Cairo University - Faculty of Engineering
Computer Engineering Department

M-ARY AMPLITIUDE SHIFT MODULATION

Subject: Digital Communication

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0.1 Part 1: Digital Communication

0.1.1 Problem 1

Figure 1 below showing the comparison between simulated BER and theoritical (analytical) BER VS the Eb/N0 in db.

Please notice, you'll have to input the no. of bits you wish to be transmitted, and it has to be divisible by 3.

0.1.2 Problem 2

The constellation of the 8-ary with decision region pf each symbol.

Boundaries are at:

$$-6\sqrt{E}, -4\sqrt{E}, -2\sqrt{E}, 0, 2\sqrt{E}, 4\sqrt{E}, 6\sqrt{E}$$

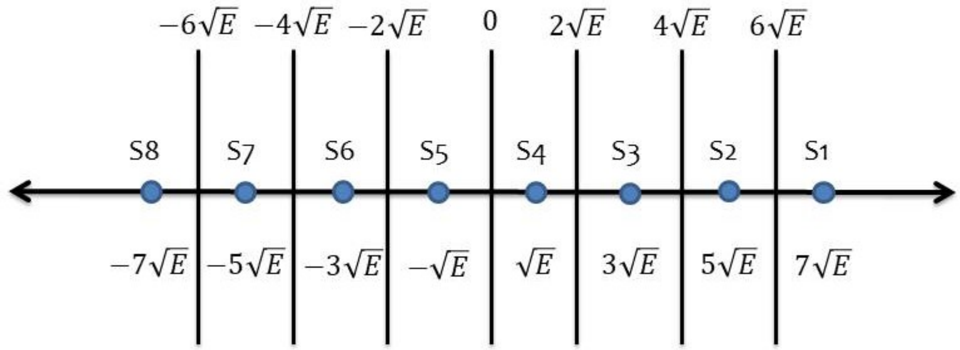

Figure 1: Symbols Boundary

0.1.3 Problem 3

The derivation of theoritical bit error rate using 8-ary.

$$Pe = \frac{1}{8}\sum_{i=0}^{7} P(e|Si) \tag{1}$$

$$Pe(e|S0) = Pe(e|S7) \tag{2}$$

$$Pe(e|S1) = Pe(e|S2) = Pe(e|S3) = Pe(e|S4) = Pe(e|S5) = Pe(e|S6) \tag{3}$$

Using Union bound S0, S7 only one neighbour and S1, S2,...S6 has two neighbours.

$$Pe(e|S0) = \frac{1}{2}erfc(\frac{\sqrt{E}}{\sqrt{N}}) \tag{4}$$

$$Pe(e|S1) = \frac{1}{2}erfc(\frac{\sqrt{E}}{\sqrt{N}}) + \frac{1}{2}erfc(\frac{\sqrt{E}}{\sqrt{N}}) \tag{5}$$

$$Pe(e|S1) = erfc(\frac{\sqrt{E}}{\sqrt{N}}) \tag{6}$$

$$Pe = \frac{1}{8*3}(2*\frac{1}{2}erfc(\frac{\sqrt{E}}{\sqrt{N}}) + 6*erfc(\frac{\sqrt{E}}{\sqrt{N}})) \quad (7)$$

$$Pe = \frac{7}{24}(erfc(\frac{\sqrt{E}}{\sqrt{N}})) \quad (8)$$

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Figure 1 below showing the comparison between simulated BER and theoritical (analytical) BER VS the Eb/N0 in db.

