VBS and VSM Internals

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whoami

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- Security Researcher
 - vulnerabilities && exploitation
 - lowlevel internals
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- Addicted to CTFs!
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Talk outline

- Windows before VBS
- VSM 101
 - VTLs, OS architecture, new components
- Attack surface
- Few bugs I found and reported ©

Motivation

- Windows kernel exposes a HUGE attack surface
 - Networking, filesystems, fonts, scrollbars, etc.
 - However, the kernel is privileged => can basically do whatever it wants
- Can't just change the entire OS architecture and design in one day...

Motivation

- Before VBS kernel is the root of trust
 - Hyper-V fully trusts the root partition
- <u>After VBS</u> we have a generic architecture to implement independent secure trusted components
 - Without a single binary for everything
 - If you own one component you still have work to do!
- Outcome: we made ring0 more restricted, and mitigated kernel exploits! (at least some of them;))

The world - after VBS

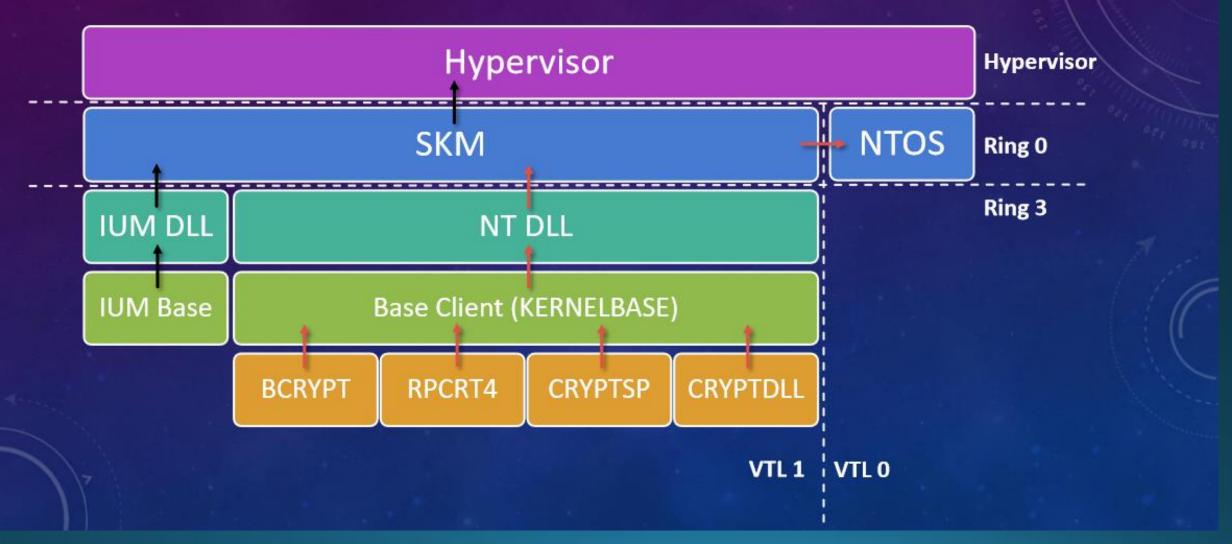
- Kernel code signing ntos can't have any +X unsigned pages
- Device Guard strict execution control
 - PowerShell scripts
 - User/kernel mode binaries
 - Everything!
- Credential Guard protects cryptographic secrets (LSA)
 - From anyone on the machine even kernelspace code!
- Extendible architecture: many more, and more to come
- Today we will focus on kernel exploit mitigations

VBS / VSM 101

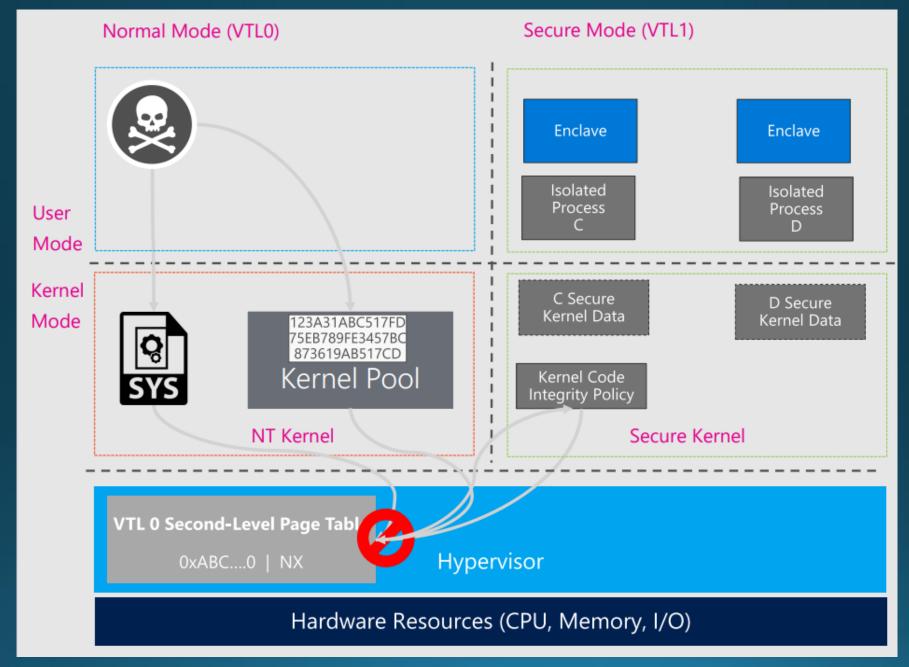
VSM 101

- Now, we have VTLs, which are orthogonal to rings
 - VTL1 is secure world, VTL0 is normal world => the higher the VTL is, the more privileged it gets
 - No more "ring -1/ring -2" confusion: more VTLs may be added in the future
 - Lets us write trusted code in userspace runs in ring3vtl1
- All managed by Hyper-V
 - securekernel runs in ring0vtl1
 - ntos runs in ring0vtl0
- Hyper-V exposes 2 hypercalls for normal call and secure call
 - Normal call services provided by NTOS to SK
 - Secure call services provided by SK to NTOS
 - Securekernel!lumInvokeSecureService

ARCHITECTURAL LAYER OVERVIEW



Source: Alex Ionescu, Blackhat 15



Source: Dave Weston && Alex Ionescu, OPCDE 2018

hv!vmcall_handler

rcx, [r15+7C0h]

