

Persistent Homology

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Persistent Homology

Filtration

Definition (filtration)

A filtration of a simplicial complex \mathbb{K} is a nested sequence of complexes,

$$\emptyset = \mathbb{K}_{-1} \subset \mathbb{K}_0 \subset \mathbb{K}_1 \subset \cdots \subset \mathbb{K}_n = \mathbb{K}.$$

Example

Suppose \mathbb{K} is a simplicial complex, we sort all the simplices in a sequence

$$\sigma_1^0, \sigma_2^0, \dots, \sigma_{n_0}^0, \sigma_1^1, \sigma_2^1, \dots, \sigma_{n_1}^1, \sigma_1^2, \sigma_2^2, \dots, \sigma_{n_2}^2.$$

where σ_i^k is the i -th k -simplex in \mathbb{K} . Then we relabel all the simplices as

$$\sigma^0, \sigma^1, \sigma^2, \dots,$$

We define \mathbb{K}_i as the union of $\sigma^0, \sigma^1, \dots, \sigma^i$.

Homology

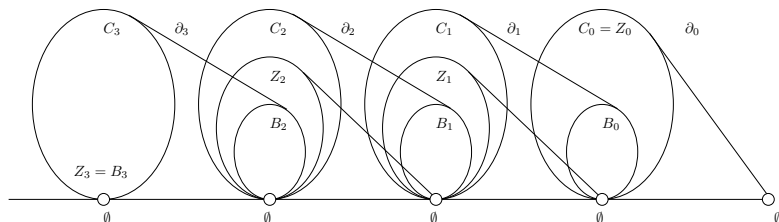


Figure: Chain, cycle, boundary groups and their images under the boundary operators.

$$H_k(\mathbb{K}, \mathbb{Z}_2) = \frac{\text{Ker} \partial_k}{\text{Im} \partial_{k+1}} = \frac{Z_k}{B_k}.$$

Persistent Homology

The inclusion map $f : \mathbb{K}_{i-1} \hookrightarrow \mathbb{K}_i$ defined by $f(x) = x$ induces a homomorphism $f_* : H_p(\mathbb{K}_{i-1}) \rightarrow H_p(\mathbb{K}_i)$. The nested sequence of complexes corresponds to a sequence of homology groups connected by the induced maps,

$$0 = H_p(\mathbb{K}_{-1}) \rightarrow H_p(\mathbb{K}_0) \rightarrow \cdots \rightarrow H_p(\mathbb{K}_n) = H_p(\mathbb{K})$$

Persistent homology studies how the homology groups change over the filtration.

Definition (positive simplex)

Given a filtration of \mathbb{K} , suppose $\mathbb{K}_i - \mathbb{K}_{i-1} = \sigma_i$, where σ_i is a $(k+1)$ -simplex. We call σ_i is **positive** if it belongs to a $(k+1)$ -cycle in \mathbb{K}_i and **negative** otherwise.

A positive simplex is also called a **generator**, a negative simplex a **killer**.

Definition (Betti Number)

Given a complex K , the i -th Betti number β_i is the rank of $H_i(K)$,

$$\beta_i = \text{Rank} H_i(K, \mathbb{Z}_2)$$

Suppose the number of positive k -simplexes is pos_k , and the number of negative k -simplexes is neg_k , then

$$\beta_k = \text{pos}_k - \text{neg}_{k+1}$$

Persistent Homology

Definition (Persistent Homology)

Define Z_k^l, B_k^l be the K -th cycle group and k -th boundary group respectively, of the l -complex K^l in a filtration. The p -persistent k -th homology group K^l is

$$H_k^{l,p} := \frac{Z_k^l}{B_k^{l+p} \cap Z_k^l}.$$

The p -persistent k -th Betti number $\beta_k^{l,p}$ of K^l is the rank of $H_k^{l,p}$.

Lemma

Consider the homomorphism $\eta_k^{l,p} : H_k^l \rightarrow H_k^{l+p}$, then

$$\text{img } \eta_k^{l,p} \cong H_k^{l,p}$$

Lemma

For each *positive* k -simplex σ^i , there exists a non-exact k -cycle c^i , c^i contains σ^i but no other positive k -simplices.

Proof.

Start with an arbitrary a k -cycle that contains σ^i and remove other positive k -simplices by adding their corresponding k -cycles. This method succeeds because each added cycle contains only one positive k -simplex by inductive assumption. \square

We use σ^i to represent c^i , and in turn the homologous class $[c^i] = c^i + B_k$.

$$\sigma^i \rightarrow c^i \rightarrow [c^i] = c^i + B_k. \quad \sigma^i \sim c^i$$

We add $[c^i]$ to the basis of $H_k(\mathbb{K}^i)$.

