

DIFFERENTIAL PRIVACY FOR FREE?

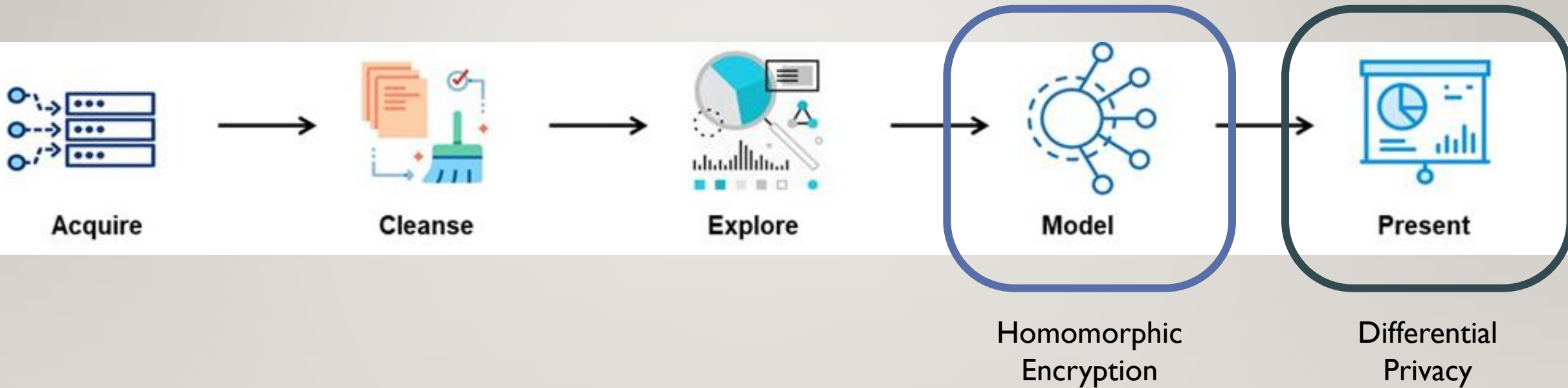
HARNESSING THE NOISE IN APPROXIMATE HOMOMORPHIC ENCRYPTION



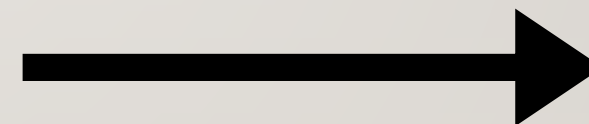
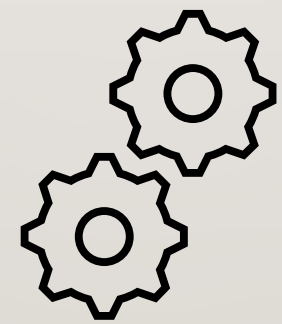
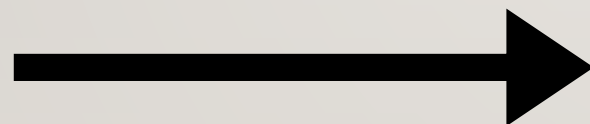
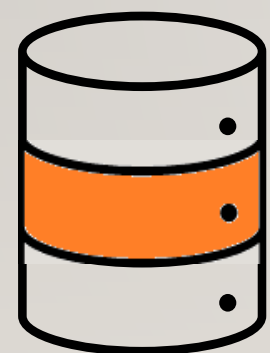
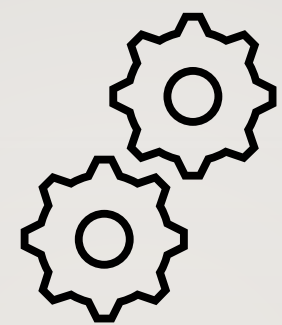
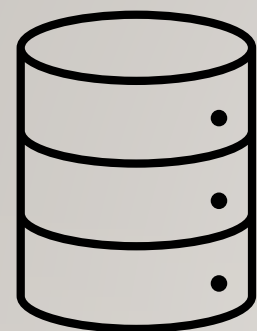
TALK PLAN