[rsp+88h+var_58], r8

call rax pause lfence

d1, r12b

r8, [rcx]

rcx, r15

vmcall_handler+14C

vmcall_handler+153

vmcall handler+156

vmcall_handler+159

vmcall_handler+15C

vmcall handler+161

mov

mov

mov

mov

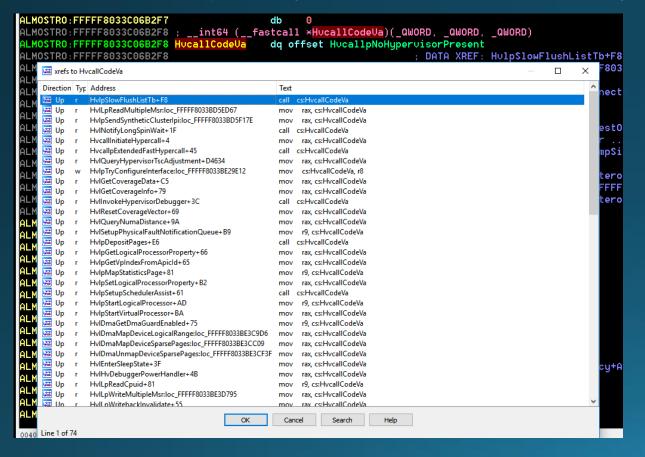
mov

call

```
🗾 🚄 🖼
vmcall_handler+11A
                                                                   umcall_handler+130
vmcall handler+11A
                                                                                     vmcall_11_flow:
                   umcall 12 flow:
                                                                   umcall handler+130
vmcall handler+11A
                           sub FFFFFAEA088901E4
                                                                   umcall handler+130
                                                                                      call
                                                                                              sub FFFFFAEA088901E4
                   call
umcall_handler+11F
                                                                   umcall handler+135
                                                                                              al. al
                   test
                           al. al
                                                                                      test
                           return HU STATUS_INVALID_HYPERCALL_CODE
                                                                   vmcall_handler+137
vmcall_handler+121
                   iz
                                                                                      iz
                                                                                              return_HU_STATUS_INVALID_HYPERCALL_CODE
 💶 🚄 🖼
 vmcall_handler+127
                            rax. HoCallUtlReturn
                                                        vmcall_handler+13D lea
                                                                                    rax, HuCallUtlCall
                                                                                                                 umcall_handler+9E
                    lea
                                                                                                                                     jz
 vmcall_handler+12E
                    ami
                            short handle_utl_transition
                                                                                                                 vmcall_handler+A1
 vmcall_handler+144
                       amarsa: here call to one of the followin
 vmcall_handler+144
                       HuCallUtlCall (if umcall 0x11)
                      HuCallUtlReturn (if umcall 0x12)
 vmcall_handler+144
 vmcall handler+144
 vmcall handler+144
                    handle utl transition:
 vmcall handler+144
                            byte ptr [r15+8], 3
                     mov
                            r12b, 1
 vmcall handler+149
                     mov
```

In ntos

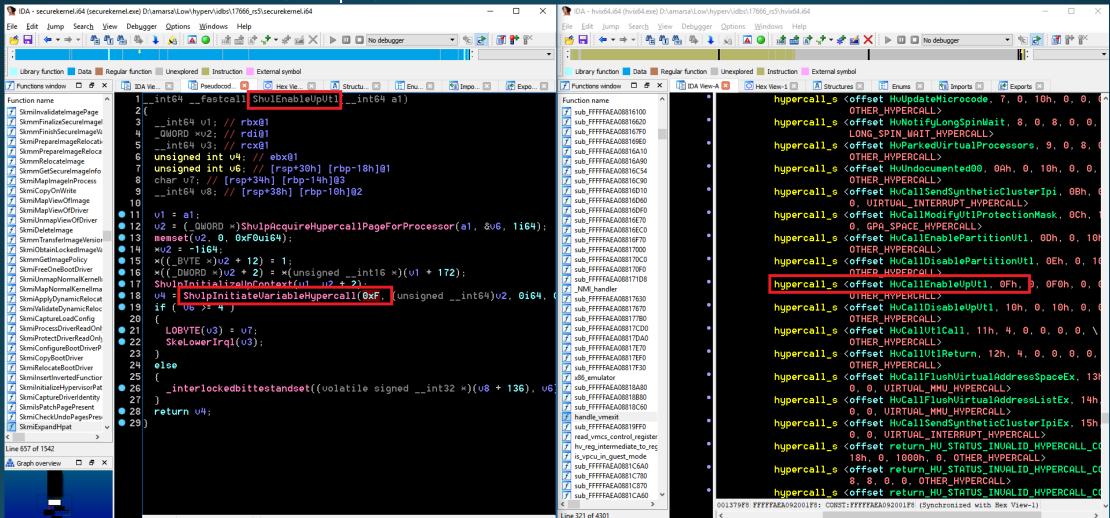
- Symbol nt!HvcallCodeVa
- Calling to hypercalls wrapped by convention nt!Hvl*



```
nt!DbgBreakPointWithStatus:
fffff803`6ca51e20 cc
                                   int
kd> dq nt!HvcallCodeVa L1
fffff803`6cd132f8 fffff803`6c77b000
kd> u poi(nt!HvcallCodeVa)
fffff803`6c77b000 0f01c1
                                   vmcal1
fffff803`6c77b003 c3
                                   ret
fffff803`6c77b004 8bc8
                                   mov
                                           ecx,eax
fffff803`6c77b006 b811000000
                                           eax,11h
                                   mov
                                   vmcal1
fffff803`6c77b00b 0f01c1
fffff803`6c77b00e c3
                                   ret
fffff803`6c77b00f 488bc1
                                   mov
                                           rax, rcx
fffff803`6c77b012 48c7c111000000
                                           rcx,11h
                                   mov
```

In securekernel

- Symbol securekernel!HvcallCodeVa
- Prefixes: shvl*, skp*, ...