Generator

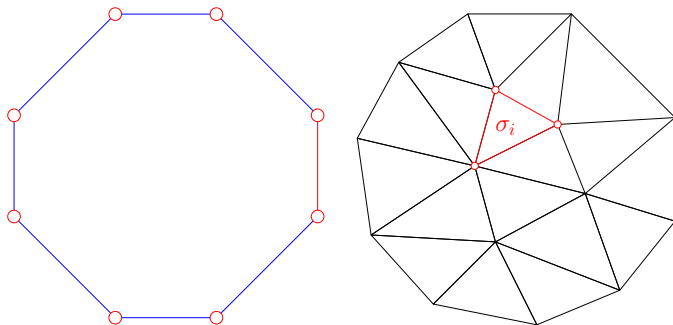


Figure: Generator, positive simplex.

For each **negative** $(k+1)$ -simplex σ^j , its boundary $d = \partial_{k+1}\sigma^j$ is a k -cycle, and can be represented as the linear combination of the basis of $H_k(\mathbb{K}_{j-1})$,

$$[d] = \sum_g [c^g], \{c^g\} \text{ basis } H_k(\mathbb{K}_{j-1}),$$

each $[c^g]$ is represented by a positive k -simplex σ^g , $g < j$, that is not yet paired. The collection of positive non-paired k -simplices is denoted as $\Gamma = \Gamma(d)$,

$$\Gamma(d) := \left\{ \sigma^g : [d] = \sum_g [c^g], \quad \sigma^g \sim c^g \right\}$$

Suppose the youngest positive simplex in $\Gamma(\partial_{k+1}\sigma^j)$ is σ^i , then we form the pair (σ^i, σ^j) , and remove $[c^i]$ from $H_k(\mathbb{K}_j)$.

$[c^i]$ is created by σ^i and killed by σ^j , the persistence life of the k -cycle $[c^i]$ is $j - i - 1$.

Example Filtration

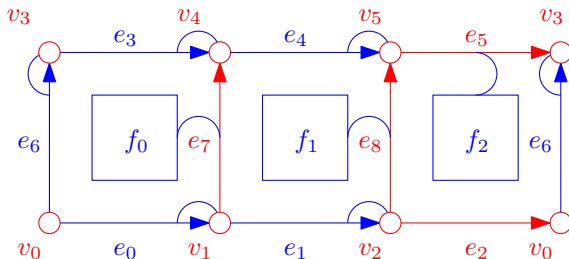


Figure: Generators and killers.

Filtration

$v_0, v_1, v_2, v_3, v_4, v_5, e_0, e_1, e_2, e_3, e_4, e_5, f_0, f_1, f_2$

Relabel them as

$\sigma^0, \sigma^1, \sigma^2, \sigma^3, \sigma^4, \sigma^5, \sigma^6, \sigma^7, \sigma^8, \sigma^9, \sigma^{10}, \sigma^{11}, \sigma^{12}, \sigma^{13}, \sigma^{14}, \sigma^{15}$

Example Generators

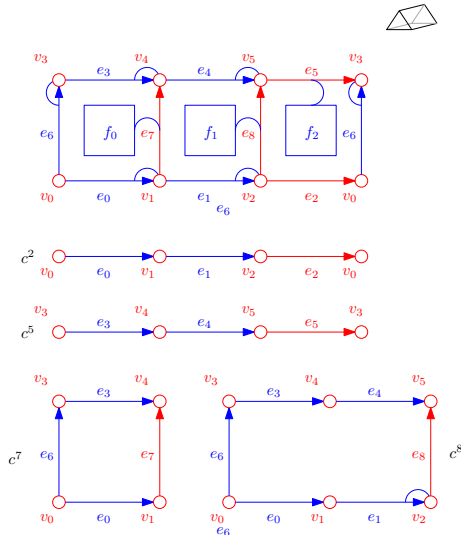


Figure: Generators.

Example Killers

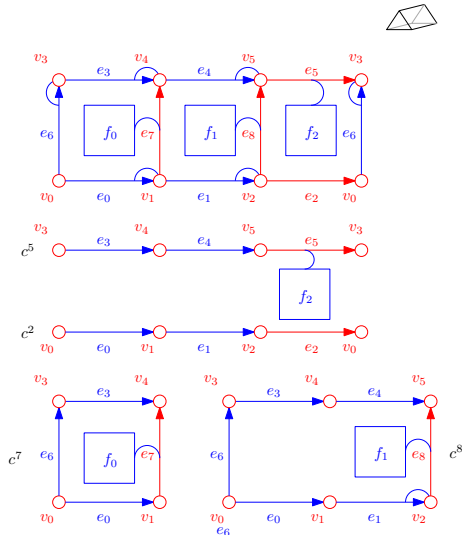


Figure: Killers.