- Motivation
- What is Differential Privacy?
- Noise in Homomorphic Encryption
- Differential Privacy for Free?
- Analysis
- Case Study Results
- Further Work

MOTIVATION

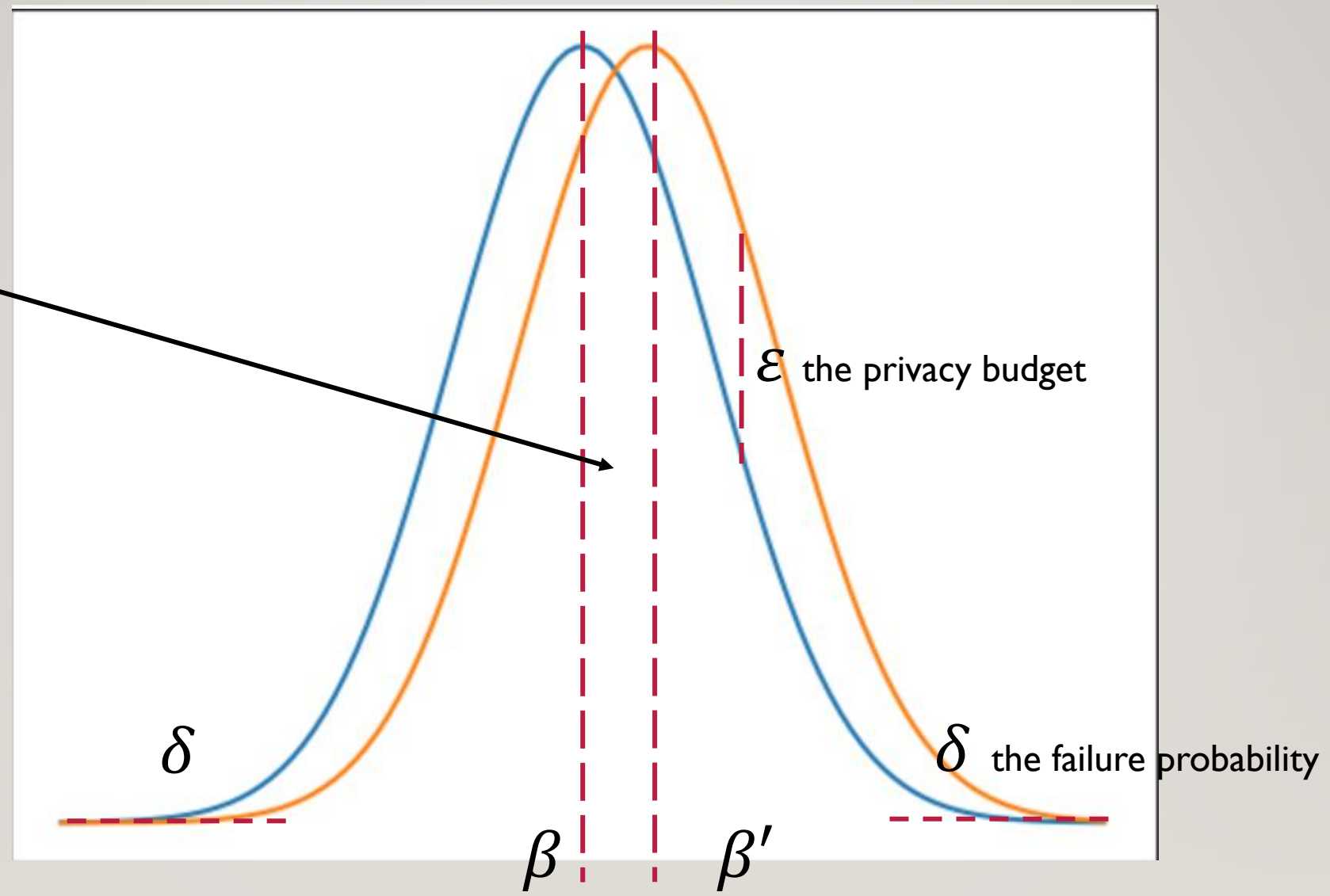


WHAT IS DIFFERENTIAL PRIVACY I

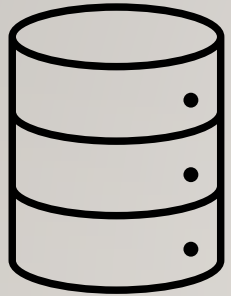


WHAT IS DIFFERENTIAL PRIVACY II

sensitivity



WHAT IS DIFFERENTIAL PRIVACY III



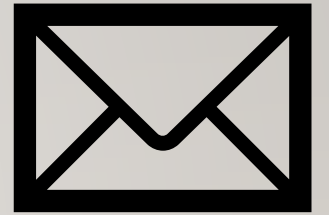
data



objective



gradient



output



$$(a, as + e)$$

- How large is the noise?
- How does the noise change when we perform homomorphic operations?
- Is the noise small enough to remove during decryption or bootstrapping?

Worst Case

1. Bound fresh sources of noise using tail bounds
2. Update after each operation according to the “worst case” growth

- ✓ Very robust
- ✓ (Relatively) easy to implement
- ✗ Loose bounds
