4. Basic building blocks of concurrency

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Concurrent programming (CC3040) 2024/2025

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https://fm-dcc.github.io/cp2425







Overview

We are here



Blocks of sequential code running concurrently and sharing memory:

- What is Scala?
- Concurrency in Java and its memory model
- Basic concurrency blocks and libraries
- Futures and Promises
- Data-Parallel Collections
- Reactive Programming (Concurrently)
- Software Transactional Memory
- Actor model

What we will see



- Tread pools: Executor and ExecutionContext
- Non-blocking synchronisation compare-and-set (CAS)
- Lazy (concurrent) values
- Concurrent collections
- Running OS processes
- Futures and Promises

Existing thread pools in Scala

Executor interface



```
Executor executor = anExecutor;
executor.execute(new RunnableTask1());
executor.execute(new RunnableTask2());
...
```

```
import scala.concurrent._
import java.util.concurrent.ForkJoinPool
object ExecutorsCreate extends App {
  val executor = new ForkJoinPool
  executor.execute(new Runnable {
    def run() = log("This; task; is; run;;
        asvnchronously.")
 })
  Thread.sleep(500) // not needed with
      fork := false in SBT
```

- Executor: can start a new thread, an existing one, or the current one
- Abstracts from the management of threads
- ExecutorService: API that extends
 Executor with shutdown
 - executor.shutdown → executes all tasks and then stops working threads
 - executor.awaitTermination(...) →
 force termination if, after a given
 time, the tasks are not completed

Scala's ExecutionContext



```
import scala.concurrent._
object ExecutionContextGlobal extends App {
  val ectx = ExecutionContext.global
  ectx.execute(new Runnable {
    def run() = log("Running_on_the_execution_context.")
  })
  Thread.sleep(500)
}
```

```
object ExecutionContextCreate extends App {
  val pool = new forkjoin.ForkJoinPool(2)
  val ectx = ExecutionContext.fromExecutorService(pool)
  ectx.execute(new Runnable {
    def run() = log("Running_on_the_execution_context_again.")
  })
  Thread.sleep(500)
```

- scala.concurrent: has
 ExecutionContext
- Similar to Executor but more Scala specific
- often used as implicit parameter
- global: default execution context (internally uses a ForkJoinPool)
 - romExecutorService:
 creates ExecutionContext
 from ExecutorService

Simplifying the execution



Similar to threads:

We now define execute

```
def execute(body: =>Unit) =
    ExecutionContext.global.execute(
     new Runnable { def run() = body }
// For example:
object ExecutionContextSleep extends App {
  for (i<- 0 until 32) execute {</pre>
    Thread. sleep (2000)
    log(s"Tasku$iucompleted.")
  Thread. sleep (10000)
```



```
object ExecutionContextSleep
        extends App {
  for (i<- 0 until 32) execute {
    Thread.sleep(2000)
    log(s"Tasku$iucompleted.")
  }
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```

- Expected: all executions terminate after 2s
- Result: only some execute after 2s



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- Using quad-core CPU with hyper threading
- global has 8 threads in the thread pool



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- global has 8 threads in the thread pool
- executes tasks in batches of 8
- after 2s, 8 tasks print "completed"
- after 2s more, 8 more print "completed"
- sleep: all enter a timed waiting state



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- global has 8 threads in the thread pool
- executes tasks in batches of 8
- after 2s, 8 tasks print "completed"
- after 2s more, 8 more print "completed"
- sleep: all enter a timed waiting state
- if T1 waits for T10 to notify: blocks indefinitely

Lock-free programming

Avoiding syncrhonized with atomic variables



- atomic variable: memory location that supports complex linearizable operations
- ... i.e., appears to occur atomically
- write of a volatile operation: simple linearizable operation
- at least two reads and/or writes: complex linearizable operation

Avoiding syncrhonized with atomic variables



- atomic variable: memory location that supports complex linearizable operations
- ... i.e., appears to occur atomically
- write of a volatile operation: simple linearizable operation
- at least two reads and/or writes: complex linearizable operation
- java.util.concurrent.atomic supports some complex ones:
 - AtomicBoolean
 - AtomicInteger
 - AtomicLong
 - AtomicReference

Variation of Example 1 (getUniqueId)

