

## Question 1

Given conditions:

1 cpu, 1 disk  $T_{\text{monitor}} = 60 \text{ min}$

# completed job 1267

# cpu accesses 2178

# disk accesses 2412

cpu busy time 2929 s

Disk busy time 2765 s

(a) Define CPU as device 1, Disk as device 2

Based on Service Demand Law  $D(j) = \frac{U(j)}{X(0)}$ ,  $U(j)$  is the utilization of device  $j$

$X(0)$  is the throughput of the system

$U(j) = \frac{B(j)}{T}$ ,  $B(j)$  is the busy time of device  $j$

output rate  $X(0) = \frac{C(0)}{T}$ ,  $C(0)$  is the number of requests completed for the system

Therefore, the service demand of a job at device  $j$   $D(j) = \frac{B(j)}{C(0)}$

⇒ The service demand at CPU is  $D(1) = \frac{2929}{1267} = 2.31 \text{ s}$

The service demand at Disk is  $D(2) = \frac{2765}{1267} = 2.18 \text{ s}$

(b) Utilisation limit:  $U(j) \leq 1 \Rightarrow X(0) \leq \frac{1}{D(j)}$

$$\Rightarrow X(0) \leq \min \frac{1}{D_j}, X_0 \leq \frac{1}{\max D_j}$$

$$R \geq \sum_{j=1}^K D_j, R = \frac{M}{X(0)} - Z, X(0) = \frac{M}{Z+R} \leq \frac{M}{Z + \sum_{j=1}^K D_j}$$

$M$ : Interface Clients  
 $Z$ : mean thinking time  
 $R$ : mean response time of the computer system  
 $X(0)$ : throughput

$$\Rightarrow X(0) \leq \min \left[ \frac{1}{\max D_i}, \frac{N}{\sum_{i=1}^K D_i} \right]$$

thinking time = 14 s Users = 20  $\Rightarrow Z=14, N=20$

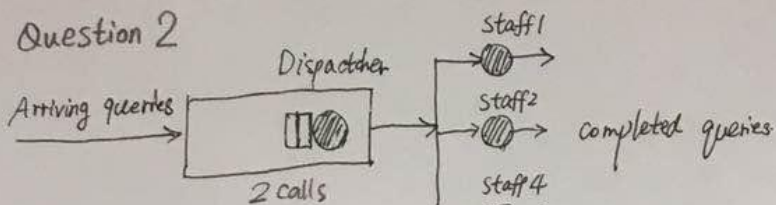
Because  $D(1) = 2.31 \text{ s}$

$D(2) = 2.18 \text{ s}$

$D_{\max} = 2.31 \text{ s} = D_{\text{cpu}} = D(1)$

$$X(0) \leq \min \left[ \frac{1}{2.31}, \frac{20}{2.31 + 2.18 + 14} \right] = \min(0.43, 1.082) = 0.43$$

## Question 2

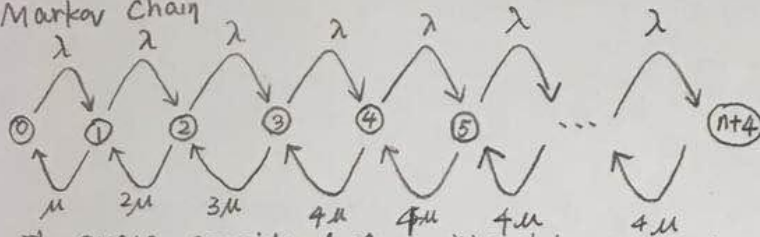


Solution: From the conditions provided

Voice-over-IP records rate  $\lambda = 15$  queries  $\mu = 3$  queries

(a)

Markov Chain



The queue consists of  $n$  waiting slots, so the format is  $M/M/4/4+n$  (Kendall's notation)

State 0: 0 calls in system (4 staffs are idle)

State 1: 1 call in system (3 staffs are idle, 1 staff has jobs)

State 2: 2 calls ----- (2 ----- idle, 2 -----)

State 3: 3 calls ----- (1 -----, 3 -----)

State 4: 4 calls ----- (All staffs has jobs)

State 5: 5 calls ----- (-----, There is 1 staff in waiting slot)

State  $n+4$ :  $n+4$  calls ----- (-----  $n$  -----)

