

Quantum Information Theory

quinten tupker

January 22 2021 - January 25, 2021

Introduction

These notes are based on the course lectured by Professor S Strelchuk in Lent 2020. This was lectured online due to measures taken to counter the spread of Covid-19 in the UK. These are not necessarily an accurate representation of what was lectures, and represent solely my personal notes on the content of the course, combining with probably, very very many personal notes and digressions... Of course, any corrections/comments would be appreciated.

Information theory is the theory of information storage and transmission. It provides the theoretical limits on what is possible with information technologies in much of our world, and a framework to study many other fields, such as animal communication as well. This is the quantum version of that theory.

1 Classical Information Theory

We begin by observing that information is closely related to uncertainty. In particular, what might be able to say it is the opposite of uncertainty, and so then, it is no surprise that we build our theory of information using concepts from probability theory. As such we define

Definition 1. The **surprisal** of random variable X taking values in discrete finite **alphabet** J according to distribution $p(x)$ is

$$\mathcal{I}(x) = -\ln(p(x))$$

Definition 2. The **Shannon entropy** of a random variable X is

$$H(X) = -\sum_{x \in J} p(x) \ln(p(x))$$

(the logarithm is base 2)

This may appear to be a somewhat arbitrary definition, but it has a strong theoretical basis given that

Theorem 1. *The **Shannon Source Coding Theorem** states (informally) that the limit that information can be compressed so that it can be reliably retrieved is the Shannon entropy of the source.*

Here a basic example of a source is a **memoryless source** which is an object producing a sequence of signals, but since it is memoryless, each signal is completely independent from any other, so $\mathbb{P}(u_1, \dots, u_n) = \mathbb{P}(u_1) \dots \mathbb{P}(u_n)$. It is also known as an **i.i.d. information source**.

But how do we actually compress information? Conceptually there are two ways

- a **variable length encoding** stores higher probability signals in shorter codes, and lower probability signals in longer codes.
- a **fixed length encoding** stores higher probability signals in unique fixed length codes, and lower probability signals in the same fixed code.

Example 1. *An example of a fixed length code for the numbers $1, \dots, 8$ are their binary representations, but if we also know that $p(1, \dots, 8) = 1/2, 1/4, 1/8, 1/16, 1/64, 1/64, 1/64, 1/64$ then the code $C(1, \dots, 8) = 0, 10, 110, 1110, 111100, 111101, 111110, 111111$ has an average length 2 compared to the fixed length of 3. Furthermore, the Shannon entropy of this source is 2 as well, so this is maximally efficient.*

1.1 Classical Data Compression

Let's start making the definitions necessary to formalise compression.

Definition 3. A **compression map** is a map $C^n : u^{(n)} = (u_1, \dots, u_n) \mapsto x(x_1, \dots, x_{nR})$ sending a **message** u to a **code** x .

Definition 4. A **decompression map** D^n sends $D^n : x \in \{0, 1\}^{[nR]} \mapsto u'^{(n)}$ with probability $\mathbb{P}(u^{(n)}|x)$.

Definition 5. A **code** of rate R and blocklength n is the triple (C^n, D^n, R) .

Here we can compute the **probability of error** as

$$P_{av}^{(n)}(C_n) = \sum_{u^{(n)} \in J^n} \mathbb{P}(u^{(n)}) \mathbb{P}(D^n(C^n(u^{(n)})) \neq u^{(n)}). \quad (1)$$

[End of lecture 1]

Definition 6. A compression/decompression scheme is **reliable** iff $\forall \epsilon > 0 \exists$ sequence of codes C_n such that $\lim_{n \rightarrow \infty} P_{av}^{(n)}(C_n) = 0$

(does this require the code to i.i.d?) As such one could define the data compression rate as

$$\inf\{R : \exists C_n = \{C^n, D^n, R\} : \lim_{n \rightarrow \infty} P_{av}^{(n)}(C_n) = 0\} \quad (2)$$

1.2 Typical Sequences

It is honestly surprising how easy it is to prove Shannon's source coding theorem, since we find we in fact only need one simple, and honestly rather crude tool to do so. That tool is