- ✗ Returns a bound on the noise only

Average Case

1. Analyse how the distribution of the noise changes over the course of a circuit
2. Use tail bounds to return a final upper bound on the noise

- ✓ Tight bounds
- ✓ Description of the noise distribution
- ✗ Requires many assumptions
- ✗ Can be difficult to deploy

can homomorphic encryption noise give
differential privacy *for free?*

- Noise is removed during decryption \longrightarrow *approximate* homomorphic encryption
- Noise remains small \longrightarrow high depth circuit
- Need to know the noise distribution to accurately specify the privacy leakage \longrightarrow use the heuristics of [1], which argues that noise in CKKS follows a normal distribution throughout a circuit

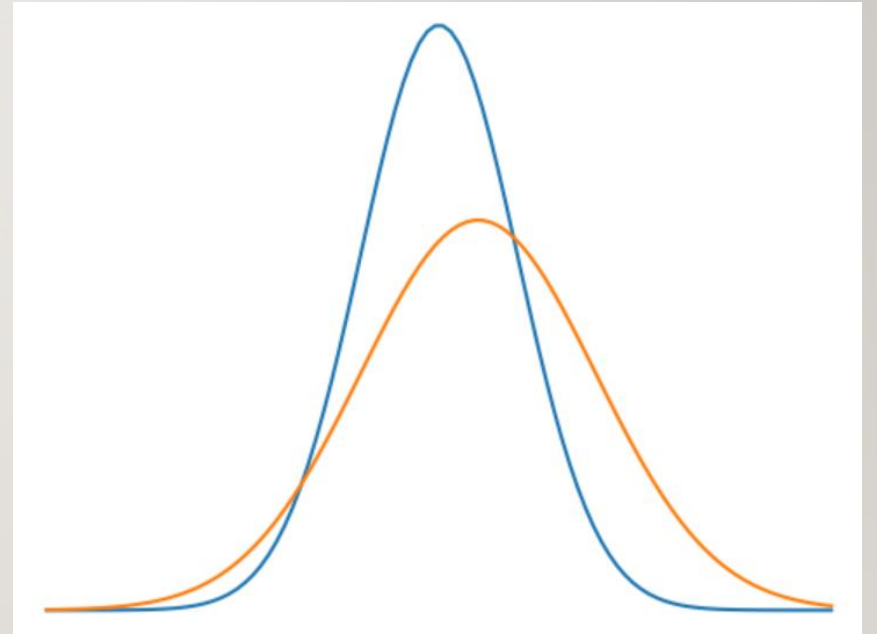
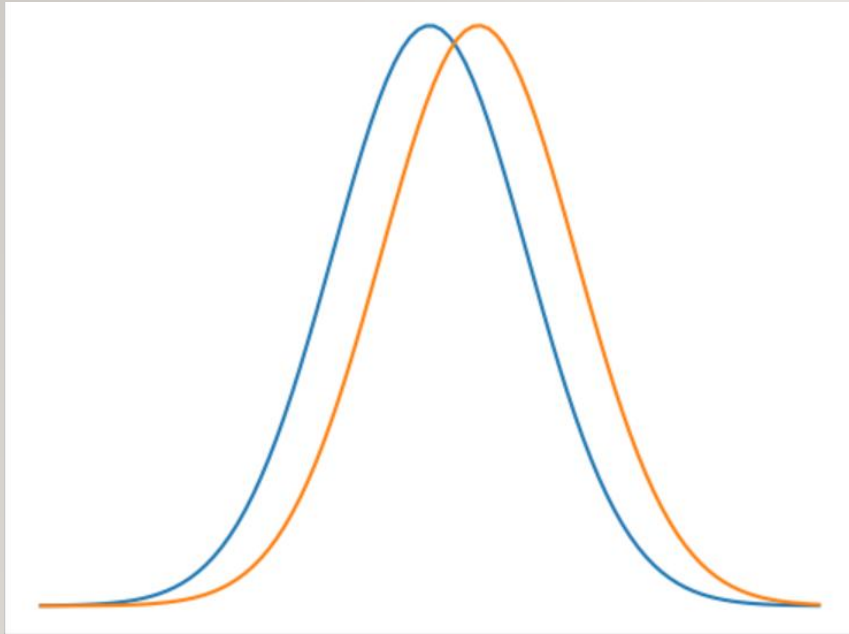
* *Can only evaluate quadratic polynomials*