The transition probabilities (with a small  $\delta$  time interval)

$$P[0 \rightarrow 1] = \lambda \delta$$

$$P[0 \rightarrow 2] = 0$$

$$P[1 \rightarrow 0] = \mu \delta$$

$$P[2 \rightarrow 0] = 0$$

$$P[2 \rightarrow 1] = 2\mu \delta$$

$\vdots$

$$P[k \rightarrow k+1] = \lambda \delta \quad (0 \leq k \leq n+3)$$

$$P[k \rightarrow k-1] = \begin{cases} k\mu \delta & (1 \leq k \leq 4) \\ 4\mu \delta & (5 \leq k \leq n+4) \end{cases}$$

(b) Balance equation for the continuous-time Markov Chain

State 0:  $\lambda p_0 = \mu p_1$

State 1:  $\lambda p_0 + 2\mu p_2 = (\lambda + \mu) p_1$

State 2:  $\lambda p_1 + 3\mu p_3 = (\lambda + 2\mu) p_2$

State 3:  $\lambda p_2 + 4\mu p_4 = (\lambda + 3\mu) p_3$

State 4:  $\lambda p_3 + 4\mu p_5 = (\lambda + 4\mu) p_4$

$\vdots$

State  $k$ :  $\lambda p_{k-1} + 4\mu p_{k+1} = (\lambda + 4\mu) p_k \quad (5 \leq k \leq n+3)$

$\vdots$

State  $n+3$ :  $\lambda p_{n+2} + 4\mu p_{n+4} = (\lambda + 4\mu) p_{n+3}$

State  $n+4$ :  $\lambda p_{n+3} = 4\mu p_{n+4} \quad (\lambda = 15, \mu = 3)$



(c) We define  $\lambda_0 = \frac{\lambda}{\mu}$

When  $1 \leq n \leq 4$   $P_n = \left(\frac{\lambda}{\mu}\right)^n \cdot P_0 \cdot \frac{1}{n!} = \lambda_0 \cdot P_0 \cdot \frac{1}{n!}$

$n > 4$   $P_n = \frac{\lambda^{n+4} \cdot P_0}{4^n \cdot 4! \cdot \mu^{n+4}} = (\lambda_0)^{n+4} \cdot \frac{P_0}{4^n \cdot 4!}$

(d) (i) Based on 6 query, we use  $n=6$   $\lambda=15$ ,  $\mu=3$  to calculate  $P_n$  by matlab.

$P_6 = 0.2935$

(ii)  $N_{avg} = X \cdot R_{avg}$  (Little's Law)   
 $n=6$   $P_6 = 0.2935$   $X = \text{rate} \times (1 - P_{reject})$    
 $\Rightarrow X = (1 - P_6) \times 15 = 10.605$  transactions/h

$R_{avg} = \frac{N_{avg}}{X} = \frac{4.35}{10.605} \approx 0.4$  hours

Mean waiting time = Mean response time - Mean service time

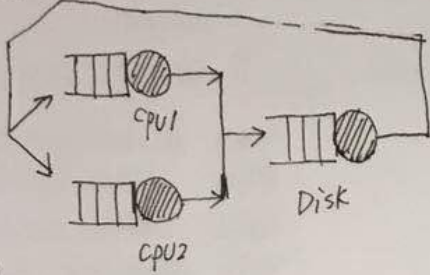
(e) We calculate these results from matlab programs  $= 0.41 - \frac{1}{4 \cdot \mu} \approx 0.32$  hour

adding  $\begin{cases} 5 \\ 10 \\ 15 \\ 20 \end{cases}$  slots  $\Rightarrow n = \begin{cases} 5+2+4=11 \\ 10+2+4=16 \\ 15+2+4=21 \\ 20+2+4=26 \end{cases}$  requires.

We run the results as  $\begin{cases} P_{11} = 0.2233 & P_{21} = 0.2071 \\ P_{16} = 0.2023 & P_{26} = 0.2007 \end{cases}$

(f) We can make a conclusion that the limit capability for system to process is adding 10 more waiting slots. Because  $\underbrace{\text{staffs}}_{\text{numbers of}}$  are only 4. When adding more than 10 slots, it has little influence on the system.  $\underbrace{\text{Under this limitation,}}_{\text{Only when we enforce work efficiency of staff can help enhance the system efficiency.}}$