MMU Security

- MMU configuration is restricted to VTL1 only
- The processor uses the SLAT (EPT) instead of PTEs
- ntos manages PTEs, Hyper-V and sk are the only ones with write access to EPT
- Result: with VBS enabled, virtual to physical translation is not managed by ntos!
 - MMU uses SLAT instead of PTEs

Implementing Mitigations

- So we split the kernel into separate spaces
- And we have IPC between VTLs.
- Now it's trivial to implement the mitigations we talked about!
- VTL1, Hyper-V and CPU can have these "contracts":
 - Nobody can read this physical page Credential Guard, Lsalso.exe
 - Block +RX/+RWX <u>Device Guard</u>
 - Don't execute code from this physical page (even when PTE has -NX, kills kernel exploits that generate code)

Execution Security

- New: Mode Based Execution Control (MBEC)
 - Even if the processor doesn't support it, Hyper-V emulates it
 - Can't disable SMEP and jump directly to userspace anymore!
- The new design lets us build really strong mitigations like that
- Implementation:
 - ntos can't allocate executable physical pages (only sk can)
 - ExAllocatePoolWithTag(poolType=0) actually causes a BSOD
 - Instead, it calls into sk for this allocation
 - sk uses skci to verify that all code in ring0 is signed

Other pieces of the puzzle

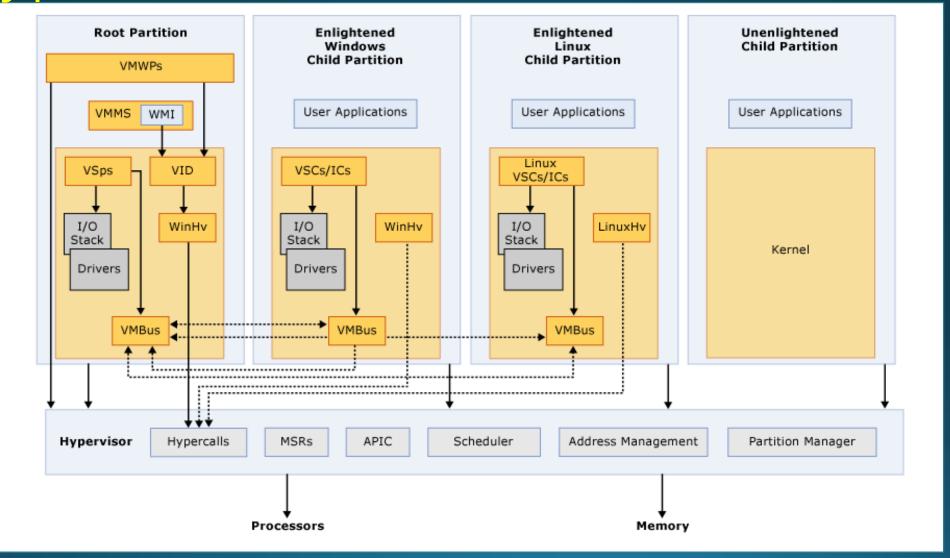
- Secureboot
 - Otherwise, can write into bootloader from the root partition, reboot and profit
- VTd
 - Otherwise, possible to overwrite HV via DMA attacks
- TPM
 - Only way to securely hold encryption keys for hiberfile across sleep states
 - Otherwise, hibernate, modify Hyper-V text section, boot and profit
 - Critical to secure S4 sleep state

VBS Attack Surface

Attack surface

- Heart of all that architecture is what we covered
 - Hypervisor platform/services, VTL1 securekernel and trustlets
- So, immediate attack surface:
 - Services exposed by VTL1: securekernel!lumInvokeSecureService
 - VTL1 calls into VLT0 for normal services and it has to sanitize responses correctly
 - RPC with trustlets
 - Hyper-V: hypercalls, MSRs, IO ports, VMBus, etc...

Hypervisor architecture



Hypervisor attack surface

- Hypercall handlers
- MSRs
 - Reads and writes MSR handlers called from vmexit loop handler
- x86_emulator
 - Called from the vmexit loop handler, emulates many interesting things:
 - Instructions
 - Triple-fault
 - etc.
- Virtual address emulation

```
case EXIT_REASON_APIC_WRITE:
    read_EXIT_QUALIFICATION(&v_exit_qualification);
    emulate_apic_write(
      υ93.
      v_exit_qualification & 0xFF0,
      \times( DWORD \times)(\times(_QWORD \times)(\times(_QWORD \times)(\cup16 + 3344) + 280i64) + (\cup_exit_qualification & 0xFF0)));
    goto LABEL 219:
  case EXIT_REASON_APIC_ACCESS[EXIT_REASON_EXTERNAL_INTERRUPT:
    read_EXIT_QUALIFICATION(&v106);
    emulate_apic_access_or_external_interrupt_handling(v17, v106);
    goto LABEL_219;
if ( (_DWORD)v_exit_reason != EXIT_REASON_APIC_ACCESS )
  if ( (_DWORD)v_exit_reason == EXIT_REASON_EPT_UIOLATION )
    handle_ept_violation(v16 + 0xD00, a1);
  else if ( (_DWORD)v_exit_reason == EXIT_REASON_EPT_MISCONFIG )
    if ( ( DWORD) RCX == 0x10000031 )
      handle_ept_misconfig((__int64)v17);
  else
    x86_emulator(
      (unsigned __int64 **)a1,
      v_exit_reason,
      a5,
      (__int64)a1,
      υ19.
      a6,
      a7,
      a8.
```

```
if ( (_DWORD)exit_reason == EXIT_REASON_TRIPLE_FAULT )
   if ( is_vpcu_in_guest_mode((__int64)*(a1 - 200)) )
     sub_FFFFFAEA088F6088((__int64)v16, 2, 0i64, 0);
   else
     check_if_keRaiseSystemError((__int64)v16, v76);
   return;
 if ( (_DWORD)v_exit_reason == EXIT_REASON_INIT )
   call_do_reboot_((__int64)a1, exit_reason);
 switch ( (_DWORD)v_exit_reason )
   case EXIT_REASON_PENDING_UIRT_NMI:
     if ( is_vpcu_in_guest_mode((__int64)*(a1 - 200)) && (unsigned __int8)does_cpu_have(v69, 0x400000i64) )
       \times((_DWORD \times) \cup 16[33] + 92) = 26;
       read_GUEST_INTERRUPTIBILITY_INFO(&v86);
       if ( U86 & 1 )
          ∪86 = ∪86 & 0xFFFFFFF | 2;
         write_GUEST_INTERRUPTIBILITY_INFO(&v86);
00018350 x86 emulator:76
```

```
case EXIT_REASON_TASK_SWITCH:
     handle_reason_task_switch(v15 + 3328, (__int64)v16);
     return:
   case EXIT_REASON_CPUID:
     if ( !is_vpcu_in_guest_mode((__int64)*(a1 - 200)) );
       handle_reason_cpuid((__int64)v16);
       return:
     \times((_DWORD \times) \cup 16[33] + 92) = 18;
lo_exit_inst_flow_return:
     do_exit_inst_flow();
     return:
   case EXIT REASON INUD:
     if ( !is_vpcu_in_guest_mode((__int64)*(a1 - 200)) )
       handle_reason_invd((__int64)v16);
       return;
     goto do_exit_inst_flow_return;
   case EXIT_REASON_INULPG:
     read_EXIT_QUALIFICATION(v16 + 2);
     if ( is_vpcu_in_quest_mode(v67) && (unsigned __int8)does_cpu_have(v68, 0x200i64) )
0018CF1 x86 emulator:133
```