Example Pairing

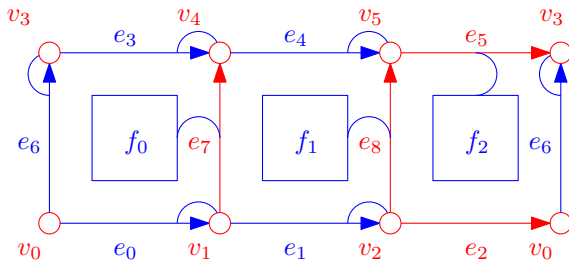


Figure: Generators and killers.

$$\begin{aligned}\partial_2 f_2 &= \mathbf{e}_2 + \mathbf{e}_5 + \mathbf{e}_6 + \mathbf{e}_8 = (\mathbf{e}_5 + 2\mathbf{e}_4 + 2\mathbf{e}_3) + (\mathbf{e}_2 + 2\mathbf{e}_1 + 2\mathbf{e}_0) + \mathbf{e}_6 + \mathbf{e}_8 \\ &= (\mathbf{e}_5 + \mathbf{e}_4 + \mathbf{e}_3) + (\mathbf{e}_2 + \mathbf{e}_1 + \mathbf{e}_0) + \partial_2(f_0 + f_1)\end{aligned}$$

Key Lemma

Definition (Collision Free Cycle)

A collision free cycle is one where the youngest positive simplex has not been paired (killed).

Lemma (Collision)

Given a filtration, $\mathbb{K}_j - \mathbb{K}_{j-1} = \sigma^j$, σ^j is the youngest positive simplex in $\Gamma(\partial_{k+1}\sigma^j)$. Let e be a collision free k -cycle in \mathbb{K}_{j-1} homologous to $\partial_{k+1}\sigma^j$. Suppose the youngest positive simplex in e is σ^g , then

$$\sigma^j = \sigma^g.$$

$$\max \Gamma(\partial_{k+1}\sigma^j) = \max(e) \quad \forall e \text{ collision free, } [e] = [\partial\sigma^j].$$

Key Lemma

Proof.

Let f be the sum of the basis cycles, homologous to $d = \partial_{k+1}\sigma^j$. By definition, f 's youngest positive simplex is σ^i , namely the youngest simplex in $\Gamma(\partial_{k+1}\sigma^j)$,

$$\sigma^i = \max \Gamma(\partial_{k+1}\sigma^j).$$

This implies that there are no cycles homologous to d in \mathbb{K}_{i-1} or earlier complexes. Let σ^g be the youngest positive simplex in e . $[e] = [d]$, therefore $g \geq i$.

If $g > i$, then $e = f + c$, where c bounds in \mathbb{K}^{j-1} . $\sigma^g \notin f$, implies $\sigma^g \in c$, and as σ^g is the youngest in e , it is also the youngest in c . \square

Key Lemma

continued.

Since e is collision free, the cycle created by σ^g , denoted as c^g , is still a non-boundary cycle in \mathbb{K}_{j-1} . Hence c^g can't be c , and can't be homologous to c when c becomes a boundary. Namely, when c is killed, σ^g is not paired yet.

It follows that the negative $(k+1)$ -simplex that kills c must pair a positive k -simplex in c , which is younger than σ^g , a contradiction.

This lemma shows, when σ^j is added to \mathbb{K}_{j-1} , we need to find any collision free cycle e homologous to $\partial_{k+1}\sigma^j$, and pair σ^j with the youngest positive simplex of e .

Pair Algorithm

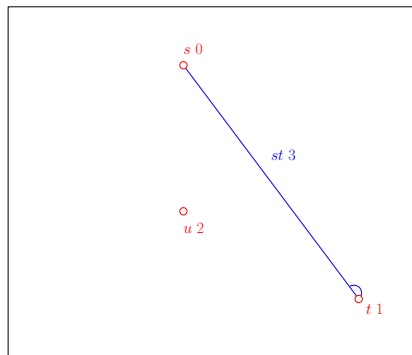
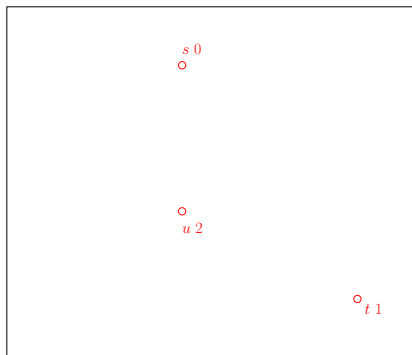
Pair(σ)

- 1 $c = \partial_p \sigma$
- 2 τ is the youngest positive $(p-1)$ -simplex in c .
- 3 **while** σ is paired and c is not empty **do**
- 4 find (τ, d) , d is the p -simplex paired with τ ;
- 5 $c \leftarrow \partial_p d + c$
- 6 Update τ to be the youngest positive $(p-1)$ -simplex in c
- 7 **end while**
- 8 **if** c is not empty **then**
- 9 σ is negative p -simplex and paired with τ
- 10 **else**
- 11 σ is a positive p -simplex
- 12 **endif**

