

Nonlinear Derivative-free Constrained Optimization with a Mixed Penalty-Logarithmic Barrier Approach and Direct Search

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Outline

- ① Introduction
- ② Algorithmic Structure and Convergence Analysis
- ③ Implementation Details
- ④ Numerical Experiments
- ⑤ Conclusions and Future Work

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Problem Features

$$\begin{array}{ll}\min & f(\mathbf{x}) \\ \text{s.t.} & g(\mathbf{x}) \leq 0 \\ & h(\mathbf{x}) = 0 \\ & \mathbf{x} \in X\end{array}$$

$$\begin{array}{l}f : X \subseteq \mathbb{R}^n \rightarrow \mathbb{R} \cup \{+\infty\} \\ g : X \subseteq \mathbb{R}^n \rightarrow \{\mathbb{R} \cup \{+\infty\}\}^m \\ h : X \subseteq \mathbb{R}^n \rightarrow \{\mathbb{R} \cup \{+\infty\}\}^p \\ X = \{\mathbf{x} \in \mathbb{R}^n \mid \mathbf{Ax} \leq \mathbf{b}\}\end{array}$$

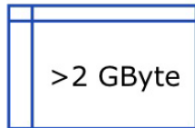
- f , g , and h are **black-box** type and continuously differentiable
- Computing f , g , and h is expensive



No derivatives
available for use



Long runtime



Large memory
requirement

* Image credits to Joseph Simonis, ISMP 2009, Chicago, US

SID-PSM: Direct Search using Simplex Derivatives

- Search step based on the minimization of some quadratic polynomial model (interpolation, minimum $\|\cdot\|_F$ or regression)



A. L. Custódio, H. Rocha, and L. N. Vicente

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Comput. Optim. Appl., 46: 265-278, 2010.

- Order of the poll vectors according to a negative simplex gradient



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- Handles constraints using an extreme barrier approach
- Efficient approaches to address general constraints (not yet!...)

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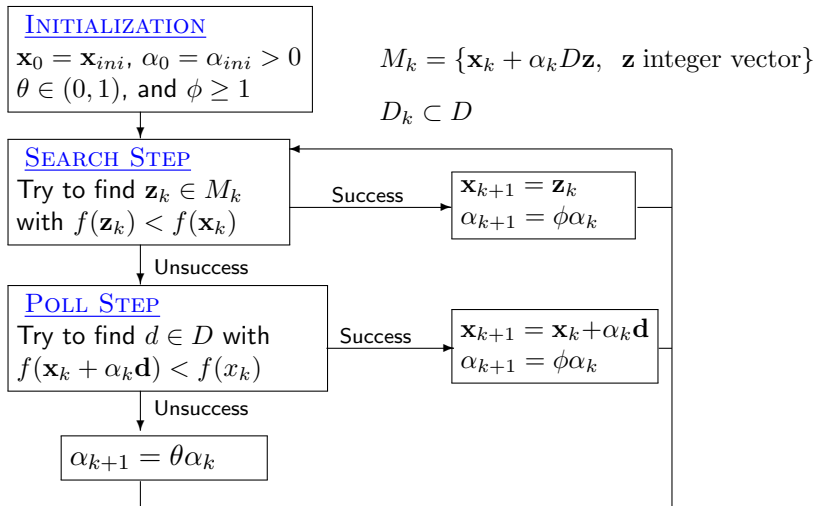


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SID-PSM – Algorithmic Structure



LOG-DFL: Logarithmic barrier penalty DFO algorithm

- New DFO method that uses a merit function with a **log-barrier for inequality** constraints and a **penalty approach for equality** constraints.



A. Brilli, G. Liuzzi, and S. Lucidi

An interior point method for nonlinear constrained derivative-free optimization

arXiv, 2108.05157 [math.OC], 2021.

Mixed Penalty-Log Barrier Approach in Direct Search

Given an initial point $\mathbf{x}_0 \in X$, define:

$$\begin{aligned}\mathcal{G}^{\log} &= \{\ell \mid g_\ell(\mathbf{x}_0) < 0\} \\ \mathcal{G}^{\text{ext}} &= \{\ell \mid g_\ell(\mathbf{x}_0) \geq 0\}\end{aligned}$$

Feasible region

$$\mathcal{F} = X \cap \Omega_{\mathcal{G}^{\log}} \cap \Omega_{\mathcal{G}^{\text{ext}}} \cap \Omega_h \neq \emptyset, \text{ a compact set.}$$

Mixed Penalty-Logarithmic Barrier function

$$Z(\mathbf{x}; \rho) = f(\mathbf{x}) - \rho \sum_{\ell \in \mathcal{G}^{\log}} \log(-g_\ell(\mathbf{x})) + \frac{1}{\rho^{\nu-1}} \left(\sum_{\ell \in \mathcal{G}^{\text{ext}}} (\max\{g_\ell(\mathbf{x}), 0\})^\nu + \sum_{j=1}^p |h_j(\mathbf{x})|^\nu \right),$$

where $\rho > 0$ and $\nu \in (1, 2]$.

