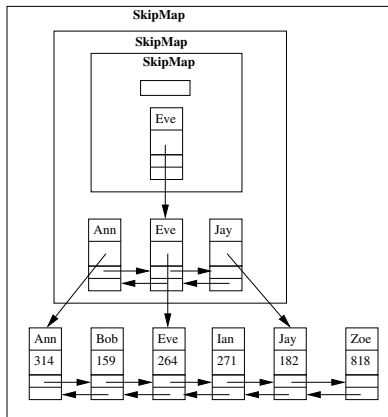
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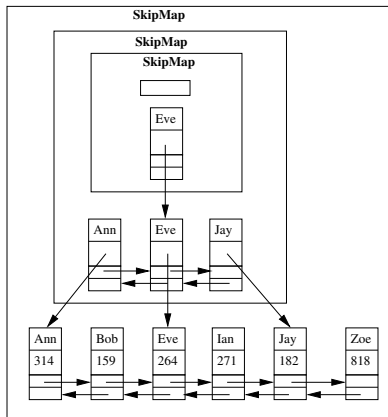
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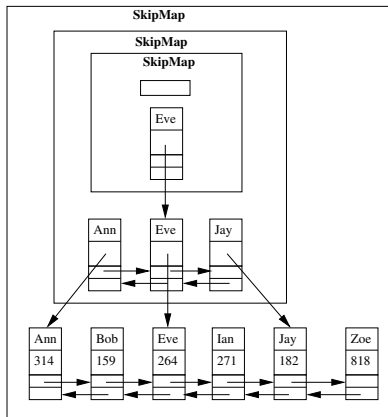
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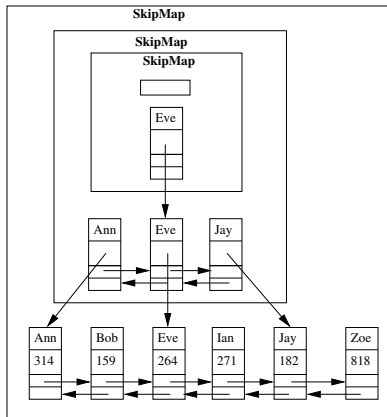
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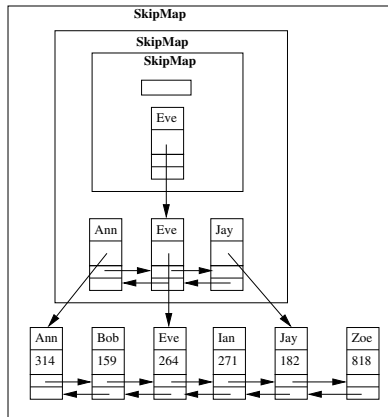
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- ▶ Finding the gold coin in 1000 required
- ▶ 1 weighing plus finding the coin in 500.
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- ▶ find in 500 plus 1 weighing (step forward)?

# SkipMap add and remove

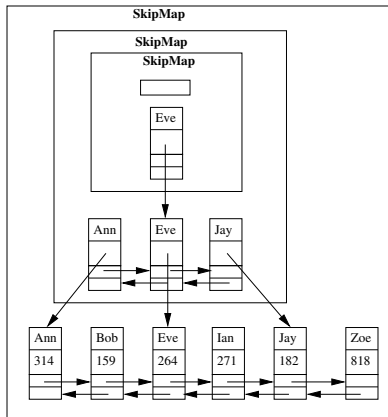


# SkipMap add and remove



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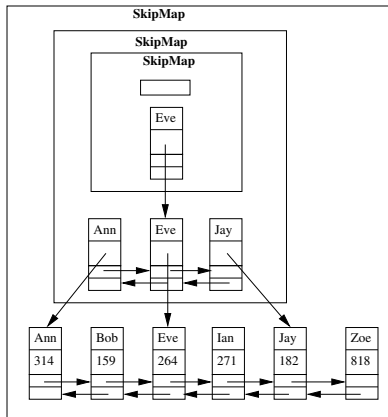
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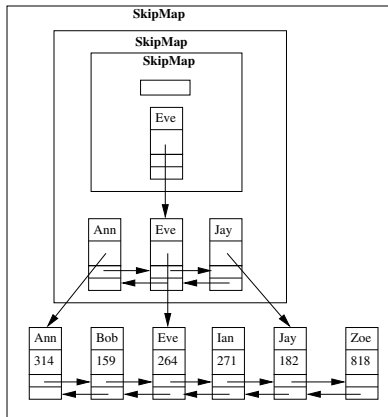


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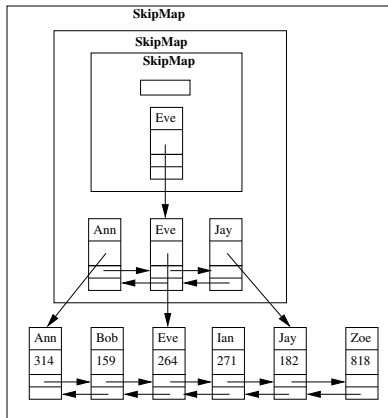
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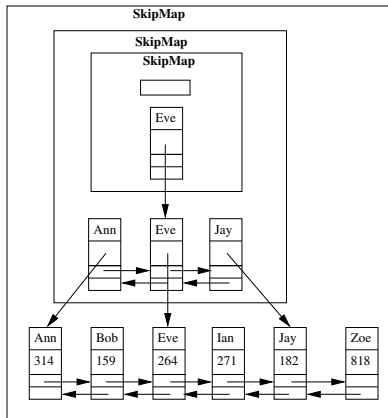
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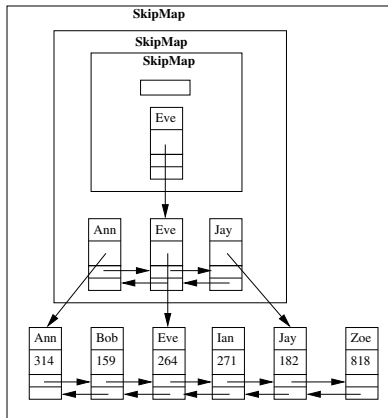
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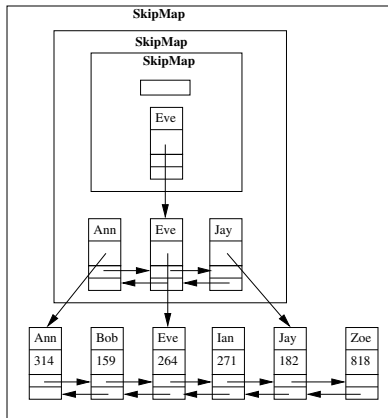
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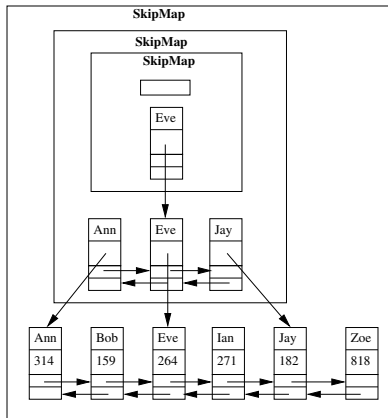


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- ▶ So on average? Two!
- ▶ add and remove are still  $O(\log n)$ .
- ▶ This also means a SkipMap only requires twice as much space as a LinkedMap.