

## EXAMPLE 2.

How many levels of indexing would be needed for the file from example 1?

- Assume: ORDERED FILE WITH PRIMARY INDEX.

$$b = \text{number of data blocks} = 3,000$$

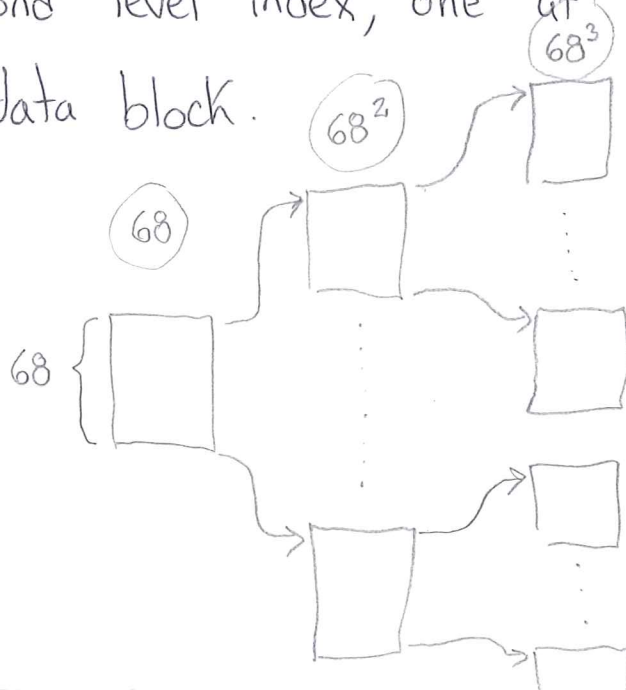
$$\text{number of index blocks} = 45$$

$$\text{BFRI} = \text{blocking factor for index block} = 68$$

We will need to add 45 second-level index entries to access 45 primary-level index blocks.

45 entries fits in one block.

The total numbers of disk blocks read is 3:  
one at second level index, one at primary level index  
and one data block.



$$\text{level 1} = 68^1 = 68$$

$$\text{level 2} = 68^2 = 4,624$$

$$\text{level 3} = 68^3 = 314,432$$

### EXAMPLE 3.

Calculate the order  $p$  of a  $B^+$  tree given:

$BS$  = block size = 512

$V$  = index size = 9  $\leftarrow$  this is the size of the Key

$P_r$  = pointer size = 7

$P$  = index tree pointer size = 6  $\leftarrow$  this is the size of the tree pointer  
 $\nwarrow$  capital  $p$

"How many tree pointers can we store?"



$$[p * 6] + [(p-1) * 9] \leq 512$$

$$6p + 9p - 9 \leq 512$$

$$15p - 9 \leq 512$$

$$15p \leq 512 + 9$$

$$15p \leq 521$$

$$p \leq 521/15$$

$$p \leq 34.7 \quad \therefore p = 34$$

Approximately how many levels in this  $B^+$  tree would we need to store 1 million records?

$$\log_{34}(1,000,000) \approx 4$$

$$34^4 = 1,336,336$$

• 512 byte/block , 64 byte record

What is the blocking factor of the disk?

$$BFR = \frac{512}{64} = 8$$

∴ 8 records per block.

• 512 byte/block , 65 byte record

$$BFR = 7.87... \rightarrow \begin{array}{l} \text{UNSPANNED} \\ \text{SPANNED} \end{array} = \begin{array}{l} 7 \\ 7.87 \end{array} \begin{array}{l} \text{records per block} \\ \text{records per block} \end{array}$$

## B+ Trees (multi-level index)

- Efficient direct access to records.

⇒ Multi-level index, so we can do a multi-level search

- Efficient ordered access through all records.

⇒ Bottom level is chain together as a linked list

- Efficient overall balanced index structure.

⇒ Expand tree upward, add values to blocks or split blocks

⇒ It is wide, we have a large BFR

⇒ It is shallow, we donot need many levels to represent it

- Efficient insertion and deletion.

⇒ It is proportional to the height of the tree.