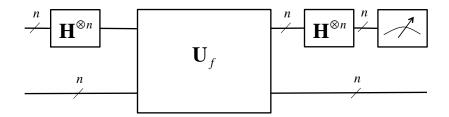
# Simon's Periodicity Algorithm



PURPOSE: Detect Patterns within a Function

# Simon's Periodicity Algorithm Problem Overview

Consider an Unknown Function of the Form:

$$f: \{0,1\}^n \to \{0,1\}^n$$

- Function is "Hidden" in a Black Box
- Known that there exists a string:

$$\mathbf{c} = c_0 c_1 c_2 ... c_{n-1}$$

• Such that for all strings:  $\mathbf{x}, \mathbf{y} \in \{0,1\}^n$  $f(\mathbf{x}) = f(\mathbf{y})$  if and only if  $\mathbf{x} = \mathbf{y} \oplus \mathbf{c}$ 

# Simon's Periodicity Algorithm Problem Overview

- XOR operation is performed bitwise on the strings y and c
- The values of f repeat themselves in some pattern c
- c is called the "period" of f
- Purpose of Simon's Algorithm is to determine c

# Periodic Function Example

- number of bits n=3
- Consider c=101

```
y c x 000 \oplus 101 = 101 \Rightarrow f(000) = f(101) this must hold if 001 \oplus 101 = 100 \Rightarrow f(001) = f(100) f(x) = f(y) over c, 010 \oplus 101 = 111 \Rightarrow f(010) = f(111) the period of f(x) = f(y) over c, the period of f(x) = f(x) = f(x) the period of f(x) = f(x) = f(x) if f(x) = f(x) the period of f(x) = f(x) if f(x) = f(x) the period of f(x) = f(x) if f(x)
```

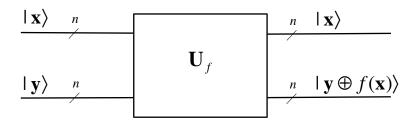
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```

#### **Function Specification**

 Unknown function specified as a unitary operation of the form:



 Setting y=0<sup>n</sup> Provides a Convenient way to evaluate f(x)

#### Classical Solution to Problem

- Evaluate *f*(**x**) on different binary strings
- After Each Evaluation, Check if Function Response has Already Been Found
- If two strings  $\mathbf{x}_1$  and  $\mathbf{x}_2$  are found such that  $f(\mathbf{x}_1) = f(\mathbf{x}_2)$  then it is assured that:

$$\mathbf{x}_1 = \mathbf{x}_2 \oplus \mathbf{c}$$

• How do we find c?

#### Classical Solution to Problem

- Evaluate  $f(\mathbf{x})$  on different binary strings
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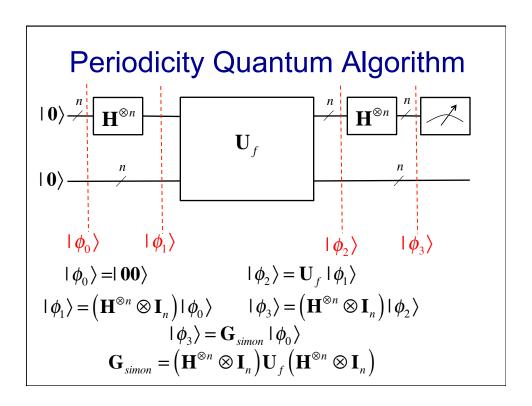
$$\mathbf{x}_{2} \oplus \mathbf{x}_{1} = \mathbf{x}_{2} \oplus \mathbf{x}_{2} \oplus \mathbf{c} = (\mathbf{x}_{2} \oplus \mathbf{x}_{2}) \oplus \mathbf{c} = 0 \oplus \mathbf{c} = \mathbf{c}$$
$$\mathbf{x}_{1} \oplus \mathbf{x}_{2} = \mathbf{c}$$

