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# Week 4

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Xiaoyuan Xie 谢晓园

[xxie@whu.edu.cn](mailto:xxie@whu.edu.cn)

计算机学院E301

## Week 4 作业

- **P130 4.2.1** Consider the context-free grammar:

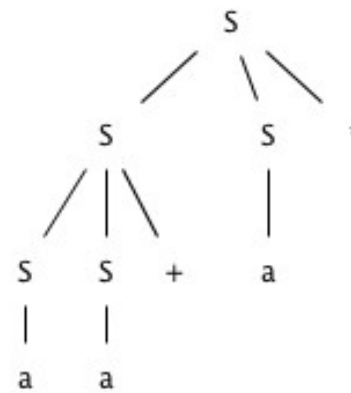
$$S \rightarrow SS + \mid SS^* \mid a$$

and the string **aa + a\***.

- 1) Give a leftmost derivation for the string.
  - $S \Rightarrow SS^* \Rightarrow SS+S^* \Rightarrow aS+S^* \Rightarrow aa+S^* \Rightarrow aa+a^*$
- 2) Give a rightmost derivation for the string.
  - $S \Rightarrow SS^* \Rightarrow Sa^* \Rightarrow SS+a^* \Rightarrow Sa+a^* \Rightarrow aa+a^*$

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- 3) Give a parse tree for the string.



- 4) Is the grammar ambiguous or unambiguous? Justify your answer.
  - Unambiguous.
- 5) Describe the language generated by this grammar.
  - The set of all postfix expressions consist of addition and multiplication.

## Week 4 作业

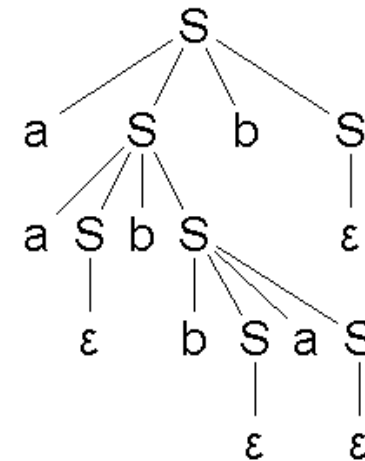
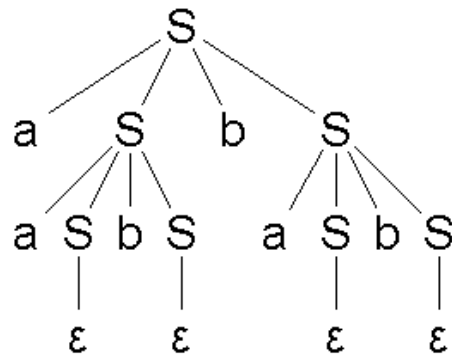
### ■ P131 4.2.2(6)

$$S \rightarrow a S b S \mid b S a S \mid \epsilon$$

with string aabbab.

- 1)  $S \Rightarrow aSbS \Rightarrow aaSbSbS \Rightarrow aabSbS \Rightarrow aabbS \Rightarrow aabbaSbS \Rightarrow aabbabS \Rightarrow aabbab$
- 2)  $S \Rightarrow aSbS \Rightarrow aSbaSbS \Rightarrow aSbaSb \Rightarrow aSbab \Rightarrow aaSbSbab \Rightarrow aaSbbab \Rightarrow aabbab$

- 3) Give a parse tree for the string.



- 4) Ambiguous.

- 5) The set of all strings of 'a's and 'b's of the equal number of 'a's and 'b's.



## Week 4 作业

### ■ P131 4.2.3 Design grammars for the following languages:

- 2) The set of all strings of 0s and 1s that are palindromes; that is, the string reads the same backward as forward.
  - $S \rightarrow 0S0 \mid 1S1 \mid 0 \mid 1 \mid \varepsilon$
- 5) The set of all strings of 0s and 1s in which 011 does not appear as a substring.
  - $S \rightarrow 1^* (0+1?)^*$

## Week 4 作业

- **P131 4.2.5** Use the braces described in Exercise 4.2.4 to simplify the following grammar for statement blocks and conditional statements:

stmt -> **if** expr **then** stmt **else** stmt  
          | **if** stmt **then** stmt  
          | **begin** stmtList **end**  
stmtList -> stmt; stmtList | stmt

- **Answer:**

stmt -> **if** expr **then** stmt [**else** stmt]  
          | **begin** stmtList **end**  
stmtList -> stmt [; stmtList]

## Week 4 作业

- **P137 4.3.3** The following grammar is proposed to remove the "dangling-else ambiguity" discussed in Section 4.3.2:

```
stmt -> if expr then stmt  
      | matchedStmt  
matchedStmt -> if expr then matchedStmt else stmt  
             | other
```

Show that this grammar is still ambiguous.



## Week 4 作业

1)

```
if expr
then
    if expr
    then matchedStmt
    else
        if expr
        then matchedStmt
else stmt
```

2)

```
if expr
then
    if expr
    then matchedStmt
    else
        if expr
        then matchedStmt
    else stmt
```



*Thank you!*