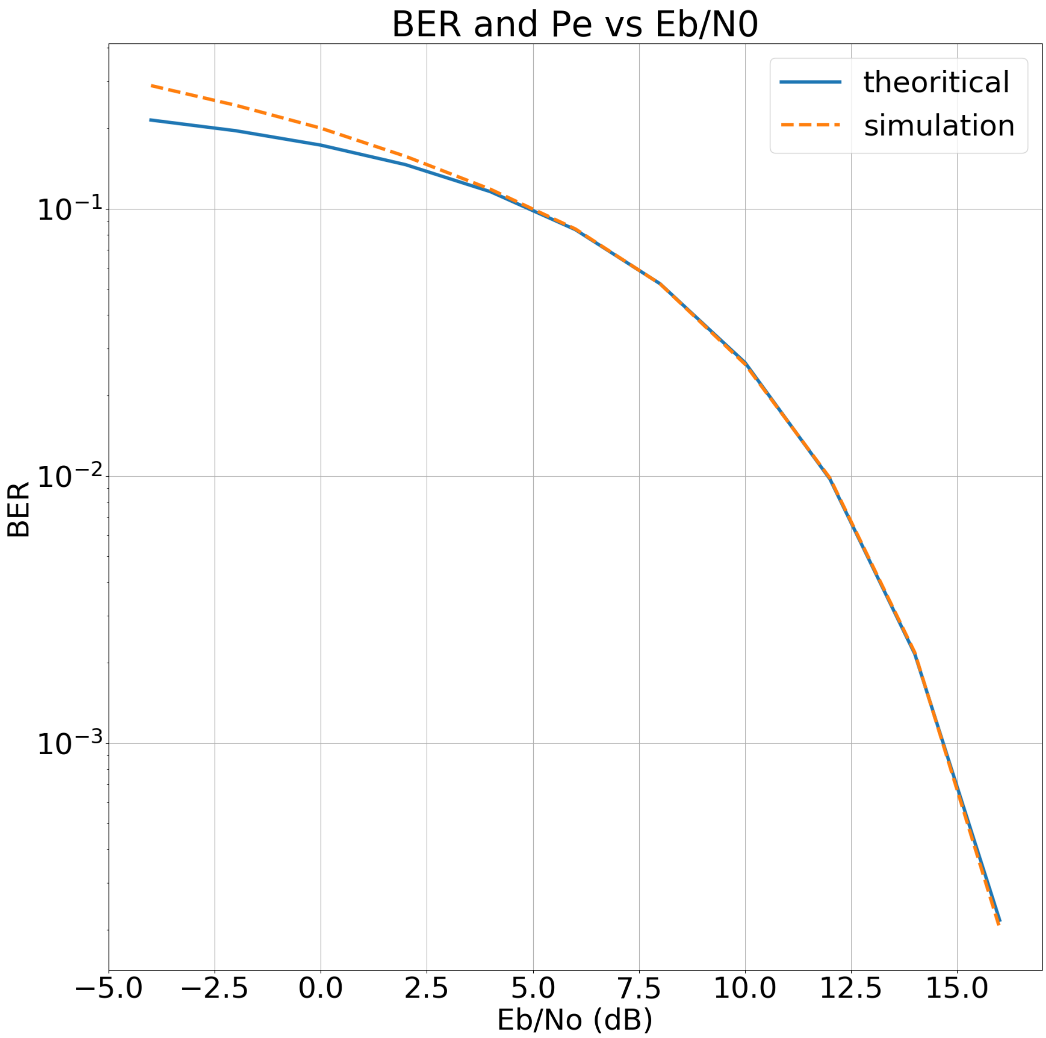

Figure 2: BER vs Eb/N0

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The answer is NO, We can't transmit at Rate 1 Mbps with bandwidth 0.5 MHz in passband transmition.

The minimum M required: 16 Only bit by bit transmittion is allowed.

GIVEN:

$$Bt = 2Rs$$

$$Rb = 1Mbps$$
$$BW = 0.5MHz$$
$$M = 3$$

REQUIRED:

$$BW =?2 * Rs$$
$$BW =?2 * \frac{Rb}{\log_2 M}$$
$$BW =?1Mbps * \frac{2}{3}$$

The Answer is: NO it can not be transmitted

$$0.5MHz =?1Mbps * \frac{2}{\log_2 M}$$
$$4MHz =?\log_2 M$$

and Minimum M alowed is:

$$M = 16$$

0.1.6 Probelm 6

Both of them satisfy the Gray Encoding criterion.

Because at the two examples only one bit is changed in each transition from symbol to the next one.

0.2 Part 2: Information Theory

Cyclic Codes.

0.2.1 Definition of Cyclic Codes

A Cyclic Codes is a block codes that follow the property of linearity where the circular i shift to the left (or n-1 shift to the right) always result in a word that belongs to the code words.