# **Classical Solution Complexity**

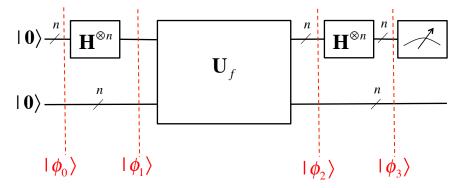
- If f(x) is not one-to-one (it is two-to-one) then a repeat will be found before half of inputs are evaluated
- If more than half of inputs checked with no match, then f(x) is one-to-one and c=0<sup>n</sup>
- Worst case number of evaluations:

$$\frac{2^n}{2} + 1 = 2^{n-1} + 1$$

Exponential Complexity



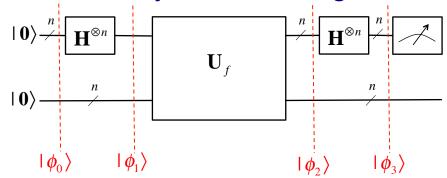
# Periodicity Quantum Algorithm



- Places Upper *n* Qubits in State of Superposition
- **EXAMPLE**: *n*=3:

$$|\phi_{1}\rangle = (\mathbf{H}^{\otimes 3} \otimes \mathbf{I}_{3})|\phi_{0}\rangle = \frac{\sum_{\mathbf{x} \in \{0,1\}^{3}} |\mathbf{x}, \mathbf{0}\rangle}{\sqrt{2}^{3}} = \frac{|000,000\rangle + |001,000\rangle + ... + |111,000\rangle}{\sqrt{8}}$$

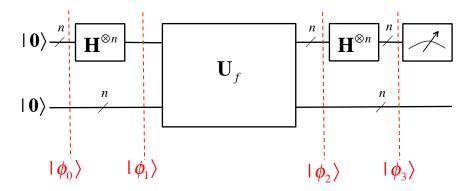
#### Periodicity Quantum Algorithm



•  $\mathbf{U}_f$  causes evaluation of f for the superimposed quantum state

$$|\phi_{2}\rangle = \mathbf{U}_{f}\left(\mathbf{H}^{\otimes 3} \otimes \mathbf{I}_{3}\right)|\phi_{1}\rangle = \frac{\sum_{\mathbf{x} \in \{0,1\}^{3}} |\mathbf{x}, f(\mathbf{x})\rangle}{\sqrt{2^{3}}}$$

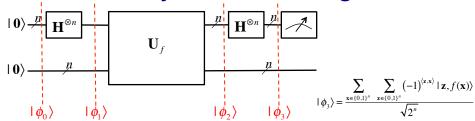
# Periodicity Quantum Algorithm



• n-Dimensional Hadamard Transform applied again

$$|\phi_{3}\rangle = (\mathbf{H}^{\otimes 3} \otimes \mathbf{I}_{3})|\phi_{2}\rangle = \frac{\sum_{\mathbf{x} \in \{0,1\}^{n}} \sum_{\mathbf{z} \in \{0,1\}^{n}} (-1)^{\langle \mathbf{z}, \mathbf{x} \rangle} |\mathbf{z}, f(\mathbf{x})\rangle}{\sqrt{2^{n}}}$$

# Periodicity Quantum Algorithm



- n-Dimensional Hadamard Transform applied again
- For each  $\mathbf{x}$  and  $\mathbf{z}$  in the Summations, Consider when:  $|\mathbf{z}, f(\mathbf{x})\rangle = |\mathbf{z}, f(\mathbf{x} \oplus \mathbf{c})\rangle$
- When this occurs, coefficient is:

$$\frac{\left(-1\right)^{\langle \mathbf{z}, \mathbf{x} \rangle} + \left(-1\right)^{\langle \mathbf{z}, \mathbf{x} \oplus \mathbf{c} \rangle}}{2} = \frac{\left(-1\right)^{\langle \mathbf{z}, \mathbf{x} \rangle} + \left(-1\right)^{\langle \mathbf{z}, \mathbf{x} \rangle \oplus \langle \mathbf{z}, \mathbf{c} \rangle}}{2}$$
$$= \frac{\left(-1\right)^{\langle \mathbf{z}, \mathbf{x} \rangle} + \left(-1\right)^{\langle \mathbf{z}, \mathbf{x} \rangle} \left(-1\right)^{\langle \mathbf{z}, \mathbf{c} \rangle}}{2}$$