Handle Loop and Tunnel Loop

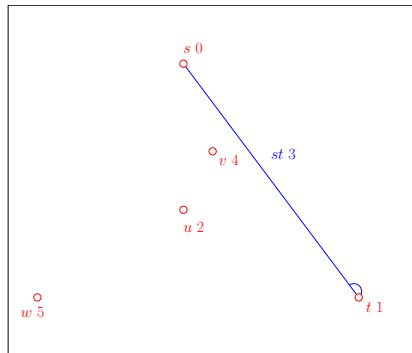
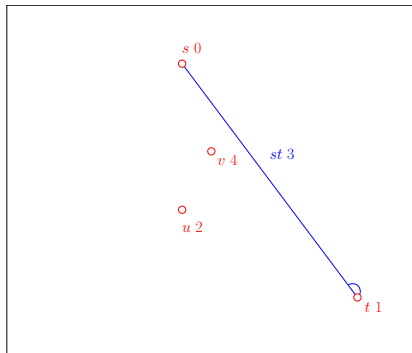
- 1 The simplices on the surface M are added into the filtration in any arbitrary order. Since $H_1(M)$ is of rank $2g$, the algorithm Pair generates $2g$ number of unpaired positive edges.
- 2 The simplices up to dimension 2 in I are added into the filtration. Since $H_1(I)$ of rank g , half of $2g$ positive edges generated in step 1 get paired with the negative triangles in I . Each pair corresponds to a killed loop, these g loops are handle loops.
- 3 Or the simplices up to dimension 2 in O are added into the filtration. Since $H_1(O)$ of rank g , half of $2g$ positive edges generated in step 2 get paired with the negative triangles in O . Each pair corresponds to a killed loop, these g loops are tunnel loops.

Example

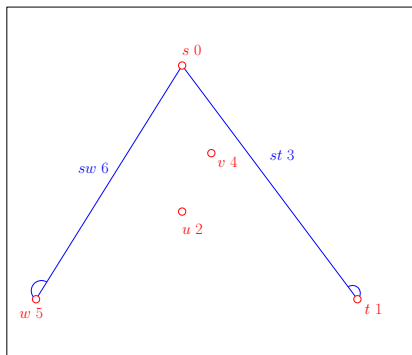


$$3. \partial st = s + t, (t_1, st_3)$$

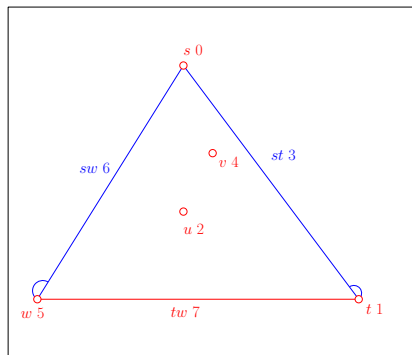
Example



Example



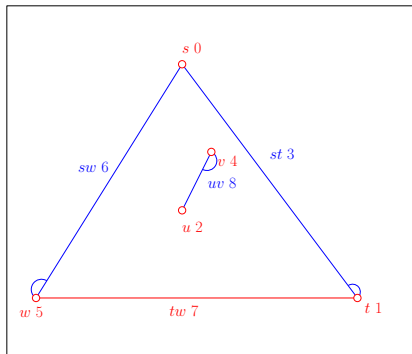
6. $\partial sw = s + w, (w_5, sw_6)$



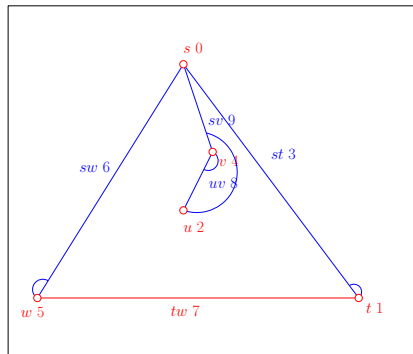
7. tw

$$\begin{aligned} 7 : \partial tw &= w + t = w + t + \partial st + \partial sw \\ &= w + t + (s + t) + (s + w) \\ &= 0. \end{aligned}$$

Example



8. $\partial uv = u + v, (v_4, uv_8)$

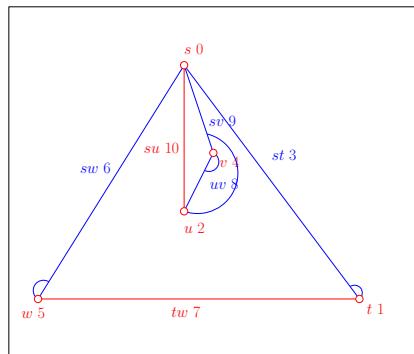


9. (u_2, sv_9)

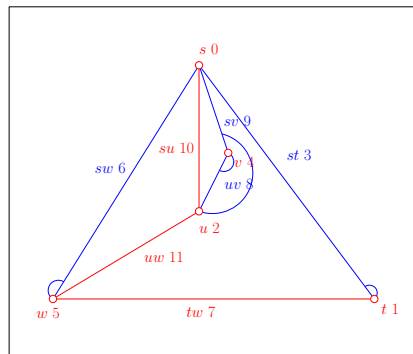
$$\begin{aligned} 9. \partial sv &= s + v = s + v + \partial uv \\ &= s + v + (u + v) \\ &= s + u \end{aligned}$$

.