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Penalized problem

$$\begin{aligned} \min \quad & Z(\mathbf{x}; \rho) \\ \text{s.t.} \quad & \mathbf{x} \in X \cap \overset{\circ}{\Omega}_{\mathcal{G}^{\log}} \end{aligned}$$

Equivalent Reformulation

$$\begin{aligned} \min \quad & Z(\mathbf{x}; \rho) \\ \text{s.t.} \quad & \mathbf{x} \in X \end{aligned}$$

- ρ must be driven to zero to solve the original problem

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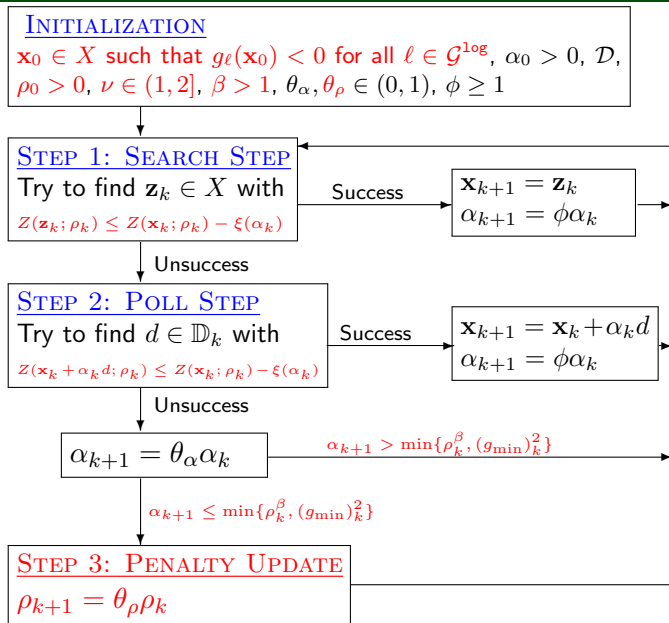
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LOG-SID-PSM – Algorithmic Structure



$$\mathcal{D} = \bigcup_{k=0}^{+\infty} \mathbb{D}_k$$

$$(g_{\min})_k = \min_{\ell \in \mathcal{G}^{\log}} \{|g_\ell(\mathbf{x}_k)|\}$$

Penalty Parameter Update

Penalty Parameter Update.

Step 3. Set $(g_{\min})_k = \min_{\ell \in \mathcal{G}^{\log}} \{|g_{\ell}(\mathbf{x}_{k+1})|\}$

If $\alpha_{k+1} < \alpha_k$ and $\alpha_{k+1} \leq \min\{\rho_k^{\beta}, (g_{\min})_k^2\}$

Then set $\rho_{k+1} = \theta_{\rho} \rho_k$

Else set $\rho_{k+1} = \rho_k$

The **measure of stationarity** α_k should go to zero faster than:

- the **measure of quality** ρ_k
- the **measure of proximity** $(g_{\min})_k$

Theorem

Let $\{\rho_k\}_{k \in \mathbb{N}}$ and $\{\alpha_k\}_{k \in \mathbb{N}}$ be sequences generated by LOG-SID-PSM. Then,

$$\lim_{k \rightarrow +\infty} \rho_k = 0 \text{ and } \lim_{k \rightarrow +\infty} \alpha_k = 0.$$

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Forcing Function $\xi : [0, +\infty) \rightarrow [0, +\infty)$

- a continuous and nondecreasing function
- $\xi(t)/t \rightarrow 0$ when $t \rightarrow 0$
- If $\xi(t) \rightarrow 0$ then $t \rightarrow 0$

Mangasarian-Fromovitz Type Constraint Qualification

MFCQ

Let $x \in X$ and $T_X(\mathbf{x})$ be the tangent cone at \mathbf{x} with respect to the linear constraints. The point \mathbf{x} is said to satisfy the MFCQ if the two following conditions are satisfied:

(a) There does not exist a nonzero vector $\alpha = (\alpha_1, \dots, \alpha_q)$ such that:

$$\left(\sum_{i=1}^q \alpha_i \nabla h_i(\mathbf{x}) \right)^\top \mathbf{d} \geq 0, \quad \forall \mathbf{d} \in T_X(\mathbf{x});$$

(b) There exists a feasible direction $\mathbf{d} \in T_X(\mathbf{x})$, such that:

$$\nabla g_\ell(\mathbf{x})^\top \mathbf{d} < 0, \quad \forall \ell \in I_+(\mathbf{x}), \quad \nabla h_j(\mathbf{x})^\top \mathbf{d} = 0, \quad \forall j = 1, \dots, p$$

where $I_+(\mathbf{x}) = \{\ell \mid g_\ell(\mathbf{x}) \geq 0\}$.