# Periodicity Quantum Algorithm

$$\frac{\left(-1\right)^{\langle z,x\rangle}+\left(-1\right)^{\langle z,x\oplus e\rangle}}{2}=\frac{\left(-1\right)^{\langle z,x\rangle}+\left(-1\right)^{\langle z,x\rangle\oplus\langle z,e\rangle}}{2}=\frac{\left(-1\right)^{\langle z,x\rangle}+\left(-1\right)^{\langle z,x\rangle}\left(-1\right)^{\langle z,e\rangle}}{2}$$

• When:  $\langle \mathbf{z}, \mathbf{c} \rangle = 1$ 

$$\frac{\left(-1\right)^{\langle \mathbf{z}, \mathbf{x} \rangle} + \left(-1\right)^{\langle \mathbf{z}, \mathbf{x} \rangle} \left(-1\right)^{\langle \mathbf{z}, \mathbf{c} \rangle}}{2} = \frac{\left(-1\right)^{\langle \mathbf{z}, \mathbf{x} \rangle} + \left(-1\right)^{\langle \mathbf{z}, \mathbf{x} \rangle} \left(-1\right)^{1}}{2} = \frac{\left(-1\right)^{\langle \mathbf{z}, \mathbf{x} \rangle} - \left(-1\right)^{\langle \mathbf{z}, \mathbf{x} \rangle}}{2} = 0$$

• When:  $\langle \mathbf{z}, \mathbf{c} \rangle = 0$ 

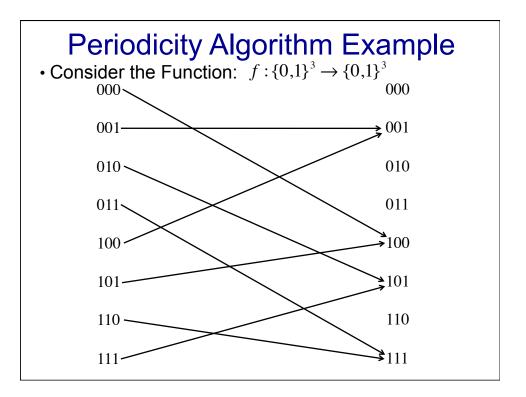
$$\frac{\left(-1\right)^{\langle z,x\rangle} + \left(-1\right)^{\langle z,x\rangle} \left(-1\right)^{\langle z,c\rangle}}{2} = \frac{\left(-1\right)^{\langle z,x\rangle} + \left(-1\right)^{\langle z,x\rangle} \left(-1\right)^{0}}{2} = \frac{\left(-1\right)^{\langle z,x\rangle} + \left(-1\right)^{\langle z,x\rangle}}{2} = 1$$

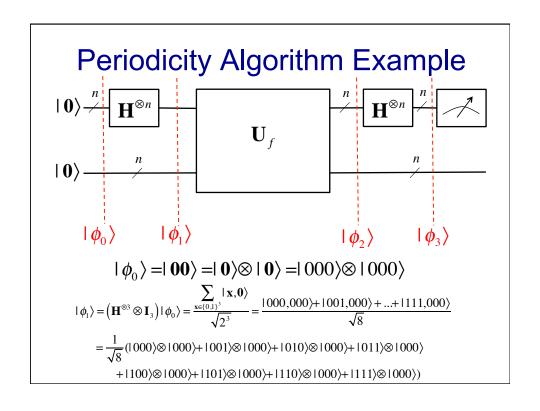
•Therefore, when measuring the top qubits, we only find those binary strings where:  $\langle \mathbf{z}, \mathbf{c} \rangle = 0$ 