```
import
    java.util.concurrent.atomic._
object AtomicUid extends App {
  private val uid =
    new AtomicLong(OL)
  def getUniqueId(): Long =
    uid.incrementAndGet()
  execute {
    log(s"Uid asynchronously: u$
        {getUniqueId()}")
  log(s"Got | a | unique | id: | $
      {getUniqueId()}")
```

Compare-And-Set (CAS) – the ♥ of complex linearizable operations



- CAS can be used to implement others:
 - getAndSet
 - decrementAndGet
 - addAndGet
- available in all atomic variables
- including AtomicReference[T]

Long-CAS conceptually equivalent to:

Ref-CAS conceptually equivalent to:

Using CAS



- Back to Example 1 (getUniqueld)
- Need to keep-on-trying
- Looks like busy-waiting, but it is much better
- Here: using (cheap) recursion instead of a loop

Lock-free programming – really?



- Lock-free programs: without locks (with synchronized)
- Achieved using atomic variables (and some re-trying)
- No locks, no deadlocks...

Lock-free programming - really?



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- No locks, no deadlocks...
- (almost):
 - lock-free ⇒ use atomic variables (for atomicity)
 - use atomic variables *>>* lock-free

Lock-free programming – really?



- Lock-free programs: without locks (with synchronized)
- Achieved using atomic variables (and some re-trying)
- No locks, no deadlocks...
- (almost):
 - lock-free ⇒ use atomic variables (for atomicity)
 - use atomic variables *>>* lock-free

```
object AtomicLock extends App {
  private val lock = new
      AtomicBoolean(false)
  def mySynchronized(body: =>Unit):
      Unit = \{
    while (!lock.compareAndSet(false,
        true)) {}
    try body finally lock.set(false)
  var count = 0
 for (i<- 0 until 10) execute {</pre>
      mySynchronized { count += 1 } }
 Thread. sleep (1000)
  log(s"Count_is:_\$count")
```

Lock-freedom definition



Lock-freedom

Given a set of threads and an operation OP.

OP is lock-free if at least one thread always completes OP after a finite number of steps, regardless of the speed at which different threads progress.

One more example: Concurrent filesystem



- Example 1: getUniqueId()
- Example 2: Logging Bank Transfers
- Example 3: Thread pool
- Example 4: Batman
- **Example 5:** Concurrent filesystem

Concurrent filesystem



Filesystem API

T1 is creating F:

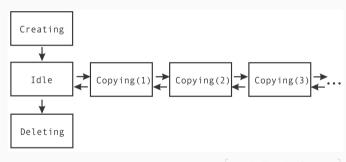
T2 cannot copy or delete F

T1 & T2 are copying F:

T3 cannot delete F

T1 is deleting F:

T2 cannot copy nor delete F



in "Learning Concurrent Programming in Scala", pg. 79

Concurrent filesystem – Scala data types



Filesystem API

```
T1 is creating F:
```

T2 cannot copy or delete F

T1 & T2 are copying F:

T3 cannot delete F

T1 is deleting F:

T2 cannot copy nor delete F

```
class Entry(val isDir: Boolean) {
  val state = new AtomicReference[State](new Idle)
}
sealed trait State
class Idle extends State
class Creating extends State
class Copying(val n: Int) extends State
class Deleting extends State
```

Deleting and Copying



Deleting: prepare (checks for permission) then delete (perform delete)

Copying: aquire (get permission); copy (perform action); then release (give permission)

Prepare for deleting



