

# Hilbert's 10<sup>th</sup> Problem, Undecidability, & The Halting Problem

A mini-lecture series

CSE498 Collaborative Design (W) - Secure and Efficient C++ Software Development

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# David Hilbert

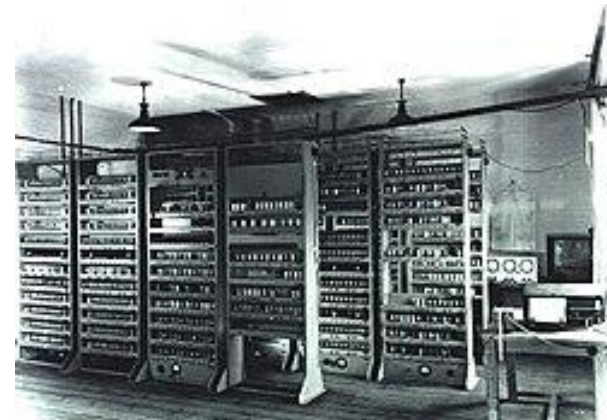
- Brilliant mathematician in the early 1900s
- During his speech to the International Congress of Mathematicians in 1900, he presented 23 unsolved problems
- His 10<sup>th</sup> problem:
  - For any given Diophantine equation, can we *decide* whether the equation has a solution?
- This problem remained unsolved for 70 years, until 1970



# Why are we even talking about this?

# Very early programming languages

- 1843: Ada Lovelace writes the first algorithm
  - Note G – used to calculate Bernoulli numbers on the analytical machine
  - Lovelace translated seminar and added her own theories
  - Lovelace is often considered the first programmer
- 1936: Alan Turing propose the concept of a universal machine
- 1950-ish: Assembly languages, EDSAC machine completed



# Decidability

- In computer theory, an undecidable problem is one which it is impossible to construct an algorithm that can always lead to a yes or no answer
- Hilbert called it ***Entscheidungsproblem*** (decision problem)
- Example: Halting problem

# Implications

- No perfect program analyzer can exist
- Can we generate a universal machine that can identify any UB in a C++ program?
- Can we automatically test whether a new update is backwards compatible with all previous versions and hardware
- Can we build an automated tool that can detect all bugs
- Can we build a program to check if another program will crash
- In other words: you can tell your boss that the task they assigned you is mathematically impossible<sup>\*</sup>
  - <sup>\*</sup>do it at your own risk

# Halting Problem

- Can you write a computer program  $P$  that:
  - For a given input, a computer program  $i$ , determine whether  $i$  will:
    - Halt (exit execution)
    - Run indefinitely
- To prove this, you must show that
  - Decide on ALL existing programs, even ones that have not existed
  - OR use prove by contradiction (show an instance of the problem that cannot logically exist)

# Alan Turing's Paper in 1936

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A. M. TURING

[Nov. 12,

## ON COMPUTABLE NUMBERS, WITH AN APPLICATION TO THE ENTSCHEIDUNGSPROBLEM

*By* A. M. TURING.

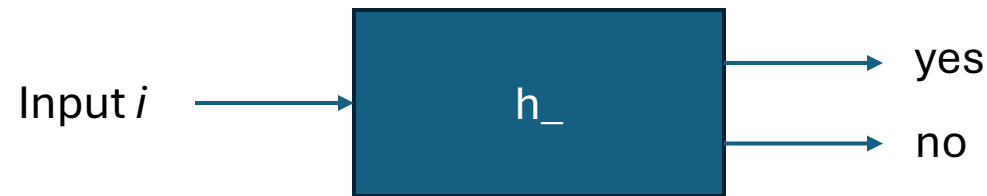
[Received 28 May, 1936.—Read 12 November, 1936.]

The “computable” numbers may be described briefly as the real numbers whose expressions as a decimal are calculable by finite means. Although the subject of this paper is ostensibly the computable *numbers*, it is almost equally easy to define and investigate computable functions of an integral variable or a real or computable variable, computable predicates, and so forth. The fundamental problems involved are, however, the same in each case, and I have chosen the computable numbers for explicit treatment as involving the least cumbrous technique. I hope shortly to give an account of the relations of the computable numbers, functions, and so forth to one another. This will include a development of the theory of functions of a real variable expressed in terms of computable numbers. According to my definition, a number is computable if its decimal can be written down by a machine.