### Question 3



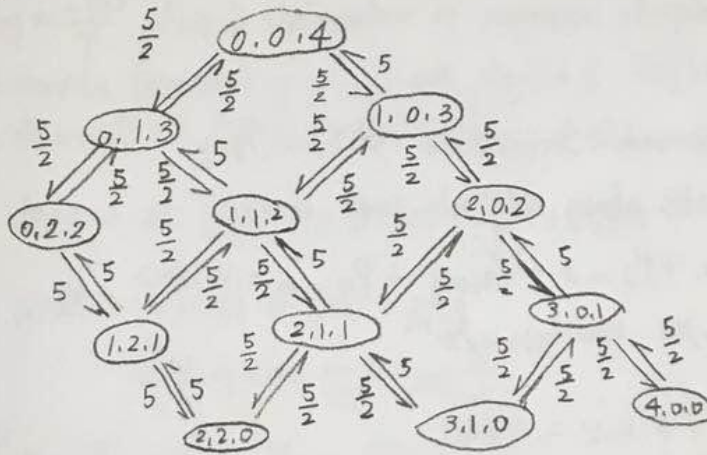
(a)

(1) Obviously We have 12 states. The list of states is  
 $(0,0,4), (1,0,3), (0,1,3), (2,0,2), (1,1,2), (0,2,2), (3,0,1), (2,1,1), (1,2,1), (4,0,0), (3,1,0), (2,2,0)$

(2) From the condition we get that mean processing time  $T_{\text{disk}} = 0.2s$ ,  $T_{\text{cpu1}} = 0.2s$ ,  $T_{\text{cpu2}} = 0.4s$

Thus, the transition rate of CPU1, CPU2, Disk is

$$R_{\text{cpu1}} = \frac{1}{0.2} = 5 \text{ transactions/s}, R_{\text{cpu2}} = \frac{1}{0.4} = \frac{5}{2} \text{ transactions/s}, R_{\text{disk}} = \frac{1}{0.2} = 5 \text{ transactions/s}$$



(b) The balance equations are based on The in & out of Nodes in Markov chain

$$\begin{aligned} & 5P_{0,0,4} - 5P_{1,0,3} - 2.5P_{0,1,3} + 0P_{2,0,2} + 0P_{1,1,2} + 0P_{0,2,2} + 0P_{3,0,1} + 0P_{2,1,1} + 0P_{1,2,1} + 0P_{4,0,0} + 0P_{3,1,0} + 0P_{2,2,0} + \\ & -\frac{5}{2}P_{0,0,4} + 10P_{1,0,3} + 0P_{0,1,3} - 5P_{2,0,2} - \frac{5}{2}P_{1,1,2} + 0P_{0,2,2} + 0P_{3,0,1} + 0P_{2,1,1} + 0P_{1,2,1} + 0P_{4,0,0} + 0P_{3,1,0} + 0P_{2,2,0} + \\ & -\frac{5}{2}P_{0,0,4} + 0P_{1,0,3} + \frac{15}{2}P_{0,1,3} + 0P_{2,0,2} - 5P_{1,1,2} - \frac{5}{2}P_{0,2,2} + 0P_{3,0,1} + 0P_{2,1,1} + 0P_{1,2,1} + 0P_{4,0,0} + 0P_{3,1,0} + 0P_{2,2,0} + \\ & 0P_{0,0,4} - \frac{5}{2}P_{1,0,3} + 0P_{0,1,3} + 0P_{2,0,2} + 0P_{1,1,2} + 0P_{0,2,2} - 5P_{3,0,1} - \frac{5}{2}P_{2,1,1} + 0P_{1,2,1} + 0P_{4,0,0} + 0P_{3,1,0} + 0P_{2,2,0} + \\ & 0P_{0,0,4} - \frac{5}{2}P_{1,0,3} - \frac{5}{2}P_{0,1,3} + 0P_{2,0,2} + \frac{25}{2}P_{1,1,2} + 0P_{0,2,2} + 0P_{3,0,1} - 5P_{2,1,1} - \frac{5}{2}P_{1,2,1} + 0P_{4,0,0} + 0P_{3,1,0} + 0P_{2,2,0} + \\ & 0P_{0,0,4} + 0P_{1,0,3} - \frac{5}{2}P_{0,1,3} + 0P_{2,0,2} + 0P_{1,1,2} + \frac{15}{2}P_{0,2,2} + 0P_{3,0,1} + 0P_{2,1,1} - 5P_{1,2,1} + 0P_{4,0,0} + 0P_{3,1,0} + 0P_{2,2,0} + \\ & 0P_{0,0,4} + 0P_{1,0,3} + 0P_{0,1,3} - \frac{5}{2}P_{2,0,2} + 0P_{1,1,2} + 0P_{0,2,2} + 0P_{3,0,1} + 0P_{2,1,1} + 0P_{1,2,1} - 5P_{4,0,0} - \frac{5}{2}P_{3,1,0} + 0P_{2,2,0} + \\ & 0P_{0,0,4} + 0P_{1,0,3} + 0P_{0,1,3} - \frac{5}{2}P_{2,0,2} - \frac{5}{2}P_{1,1,2} + 0P_{0,2,2} + 0P_{3,0,1} + \frac{25}{2}P_{2,1,1} + 0P_{1,2,1} + 0P_{4,0,0} + 0P_{3,1,0} - 5P_{2,2,0} + \\ & 0P_{0,0,4} + 0P_{1,0,3} + 0P_{0,1,3} + 0P_{2,0,2} + 0P_{1,1,2} + 0P_{0,2,2} - \frac{5}{2}P_{3,0,1} + 0P_{2,1,1} + 0P_{1,2,1} + 5P_{4,0,0} + 0P_{3,1,0} + 0P_{2,2,0} + \\ & 0P_{0,0,4} + 0P_{1,0,3} + 0P_{0,1,3} + 0P_{2,0,2} + 0P_{1,1,2} + 0P_{0,2,2} - \frac{5}{2}P_{3,0,1} - \frac{5}{2}P_{2,1,1} + 0P_{1,2,1} + 0P_{4,0,0} + \frac{15}{2}P_{3,1,0} + 0P_{2,2,0} + \\ & 0P_{0,0,4} + 0P_{1,0,3} + 0P_{0,1,3} + 0P_{2,0,2} + 0P_{1,1,2} + 0P_{0,2,2} + 0P_{3,0,1} - \frac{5}{2}P_{2,1,1} - 5P_{1,2,1} + 0P_{4,0,0} + 0P_{3,1,0} + \frac{15}{2}P_{2,2,0} + \\ & = 1 \end{aligned}$$



