

ECE374 Assignment 2

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T3: Proof of closure

3. Let B and C be languages over $\Sigma = \{0, 1\}$. Define:

$$B \xrightarrow{0} C = \{w \in C \mid \text{for some } x \in B \text{ strings } w \text{ and } x \text{ contain an equal number of } 0\text{'s}\} \quad (1)$$

Show that the class of regular languages is closed under the $\xrightarrow{0}$ operation.

The class of regular languages is indeed closed under the " $\xrightarrow{0}$ " operation. Given a DFA **D1** that recognizes B and a DFA **D2** that recognizes C , we can prove B can be adopted into a NFA **N1** = $(0+1)^*$, to include all possible string in $D2$.

Here's how it works:

1. Based on the question, the important thing is number of 0, so **all 1-transition function can be treated as ϵ - reach**, modifying $D1$ into NFA $N1$.

2. Then for all string in $B1$, we have 1 to be 1^* . For example,

$$101001 \rightarrow 1^*01^*001^* \quad (1)$$

3. Therefore, the 0's in $N1$ can be treated as subsequence, and the language C , $D2$ is included in $N1$. There must exists some x in B can also be recognized as string in C .

By following this construction, it can be shown that the language recognized by the new NFA is exactly $B \xrightarrow{0} C$. Vice versa, $C \xrightarrow{0} B$ also established, this shows that the class of regular languages is closed under the " $\xrightarrow{0}$ " operation.