Hypervisor attack surface

- The root partition can access almost the entire physical address space, including:
 - MMIO
 - exceptions being LAPIC
 - VTd bars
- Some pages are shared with VTL1
 - Libraries, like iumdll.dll
 - Marked as RO in the root partition's EPT
- With HVCI enabled:
 - UEFI runtime pages are marked as RO
 - UEFI runtime executes in VTL1 context
- Without HVCI
 - UEFI runtime pages are marked as writable!
 - UEFI runtime executes in root partition kernel context

Example: HVCI bypass

- CVE-2016-0181 by Rafal Wojtczuk
- There were some pages with RWX permission in ring0 EPT
 - Likely artifacts of early boot phase / EFI code
- Trivial HVCI bypass execute unsigned code in ring0vtl0
 - With VBS, it is a security boundary ©

```
C:\Users\testuser\probex>probex 0xf000000 2000000
ntoskrnl.exe at FFFFF802EDA82000
tryexcept at FFFFF802EDE5E292
kthread at FFFFE000E6CBC080
stack_base=FFFFD000769C5000 limit=FFFFD000769CB000
starting phys mem probe:
0xf000000-0x100000000 (size 0x1000000): not rwx
0x10000000-0x10157000 (size 0x157000): rwx
0x10157000-0x11000000 (size 0xea9000): not rwx
```

Hypervisor attack surface

- PCIe config space (accessible via MMCFG)
 - Device-specific registers
 - Memory bars locations
 - REMAP_LIMIT/REMAP_BASE are locked (trivially)
 - What can we overlap?
- 10 port
 - All are directly accessible, with some exceptions
 - · Easy to see them, configured in the root partition's VMCS IO bitmap
- Finally, in RS4: protect vtl0 from DMA attacks
 - not only vtl1:)
 - Using IOMMÚ remapping (hal!dmr*)
 - https://twitter.com/AmarSaar/status/985618204184768513

VTL1 attack surface

- Services exposed by VTL1
 - Lsalso RPC services
 - Lots of RPC demarshalling code (vtl1ring3)
 - >50 services implemented in securekernel!lumInvokeSecureService (vtl1ring0)
- VTL1 extensively calls to the root partition for some services
 - We want to maintain lots of the current mechanisms in VTL0
 - To keep VTL1 attack surface as small as possible
 - So, lots of sanitizing/parsing of responses from VTL0 occurs in VTL1

VMBus && vmswitch

- VMBus has a serious part in this design
 - IPC between partitions
 - Really important in Azure
- VMBus has to do lots of parsing, marshalling, etc...
 - Can be complex and sensitive code
- Same goes for vmswitch
 - Google Project Zero found a potential g2h there
 - https://bugs.chromium.org/p/project-zero/issues/detail?id=688

Hyper-V vmswitch.sys VmsMpCommonPvtHandleMulticastOids Guest to Host Kernel-Pool Overflow

Reported by kost...@google.com, Jan 4 2016

This bug is subject to a 90 day disclosure deadline. If 90 days elapse without a broadly available patch, then the bug report will automatically become visible to the public.

This function is reachable by sending a RNDIS Set request with OID 0x01010209 (OID_802_3_MULTICAST_LIST) from the Guest to the Host.

This function potentially allocates a buffer based on the addresses sent. The number of entries is determined by dividing the length of the data by 6:

```
.text:000000000001D717 mov eax, 0AAAAAAABh
 .text:000000000001D71C mov r13b, 1
 .text:000000000001D71F mul r14d
 .text:000000000001D722 mov ebp, edx
 .text:000000000001D724 shr ebp, 2
 .text:000000000001D727 test ebp, ebp ; ebp=r14d//6
 .text:000000000001D729 jz loc 31B04
 .text:000000000001D72F
 .text:0000000001D72F loc 1D72F: ; CODE XREF: VmsMpCommonPvtHandleMulticastOids+144CEj
 .text:00000000001D72F cmp ebp, [rbx+0EE8h]
 .text:000000000001D735 jz loc 31B2B
 .text:00000000001D73B mov r8d, 'mcMV'; Tag
 .text:000000000001D741 mov rdx, r14; NumberOfBytes
 .text:00000000001D744 mov ecx, 200h ; PoolType
 .text:000000000001D749 mov r12, r14
 .text:00000000001D74C call cs: imp ExAllocatePoolWithTag .text:00000000001D752 mov r14, rax
 .text:000000000001D755 test rax, rax
.text:000000000001D758 jz loc 1D7E8
 .text:000000000001D75E mov r8, r12; Size
 .text:000000000001D761 mov rdx, r15 ; Src
 .text:000000000001D764 mov rcx, rax ; Dst
 .text:000000000001D767 call memmove
An interesting test is located at 0x1D72F.
If the number of entries is identical to the currently stored one, then we jump to this piece of code:
```

So... what's new?

Last year...

- MS keeps adding code, features and patches to these components
- It's REALLY important to keep track of that...