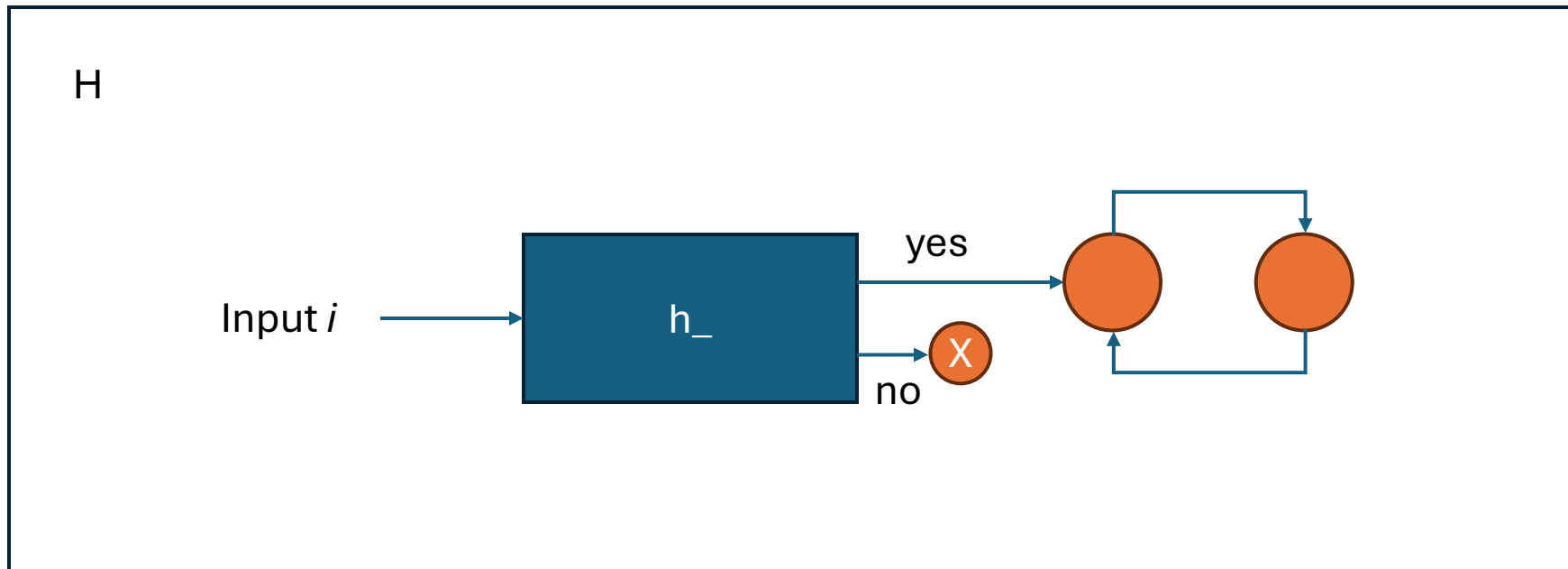
# Halting Problem

- Suppose there exist a program  $h_*$  that can magically decide whether a program halts or not



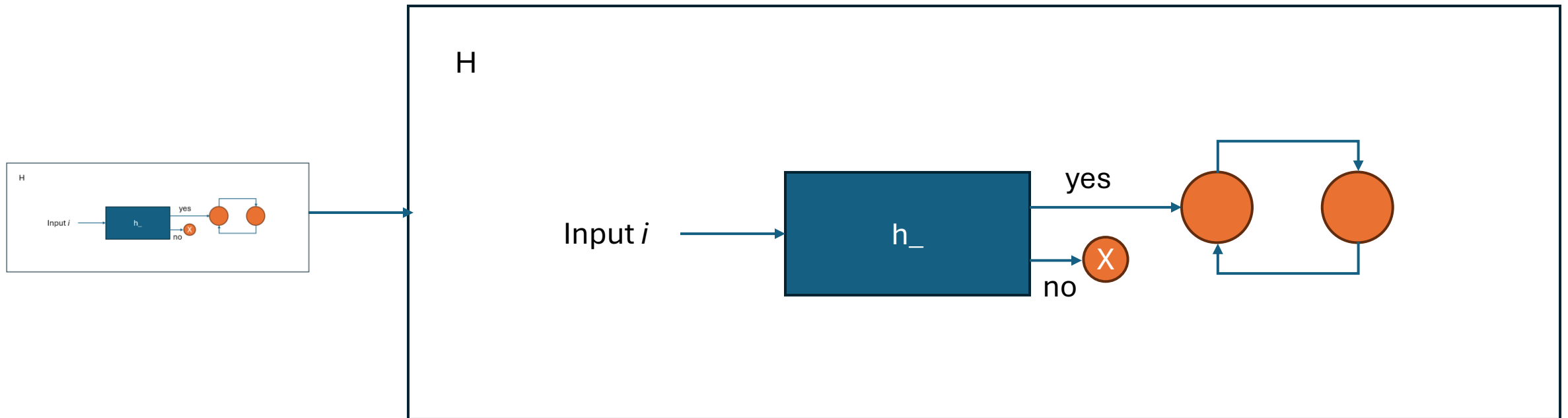
# Halting Problem

- Let's construct a new program  $H$ , based on  $h_*$  as follows
  - If  $h_*$  says that the problem halts, then run forever
  - Else, halt



# Contradiction

- What happens when we feed the program an instance of itself?
  - If  $H$  halts, then  $h_+$  returns yes; but it runs forever
  - If  $H$  does not halt, then  $h_+$  returns no; but it does halt



# Halting Problem - Results

- The halting problem showed an instance of a problem that *cannot be solved by computers*
- This problem showed that some functions are mathematically definable, but not computable
- Like most other proofs in computer science, this extends beyond just the Halting problem
  - Mapped and demonstrated to be the same problem as computability, Hilbert's 10<sup>th</sup> problem, the other problems we discussed earlier

# Implications

- Actually equivalent (can be reduced) to Hilbert's 10<sup>th</sup> problem
  - Proved by Matiyasevich
- Wednesday, we will talk about testing in details
  - Exhaustive testing is mathematically infeasible
- Microsoft rolls out a new update, and they want to make sure that it does not break backwards compatibility at all. This is actually just the halting problem

# XKCD - 1425



IN CS, IT CAN BE HARD TO EXPLAIN  
THE DIFFERENCE BETWEEN THE EASY  
AND THE VIRTUALLY IMPOSSIBLE.

# Summary

- One of the most important results in our field
- Halting problem showed that there exist problems in which computers cannot compute
- Some what defines the limits of what programming can do

# Person of the day

## David Hilbert

- Brilliant mathematician in the early 1900s
- Visionary for “math” and how to the next sets of solve problems (agenda)
- *Hilbert spaces* are named after him
- *Wir müssen wissen.*  
*Wir werden wissen.*
- Translation:
  - We must know.  
We shall know.





