



# KSFUPRO1K1U – Functional Programming

## Lecture 2: Values, operators, expressions and functions

Niels Hallenberg

These slides are based on original slides by Michael R. Hansen, DTU. Thanks!!!



The original slides has been used at a course in functional programming at DTU.



# WELCOME to KSFUPRO1K1U – Functional Programming

**Teacher:** Niels Hallenberg, [nh@itu.dk](mailto:nh@itu.dk)  
Peter Sestoft, [sestoft@itu.dk](mailto:sestoft@itu.dk)  
IT University of Copenhagen

**Teaching assistant:** Jonas Christiansen Røyggaard (TA), [jocr@itu.dk](mailto:jocr@itu.dk)  
Jonas Ishøj Nielsen (TA), [join@itu.dk](mailto:join@itu.dk)  
IT University of Copenhagen

**Homepage:** <https://learnit.itu.dk/course/view.php?id=3018548>



Exam information is kept updated here:

<https://learnit.itu.dk/mod/wiki/create.php?wid=82&group&uid=12544&title=Exam>.

- **Date and Time:** May 25, 2019
- Exam Syllabus: Functional Programming using F#, Michael R. Hansen and Hans Rischel, ISBN 9781107684065, Chapter 1 - 13.
- **More to come . . .**

**You must pass the Mandatory assignments in order to attend exam.**

See course homepage in LearnIt for the details.

- Mandatory assignments does improve exam results!
- Main purpose of feedback on assignments is learning - not precise counting of points.
- If you are not satisfied with the grading, then just handin again with your corrections
- Remember the cut-off date - it is very important.



- A further look at functions, including higher-order (or curried) functions
- A further look at basic types, including characters, equality and ordering
- A first look at polymorphism
- A further look at tuples and patterns
- A further look at lists and list recursion

Goal: By the end of the day you are acquainted with a major part of the F# language.



Function expressions with general patterns, e.g.

```
function
| 2          -> 28  // February
| 4|6|9|11   -> 30  // April, June, September, November
| _          -> 31  // All other months
;;
```

Simple function expressions, e.g.

```
fun r -> System.Math.PI * r * r ;;
val it : float -> float = <fun:clo@10-1>

it 2.0 ;;
val it : float = 12.56637061
```



Simple function expressions with *currying*

$$\text{fun } x \ y \ \cdots \ z \rightarrow e$$

with the same meaning as

$$\text{fun } x \rightarrow (\text{fun } y \rightarrow (\cdots (\text{fun } z \rightarrow e) \cdots))$$

For example: The function below takes an integer as argument and returns a function of type `int -> int` as value:

```
fun x y -> x + x*y;;  
val it : int -> int -> int = <fun:clo@2-1>  
  
let f = it 2;;  
val f : (int -> int)  
  
f 3;;  
val it : int = 8
```

Functions are **first class citizens**:  
the argument and the value of a function may be functions



A simple function declaration:

`let  $f$   $x$  =  $e$`  means `let  $f$  = fun  $x$  →  $e$`

For example: `let circleArea r = System.Math.PI *r*r`

A declaration of a **curried function**

`let  $f$   $x$   $y$  ...  $z$  =  $e$`

has the same meaning as:

`let  $f$  = fun  $x$  → (fun  $y$  → ( ... (fun  $z$  →  $e$ ) ... ))`

For example:

```
let addMult x y = x + x*y;;  
val addMult : int -> int -> int
```

```
let f = addMult 2;;  
val f : (int -> int)
```

```
f 3;;  
val it : int = 8
```



Suppose that we have a cube with side length  $s$ , containing a liquid with density  $\rho$ . The weight of the liquid is then given by  $\rho \cdot s^3$ :

```
let weight ro s = ro * s ** 3.0;;  
val weight : float -> float -> float
```

We can make *partial evaluations* to define functions for computing the weight of a cube of either water or methanol:

```
let waterWeight = weight 1000.0;;  
val waterWeight : (float -> float)  
  
waterWeight 2.0;;  
val it : float = 8000.0  
  
let methanolWeight = weight 786.5 ;;  
val methanolWeight : (float -> float)  
  
methanolWeight 2.0;;  
val it : float = 6292.0
```





A *closure* is a value in F# representing a function and defined as a triple

$$(x, \text{exp}, \text{env})$$

where  $x$  is the argument to the function,  $\text{exp}$  is the expression to calculate (i.e., code) and  $\text{env}$  is the environment holding free variables.

