# RW771 Tutorial 1

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# 1 Problem Description

Write a program that can parse and simulate any DFA that is in the correct EBNF<sup>1</sup>, the EBNF is show in figure one. The parser needs to be able to do the following: detect incorrect format, detect inconsistencies and extract and DFA 5-tuple<sup>2</sup>. The inconsistencies that should be detected are the following: The transition function specifies a mapping between two states and one of those state do not exist, when a symbol is used to signal a transition between two states that is not part of the specified alphabet. The program needs to take in two argument. The first argument specifies the file where the DFA is located and the second argument specifies the input to that DFA. When a DFA executes the input and it finishes in an accept state the word "accept" should be printed else the word "reject" should be printed and followed by the word "reject".

# 2 Design and Implementation

The program design is divide up into 3 section, a scanner, a parser and simulater.

#### 2.1 Scanner

The task of this section is to build up tokens that is reconized by the specified EBNF. This will be achived in the follow way.

- Read a character.
- Is the character a letter. If so build up a string of characters till a white space is encountered and start again. If not go to the next step.
- Is the character a number. If so Create a symbol token that has the value of the number and start again.
- Is the character one of the following symbols,  $\{\ \} = (\ )$  and ,
- In the case where a symbol is read. Read and check if the next symbol is a > sign if so create the appropriate token and start again else throw an error.

 $<sup>^{1}</sup>$ Extended Backus-Naur Form

<sup>&</sup>lt;sup>2</sup>(states, alphabet, transition function, starting state, accepting states)

- If the EOF character is read stop. And test if the current symbol is valid. If so then create the appropriate token and stop else if not error.
- If the character is not identified stop processing the file and throw an error.

### 2.2 Parser

The task of this section is to analise the syntax of the specified DFA to make sure it is in the correct format and extract the DFA 5-tuple. This is achived by implementing a recursive desent parser in the following manner. Variables<sup>3</sup> become functions. The left had side of the variables are then checked symbol by symbol. If the left hand side contains a variable then the appropriate function must be called. The 5-tuple that a DFA consists of is describe in the first production<sup>4</sup> of the EBNF thus as each part of the top production is finished. Its result can be stored into the corresponding variable.

#### 2.3 Simulater

The task of this section is to simulate a DFA. This is achive by using the following algorithm.

- Step 1 set currentState = startingState
- Step 2 Read a character from the input and put it equal to currentCharacter
- Step 3 Concatenate the currentState and the currentCharacter to form a key.
- Step 4 set currentState = hashmap(key). If there are no more characters go to step 5 else got to step 2.
- Step 5 Check if the currentState is indeed a accepting state if so display "accept" else "reject".

### 3 Test Cases

#### 3.1 Test Case 1

Test: This Test case tests if the starting state is not the final state and there is no input string.

DFA:automaton = 
$$(\{q1,q2\}, \{0, 1\}, \{(q1,0) -> q1, (q1,1) -> q2\}, q1, \{q2\})$$
  
Input String:  
Result:reject

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<sup>&</sup>lt;sup>3</sup>The words to the left of the equals sign that doesn't equate to only one token.

<sup>&</sup>lt;sup>4</sup>A production looks like a mathematical equation

### 3.2 Test Case 2

Test: This Test case tests if the starting state is the final state and there is no input string.

 $DFA: automaton = (\{q1,q2\}, \{0,1\}, \{(q1,0) -> q1, (q1,1) -> q2\}, q2, \{q2\})$ 

Input String: Result:accept

## 3.3 Test Case 3

- 3.4 Test Case 4
- 3.5 Test Case 5
- 3.6 Test Case 6