Microcode update && Hotpatch

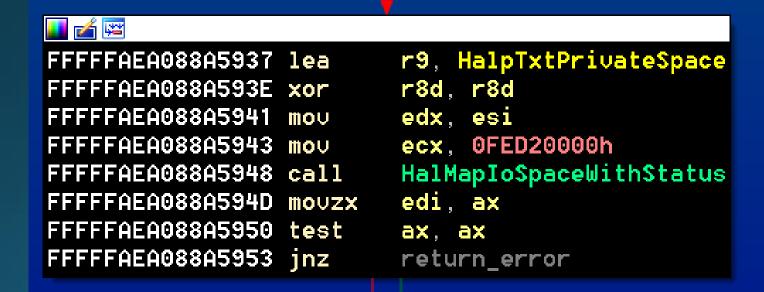
- Microcode update and hotpatch mechanisms were added to the hypervisor in RS4
- Microcode update:
 - look out for a new hypercall: 0x7
 - Called from ntos: nt!HvlUpdateMicrocode
- Easy flow to reverse, see known MSRs:

IDA View-A	Hex View-1 🔲 🗚 Structures 🖸	Enums	E3	Imports	1	Exports	Matched Functions	Primary Unmatched
EA	Name	Basicblock	Instruc	tions Edges				
0000000140019E10	ShvlGetInterceptData	1	20	0				
0000000140019E68	ShvlCompleteIntercept	9	50	11				
000000014001A2E8	ShvIPrepareHypervisorPatch	18	78	24				
000000014001A414	ShvlLoadHypervisorPatch	22	174	32				
000000014001F080	SkmiConvertImagePfn	4	19	4				
000000014001F0D4	SkmiReplicateGapFrame	5	59	6				
0000000140025A10	SkmiReserveHpatAddresses	12	61	18				
0000000140034F44	SkmilnitializeHypervisorPatching	17	70	24				
000000014003505C	SkmiCaptureDriverIdentity	8	38	10				
00000001400350E8	SkmilsPatchPagePresent	8	29	10				
0000000140035148	SkmiCheckUndoPagesPresent	19	61	27				
000000014003520C	SkmiExpandHpat	24	197	34				
0000000140035534	SkmiFreeDummyPatchMapping	9	56	13				
0000000140035618	SkmiCreateDummyPatchMapping	67	301	103				
0000000140035A50	SkmiPrepareNormalDummyPatchContext	18	165	26				
0000000140035CE8	SkmiPrepareGuardHotPatch	10	65	15				
0000000140035DDC	SkmiPreApplyHotPatch	16	112	22				
0000000140035FA4	SkmiApplyPatchesToPatchImage	4	39	4				
000000014003604C	SkmiApplyHotPatch	88	419	142				
0000000140036640	SkmiRevertHotPatch	14	110	19				
0000000140036800	SkmiPrepareHotPatchBaseImage	11	55	15				
00000001400368C4	SkmiMatchPatchImage	10	32	14				
0000000140036934	SkmiFreeSecureImagePages	10	50	14				
0000000140036A04	SkmiCopyPatchImage	25	174	38				
0000000140036CB4	SkmiCopySecurePatchImage	8	71	11				
0000000140036DCC	SkmiUnloadPatchImage	12	60	16				
0000000140036EB8	SkmiPrepareSecurePatchContext	18	95	25				
000000014003703C	SkmiCleanupSecurePatchContext	6	31	8				
00000001400370BC	SkmiPatchSecureImage	20	138	29				
00000001400372D4	SkmiApplyHypervisorHotPatch	28	146	42				
00000001400374EC	SkmiApplySecureHotPatch	12	54	17				
00000001400376F8	SkmmDetermineHotPatchUndoTableSize	16	62	24				
00000001400377FC	SkmmObtainHotPatchUndoTable	34	158	53				
0000000140037A5C	SkmiPrepareDriverPatchContext	6	67	8				
0000000140037B60	SkmmApplyHotPatch	74	325	118				
00000001400380E8	SkmmRevertHotPatch	19	77	29				
00000001400404A0	SkpgEnumerateSections	22	86	33				
0000000140042AF0	SkpgVerifyKernelVaProtectionExtent	19	87	27				
000000014004380C	SkpgCompareExtensionTableEntries	3	12	2				
0000000140043834	SkpgCompareVas	3	10	2				
0000000140043854	SkpgLocatelmageExtents	23	133	35				
0000000140043A3C	SkpgHotpatchApply	8	32	11				
2000000110010100	and the state of t							

Trusted Boot

- SMX support was added to the hypervisor in RS4
 - For example, new MSR, 0x40000116
- According to the documentation:
 - SMX registers are memory mapped to 0xFED20000 FED3FFFF
 - TPM registers are mapped to 0xFED40000 0xFED4FFFF
 - These regions are listed as allocated resources in the Device Manager application, so the system is finding them
 - https://software.intel.com/en-us/forums/virtualization-softwaredevelopment/topic/297709

```
💶 🚄 🖼
FFFFFAEA088A5910 mov
                        rcx, [rbx+2380h]
                        r9, HalpTxtProcStartup
FFFFFAEA088A5917 lea
                        r8d. 6
FFFFFAEA088A591E mov
FFFFFAEA088A5924 mov
                        edx, esi
                        HalMapIoSpaceWithStatus
FFFFFAEA088A5926 call
FFFFFAEA088A592B movzx
                        edi, ax
FFFFFAEA088A592E test
                        ax, ax
                        return_error
FFFFFAEA088A5931 jnz
```



Hyper-V ASLR

ASLR

- Usually we need to gain relative/arbitrary read primitive
- Fixed addresses are very helpful...
 - PTE base (and then... Anniversary update!)
 - HAL's heap (and then... Creators update!)
 - sharedpage
 - Now we even have a hypervisor sharedpage: ntdll!RtlphypervisorSharedUserVa
 - Not fixed, but can be predicted...
 - etc...
- Some pages in hv are not fully randomized...
 - PTE base (0xffffff7f80000000)
 - hv_lowstub (0x1000 / 0x2000)
 - self-referenced PTE entry 0x1fe (0xffffff7fbfdfe000)
 - And I keep finding stuff...