1. Use CKKS
2. Choose a high depth application
3. Allow the noise to grow large enough

$$N\rho_1^2\rho_2^2 + \rho_1^2 \underline{|m_2|^2} + \rho_2^2 \underline{|m_1|^2}$$

...the shape of the distribution depends on the input data

ANALYSIS II



Let $\kappa = \text{sensitivity/standard deviation}$, $\tau = \text{standard deviation/standard deviation}'$

Standard case

$$\varepsilon \geq \sqrt{c} \kappa$$

Our case, ID

$$\varepsilon \geq \tau^2 \sqrt{c'} \kappa + \frac{1}{2} \tau^2 \kappa^2 + \frac{1}{2} (\tau^2 - 1) c' + \ln \tau \quad \longrightarrow \quad \varepsilon \geq \sqrt{c'} \kappa + \frac{1}{2} \kappa^2$$

Our case, n-D

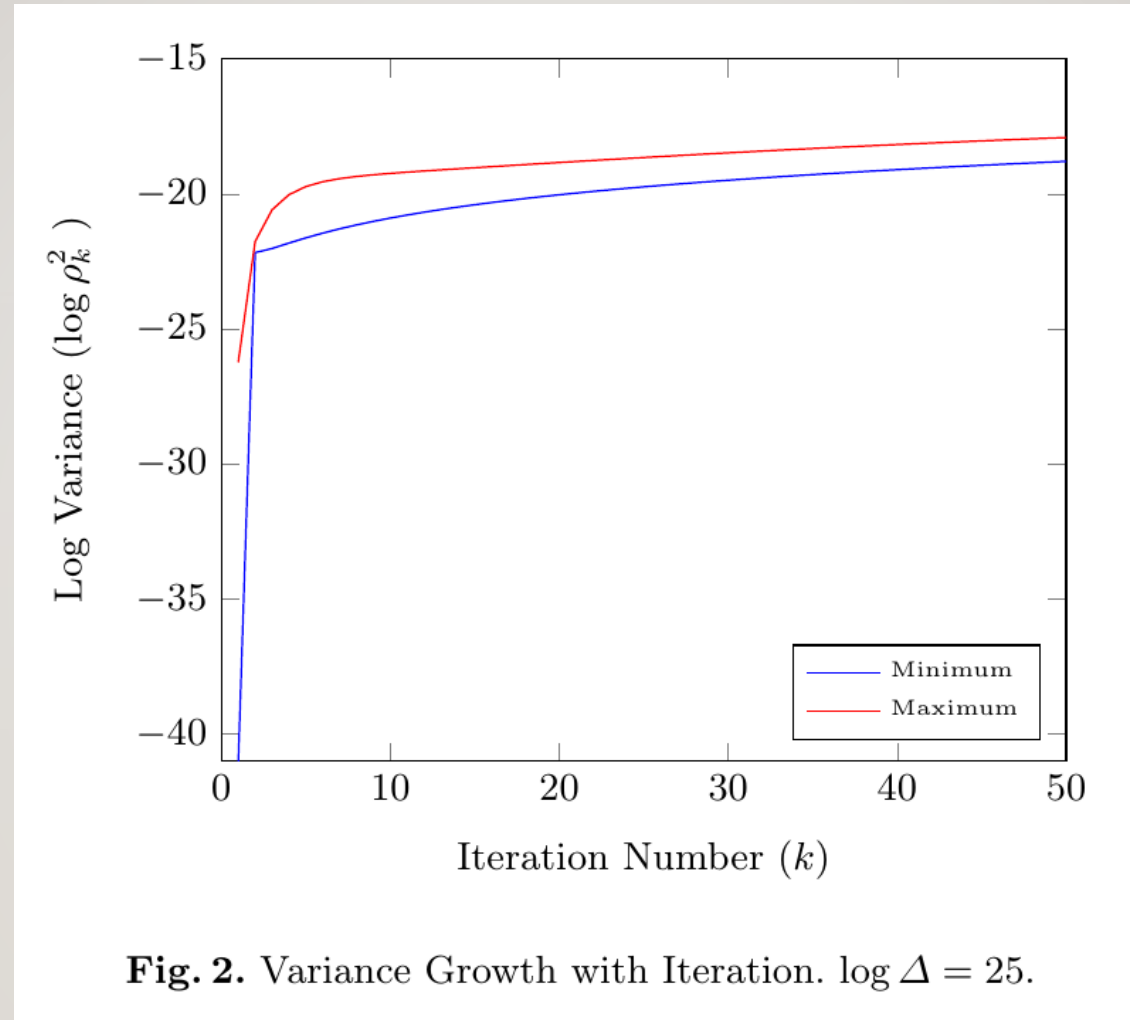
$$\varepsilon \geq \sqrt{c'' \left(\tau^4 \kappa^2 + \frac{n(\tau^2 - 1)}{2} \right)} + \frac{1}{2} \tau^2 \kappa^2 + \frac{1}{2} (\tau^2 - 1) (c'' + n) + \ln \tau \quad \longrightarrow \quad \varepsilon \geq \sqrt{c''} \kappa + \frac{1}{2} \kappa^2$$

case study requirements:

- CKKS
- (arbitrarily) high depth
- quadratic



Ridge regression training
using gradient descent [2]