Active Constraints and Tangent Cone

For every $\mathbf{x} \in X$, i.e., such that $\mathbf{Ax} \leq \mathbf{b}$:

$$I_X(\mathbf{x}) = \{i \mid \mathbf{a}_i^\top \mathbf{x} = b_i\} \quad (\text{set of indices of active constraints})$$

$$T_X(\mathbf{x}) = \{\mathbf{d} \in \mathbb{R}^n \mid \mathbf{a}_i^\top \mathbf{d} \leq 0, i \in I_X(\mathbf{x})\} \quad (\text{tangent cone at } \mathbf{x})$$

ε -Active Constraints and Tangent Cone

For every $\mathbf{x} \in X$, i.e., such that $\mathbf{Ax} \leq \mathbf{b}$:

$$I_X(\mathbf{x}_k, \varepsilon) = \{i \mid \mathbf{a}_i^\top \mathbf{x}_k \geq b_i - \varepsilon\} \quad (\text{set of indices of } \varepsilon\text{-active constraints})$$

$$T_X(\mathbf{x}_k, \varepsilon) = \{\mathbf{d} \in \mathbb{R}^n \mid \mathbf{a}_i^\top \mathbf{d} \leq 0, \ i \in I_X(\mathbf{x}_k, \varepsilon)\} \quad (\varepsilon\text{-tangent cone at } \mathbf{x}_k)$$

Proposition

Let $\{\mathbf{x}_k\}_{k \in \mathbb{N}}$ be a sequence of points in X converging to $\mathbf{x}^* \in X$. Then, there exists an $\varepsilon^* > 0$ (depending only on \mathbf{x}^*) such that for any $\varepsilon \in (0, \varepsilon^*]$ there exists $k_\varepsilon \in \mathbb{N}$ such that

$$I_X(\mathbf{x}^*) = I_X(\mathbf{x}_k, \varepsilon)$$

$$T_X(\mathbf{x}^*) = T_X(\mathbf{x}_k, \varepsilon)$$

for all $k \geq k_\varepsilon$.

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Geometry Assumption

Let $\{\mathbf{x}_k\}_{k \in \mathbb{N}}$ be a sequence of points such that $\mathbf{x}_k \in X$. The sequence \mathbb{D}_k of poll directions satisfies:

$$\mathbb{D}_k = \{\mathbf{d}_k^i \mid \|\mathbf{d}_k^i\| = 1, i = 1, \dots, |\mathbb{D}_k|\}$$

and for some $\bar{\varepsilon} > 0$,

$$\text{cone}(\mathbb{D}_k \cap T_X(\mathbf{x}_k, \varepsilon)) = T_X(\mathbf{x}_k, \varepsilon), \quad \forall \varepsilon \in (0, \bar{\varepsilon}].$$

Furthermore, $\mathcal{D} = \bigcup_{k=0}^{+\infty} \mathbb{D}_k$ is a finite set, and $|\mathbb{D}_k|$ is bounded.

Convergence of LOG-SID-PSM

Lagrange Multipliers

$$\begin{aligned}\nabla Z(\mathbf{x}; \rho_k) = & \nabla f(\mathbf{x}) + \sum_{\ell \in \mathcal{G}^{\log}} \frac{\rho_k}{-g_\ell(\mathbf{x})} \nabla g_\ell(\mathbf{x}) + \sum_{\ell \in \mathcal{G}^{\text{ext}}} \nu \left(\frac{\max\{g_\ell(\mathbf{x}), 0\}}{\rho_k} \right)^{\nu-1} \nabla g_\ell(\mathbf{x}) + \\ & + \sum_{j=1}^p \nu \left(\frac{|h_j(\mathbf{x})|}{\rho_k} \right)^{\nu-1} \nabla h_j(\mathbf{x})\end{aligned}$$

$$\lambda_\ell(\mathbf{x}; \rho) = \begin{cases} \frac{\rho}{-g_\ell(\mathbf{x})}, & \text{if } \ell \in \mathcal{G}^{\log} \\ \nu \left(\frac{\max\{g_\ell(\mathbf{x}), 0\}}{\rho} \right)^{\nu-1}, & \text{if } \ell \in \mathcal{G}^{\text{ext}} \end{cases}$$

$$\mu_j(\mathbf{x}; \rho) = \nu \left(\frac{|h_j(\mathbf{x})|}{\rho} \right)^{\nu-1}, \quad j = 1, \dots, p$$