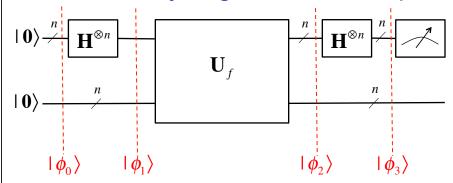
# **Example Function**

- Consider the Function:  $f: \{0,1\}^3 \rightarrow \{0,1\}^3$
- ulletTruth Table Representation Embedded Inside  $\mathbf{U}_f$

$x_1 x_2 x_3$	f
0 0 0	100
0 0 1	001
0 1 0	101
0 1 1	111
1 0 0	001
1 0 1	100
1 1 0	111
1 1 1	101







$$|\phi_2\rangle = \mathbf{U}_f \left( \mathbf{H}^{\otimes 3} \otimes \mathbf{I}_3 \right) |\phi_0\rangle = \frac{\sum_{\mathbf{x} \in \{0,1\}^3} |\mathbf{x}, f(\mathbf{x})\rangle}{\sqrt{2^3}} = \frac{\sum_{\mathbf{x} \in \{0,1\}^3} |\mathbf{x}\rangle \otimes |f(\mathbf{x})\rangle}{\sqrt{8}}$$

# Periodicity Algorithm Example

$$|\phi_2\rangle = \mathbf{U}_f \left( \mathbf{H}^{\otimes 3} \otimes \mathbf{I}_3 \right) |\phi_0\rangle = \frac{\sum_{\mathbf{x} \in \{0,1\}^3} |\mathbf{x}, f(\mathbf{x})\rangle}{\sqrt{2^3}} = \frac{\sum_{\mathbf{x} \in \{0,1\}^3} |\mathbf{x}\rangle \otimes |f(\mathbf{x})\rangle}{\sqrt{8}}$$

$$\begin{split} |\phi_2\rangle &= \frac{1}{\sqrt{8}}(|000\rangle\otimes|100\rangle + |001\rangle\otimes|001\rangle + |010\rangle\otimes|101\rangle + |011\rangle\otimes|111\rangle \\ &+ |100\rangle\otimes|001\rangle + |101\rangle\otimes|100\rangle + |110\rangle\otimes|111\rangle + |111\rangle\otimes|101\rangle) \end{split}$$

•In the Next Stage of the Cascade:

$$|\phi_{3}\rangle = \frac{\sum_{\mathbf{x} \in \{0,1\}^{3}} \sum_{\mathbf{z} \in \{0,1\}^{3}} (-1)^{\langle \mathbf{z}, \mathbf{x} \rangle} |\mathbf{z}\rangle \otimes |f(\mathbf{x})\rangle}{(\sqrt{8})(\sqrt{8})}$$