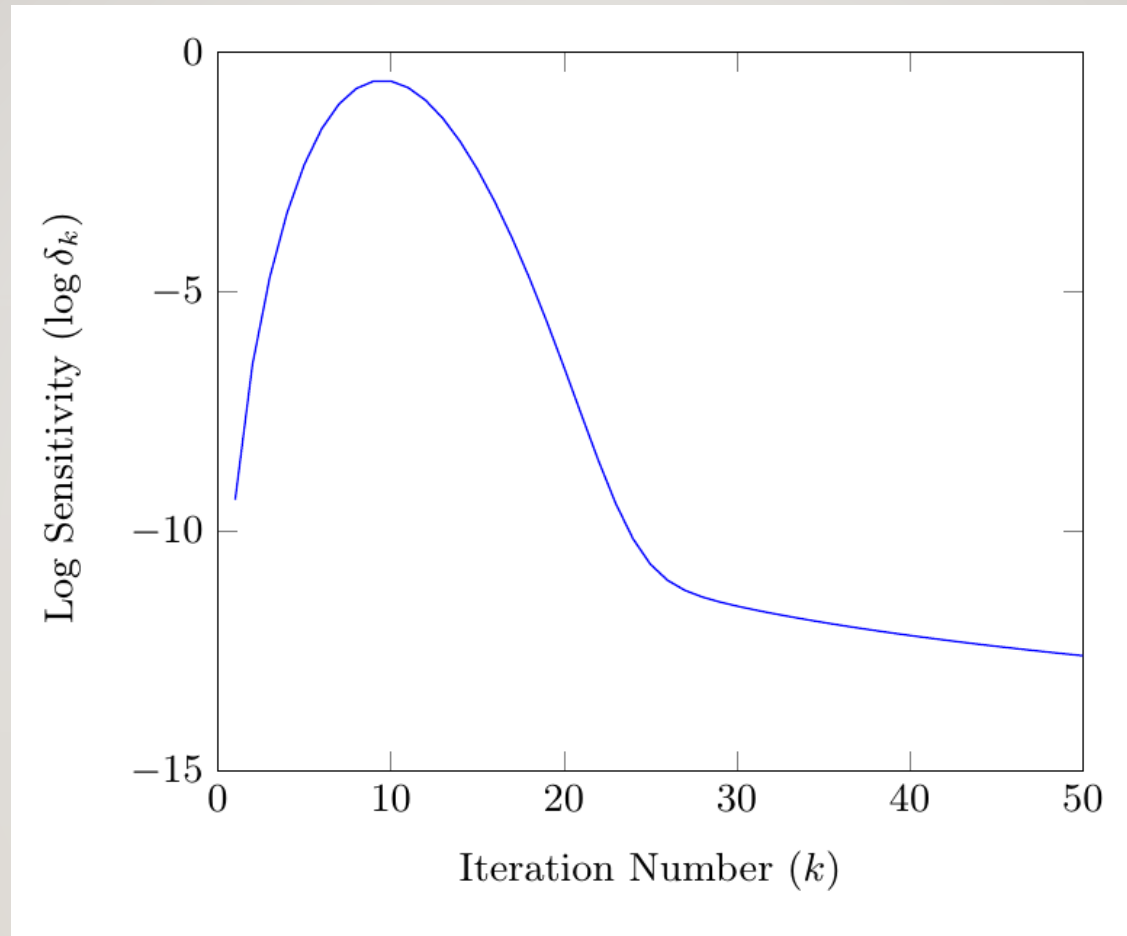
