

DIJKSTRA'S ALGORITHM

Dijkstra's algorithm is based on a cost/value based system where each node is assigned a minimum by a set of rules (known as relaxation).

It is represented using graphs and tables.

It is applicable to both direct and indirect paths.

Algorithm is as follows-

$$D(u) + c(u, v) < d(v)$$

Where u, v are nodes.

$D(u)$ represents the distance from the initial node to node ' u '

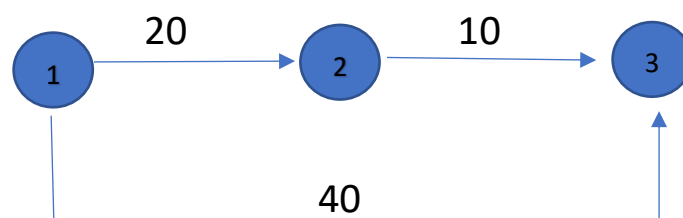
$C(u, v)$ represents the cost to travel to final node v from initial node u .

RULES OF THE ALGORITHM

- Each node's value is initially initialized to infinity.
- Later when paths between the nodes is established the value (infinity) is reduced to a minimum value.
- Each set of nodes has an initial start point is initialised to zero.

Let's understand this concept using two examples

1.



Here, $D(u) + c(u, v) < d(v)$ – eqn 1

$D(v) = d(u) + c(u, v)$ – eqn 2

Node 1 is the initial node so we set its value to zero, node 2, 3 are set to infinity.

Putting values in eqn 1 for moving from 1 to 2

$$0+20 < \text{infinity}$$

As this is true we can set the value of 2 as 20 instead of infinity.

Putting values in eqn 1 for moving from 2 to 3

$$0+40 < \text{infinity}$$

As this is true we can set the value of 3 as 40 instead of infinity.

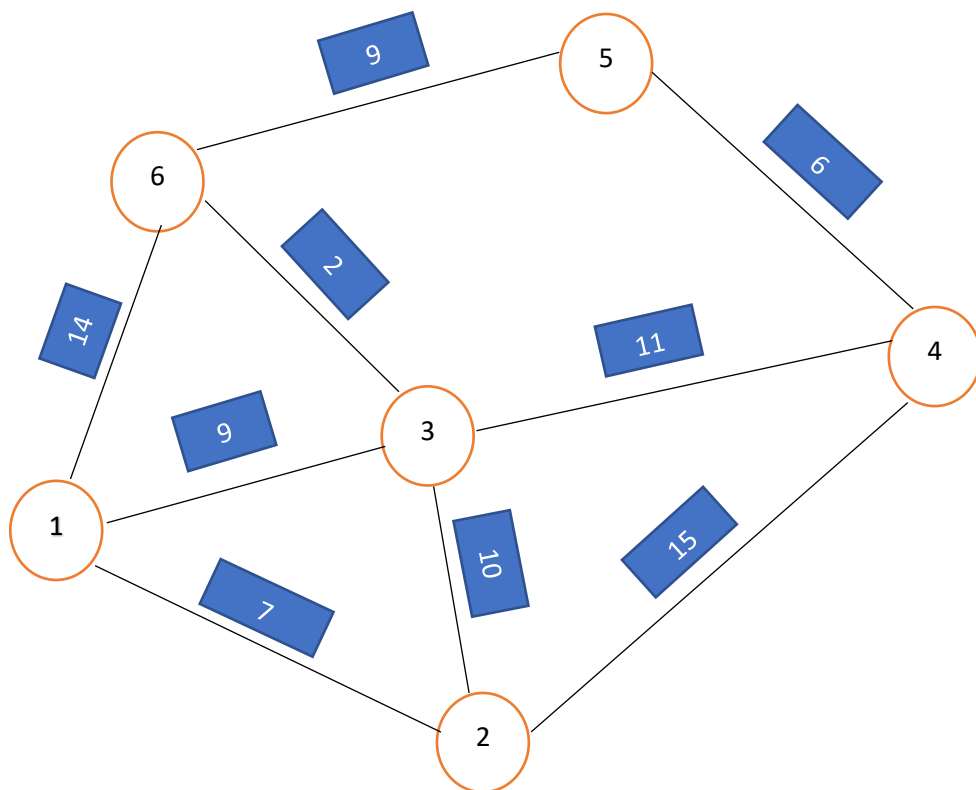
Now taking distance between 1 and 3 which is 40, and substituting it in eqn 1, we get

$$=20+10 < 40$$

$$=30 < 40 \text{ which is true}$$

So we can once again we can reduce (relaxation) the value of node 3 as 30.