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Theorem

Let $\{\mathbf{x}_k\}_{k \in \mathbb{N}}$ be the sequence of iterates generated by LOG-SID-PSM. Consider the set $K = \{k \in \mathbb{N} : \rho_{k+1} < \rho_k\}$, assume that the sets of directions $\{\mathbb{D}_k\}_{k \in \mathbb{N}}$, used by the algorithm, satisfy the Geometry Assumption and let x^* be a limit point of $\{\mathbf{x}_k\}_{k \in \hat{K}}$, $\hat{K} \subseteq K$, that satisfies the MFCQ. Then

- (i) The sequences of **Lagrange multipliers** $\{\lambda_\ell(\mathbf{x}_k; \rho_k)\}_{k \in \hat{K}}$, $\ell = 1, \dots, m$ and $\{\mu_j(\mathbf{x}_k; \rho_k)\}_{k \in \hat{K}}$, $j = 1, \dots, p$ are **bounded**.
- (ii) x^* is a **stationary point**

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- Attempts a **search step** based on the **minimization of quadratic models**
- **Orders** the **poll vectors** according to a negative **simplex gradient**
- Handles constraints using a mixed **Penalty-Logarithmic Barrier**

Sufficient decrease condition

$$Z(\mathbf{x}_{k+1}; \rho_k) \leq Z(\mathbf{x}_k; \rho_k) - \gamma \alpha_k^2,$$

where $\gamma = 10^{-9}$

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Search Step in LOG-SID-PSM – Model Building

- Reuses previous function evaluations and select points in $B(\mathbf{x}_k; \Delta_k)$, with

$$\Delta_k = \sigma \alpha_k \max_{\mathbf{d} \in D_{k-1}} \|\mathbf{d}\|$$

- Builds quadratic models for each function (f , g , and h) using the selected points:

points in $[n + 2, (n + 1)(n + 2)/2[\Rightarrow$ MFN model

points = $(n + 1)(n + 2)/2 \Rightarrow$ Determined interpolation model

points in $] (n + 1)(n + 2)/2, (n + 1)(n + 2)] \Rightarrow$ Regression model

- 80% of the points selected nearest to the current iterate
- 20% of the points selected farthest away from the current iterate

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$$\Delta_k = \sigma \alpha_k \max_{\mathbf{d} \in D_{k-1}} \|\mathbf{d}\|$$

- Builds quadratic models for each function (f , g , and h) using the selected points:

points in $[n + 2, (n + 1)(n + 2)/2[\Rightarrow$ MFN model

points = $(n + 1)(n + 2)/2 \Rightarrow$ Determined interpolation model

points in $] (n + 1)(n + 2)/2, (n + 1)(n + 2)] \Rightarrow$ Regression model

- 80% of the points selected nearest to the current iterate
- 20% of the points selected farthest away from the current iterate

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Search Step in LOG-SID-PSM

- Computes \mathbf{z}_k as the solution to the following problem:

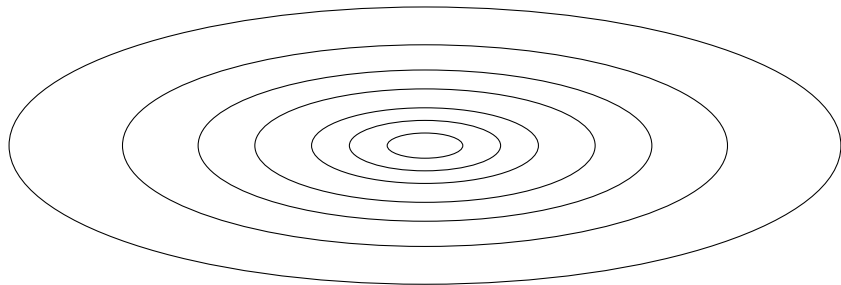
$$\begin{aligned} \min \quad & Z_m(\mathbf{z}; \rho_k) \\ \text{s.t. } \quad & \mathbf{z} \in X \cap B(x_k; \Delta_k) \end{aligned}$$