```
@tailrec
private def prepareForDelete(entry: Entry): Boolean = {
  val s0 = entry.state.get
  s0 match {
    case i: Idle =>
      if (entry.state.compareAndSet(s0, new Deleting)) true
      else prepareForDelete(entry)
    case c: Creating =>
      logMessage("File_currently_created,cannot_delete."); false
    case c: Copving =>
      logMessage("File | currently | copied , | cannot | delete."); false
    case d: Deleting =>
      false
```

logMessage: presented later - similar to log, but stores the log message

Bad copy - the ABA problem



"ABA" problem: two readings of the same value A lead to believe that B was never present (type of race condition)

Illustrated by a bad acquire-release for Copying, using a mutable n in:

Copying(var n: Int)

Bad code – acquire/release Copying



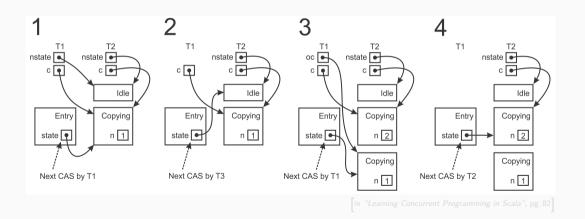
```
def releaseCopy(e: Entry): Copying = e.state.get match {
  case c: Copying =>
    val nstate = if (c.n == 1) new Idle else new Copying(c.n - 1)
    if (e.state.compareAndSet(c, nstate)) c
    else releaseCopy(e)
}
```

```
def acquireCopy(e: Entry, c: Copying) = e.state.get match {
  case i: Idle =>
    c.n = 1
    if (!e.state.compareAndSet(i, c)) acquireCopy(e, c)
  case oc: Copying =>
    c.n = oc.n + 1
    if (!e.state.compareAndSet(oc, c)) acquireCopy(e, c)
}
```

Optimization: reusing previous Copying if possible

Bad trace – T1 release (T2 starts rel.); T3 acquires; T1 acquire; T2 releases

wrongly



Some guidelines to avoid the ABA problem



- use fresh objects in AtomicReference
- use immutable objects in AtomicReference
- avoid re-assigning the same value to an atomic variable
- only increment or decrement values of numeric atomic variables (if possible)

Lazy values

Lazy values can cause deadlocks



- lazy values: initialized when read for the first time
- these should not depend-on/modify state (non-determinism)
- code in singleton objects: lazy execution
- under the hood: first write uses a lock – to ensure at most a thread initialises a lazy value
- stack overflow (sequential code)
 can become
 deadlock (concurrent code)

```
object LazyValsCreate extends App {
  var x = 5
  lazy val y = x+2
  execute {log(s"Wrk: \( \pu y \) \( \pu \) = 10
  log(s"Main: \( \pu y \) \( \pu \) = 7 or 12 in both cases
  Thread. sleep(500)
}
```

```
object LazyValsDeadlock extends App {
  object A { lazy val x: Int = B.y }
  object B { lazy val y: Int = A.x }
  execute { B.y }
  A.x
}
```

Concurrent (mutable) collections

Default mutable collections are not concurrent



- Naive code: arbitrarily returns different results and exceptions
- Corruption of the internal state
- Possible fixes:
 - immutable collections + atomic variables
 - mutable collections + synchronized
 - dedicated libraries

```
import scala.collection.
object CollectionsBad extends App {
  val buffer =
      mutable.ArrayBuffer[Int]()
  def asyncAdd(numbers: Seq[Int]) =
      execute {
    buffer ++= numbers
    log(s"buffer_=== $buffer")
  asyncAdd(0 until 10)
  asyncAdd(10 until 20)
  Thread.sleep (500)
```

Default mutable collections are not concurrent



- Naive code: arbitrarily returns different results and exceptions
- Corruption of the internal state
- Possible fixes:
 - immutable collections + atomic variables (does not scale)
 - mutable collections + synchronized (assuming collections do not block; may not scale)
 - dedicated libraries
 (far better performance and scalability)