2.



This is an example of indirect path approach which can be solved using the help of a table

Source	Destinations				
1	2	3	4	5	6
	infinity	infinity	infinity	infinity	infinity
1,2	7	9	infinity	infinity	14
1,2,3	7	9	22	infinity	14
1,2,3,6	7	9	20	infinity	11
1,2,3,6,4	7	9	20	20	11
1,2,3,6,4,5	7	9	20	20	11

Here, all nodes except 1 are set to infinity
We see that distance of 2 from 1 is 7,

Distance of 3 from 1 is 9,
Distance of 6 from 1 is 14.

Direct path

Distance of 4 from 1 via 2 is $7+15=22$,
Distance of 6 from 1 via 3 is $9+2=11 < 14$,
So we now update the value of node 6 as 11.

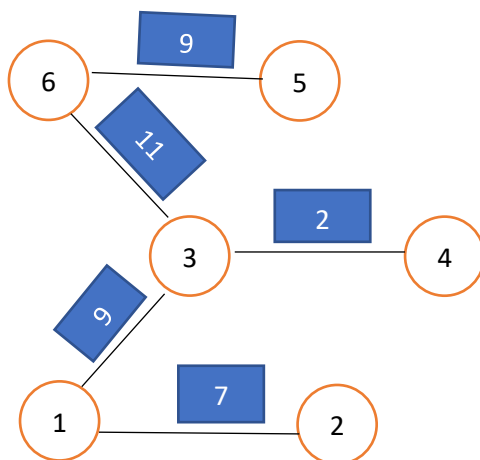
Distance of 4 from 1 via 3 is $9+11=20$.

Note there is no direct or indirect path from 1 to 5 as of now.

Indirect path

Distance of 5 from 1 via updated value of 6 (ie, 11) is $11+9=20$.

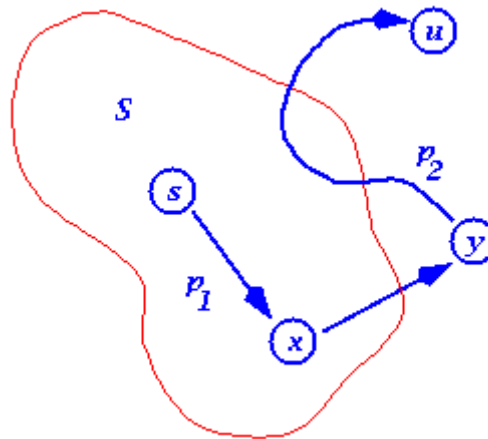
Here is the updated schematic diagram after applying the dijkstra's algorithm.



OTHER APPLICATIONS OF THE ALGORITHM

- Used for navigation purpose in google maps.
- Used in telephone networking.

PROOF OF DIJSTRA'S ALGORITHM



Proof using contradiction method:-

Let 's' be the source.

P_1 is the path from s to x ;

P_2 is the path from y to u .

Let us denote the shortest distance between any points as $\text{del}()$, and the distance from the initial node to a particular node say 'a', is given by $d(a)$.

Let's assume 'u' to be the initial node then,

1. 'u' can't be 's', as $d[s]=0$.
2. There also should be a path from s to u , otherwise $d[u]=\text{infinity}$ (by the property of the algorithm).
3. If there is a path then there must also be a shortest path.

When 'x' is inserted into 's' then $d[x]=\text{del}(x,s)$ {since we assume that u is the first node for which this is not true}

Now (x,y) is relaxed so,

$$D[y]=\text{del}(s,y)\leq\text{del}(s,u)\leq d[u]$$

But since u is the initial node so $d[u]\leq d[y]$

Thus the two inequalities must be equalities

$$D[y]=\text{del}(s,y)=\text{del}(s,u)=d[u]$$

So $d[u] = \text{del}(s, u)$ is contradicting our assumption

Therefore when ant 'u' is added we get, $d[u] = \text{del}(S, u)$