Example



10. *su*



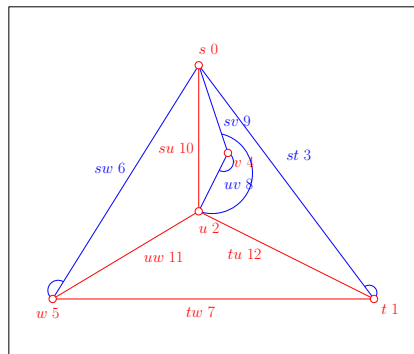
11. *uw*

Example

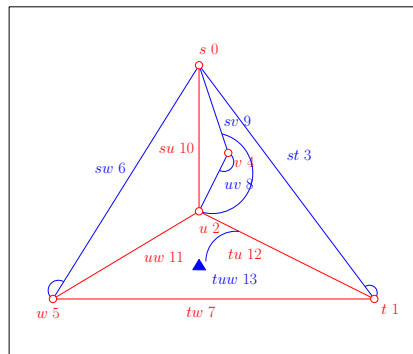
$$\begin{aligned} 10. \partial su &= s + u = s + u + \partial sv \\ &= s + u + (s + v) \\ &= u + v = u + v + \partial uv \\ &= 0. \end{aligned}$$

$$\begin{aligned} 11. \partial uw &= u + w = u + w + \partial sw \\ &= u + w + (s + w) \\ &= s + u = s + u + \partial sv \\ &= s + u = s + u + (s + v) \\ &= u + v = u + v + \partial uv \\ &= u + v + (u + v) \\ &= 0. \end{aligned}$$

Example



12. tu

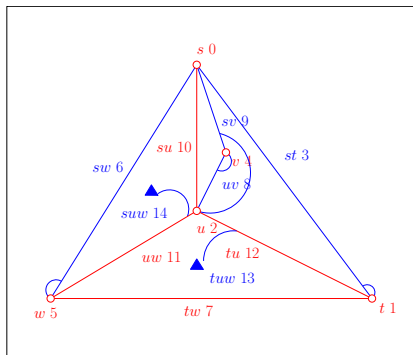


13. $tuw, (tu_{12}, tuw_{13})$

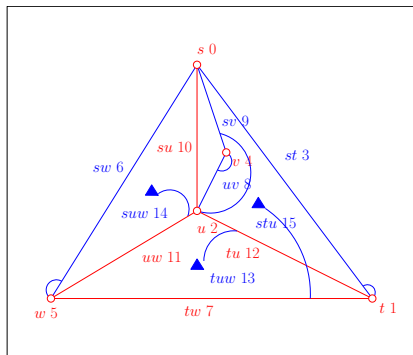
Example

$$\begin{aligned} 12. \partial tu &= t + u = t + u + \partial sv \\ &= t + u + (s + v) \\ &= t + u + s + v + \partial uv \\ &= t + u + s + v + (u + v) \quad 13. \partial tuw = tu + uw + wt \\ &= t + s \quad (tuw, tu) \\ &= s + t + \partial st \\ &= s + t + (t + s) \\ &= 0. \end{aligned}$$

Example



$$14. \partial suw = uw + su + sw \\ (uw_{11}, suw_{14})$$



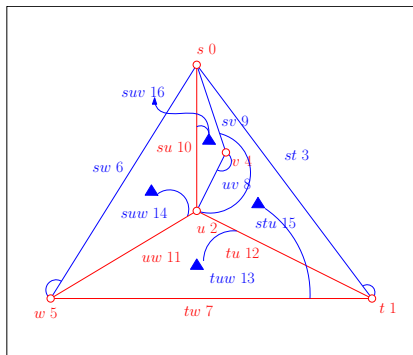
$$15. stu, (tw_7, stu_{15})$$

Example

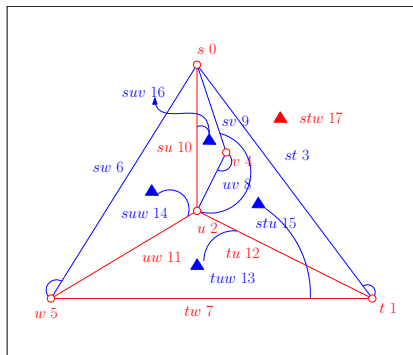
$$\begin{aligned} 15. \partial stu &= su + tu + st \\ &= su + st + tu + \partial tuw \\ &= su + st + tu + (tu + uw + tw) \\ &= su_{10} + st + uw_{11} + tw_7 \\ &= su_{10} + st + uw_{11} + tw_7 + \partial suw \\ &= su_{10} + st + uw_{11} + tw_7 + (sw + su_{10} + uw_{11}) \\ &= st + tw_7 + sw \end{aligned}$$

Hence we obtain the pair (stu, tw) .