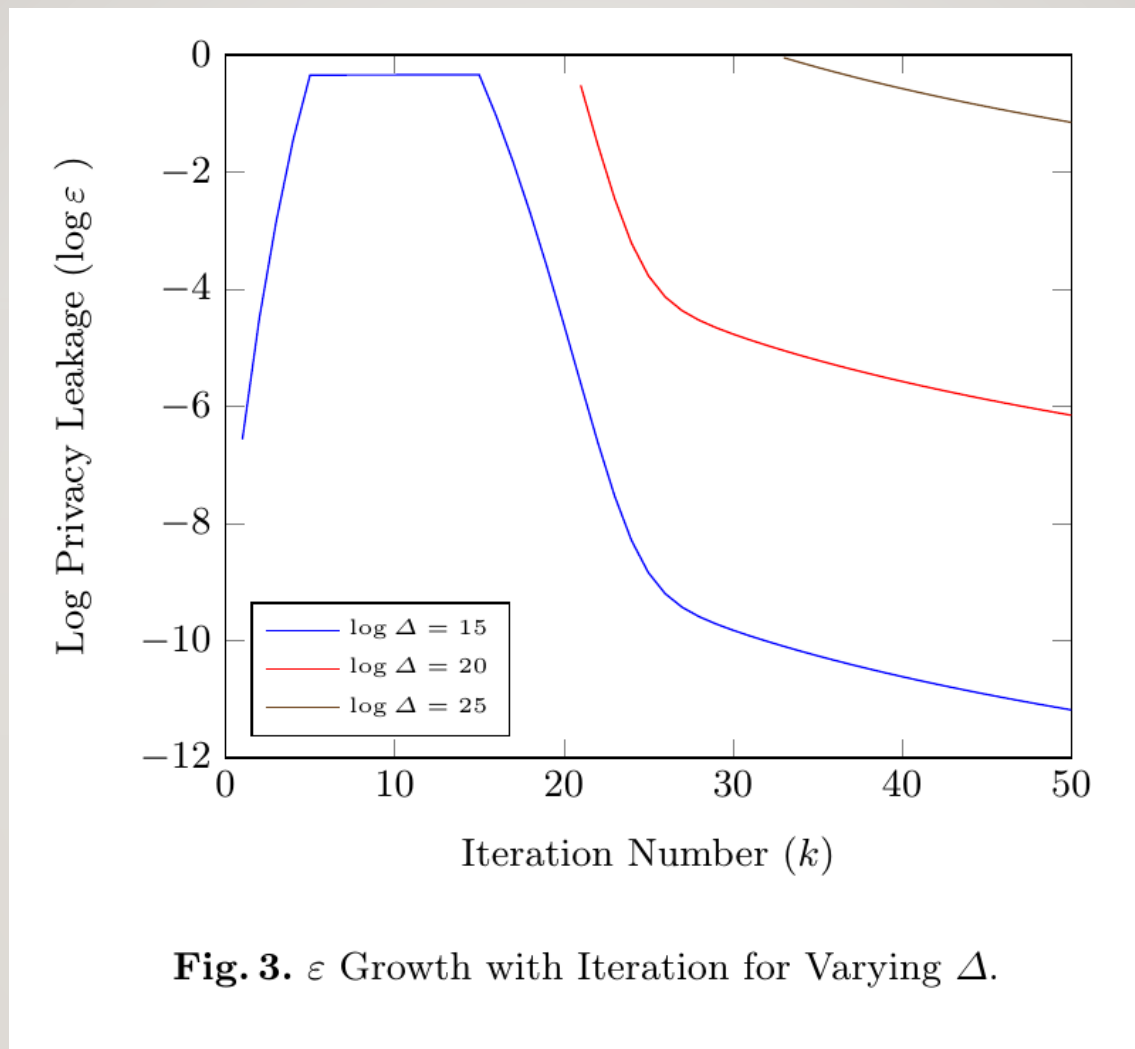