(c) We use Matlab to solve the question. The solution is

$$\begin{cases} P_{0,0,4} = 0.1711 \\ P_{1,0,3} = 0.0912 \\ P_{0,1,3} = 0.1598 \\ P_{2,0,2} = 0.0501 \\ P_{1,1,2} = 0.0935 \\ P_{0,2,2} = 0.1213 \\ P_{3,0,1} = 0.0259 \\ P_{2,1,1} = 0.0572 \\ P_{1,2,1} = 0.1021 \\ P_{4,0,0} = 0.0130 \\ P_{3,1,0} = 0.0277 \\ P_{2,2,0} = 0.0871 \end{cases}$$

(d) Throughput = Utilisation  $\times$  Service Rate  $\Rightarrow Th = U \times r$

We compute the rate where only the part works.

$$\text{Thus, } r_{cpu1} = P_{1,0,3} + P_{2,0,2} + P_{3,0,1} + P_{4,0,0} = 0.1802$$

$$Th_{cpu1(only)} = 5 \times 0.1802 = 0.901 \text{ transactions/s}$$

$$r_{cpu2} = P_{0,1,3} + P_{0,2,2} = 0.2811$$

$$Th_{cpu2(only)} = \frac{5}{2} \times 0.2811 = 0.7028 \text{ transactions/s}$$

$$Th_{cpu1+cpu2} = r_{cpu1+cpu2} \cdot U_{1+2} = U(P_{1,1,2} + P_{2,1,1} + P_{1,2,1} + P_{3,1,0} + P_{2,2,0}) = 0.368 \times 7.5 = 2.76$$

$$Th_{total} = 0.901 + 0.7028 + 2.76 = 4.364 \text{ transactions/s}$$

(e) Mean number of jobs:

$$N_{cpu1} = \bar{m} = \sum_{i=1}^n n_i p_i = (P_{1,2,1} + P_{1,0,3} + P_{1,1,2}) \times 1 + (P_{2,0,2} + P_{2,1,1} + P_{2,2,0}) \times 2 + (P_{3,0,1} + P_{3,1,0}) \times 3 + P_{4,0,0} \times 4 = 0.89$$

$$(f) \text{ Generally, } Th_{cpu1} = Th_{cpu1(only)} + Th_{total} U_1 = 0.901 + 0.368 \times 5 = 2.74 \text{ transactions/s}$$

Based on Little's Law

$$R = \frac{N}{X} = \frac{N_{cpu1}}{Th_{cpu1}} = \frac{0.89}{2.739} = 0.33 \text{ seconds}$$