

$$13. L = \{a^n b^n c^k \mid k \geq n \geq 0\}$$

~~if L is regular context-free,~~

~~then $L = C^{k=n}$~~

Two cases:

$$(1) L_1 = \{a^n b^n c^n \mid n \geq 0\} \quad (k=n)$$

shown in class as non-context-free,

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$$(2) L_2 = \{a^n b^n c^k \mid k > n \geq 0\} \quad (k > n)$$

if we assume L_2 is ^{context-free} ~~regular~~, then
there is pumping $\#$ P.

$$w = a^p b^p c^{p+1}, w \in L_2 \text{ s.t. } |w| \geq p.$$

~~Then if we let~~

If we pick either a or b for v or y ,
then we break either the a/b parity, or
we increase a/b beyond the count of c .

If we pick c , then we can pump down,
resulting in $k \leq n$, which
is not a word in L_2 .

Contradiction. L_2 cannot be context-free.