Fig. 2. Sensitivity Growth with Iteration.



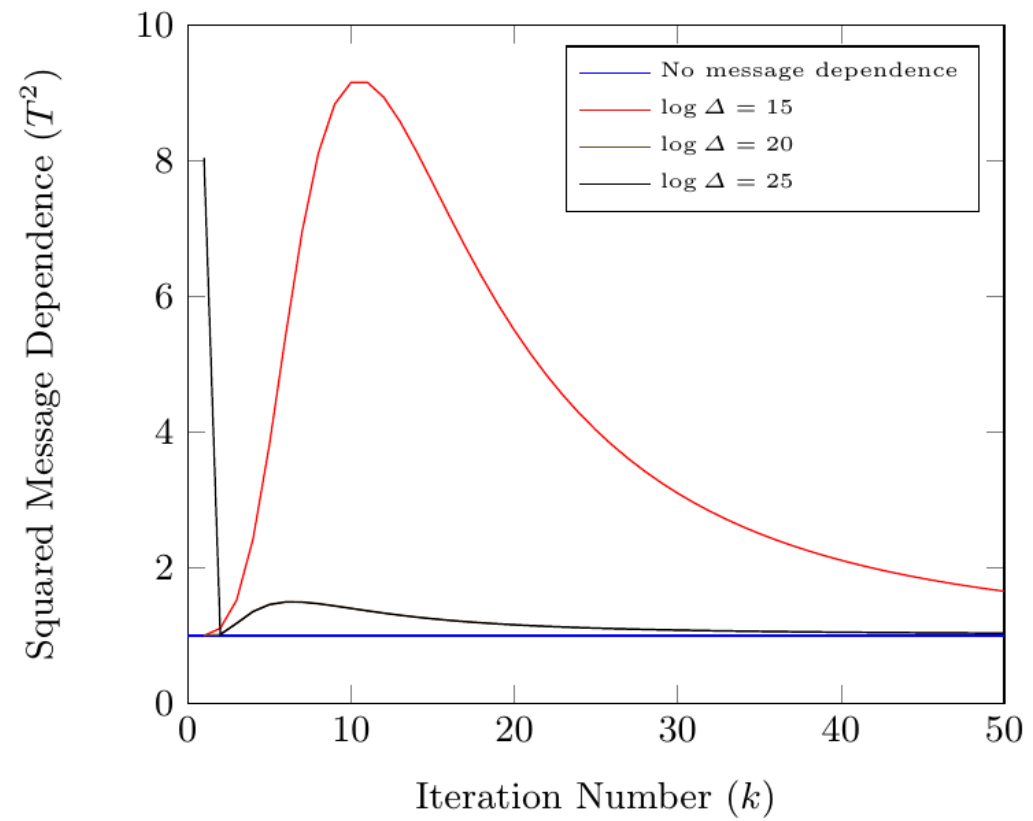


Fig. 6. Message Dependence Change with Iteration.

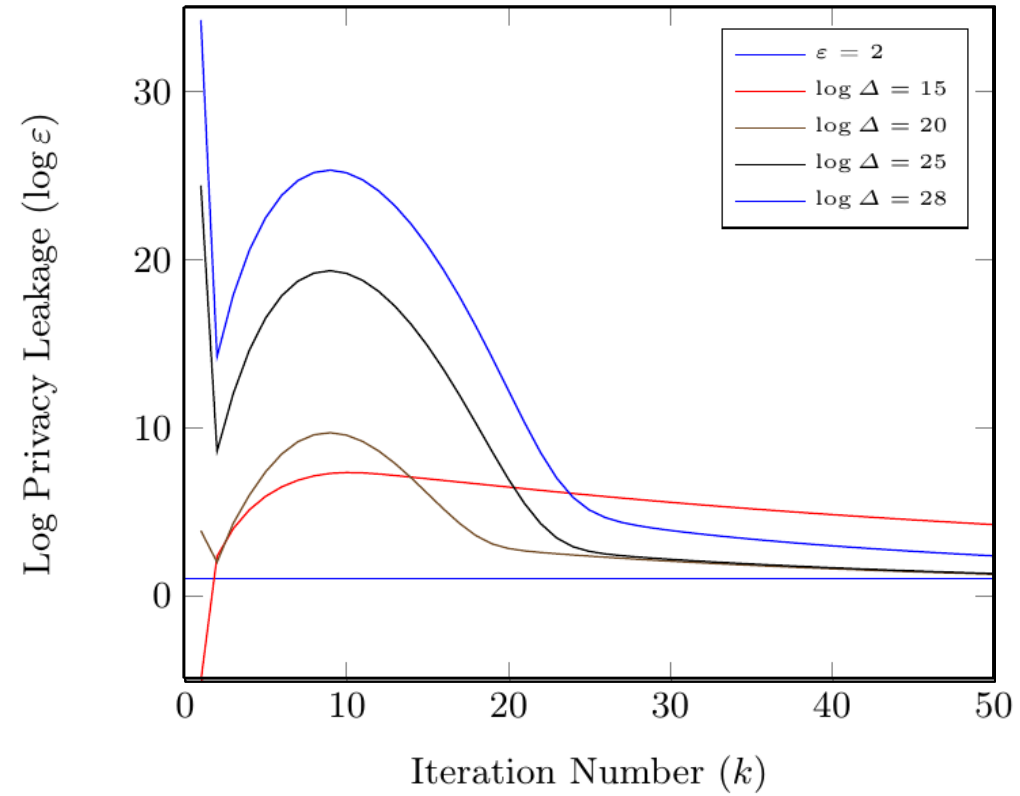


Fig. 7. Change in Log Privacy Leakage with Iteration.

- We investigated the extent to which HE noise can provide differential privacy “for free”
- Identified message dependence as a key barrier
- Derived new results on the Differential Privacy in this setting
- Explored our results with a case study, and found a privacy budget of $\epsilon \approx 2$ achievable with 50 iterations

further work

1. Further Noise Analysis
2. From Heuristic to Guarantee
3. Alternative Applications and Schemes
4. Beyond Output Perturbation
5. Differential Privacy “At A Discount”

Thanks for listening!

tabitha.l.ogilvie@gmail.com