where

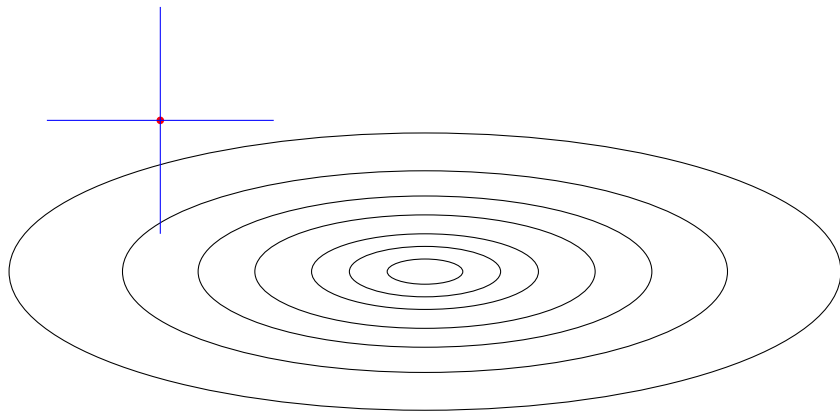
$$\begin{aligned} Z_m(\mathbf{z}; \rho_k) = & \textcolor{red}{f}^m(\mathbf{z}) - \textcolor{blue}{\rho}_k^{\log} \sum_{\ell \in \mathcal{G}^{\log}} \log(-\textcolor{red}{g}_\ell^m(\mathbf{z})) + \\ & + \frac{1}{\textcolor{blue}{\rho}_k^{\text{ext}}} \left(\sum_{\ell \in \mathcal{G}^{\text{ext}}} (\max\{\textcolor{red}{g}_\ell^m(\mathbf{z}), 0\})^\nu + \sum_{j=1}^p |\textcolor{red}{h}_j^m(\mathbf{z})|^\nu \right), \end{aligned}$$

with $\nu \in (1, 2]$.

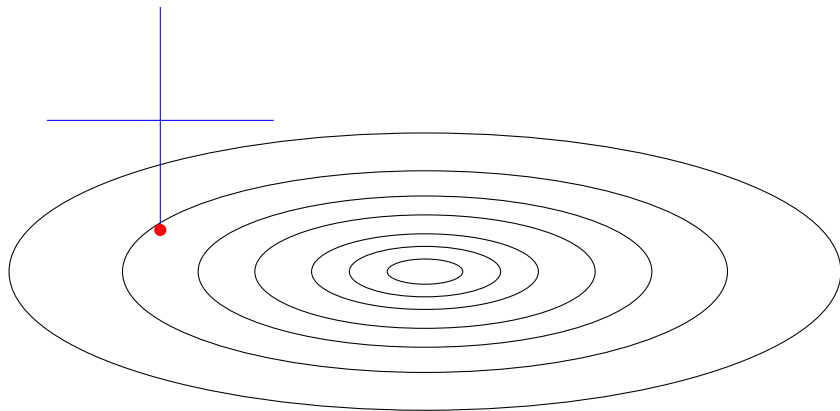
Poll Step in LOG-SID-PSM – Ordering



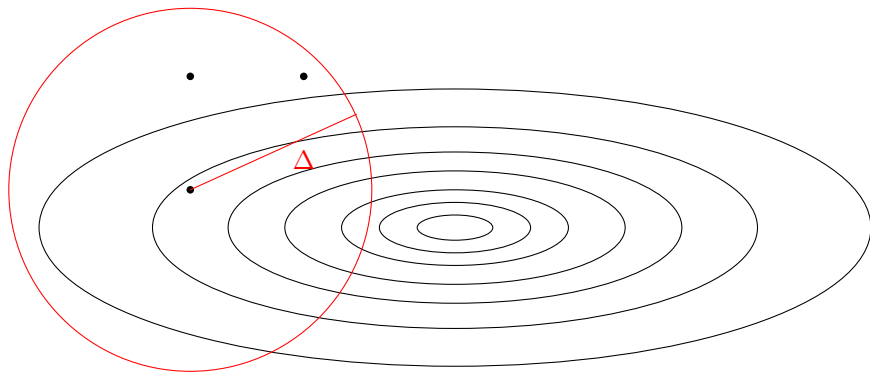
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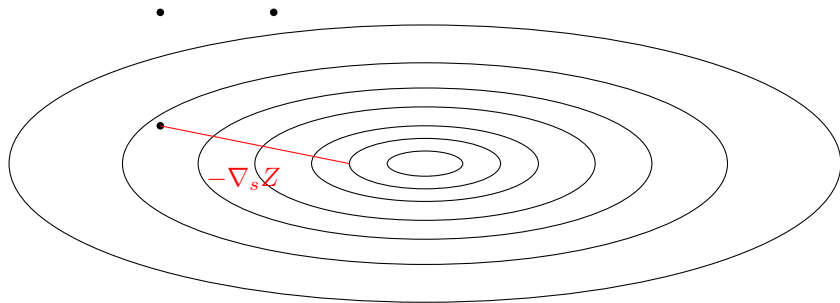
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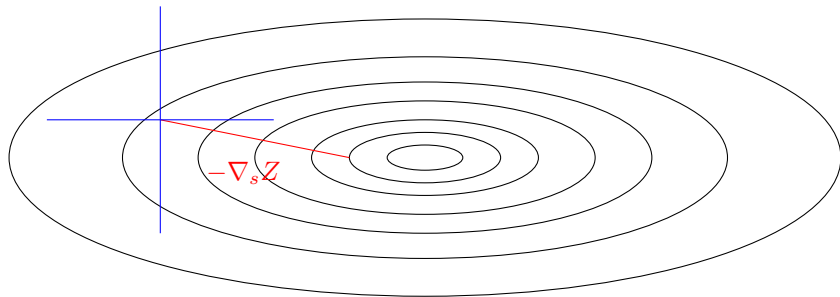
Poll Step in LOG-SID-PSM – Ordering



