

习题 6.1

1. 算法: find MinDifference(A[0..n-1])
 // 先对数组排序
 // 输入 n 个数字构成的一数组
 // 输出两个最近数的距离
 对数组排序
 $\text{min} = A[1] - A[0]$
 for $i \leftarrow 1$ to $n-2$ do
 if $A[i+1] - A[i] < \text{min}$ $\text{min} = A[i+1] - A[i]$
 return min
 排序: $O(n \log n)$, 找最小距离: $O(n)$
 效率类型: $O(n \log n)$

b. 暴力算法: $O(n^2)$
 ∴ 该算法效率类型“更高效率”

习题 6.2

1. 解:

$$\begin{bmatrix} 1 & 1 & 1 & 2 \\ 2 & 1 & 1 & 3 \\ 1 & -1 & 3 & 8 \end{bmatrix} \xrightarrow[\substack{R_2 - R_1 \\ R_3 - R_1}]{R_1 - 2R_1} \begin{bmatrix} 1 & -1 & -1 & -2 \\ 0 & -1 & -1 & -1 \\ 0 & -2 & 2 & 6 \end{bmatrix} \xrightarrow{R_3 - 2R_2} \begin{bmatrix} 1 & -1 & -1 & -2 \\ 0 & -1 & -1 & -1 \\ 0 & 0 & 4 & 8 \end{bmatrix}$$

∴ $4x_3 = 8, x_3 = 2$
 $-x_2 - x_3 = -1, x_2 = 1$
 $x_1 + x_2 + x_3 = 2, x_1 = -1$

2. 解:

a. 由 1 知 $U = \begin{bmatrix} 1 & 1 & 1 \\ 0 & -1 & -1 \\ 0 & 0 & 4 \end{bmatrix}$, $b = \begin{bmatrix} 2 \\ 3 \\ 8 \end{bmatrix}$
 $L = \begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ 1 & 2 & 1 \end{bmatrix}$

解 $Ly = b, y_1 = 2, y_2 = -1, y_3 = 8$

解 $Ux = y, x_1 = 1, x_2 = -1, x_3 = 2$

∴ $x_1 = 1, x_2 = -1, x_3 = 2$

b. 应归为递归, 将一个问题化简为两个更简单的问题, 但不能用递归处理

习题 6.3

2. 解:

a. $n=1$

①

$n=2$



$n=3$

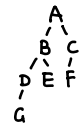


$n=4$

$n=5$

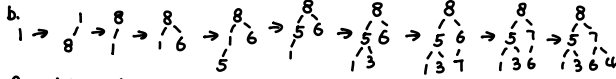
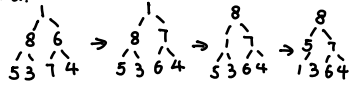


b. $h=4$

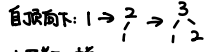
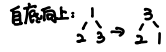


习题6.4

1. a.

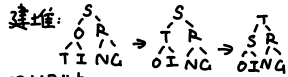


c. 不成立, 如输入为 1, 2, 3 时

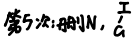
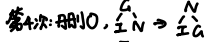
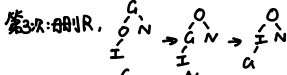
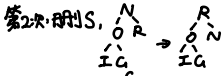
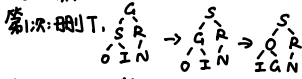


这两者不一样

7. c. 解: 堆排序 SORTING



删除根节点:



第六次: 删除 G

∴ 排序: G I N O R S T

习题6.5

4. a. 计算多项式:

系数 3 -1 0 2 5
x = -2 3 -7 14 -26 57

$$\therefore p(x = -2) = 57$$

b. 由 a 可知, 结果为 $3x^5 - 7x^4 + 14x^3 - 26x^2 + 57x - 26$,
余数为 57

7. a. 7 = 10001

$$1 \ 0 \ 0 \ 0 \ 1 \\ a^6 \ a^5 \ a^4 \ a^3 \ a^2 \ a^1 \ a^0$$

b. 可以, 将整数分解为二进制数, 逐位处理

8. 1 0 0 0 1

$$a^6 \ a^5 \ a^4 \ a^3 \ a^2 \ a^1 \ a^0 \\ a \times a^6 = a^7 \ a$$

题 6.6

1. a. 解: 假设 $\gcd(m, n) = x$

$$m = k_1 x, n = k_2 x \quad k_1, k_2 \text{ 互质}$$

$$\therefore \text{lcm}(m, n) = k_1 k_2 x = \frac{m \times n}{x} = k_1 k_2 x$$

b. $m, n \in O(1), \gcd(m, n) \in O(\log n)$

$$\therefore \text{lcm}(m, n) \text{ 类型为 } O(\log n)$$

2. 将每一个值变为值相反数, 构造最大堆
之后再还原为原值