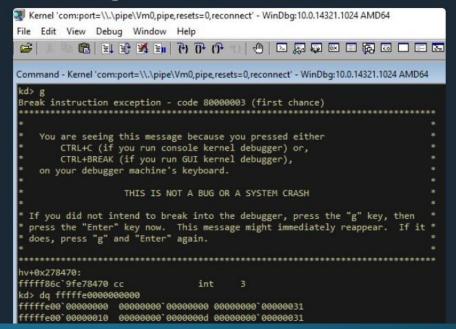


In continuation to my "fixed addresses in hv" tweet: time to talk about new structure, which I saw at address 0x40000 in few boots in a row. Details will come soon!

```
on your acoupper macritine a regularia.
                  THIS IS NOT A BUG OR A SYSTEM CRASH
 If you did not intend to break into the debugger, press the "g"
 press the "Enter" key now. This message might immediately reapp
 does, press "g" and "Enter" again.
hv+0x27bfe0:
fffffb06`aa47bfe0 cc
0: kd> dd 40000
00000000`00040000 00380001 00000000 00000009 00000000
000000000`00040010
00000000`00040020
00000000`00040030 000411e0 00000000 01080007 00000000
00000000 00040040
                ffe11290 ffe1129c ffe1120c ffe19168
00000000`00040050
                 ffe19174 ffe19180 ffe1918c ffe21058
000000000`00040060
0: kd> g
```



beside to my new 0x40000 structure (twitter.com/AmarSaar/statu ...), you can count in 0xfffffe0000000000 as well (hardcoded in binary). Looks the same between boots (more details about it will be coming soon :))



lowstub

- The reason I wake up every morning
- Extremely interesting structure, contains lots of interesting stuff
 - Located in a predictable physical address
 - 16bit code stub
 - Self referenced in virtual memory
 - KPROCESSOR_STATE (cr3, rip, rsp, rcx)
- Has 2 responsibilities:
 - On boot time— bring up from real mode to protected mode, and then to long mode
 - When processor resumes from S2/S3 sleep states
- Its (physical) address can be predicted and we have +X memory there!
- Question: What bothers me about this?



My last tweets made me and @Liran_Alon wonder: we need the lowstub in 2 states:

- in boot.
- resuming from S2/S3 sleep states.

So... why not to destroy it afterwards, and create it before each sleep? @epakskape

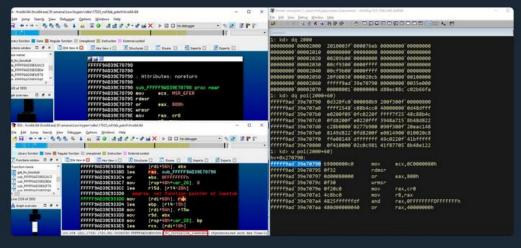
6:54 AM - 8 Mar 2018

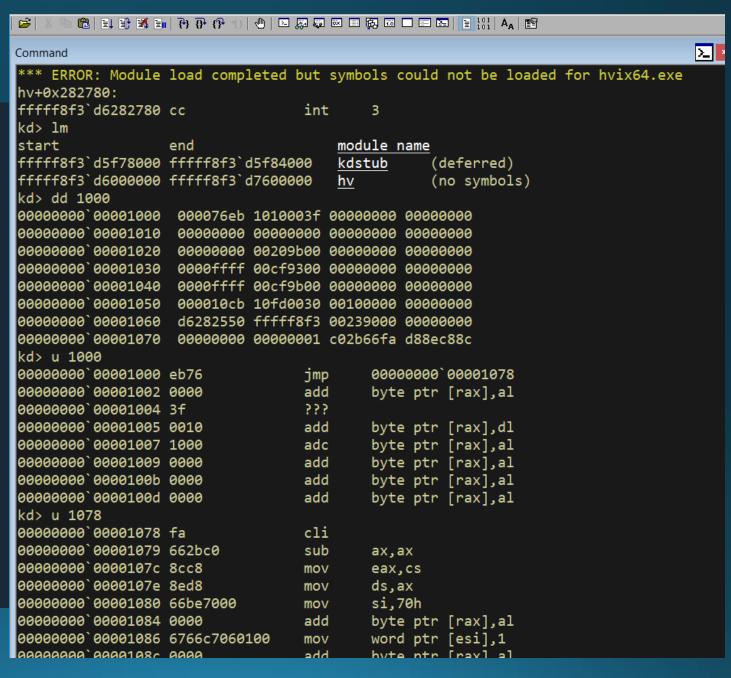






Little bonus to last tweet - the lowstub has a pointer to code, so if you know its' addr (and it's pretty much fixed), you can break randomization over hv base (credit goes to @aionescu and his talk about type-c attacks and lowstub internals)





The lowstub MSR

- Apparently, in RS4 we got a new MSR: 0x40000200
- This MSR simply reads (not writes) the address of hv_lowstub
 - Accessible from the root partition

```
case 0x40000200:
    v_partitionPermissionOK = *(__readgsqword(CurrentPartition) + offsetof(partition_s, partitionPrivilegeFlags)) &
    v_retval = HV_STATUS_ACCESS_DENIED;
    if ( v_partitionPermissionOK )
        v_retval = HV_STATUS_SUCCESS;
    if ( v_partitionPermissionOK )
        *pOut = hv_lowstub;
```

```
Break instruction exception - code 80000003 (first chance)
You are seeing this message because you pressed either
      CTRL+C (if you run console kernel debugger) or,
      CTRL+BREAK (if you run GUI kernel debugger),
   on your debugger machine's keyboard.
                THIS IS NOT A BUG OR A SYSTEM CRASH
 If you did not intend to break into the debugger, press the "g" key, then
 press the "Enter" key now. This message might immediately reappear. If it *
 does, press "g" and "Enter" again.
**************************
nt!DbgBreakPointWithStatus:
ffffff802 e1fb8e20 cc
                            int
kd> rdmsr 40000000+200
msr[40000200] = 00000000 00001000
kd> g
*BUSY* Debuggee is running...
```

```
Break instruction exception - code 80000003 (first chance)
You are seeing this message because you pressed either
      CTRL+C (if you run console kernel debugger) or,
      CTRL+BREAK (if you run GUI kernel debugger),
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 If you did not intend to break into the debugger, press the "g" key, then
 press the "Enter" key now. This message might immediately reappear. If it *
 does, press "g" and "Enter" again.
*** ERROR: Module load completed but symbols could not be loaded for hvix64.exe
hv+0x282780:
fffffbe2`56482780 cc
                            int
kd> dd 1000
 0000000`00001000  000076eb 1010003f 00000000 00000000
  300000`00001020   00000000 00209b00 00000000 00000000
  000000`00001030  0000ffff 00cf9300 00000000 00000000
  00000`00001040 0000ffff 00cf9b00 00000000 00000000
  000000`00001050 000010cb 10fd0030 00100000 00000000
00000000`00001060 56482550 fffffbe2 00239000 00000000
00000000`00001070 00000000 00000001 c02b66fa d88ec88c
kd> g
*BUSY* Debuggee is running...
```