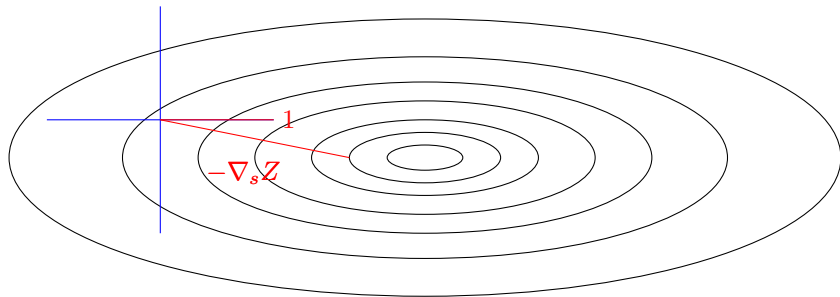
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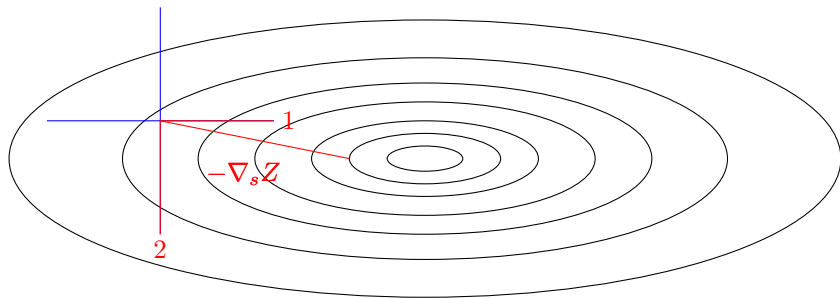
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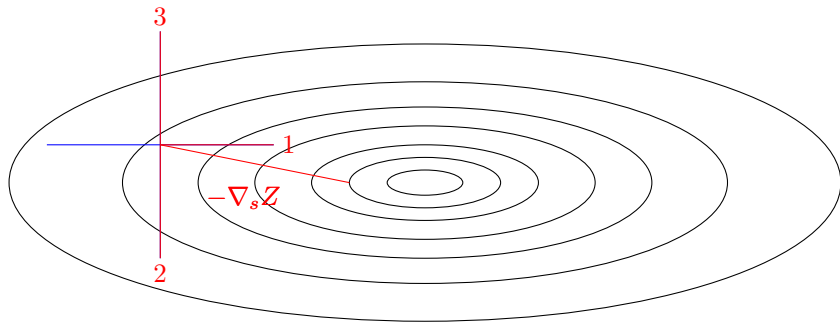
Poll Step in LOG-SID-PSM – Ordering



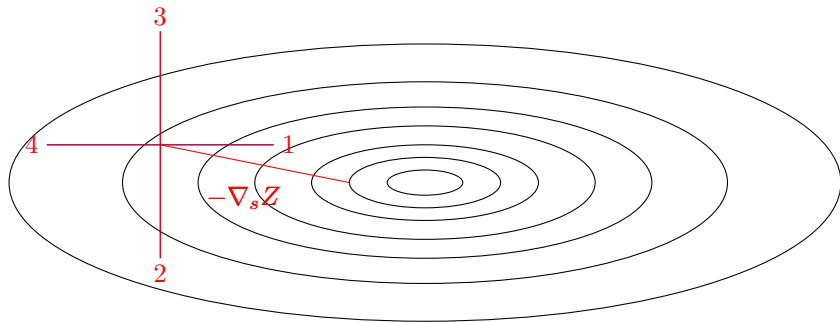
Poll Step in LOG-SID-PSM – Ordering



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LOG-SID-PSM Penalty Parameters Rule