Example



$$16. \partial \text{blue} = \text{red} + \text{blue} + \text{blue} \\ (\text{red}_{10}, \text{blue}_{16})$$

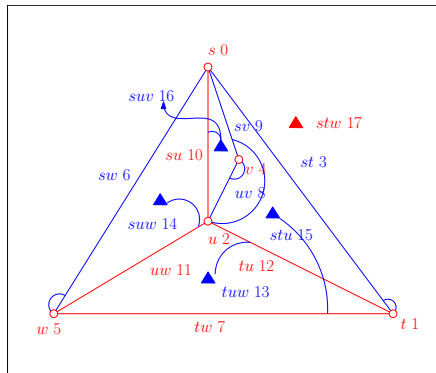


$$17. \text{red}$$

Example

$$\begin{aligned} 17. \partial stw &= tw + sw + st \\ &= sw + st + tw + \partial stu \\ &= sw + st + tw_7 + (st + tu_{12} + us_{10}) \\ &= sw + st + tw_7 + (st + tu_{12} + us_{10}) + \partial tuw \\ &= sw + st + tw_7 + (st + tu_{12} + us_{10}) + (tu_{12} + uw_{11} + wt_7) \\ &= sw + us_{10} + uw_{11} \\ &= sw + us_{10} + uw_{11} + \partial suw \\ &= sw + us_{10} + uw_{11} + (su_{10} + uw_{11} + ws) \\ &= 0. \end{aligned}$$

Example



<i>Creator</i>	<i>Killer</i>
t_1	st_3
u_2	sv_9
v_4	uv_8
w_5	sw_6
tw_7	stu_{15}
su_{10}	suv_{16}
uw_{11}	suw_{14}
tu_{12}	tuw_{13}

Incidence Matrix

Assuming an ordering of the $(p - 1)$ simplices and of the p -simplices, the boundary of a p -chain can be obtained by multiplication of the corresponding vector with the incidence matrix,

$$\partial(c_p) = D_p c_p.$$

The incidence matrix is defined as

$$D_p[i, j] = \begin{cases} 1 & \sigma_i^{p-1} \in \sigma_j^p \\ 0 & \sigma_i^{p-1} \notin \sigma_j^p \end{cases}$$

Incidence Matrix and Betti Number

A classic algorithm computes the Betti numbers of K by reducing its incidence matrices to Smith normal form. It uses row and column operations to zero out all entries except along an initial portion of the diagonal.

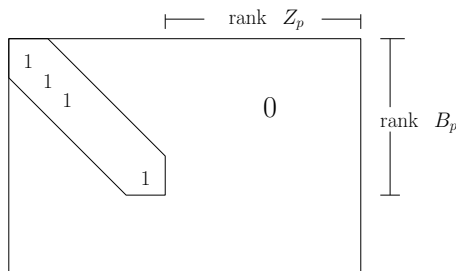


Figure: Smith norm of incidence matrix in \mathbb{Z}_2 .

The Betti number

$$\beta_p = \text{rank} Z_p - \text{rank} B_p.$$

Pairing Algorithm

Definition (Monotonous Filtering)

A filtering is monotonous, if in the ordering of K , any simplex σ is preceded by its faces.

An algorithm computes the persistence diagrams by pairing the simplices, which uses column operator to reduce D and another $0 - 1$ matrix R . Let $\text{low}_R(j)$ be the row index of the last 1 in column j of R , and (undefined if the column is zero).

Definition (Reduced Matrix and Pairing)

We call R reduced and low_R a pairing function, if

$$\text{low}_R(j) \neq \text{low}_R(j'),$$

whenever $j \neq j'$ specify two non-zero columns.

Pairing Algorithm

Algorithm: Incidence matrix reduction

- ① $R \leftarrow D$
- ② **for** $j = 1$ **to** n **do**
- ③ **while** $\exists j' < j$ **with** $\text{low}_R(j') = \text{low}_R(j)$ **do**
- ④ add column j' to column j
- ⑤ **endwhile**
- ⑥ **endfor.**

The pairing is given by

$$(\sigma_i, \sigma_j) \iff i = \text{low}_R(j).$$

σ_i is positive, it generates a homology class; σ_j is negative, it kills a homology class.

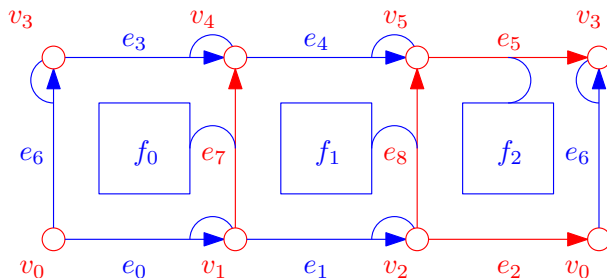


Figure: Generators and killers.

Pairing by matrix induction

Boundary operator ∂_1 , incidence matrix D_1 ,

	e_0	e_1	e_2	e_3	e_4	e_5	e_6	e_7	e_8
v_0	1	0	1	0	0	0	1	0	0
v_1	1	1	0	0	0	0	0	1	0
v_2	0	1	1	0	0	0	0	0	1
v_3	0	0	0	1	0	1	1	0	0
v_4	0	0	0	1	1	0	0	1	0
v_5	0	0	0	0	1	1	0	0	1

Pairing by matrix induction

$1 + 2, 4 + 5, 3 + 7, 4 + 8$

	e_0	e_1	e_2	e_3	e_4	e_5	e_6	e_7	e_8
v_0	1	0	1	0	0	0	1	0	0
v_1	1	1	1	0	0	0	0	1	0
v_2	0	1	0	0	0	0	0	0	1
v_3	0	0	0	1	0	1	1	1	0
v_4	0	0	0	1	1	1	0	0	1
v_5	0	0	0	0	1	0	0	0	0