Consider the function application `weight 786.5`. The result value is a closure

```
(s, ro*s**3.0, [ro->786.5,  
                *->the multiplication function,  
                **->the power function])
```



We have in previous examples exploited the pattern matching in function expression:

```
function
|  $pat_1$  →  $e_1$ 
  ⋮
|  $pat_n$  →  $e_n$ 
```

A **match expression** has a similar pattern matching feature:

```
match  $e$  with
|  $pat_1$  →  $e_1$ 
  ⋮
|  $pat_n$  →  $e_n$ 
```

The value of  $e$  is computed and the expression  $e_i$  corresponding to the first matching pattern is chosen for further evaluation.



Alternative declarations of the power function:

```
let rec power = function
  | (_,0) -> 1.0
  | (x,n) -> x * power(x,n-1);;
```

are

```
let rec power a = match a with
  | (_,0) -> 1.0
  | (x,n) -> x * power(x,n-1);;
```

and

```
let rec power(x,n) = match n with
  | 0 -> 1.0
  | n' -> x * power(x,n'-1);;
```



The prefix version ( $\oplus$ ) of an infix operator  $\oplus$  is a curried function.

For example:

```
(+);;  
val it : (int -> int -> int) = <fun:it@1>
```

Arguments can be supplied one by one:

```
let plusThree = (+) 3;;  
val plusThree : (int -> int)  
  
plusThree 5;;  
val it : int = 8
```

## Function composition: $(f \circ g)(x) = f(g(x))$



For example, if  $f(y) = y + 3$  and  $g(x) = x^2$ , then  $(f \circ g)(z) = z^2 + 3$ .

The infix operator `<<` in F# denotes functional composition:

```
let f y = y+3;;           // f(y) = y+3

let g x = x*x;;           // g(x) = x*x

let h = f << g;;          // h = (f o g)
val h : int -> int

h 4;;                     // h(4) = (f o g) (4)
val it : int = 19
```

Using just anonymous functions:

```
((fun y -> y+3) << (fun x -> x*x)) 4;;
val it : int = 19
```

Type of `(<<)` ?



The operator  $|>$  means “*send the value as argument to the function at the right*”

*arg*  $|>$  *fct* is equivalent to *fct arg*

The operator  $<|$  means “*send the value as argument to the functions at the left*”

*fct*  $|<$  *arg* is equivalent to *fct arg*

For example  $a+b |> \sin$  and  $\sin <| a+b$  means  $\sin(a+b)$ .

The two functions are predefined:

```
let (.<|. ) f a = f a
let (.|>. ) a f = f a
```

Type of  $(|>)$  ?



The basic types: integers, floats, booleans, and strings type were covered last week. Characters are considered on the next slide. For these types (and many other) equality and ordering are defined.

In particular, there is a function:

$$\text{compare } x \ y = \begin{cases} > 0 & \text{if } x > y \\ 0 & \text{if } x = y \\ < 0 & \text{if } x < y \end{cases}$$

For example:

```
compare 7.4 2.0;;  
val it : int = 1
```

```
compare "abc" "def";;  
val it : int = -3
```

```
compare 1 4;;  
val it : int = -1
```



It is often useful to have **when guards** in patterns:

```
let ordText x y = match compare x y with
    | t when t > 0 -> "greater"
    | 0             -> "equal"
    | _             -> "less";;

ordText "abc" "Abc";;
val it : string = "greater"
```

The first clause is only taken when `t > 0` evaluates to true.





The type of `ordText`

```
val ordText : 'a -> 'a -> string when 'a : comparison
```

contains

- a **type variable** `'a`, and
- a **type constraint** `'a : comparison`

The type variable can be instantiated to any type **provided** comparison is defined for that type. It is called a **polymorphic type**.