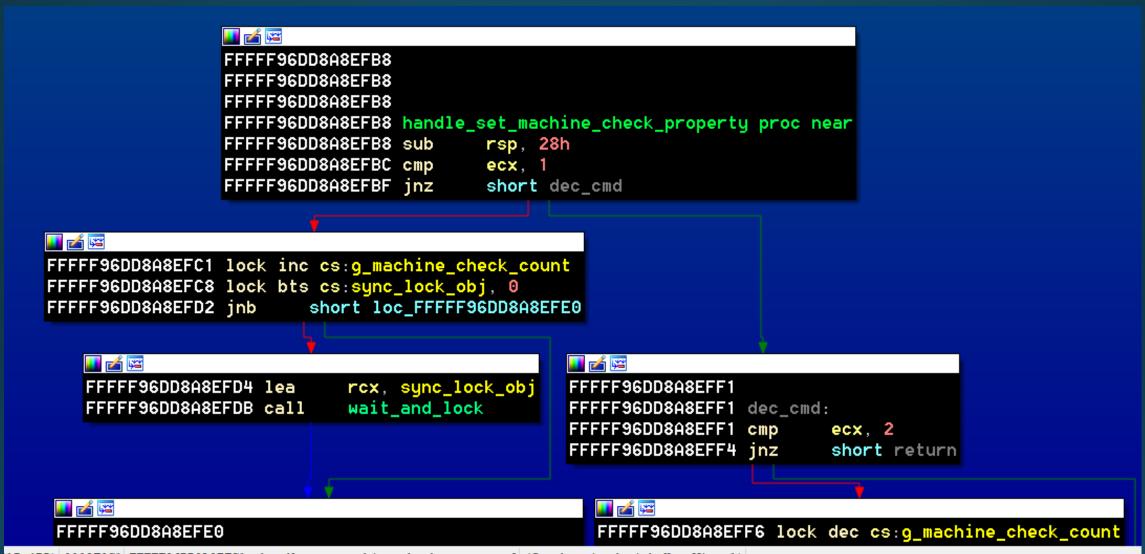
So... let's talk about NEW bugs

Bootlib

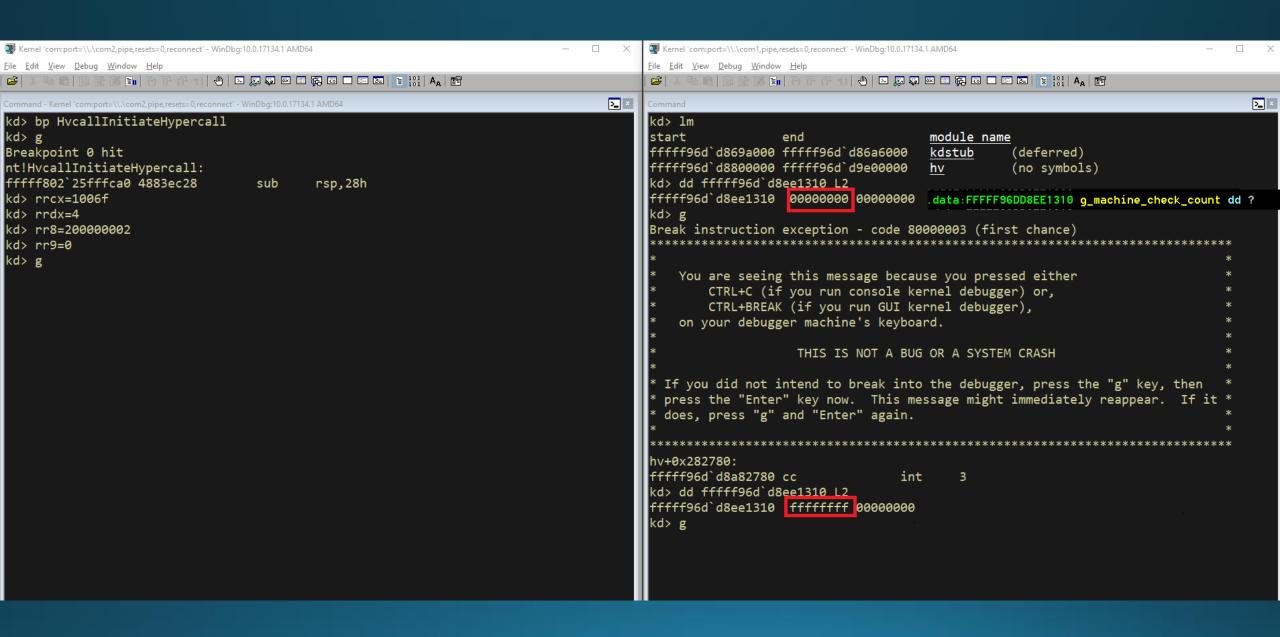
- Bootlib.dll, statically linked to winload, ci, skci, etc...
- Contains many interesting pieces of code
 - Including DER parsing
- I found 2 bugs in MinAsn1StringToOid
 - One is wrong return value
 - Second is off-by-one on a static buffer passed from the caller
- None of them is a security issue, but MS fixed them after all @
 - https://twitter.com/AmarSaar/status/976797760308633602

hv!handle_machine_check_property

- Really cute bug, accessible from a trivial hypercall flow
 - HvCallSetSystemProperty
- Integer overflow/underflow of global counts vars
- Can't crash, so it isn't security issue
- Should be fixed in the next version ©



17,475) 0008F3C1 FFFFF96DD8A8EFC1: handle_set_machine_check_property+9 (Synchronized with Hex View-1)