Pairing by matrix induction

$$1 + 2, 3 + 5, 3 + 8$$

	e_0	e_1	e_2	e_3	e_4	e_5	e_6	e_7	e_8
v_0	1	0	0	0	0	0	1	0	0
v_1	1	1	0	0	0	0	0	1	0
v_2	0	1	0	0	0	0	0	0	1
v_3	0	0	0	1	0	0	1	1	1
v_4	0	0	0	1	1	0	0	0	0
v_5	0	0	0	0	1	0	0	0	0

Pairing by matrix induction

$6 + 7, 6 + 8$

	e_0	e_1	e_2	e_3	e_4	e_5	e_6	e_7	e_8
v_0	1	0	0	0	0	0	1	1	1
v_1	1	1	0	0	0	0	0	1	0
v_2	0	1	0	0	0	0	0	0	1
v_3	0	0	0	1	0	0	1	0	0
v_4	0	0	0	1	1	0	0	0	0
v_5	0	0	0	0	1	0	0	0	0

Pairing by matrix induction

$$0 + 7, 1 + 8, 0 + 8$$

	e_0	e_1	e_2	e_3	e_4	e_5	e_6	e_7	e_8
v_0	1	0	0	0	0	0	1	0	0
v_1	1	1	0	0	0	0	0	0	0
v_2	0	1	0	0	0	0	0	0	0
v_3	0	0	0	1	0	0	1	0	0
v_4	0	0	0	1	1	0	0	0	0
v_5	0	0	0	0	1	0	0	0	0

Generators e_2, e_5, e_7, e_8 , corresponding to 0 columns. Killers corresponds to non-zero columns. Pairing

$$(e_0, v_1), (e_1, v_2), (e_3, v_4), (e_4, v_5), (e_6, v_3)$$

Pairing by matrix induction

	f_0	f_1	f_2
e_0	1	0	0
e_1	0	1	0
e_2	0	0	1
e_3	1	0	0
e_4	0	1	0
e_5	0	0	1
e_6	1	0	1
e_7	1	1	0
e_8	0	1	1

$$= \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \\ 0 & 1 & 1 \end{bmatrix} \xrightarrow{2+3} \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 1 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 1 \\ 0 & 0 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 1 \\ 0 & 1 & 0 \end{bmatrix} \xrightarrow{1+3} \begin{bmatrix} 1 & 0 & 1 \\ 0 & 1 & 1 \\ 0 & 0 & 1 \\ 1 & 0 & 1 \\ 0 & 1 & 1 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \\ 1 & 1 & 0 \\ 0 & 1 & 0 \end{bmatrix}$$

The pairing is

$$(f_0, e_7), (f_1, e_8), (f_2, e_5)$$

Topological Annulus

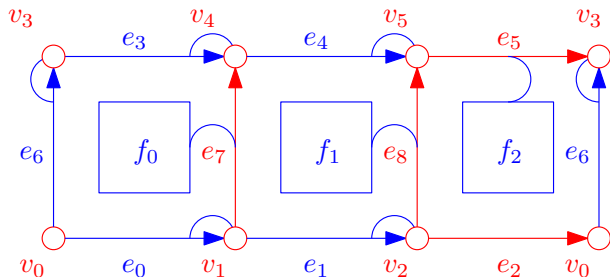
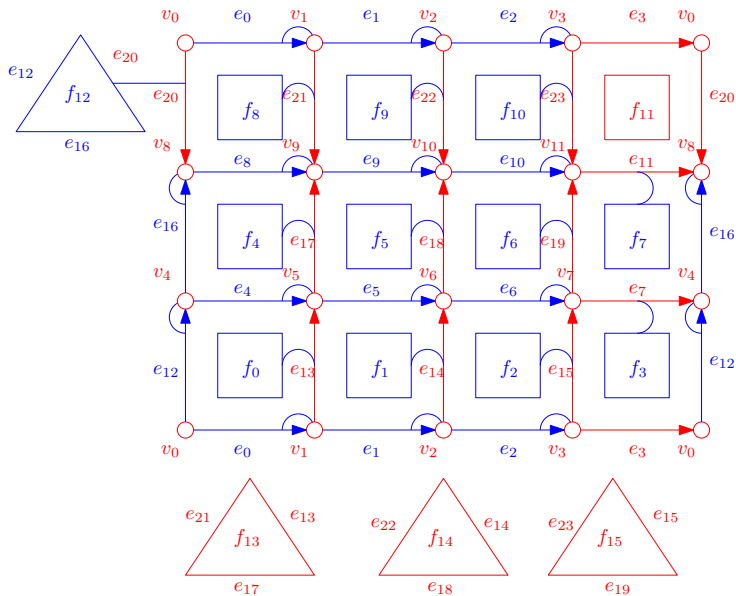
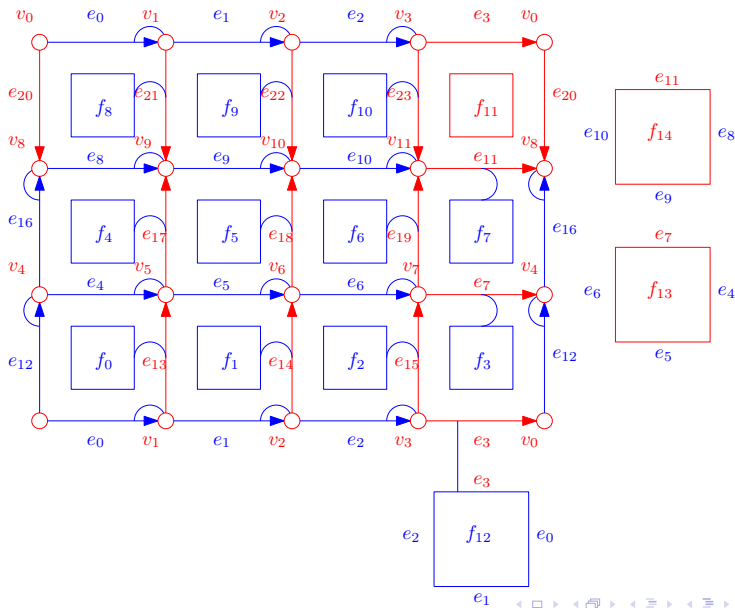