Penalty Function

$$Z(\mathbf{x}; \rho_k) = f(\mathbf{x}) - \rho_k^{\text{log}} \sum_{\ell \in \mathcal{G}^{\text{log}}} \log(-g_\ell(\mathbf{x})) + \frac{1}{\rho_k^{\text{ext}}} \left(\sum_{\ell \in \mathcal{G}^{\text{ext}}} (\max\{0, g_\ell(\mathbf{x})\})^\nu + \sum_{j=1}^p |h_j(\mathbf{x})|^\nu \right)$$

- The **updating rule** splits into

Updating ρ_k^{log}

$$\alpha_{k+1} \leq \min\{(\rho_k^{\text{log}})^\beta, (g_{\min})_k^2\}$$



$$\rho_{k+1}^{\text{log}} = \rho_k^{\text{log}} \min\{\eta, \max\{(g_{\min})_k^2, \zeta\}\}$$

Updating ρ_k^{ext}

$$\alpha_{k+1} \leq (\rho_k^{\text{ext}})^\beta$$



$$\rho_{k+1}^{\text{ext}} = \min\left\{\zeta \rho_k^{\text{ext}}, \frac{\sqrt{\alpha_{k+1}}}{10}\right\}.$$

where $\beta = 1 + 10^{-9}$, $\nu = 1.1$, and $\zeta = 10^{-2}$.

Outline

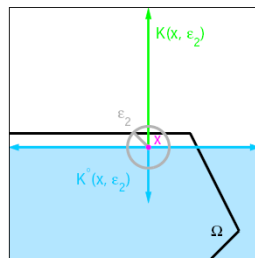
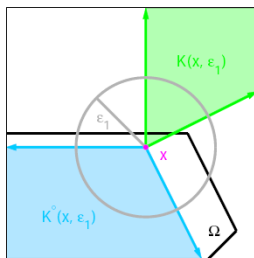
- ① Introduction
- ② Algorithmic Structure and Convergence Analysis
- ③ Implementation Details
- ④ Numerical Experiments
- ⑤ Conclusions and Future Work

- 96 problems with nonlinear constraints and bounds from the CUTEst collection
 - number of variables between 1 and 50
 - number of nonlinear inequality constraints between 1 and 144
 - number of nonlinear equality constraints between 0 and 30
 - number of linear inequality constraints (other than bounds) between 0 and 123
- Initialization: provided in CUTEst collection
- Stopping criterion:
 - $\alpha_k < 10^{-8}$
 - Maximum of 2000 function evaluations

General Linear Constraints

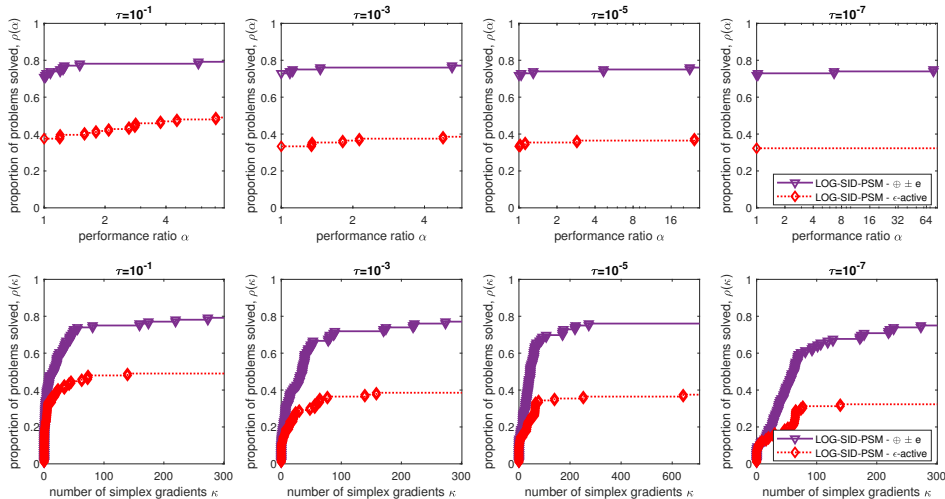
Set of poll directions **conforms to the geometry** of nearby constraints.

Implementation based in Abramson, Brezhneva, Dennis, and Pingel [2008].