Writing this out Term by Term Yields:

```
\begin{split} |\phi_3\rangle &= \frac{1}{8}((+1)|000\rangle\otimes|f(000)\rangle + (+1)|000\rangle\otimes|f(001)\rangle + (+1)|000\rangle\otimes|f(010)\rangle + (+1)|000\rangle\otimes|f(011)\rangle \\ &+ (+1)|000\rangle\otimes|f(100)\rangle + (+1)|000\rangle\otimes|f(101)\rangle + (+1)|000\rangle\otimes|f(110)\rangle + (+1)|000\rangle\otimes|f(111)\rangle \\ &+ ((+1)|001\rangle\otimes|f(000)\rangle + (-1)|001\rangle\otimes|f(001)\rangle + (+1)|001\rangle\otimes|f(010)\rangle + (-1)|001\rangle\otimes|f(011)\rangle \\ &+ ((+1)|001\rangle\otimes|f(100)\rangle + (-1)|001\rangle\otimes|f(101)\rangle + (+1)|001\rangle\otimes|f(110)\rangle + (-1)|001\rangle\otimes|f(111)\rangle \\ &+ ((+1)|010\rangle\otimes|f(000)\rangle + (-1)|010\rangle\otimes|f(001)\rangle + (-1)|010\rangle\otimes|f(010)\rangle + (-1)|010\rangle\otimes|f(011)\rangle \\ &+ ((+1)|010\rangle\otimes|f(100)\rangle + (+1)|010\rangle\otimes|f(101)\rangle + (-1)|010\rangle\otimes|f(110)\rangle + (-1)|010\rangle\otimes|f(111)\rangle \\ &+ ((+1)|011\rangle\otimes|f(000)\rangle + (-1)|011\rangle\otimes|f(001)\rangle + (-1)|011\rangle\otimes|f(010)\rangle + (+1)|011\rangle\otimes|f(011)\rangle \\ &+ ((+1)|011\rangle\otimes|f(100)\rangle + (-1)|011\rangle\otimes|f(101)\rangle + (-1)|011\rangle\otimes|f(110)\rangle + (+1)|011\rangle\otimes|f(111)\rangle \end{split}
```

```
 + (+1)|100\rangle\otimes|f(000)\rangle + (+1)|100\rangle\otimes|f(001)\rangle + (+1)|100\rangle\otimes|f(010)\rangle + (+1)|100\rangle\otimes|f(011)\rangle \\ + (-1)|100\rangle\otimes|f(100)\rangle + (-1)|100\rangle\otimes|f(101)\rangle + (-1)|100\rangle\otimes|f(110)\rangle + (-1)|100\rangle\otimes|f(110)\rangle \\ + ((+1)|101\rangle\otimes|f(000)\rangle + (-1)|101\rangle\otimes|f(001)\rangle + (+1)|101\rangle\otimes|f(010)\rangle + (-1)|101\rangle\otimes|f(011)\rangle \\ + (-1)|101\rangle\otimes|f(100)\rangle + (+1)|101\rangle\otimes|f(101)\rangle + (-1)|101\rangle\otimes|f(110)\rangle + (+1)|101\rangle\otimes|f(111)\rangle \\ + ((+1)|110\rangle\otimes|f(000)\rangle + (+1)|110\rangle\otimes|f(001)\rangle + (-1)|110\rangle\otimes|f(100)\rangle + (-1)|110\rangle\otimes|f(011)\rangle \\ + (-1)|110\rangle\otimes|f(100)\rangle + (-1)|110\rangle\otimes|f(101)\rangle + (+1)|110\rangle\otimes|f(110)\rangle + (+1)|111\rangle\otimes|f(111)\rangle \\ + ((+1)|111)\otimes|f(000)\rangle + (-1)|111\rangle\otimes|f(001)\rangle + (-1)|111\rangle\otimes|f(010)\rangle + (+1)|111\rangle\otimes|f(011)\rangle \\ + ((-1)|111)\otimes|f(000)\rangle + (-1)|111\rangle\otimes|f(001)\rangle + (-1)|111\rangle\otimes|f(100)\rangle + (-1)|111\rangle\otimes|f(110)\rangle + (-1)|111\rangle\otimes|f(110)\rangle + (-1)|111\rangle\otimes|f(111)\rangle)
```

- Coefficients in Red are the Elements of:  $\mathbf{H}^{\otimes 3}$
- Evaluating the Function *f* in the Previous Equation Yields the Following:

 Evaluating the Function f in the Previous Equation Yields the Following:

```
\begin{split} |\phi_3\rangle &= \frac{1}{8}((+1)|000\rangle\otimes|100\rangle + (+1)|000\rangle\otimes|001\rangle + (+1)|000\rangle\otimes|101\rangle + (+1)|000\rangle\otimes|111\rangle \\ &+ (+1)|000\rangle\otimes|001\rangle + (+1)|000\rangle\otimes|100\rangle + (+1)|000\rangle\otimes|111\rangle + (+1)|000\rangle\otimes|101\rangle \\ &+ ((+1)|001\rangle\otimes|100\rangle + (-1)|001\rangle\otimes|1001\rangle + (+1)|001\rangle\otimes|101\rangle + (-1)|001\rangle\otimes|111\rangle \\ &+ (+1)|001\rangle\otimes|001\rangle + (-1)|001\rangle\otimes|100\rangle + (+1)|001\rangle\otimes|111\rangle + (-1)|001\rangle\otimes|111\rangle \\ &+ ((+1)|010\rangle\otimes|100\rangle + (+1)|010\rangle\otimes|100\rangle + (-1)|010\rangle\otimes|111\rangle + (-1)|010\rangle\otimes|111\rangle \\ &+ ((+1)|010\rangle\otimes|100\rangle + (+1)|010\rangle\otimes|100\rangle + (-1)|010\rangle\otimes|111\rangle + (-1)|010\rangle\otimes|101\rangle \\ &+ ((+1)|011\rangle\otimes|100\rangle + (-1)|011\rangle\otimes|100\rangle + (-1)|011\rangle\otimes|101\rangle + (+1)|011\rangle\otimes|111\rangle \\ &+ ((+1)|011\rangle\otimes|001\rangle + (-1)|011\rangle\otimes|100\rangle + (-1)|011\rangle\otimes|111\rangle + (+1)|011\rangle\otimes|101\rangle \end{split}
```

# Periodicity Algorithm Example

 Evaluating the Function f in the Previous Equation Yields the Following (continued):

```
 + (+1) | 100 \rangle \otimes | 100 \rangle + (+1) | 100 \rangle \otimes | 001 \rangle + (+1) | 100 \rangle \otimes | 101 \rangle + (+1) | 100 \rangle \otimes | 111 \rangle \\ + (-1) | 100 \rangle \otimes | 001 \rangle + (-1) | 100 \rangle \otimes | 100 \rangle + (-1) | 100 \rangle \otimes | 111 \rangle + (-1) | 100 \rangle \otimes | 101 \rangle \\ + ((+1) | 101 \rangle \otimes | 100 \rangle + (-1) | 101 \rangle \otimes | 100 \rangle + (+1) | 101 \rangle \otimes | 101 \rangle + (-1) | 101 \rangle \otimes | 111 \rangle \\ + (-1) | 101 \rangle \otimes | 100 \rangle + (+1) | 101 \rangle \otimes | 100 \rangle + (-1) | 101 \rangle \otimes | 111 \rangle + (+1) | 101 \rangle \otimes | 101 \rangle \\ + ((+1) | 110 \rangle \otimes | 100 \rangle + (+1) | 110 \rangle \otimes | 100 \rangle + (-1) | 110 \rangle \otimes | 101 \rangle + (-1) | 110 \rangle \otimes | 111 \rangle \\ + (-1) | 110 \rangle \otimes | 100 \rangle + (-1) | 110 \rangle \otimes | 100 \rangle + (+1) | 110 \rangle \otimes | 111 \rangle + (+1) | 111 \rangle \otimes | 111 \rangle \\ + ((+1) | 111 \rangle \otimes | 100 \rangle + (-1) | 111 \rangle \otimes | 100 \rangle + (-1) | 111 \rangle \otimes | 111 \rangle + (+1) | 111 \rangle \otimes | 111 \rangle \\ + (-1) | 111 \rangle \otimes | 100 \rangle + (+1) | 111 \rangle \otimes | 100 \rangle + (+1) | 111 \rangle \otimes | 111 \rangle + (-1) | 111 \rangle \otimes | 101 \rangle )
```