They are defined as $C(x) = (n, k)$ it means the message has k bits and the code vector is n bits.

$C(x)$ are defined with the help of generator polynomial $g(x)$.

The degree of $g(x)$ is equal to the number of parity-check digits of the code.

There are total of 2^k code polynomials in $C(x)$.

They are error-correction codes, used earlier to transmit images of planets.

They have algebric properties that help detecting and correcting errors.

Cyclic codes can be used to correct errors, it can be generalized to correct burst of errors, not just one bit.

For example the set of [000, 1111, 0110, 1001]:

They follow the Linearity property, but not the Cyclic Shift property, therefore they can't be considered a valide Cyclic Codes.

Another Example the set of [0000, 1111, 0101, 1010]:

They follow both the linearity and cyclic shift property. Because, when adding any two of them it will result in another (third) codeword that lies in the finite list.

0.2.2 Systematic CodeWords:

In Systematic Codewords $C(x) = [message, parity]$ this means, each generated codeword's first k bits are the message that got us that codeword while the remaining $n - k$ bits are the parity check of the codeword.

Systematic Codewords follow the following conditions:

$$C(x) = x^{n-k}m(x) + p(x)$$

Where:

$$p(x) = Rem\left[\frac{x^{n-k}m(x)}{g(x)}\right]$$

and,
$C(x)$ is codeword polynomial
$m(x)$ is message polynomial
$g(x)$ is generator polynomial

Example: Constructing a Systematic Cyclic Codes (7,4), using generator polynomial $g(x) = x^3 + x^2 + 1$, with message [1010].

Answer:

$m(x) = x^3 + x$
$p(x) = Rem\left[\frac{x^3*(x^3+x)}{(x^3+x^2+1)}\right]$
From division:
$p(x) = 1$
therefore,
$C(x) = x^3 * (x^3 + x) + 1$
$C(x) = x^6 + x^4 + 1$
$C(x) = [1010001]$

You may notice that the message $m(x) = 1010$ and the first k bits in the generated codeword is also 1010

Will the rest of the codeword is the parity check.

0.2.3 Relation between generator polynomial and generator matrix:

At the last section we've introduced the difinition of generator polynomial such that $g(x) = x^3 + x^2 + 1$, this polynomial is unique and we can derive all the other codewords from it by multiplying with the various messages allowed.

In cyclic code C, the generator matrix's dimentsions is $[n, k]$ n number of columns and k number of rows.

Generator Matric $[G]$ is composed of $[I, P]$, I is the identity matrix with dimentsions of $[k, k]$ while P is the parity matrix with the dimentsions of $[k, (n-k)]$

So, the main core of our problem is identifying the parity matrix, and it is identified as follows:

$$k^{th} Row = Rem\left[\frac{x^{n-k}}{g(x)}\right]$$

continuing on the previous example the generator matrix will be something like this:

$$G(x) = \begin{bmatrix} 1 & 0 & 0 & 0 & - & - & - \\ 0 & 1 & 0 & 0 & - & - & - \\ 0 & 0 & 1 & 0 & - & - & - \\ 0 & 0 & 0 & 1 & - & - & - \end{bmatrix}$$

so, to calculate the parity matrix we'll solve the following equations:

$$1^{st} Row = Rem\left[\frac{x^6}{x^3 + x^2 + 1}\right]$$
$$1^{st} Row = x^2 + 1 = [101]$$

$$2^{nd} Row = Rem\left[\frac{x^5}{x^3 + x^2 + 1}\right]$$
$$2^{st} Row = x^2 + x + 1 = [111]$$

$$3^{rd} Row = Rem\left[\frac{x^4}{x^3 + x^2 + 1}\right]$$

$$3^{rd}Row = x^2 + 1 = [101]$$

$$4^{th}Row = Rem\left[\frac{x^3}{x^3 + x^2 + 1}\right]$$

$$4^{th}Row = x + 1 = [011]$$

therefore the final matrix will be:

$$G(x) = \begin{bmatrix} 1 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 & 1 \end{bmatrix}$$