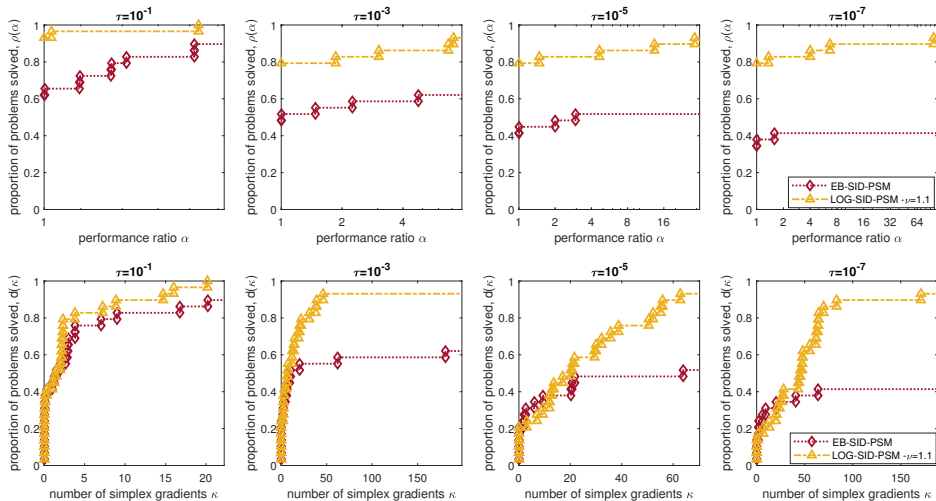
(in Kolda, Lewis, and Torczon [2003])

Comparison between strategies for linear constraints



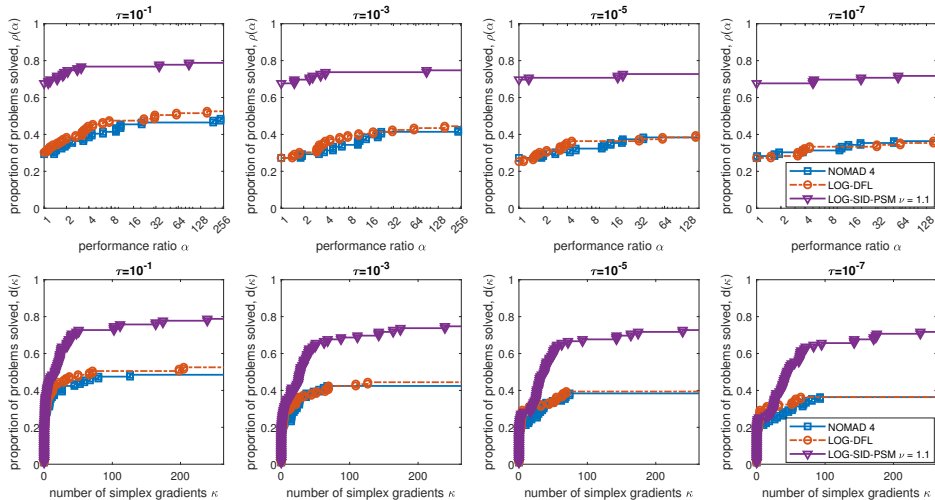
$$*\oplus \pm e \equiv [e - eI - I]$$

LOG-SID-PSM vs SID-PSM

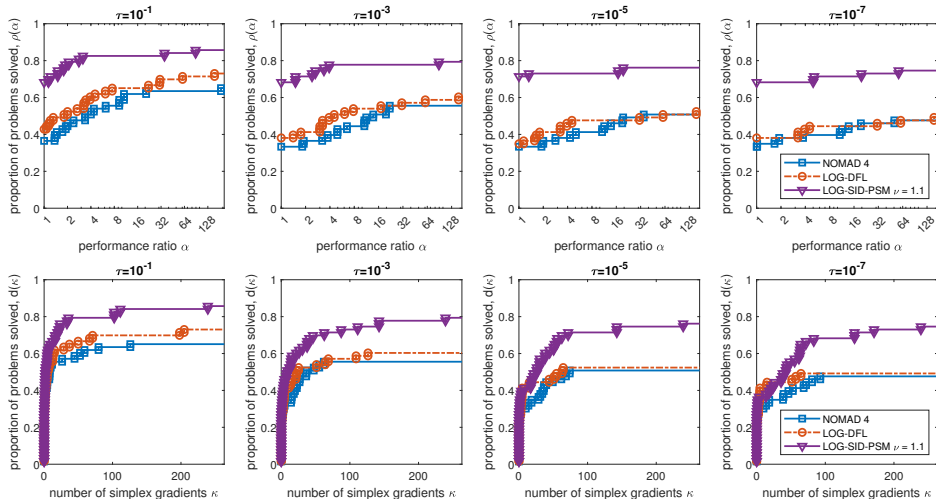


*28 problems from CUTEst

Comparison with other solvers (all problems)



Comparison with other solvers (only inequality constraints)



Outline

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Future Work

- Extension of the theoretical analysis for nonsmooth functions (*A. Brilli, A.L. Custódio, G. Liuzzi, and E.J. Silva*) ongoing work
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Thank you for your attention!

Technical report will appear soon!