 Combining Like Terms and Cancelling Out Where Possible Yields the Following:

$$\begin{split} |\phi_{3}\rangle &= \frac{1}{8}((+2)|000\rangle\otimes|100\rangle + (+2)|000\rangle\otimes|001\rangle + (+2)|000\rangle\otimes|101\rangle + (+2)|000\rangle\otimes|111\rangle \\ &+ (+2)|010\rangle\otimes|100\rangle + (+2)|010\rangle\otimes|001\rangle + (-2)|010\rangle\otimes|101\rangle + (-2)|010\rangle\otimes|111\rangle \\ &+ (+2)|101\rangle\otimes|100\rangle + (-2)|101\rangle\otimes|001\rangle + (+2)|101\rangle\otimes|101\rangle + (-2)|101\rangle\otimes|111\rangle \\ &+ (+2)|111\rangle\otimes|100\rangle + (-2)|111\rangle\otimes|001\rangle + (-2)|111\rangle\otimes|101\rangle + (+2)|111\rangle\otimes|111\rangle \\ &+ (+2)|111\rangle\otimes|100\rangle + (-2)|111\rangle\otimes|001\rangle + (100)| + (101)| + (111)| + (1111)| \\ &+ (+2)|101\rangle\otimes(|100\rangle + |1001\rangle + |101\rangle + |111\rangle) \\ &+ (+2)|101\rangle\otimes(|100\rangle - |1001\rangle + |101\rangle - |111\rangle) \\ &+ (+2)|111\rangle\otimes(|100\rangle - |1001\rangle - |101\rangle + |111\rangle)) \end{split}$$

$$\begin{split} |\phi_{3}\rangle &= \frac{1}{8}((+2)|000\rangle \otimes (|100\rangle + |001\rangle + |101\rangle + |111\rangle) \\ &+ (+2)|010\rangle \otimes (|100\rangle + |001\rangle - |101\rangle - |111\rangle) \\ &+ (+2)|101\rangle \otimes (|100\rangle - |001\rangle + |101\rangle - |111\rangle) \\ &+ (+2)|111\rangle \otimes (|100\rangle - |001\rangle - |101\rangle + |111\rangle)) \end{split}$$

- Measuring the Top 3 Qubits Gives (with equal probability): 1000>,1010>,1101>,1111>
- For Each of These Measured Quantum States, it is True that the Inner Product with the Period Bitstring c is Zero

- For Each of These Measured Quantum States, it is True that the Inner Product with the Period Bitstring  $\varepsilon$  is Zero
- The Algorithm/Circuit is Measured a Sufficient Number of Times to Ensure All Possible Measurements are Obtained
- This Yields a Set of Simultaneous Equations:
  - (i)  $\langle 000, \mathbf{c} \rangle = 0$
  - (ii)  $\langle 010, \mathbf{c} \rangle = 0$
  - (*iii*)  $\langle 101, \mathbf{c} \rangle = 0$
  - $(iv) \langle 111, \mathbf{c} \rangle = 0$

$$\langle 000, c_1 c_2 c_3 \rangle = 0$$
$$\langle 010, c_1 c_2 c_3 \rangle = 0$$

$$\langle 101, c_1 c_2 c_3 \rangle = 0$$

$$\langle 111, c_1 c_2 c_3 \rangle = 0$$

$$(0 \wedge c_1) \oplus (0 \wedge c_2) \oplus (0 \wedge c_3) = 0$$

$$(0 \wedge c_1) \oplus (1 \wedge c_2) \oplus (0 \wedge c_3) = 0$$

$$(1 \land c_1) \oplus (0 \land c_2) \oplus (1 \land c_3) = 0$$

$$(1 \land c_1) \oplus (1 \land c_2) \oplus (1 \land c_3) = 0$$

$$\begin{aligned} c_2 &= 0 \\ c_1 \oplus c_3 &= 0 \\ c_1 \oplus c_2 \oplus c_3 &= 0 \end{aligned}$$

- This Means that  $c_1$ = $c_3$ =0 or that  $c_1$ = $c_3$ =1
- ullet We Know that  ${f c}$  is NOT EQUAL to 000 Since Function was Found not to be One-to-One
- Therefore  $c_1 = c_3 = 1$
- Period of Function is:

$$\mathbf{c} = c_1 c_2 c_3 = 101$$

- Must Run Simon's Algorithm Several Times to Measure n Different z Bitstrings
- Next Use a Classical Computer for Solving n
   Different Linear Equations