Definition 7. a **typical set**, denoted $T_\epsilon^{(n)}$, which is the set of sequences $u = (u_1, \dots, u_n)$ satisfying

$$2^{-n(H(u)+\epsilon)} \leq \mathbb{P}(u) \leq 2^{-n(H(u)-\epsilon)} \quad (3)$$

Why "typical"? Because if we have a memoryless source generating a sequence of length n , then a "typical" sequence occurs with probability

$$\prod_{u \in J} \mathbb{P}(u)^{n\mathbb{P}(u)} = 2^{-nH(u)} \quad (4)$$

(a typical sequence would have the expected value as the number of occurrences of each letter in the alphabet). We consequently find the following theorem (not proven)

Theorem 2. $\forall \epsilon, \delta > 0, \exists n$ such that

1. $u \in T_\epsilon^{(n)} \implies H(u) - \epsilon \leq \frac{-1}{n} \log(\mathbb{P}(u)) \leq H(u) + \epsilon$
2. $\mathbb{P}(T_\epsilon^{(n)}) > 1 - \delta$
3. $|T_\epsilon^{(n)}| \leq 2^{n(H(u)+\epsilon)}$
4. $|T_\epsilon^{(n)}| > (1 - \delta)2^{n(H(u)-\epsilon)}$

and so

Corollary 1. $\forall \epsilon, \delta > 0, \exists n_0, \forall n > n_0, J^n$ decomposes into the disjoint **atypical** and **typical sets** $A_\epsilon^{(n)}, T_\epsilon^{(n)}$ satisfying

1. $\mathbb{P}(A_\epsilon^{(n)}) < \delta$
2. $2^{-n(H(u)+\epsilon)} \leq \mathbb{P}(u) \leq 2^{-n(H(u)-\epsilon)}$ (on the typical set?)

That allows us to formally state

Theorem 3. Shannon's source coding theorem, which claims that for a i.i.d. source U , if $R > H(U)$, then we can find a reliable compression scheme and if $R < H(U)$ there are no reliable compression schemes.

Proof. If $R > H(U)$ then pick $\epsilon > 0$ such that $H(U) + \epsilon < R$ and n such that $T_\epsilon^{(n)}$ satisfies our typical set theorem conditions. Then for $\delta > 0$ there are at most $2^{n(H(U)+\epsilon)} < 2^{nR}$ ϵ -typical sequences. Our compression scheme then works according to

1. Split J^n into typical and atypical sequences
2. order typical sequences somehow (say lexicographically), assigning each sequence a binary index.
3. typical sequences are sent to their binary code, prefixed with a 1, leading to a total length of $\lceil nR \rceil + 1$
4. atypical sequences are all sent to the fixed string $00 \dots 0$ of length $\lceil nR \rceil + 1$

If $R < H(U)$ there does not exist a reliable compression scheme (see lemma below). \square

Lemma 1. *For a collection of strings of length n , $S(n)$, with $|S(n)| \leq 2^{nR}$, $R < H(U)$. Then, $\forall \delta > 0, \exists n, \sum_{u \in S(n)} \mathbb{P}(u) < \delta$*

Basically, when $R < H(U)$ we will always get some typical sets that are indistinguishable after compression. Assume $S(n)$ is a set of such indistinguishables. This lemma states that these become less significant?

Proof. Again, we split $S(n)$ into its typical and atypical part and then observe

$$\begin{aligned}
\mathbb{P}(S(n)) &= \sum_{u \in S(n)} \mathbb{P}(u) \\
&= \sum_{u \in S(n) \cap T_\epsilon^{(n)}} \mathbb{P}(u) + \sum_{u \in S(n) \cap A_\epsilon^{(n)}} \mathbb{P}(u) \\
&\leq |S(n)| 2^{-n(H(n) - \epsilon)} + \mathbb{P}(A_\epsilon^{(n)}) \\
&\leq 2^{-n(H(n) - R)} + A_\epsilon^{(\kappa)}
\end{aligned}$$

\square

which completes the proof of Shannon's source coding theorem. Finally, we note that using a variable length encoding scheme or something like that does not fundamentally change this result. [End of lecture 2]