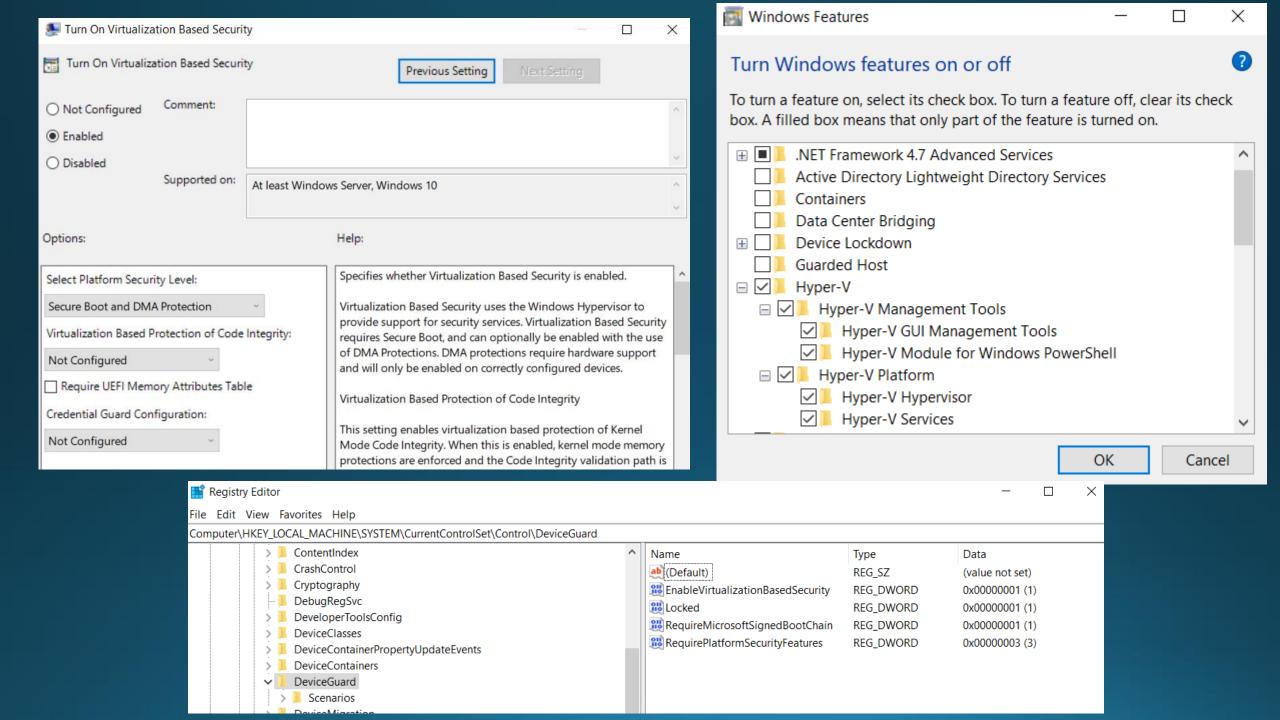
Segfault in HV!

Got fixed in build 17711 (RS5)

```
System Uptime: not available
Access violation - code c0000005 (!!! second chance !!!)
*** ERROR: Module load completed but symbols could not be loaded for hvix64.exe
hv+0x216837:
fffff96f 30616837 48396b10
                                           qword ptr [rbx+10h],rbp
kd> dq @rdi
fffff96f`30ae2000
                   fffff96f 30ae2000 00000000 000000000
fffff96f`30ae2010
                   00000020 d3abc4b0 00000000 204f33cd
fffff96f`30ae2020
                   00000000 00000000 00fbf6f1 ece7e2dd
fffff96f`30ae2030
                   00000000 9b081c40 ffffe800 0042b000
fffff96f`30ae2040
                   ffffe800`0042b000 00000000`00000021
fffff96f`30ae2050
                   00000000 25fbd736 00000000 00000000
fffff96f 30ae2060
                   00000000 00020000 00000000 00000001
fffff96f`30ae2070
                   00000000,00000000 00000000,00000000
kd> u @rip
hv+0x216837:
                                           qword ptr [rbx+10h],rbp
fffff96f`30616837 48396b10
                                          hv+0x216888 (fffff96f 30616888)
fffff96f`3061683b 774b
                                   ja
fffff96f`3061683d 0f1f00
                                           dword ptr [rax]
                                  nop
                                           rax, qword ptr [rbx]
fffff96f`30616840 488b03
                                  mov
                                           gword ptr [rax+8],rbx
fffff96f`30616843 48395808
                                   cmp
fffff96f`30616847 0f8515010000
                                           hv+0x216962 (fffff96f~30616962)
                                   ine
fffff96f`3061684d 488b4b08
                                           rcx, qword ptr [rbx+8]
                                  mov
                                           gword ptr [rcx],rbx
fffff96f`30616851 483919
                                  cmp
kd> u @rip - a
hv+0x21682d:
                                           rbx, qword ptr [rdi+14690h]
fffff96f`3061682d 488b9f90460100
                                           r15d, r15d
fffff96f`30616834 4533ff
                                   xor
                                           gword ptr [rbx+10h],rbp
fffff96f`30616837 48396b10
                                   cmp
                                  ja
                                           hv+0x216888 (fffff96f~30616888)
fffff96f`3061683b 774b
                                           dword ptr [rax]
fffff96f`3061683d 0f1f00
                                  nop
                                           rax, qword ptr [rbx]
fffff96f`30616840 488b03
                                  mov
                                           gword ptr [rax+8],rbx
fffff96f`30616843 48395808
                                   cmp
                                           hv+0x216962 (fffff96f 30616962)
fffff96f 30616847 0f8515010000
kd> dq @rdi+14690
ffftt96t 30at6690
                   00000100`00003c40 Fffff96f`30af6740
fffff96f`30af66a0
                   00000000 00000001 00000000 25fb5a58
fffff96f`30af66b0
                   00000000 25fb7ec2 00000000 25fb7ec2
fffff96f`30af66c0
                   fffff96f 30af6680 fffff96f 30af6780
fffff96f`30af66d0
                   ffffffff fffffff 00000000 000000000
fffff96f`30af66e0
                   00000000 00000001 fffff96f 306b5370
fffff96f`30af66f0
                   00000000,00000000 00000000,00000000
fffff96f`30af6700 fffff96f`30ab5b80 fffff96f`30af6680
kd> rrbx
rbx=0000010000003c40
```

Public service announcement

- Enable VSM, HVCI, Device Guard and Credential Guard
 - it really changes the rules of the game
 - And... you get kCFG for free ◎
- Supposed to be on by default in RS4+ clean installs



References

- http://www.alex-ionescu.com/syscan2015.pdf
- http://www.alex-ionescu.com/blackhat2015.pdf
- https://www.blackhat.com/docs/us-16/materials/us-16-Wojtczuk-Analysis-Of-The-Attack-Surface-Of-Windows-10-Virtualization-Based-Security-wp.pdf
- http://alex-ionescu.com/Publications/OPCDE/octagon.pdf
- https://bugs.chromium.org/p/projectzero/issues/detail?id=688

close()

- Thanks to the great folks @BluehatIL
- Follow me on twitter for lowlevel internals
 - @AmarSaar