For example:

```
ordText true false;;  
val it : string = "greater"
```

```
ordText (1,true) (1,false);;  
val it : string = "greater"
```

```
ordText sin cos;;  
... '('a -> 'a)' does not support the 'comparison' ...
```

**Comparison is not defined for types involving functions.**



Type name: `char`

Values `'a'`, `' '`, `'\'` (escape sequence for `'`)

## Examples

```
let isLowerCaseVowel ch =  
    System.Char.IsLower ch &&  
    (ch='a' || ch='e' || ch = 'i' || ch='o' || ch = 'u');;  
val isLowerCaseVowel : char -> bool
```

```
isLowerCaseVowel 'i';;  
val it : bool = true
```

```
isLowerCaseVowel 'I';;  
val it : bool = false
```

The *i*'th character in a string is achieved using the "dot"-notation:

```
"abc".[0];;  
val it : char = 'a'
```



## A squaring function on integers:

Declaration	Type	
<code>let square x = x * x</code>	<code>int -&gt; int</code>	Default

## A squaring function on floats: `square: float -> float`

Declaration	
<code>let square(x:float) = x * x</code>	Type the argument
<code>let square x:float = x * x</code>	Type the result
<code>let square x = x * x: float</code>	Type expression for the result
<code>let square x = x:float * x</code>	Type a variable

You can mix these possibilities



An ordered collection of  $n$  values  $(v_1, v_2, \dots, v_n)$  is called an  $n$ -tuple

## Examples

<pre>(3, false); val it = (3, false) : int * bool</pre>	2-tuples (pairs)
<pre>(1, 2, ("ab", true)); val it = (1, 2, ("ab", true)) : ?</pre>	3-tuples (triples)

Equality defined componentwise, ordering lexicographically

```
(1, 2.0, true) = (2-1, 2.0*1.0, 1<2);;  
val it = true : bool
```

```
compare (1, 2.0, true) (2-1, 3.0, false);;  
val it : int = -1
```

provided = is defined on components



## Extract components of tuples

```
let ((x,_), (_,y,_)) = ((1,true), ("a","b",false));;  
val x : int = 1  
val y : string = "b"
```

## Pattern matching yields bindings

### Restriction

```
let (x,x) = (1,1);;  
...  
... ERROR ... 'x' is bound twice in this pattern
```



## Examples

```
let g x =  
  let a = 6  
  let f y = y + a  
  x + f x;;  
val g : int -> int
```

```
g 1;;  
val it : int = 8
```

Note: **a** and **f** are not visible outside of **g**



Example: Solve  $ax^2 + bx + c = 0$

```
type Equation = float * float * float
type Solution = float * float
exception Solve; (* declares an exception *)
```

```
let solve(a,b,c) =
  if b*b-4.0*a*c < 0.0 || a = 0.0 then raise Solve
  else ((-b + sqrt(b*b-4.0*a*c))/(2.0*a),
        (-b - sqrt(b*b-4.0*a*c))/(2.0*a));;
val solve : float * float * float -> float * float
```

The type of the function `solve` is (the expansion of)

`Equation -> Solution`

`d` is declared once and used 3 times

readability, efficiency



```
let solve(a,b,c) =  
  let d = b*b-4.0*a*c  
  if d < 0.0 || a = 0.0 then raise Solve else  
    ((-b + sqrt d)/(2.0*a), (-b - sqrt d)/(2.0*a));;
```

```
let solve(a,b,c) =  
  let sqrtD =  
    let d = b*b-4.0*a*c  
    if d < 0.0 || a = 0.0 then raise Solve  
    else sqrt d  
  ((-b + sqrtD)/(2.0*a), (-b - sqrtD)/(2.0*a));;
```

Indentation matters



## Example: Rational Numbers



Consider the following **signature**, specifying operations and their types:

Specification	Comment
<code>type qnum = int * int</code>	rational numbers
<code>exception QDiv</code>	division by zero
<code>mkQ: int * int → qnum</code>	construction of rational numbers
<code>+. : qnum * qnum → qnum</code>	addition of rational numbers
<code>-. : qnum * qnum → qnum</code>	subtraction of rational numbers
<code>.*. : qnum * qnum → qnum</code>	multiplication of rational numbers
<code>./. : qnum * qnum → qnum</code>	division of rational numbers
<code>.=. : qnum * qnum → bool</code>	equality of rational numbers
<code>toString: qnum → string</code>	String representation of rational numbers



```
let q1 = mkQ(2,3);;
```

$$q_1 = \frac{2}{3}$$

```
let q2 = mkQ(12, -27);;
```

$$q_2 = -\frac{12}{27} = -\frac{4}{9}$$

```
let q3 = mkQ(-1, 4) .* q2 ./ q1;;
```

$$q_3 = -\frac{1}{4} \cdot q_2 - q_1 = -\frac{5}{9}$$

```
let q4 = q1 ./ q2 ./ q3;;
```

$$q_4 = q_1 - q_2 / q_3 = \frac{2}{3} - \frac{-4}{9} / \frac{-5}{9}$$

```
toString q4;;
```

```
val it : string = "-2/15"
```

$$= -\frac{2}{15}$$

Operators are infix with usual precedences

Note: Without using infix:

```
let q3 = (.-.)((.*.) (mkQ(-1,4)) q2) q1;;
```

Representation:  $(a, b)$ ,  $b > 0$  and  $\text{gcd}(a, b) = 1$



Example  $-\frac{12}{27}$  is represented by  $(-4, 9)$

### Greatest common divisor (Euclid's algorithm)

```
let rec gcd = function
  | (0,n) -> n
  | (m,n) -> gcd(n % m,m);;
val gcd : int * int -> int
```

```
- gcd(12,27);;
val it : int = 3
```

Function to **cancel** common divisors:

```
let canc(p,q) =
  let sign = if p*q < 0 then -1 else 1
  let ap = abs p
  let aq = abs q
  let d = gcd(ap,aq)
  (sign * (ap / d), aq / d);;
```

```
canc(12,-27);;
val it : int * int = (-4, 9)
```



Declaration of the constructor:

```
exception QDiv;;  
let mkQ = function  
  | (_, 0)  -> raise QDiv  
  | pr      -> cancel pr;;
```

Rules of arithmetic:

$$\begin{aligned}\frac{a}{b} + \frac{c}{d} &= \frac{ad+bc}{bd} & \frac{a}{b} - \frac{c}{d} &= \frac{ad-bc}{bd} \\ \frac{a}{b} \cdot \frac{c}{d} &= \frac{ac}{bd} & \frac{a}{b} / \frac{c}{d} &= \frac{a}{b} \cdot \frac{d}{c} \quad \text{when } c \neq 0 \\ \frac{a}{b} = \frac{c}{d} &= ad = bc\end{aligned}$$

Program corresponds directly to these rules

```
let (+. ) (a,b) (c,d) = cancel(a*d + b*c, b*d);;  
let (-. ) (a,b) (c,d) = cancel(a*d - b*c, b*d);;  
let (*. ) (a,b) (c,d) = cancel(a*c, b*d);;  
let (./.) (a,b) (c,d) = (a,b) *. mkQ(d,c);;  
let (.=.) (a,b) (c,d) = (a,b) = (c,d);;
```

Note: Functions must preserve the **invariant** of the representation



Consider `unzip` that maps a list of pairs to a pair of lists:

$$\begin{aligned} \text{unzip}([ (x_0, y_0); (x_1, y_1); \dots; (x_{n-1}, y_{n-1}) ]) \\ = ([x_0; x_1; \dots; x_{n-1}], [y_0; y_1; \dots; y_{n-1}]) \end{aligned}$$

with the declaration:

```
let rec unzip = function
  | []          -> ([], [])
  | (x,y)::rest -> let (xs,ys) = unzip rest
                   (x::xs,y::ys);;

unzip [(1,"a"); (2,"b")];;
val it : int list * string list = ([1; 2], ["a"; "b"])
```

Notice

- pattern matching on result of recursive call
- `unzip` is polymorphic. Type?
- `unzip` is available in the `List` library.



You are acquainted with a major part of the F# language.

- Higher-order (or curried) functions
- Basic types, equality and ordering
- Polymorphism
- Tuples
- Patterns
- A look at lists and list recursion