Figure: Topological Annulus.

Topological Torus



Topological Torus



Topological Torus

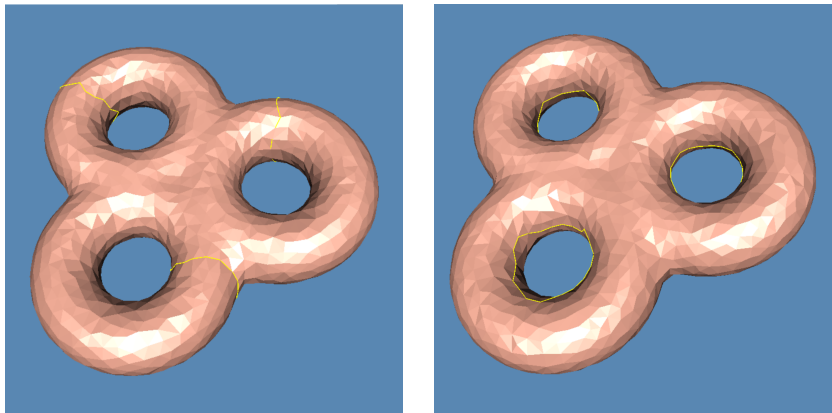


Figure: Handle and tunnel loops.

Topological Torus

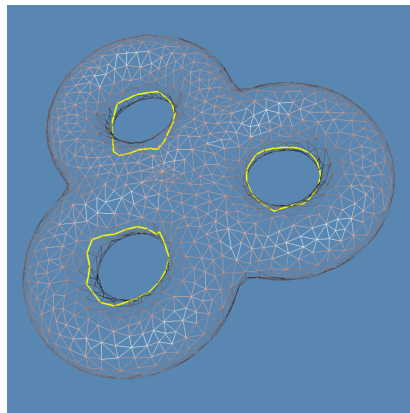
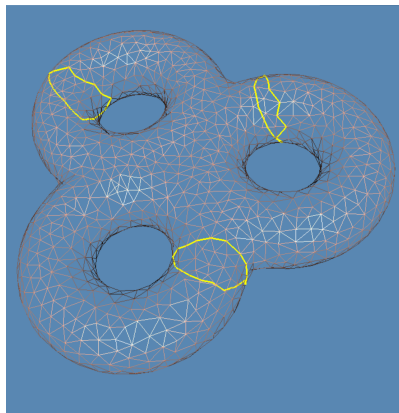


Figure: Handle and tunnel loops.

Topological Torus

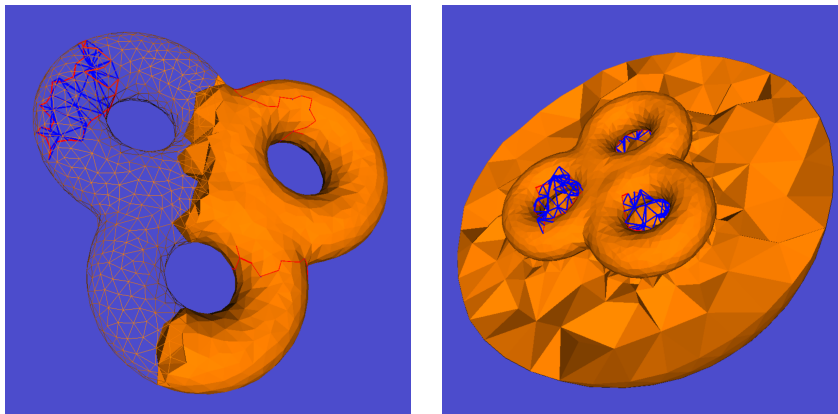


Figure: Interior and exterior volumes.