```
import scala.collection.
object CollectionsBad extends App {
  val buffer =
      mutable.ArrayBuffer[Int]()
  def asyncAdd(numbers: Seq[Int]) =
      execute {
    buffer ++= numbers
    log(s"buffer_=== $buffer")
  asyncAdd(0 until 10)
  asyncAdd(10 until 20)
  Thread.sleep (500)
```

Some concurrent collections



- Concurrent queues
 - java.util.concurrent.BlockingQueue interface
 - ...ArrayBlockingQueue class (bounded)
 - ...LinkedBlockingQueue class (unbounded)
- Concurrent Sets and Maps
 - scala.collection.concurrent.Map trait
 - java.util.concurrent.ConcurrentHashMap class

BlockingQueue API



Operation	Exception	Special value	Timed	Blocking
Dequeue	remove(): T	poll(): T	poll(t: Long, u: TimeUnit): T	take(): T
Enqueue	add(x: T)	offer(x: T): Boolean	offer(x: T, t: Long, u: TimeUnit)	put(x: T)
Inspect	element: T	peek: T	N/A	N/A

in *"Learning Concurrent* Programming in Scala", pg. 90

Back to Example 5: logging in our concurrent filesystem



We will compile a queue of messages when logging messages in our file system

```
class FileSystem(...) {
    ...
    private val messages = new LinkedBlockingQueue[String]
    val logger = new Thread {
        setDaemon(true)
        override def run() = while (true) log(messages.take())
    }
    logger.start()
    def logMessage(msg: String): Unit = messages.offer(msg)
}
```

```
...

val fileSystem = new FileSystem(".") // to be defined later

fileSystem.logMessage("Testingulog!")
```

Note on iterators



- concurrentQueue.iterator
- can produce inconsistent results
- while traversing and modifying, the iterator can be updated
- (heavier) exceptions create a copy when producing an iterator

Example 5: files as a Map in our FileSystem



```
import scala.collection.convert.decorateAsScala._
import java.io.File
import org.apache.commons.io.FileUtils // needs "commons-io" in build.sbt
class FileSystem(val root: String) {
 val rootDir = new File(root)
 val files: concurrent.Map[String, Entry] =
    new ConcurrentHashMap().asScala
 for (f <- FileUtils.iterateFiles(rootDir, null, false).asScala)</pre>
    files.put(f.getName, new Entry(false))
  . . .
```

Deleting a file



Recall the prepareForDelete(entry)

```
def deleteFile(filename: String): Unit = {
 files.get(filename) match {
    case None =>
      logMessage(s"Pathu',$filename,udoesunotuexist!")
    case Some(entry) if entry.isDir =>
      logMessage(s"Pathu', filename, uisuaudirectory!")
    case Some(entry) => execute {
      if (prepareForDelete(entry))
        if (FileUtils.deleteQuietly(new File(filename)))
          files . remove (filename)
```

Some complex linearizable methods of concurrent Map



Signature	Description		
putIfAbsent (k: K, v: V): Option[V]	This atomically assigns the value v to the key k if k is not in the map. Otherwise, it returns the value associated with k .		
remove (k: K, v: V): Boolean	This atomically removes the key k if it is associated to the value equal to v and returns true if successful.		
replace (k: K, v: V): Option[V]	This atomically assigns the value ν to the key k and returns the value previously associated with k .		
replace (k: K, ov: V, nv: V): Boolean	This atomically assigns the key k to the value nv if k was previously associated with ov and returns true if successful.		

- These use "equals" instead of the reference (which CAS does)
- Avoid null as key or valye (often used as special values)

in "Learning Concurrent Programming in Scala", pg. 95

Methods +=, -=, put, update, get, apply, remove are (non-complex) linearizable

Wrapping up our Filesystem

(Example 5)

Copying in our Filesystem



Recall our broken aquireCopy/releaseCopy methods (ABA problem) - slide19

```
0tailrec
private def acquire(entry: Entry): Boolean = {
  val s0 = entry.state.get
  s0 match {
    case _: Creating | _: Deleting =>
      logMessage("File inaccessible, cannot copy."); false
    case i: Idle =>
      if (entry.state.compareAndSet(s0, new Copying(1))) true
      else acquire(entry)
    case c: Copying =>
      if (entry.state.compareAndSet(s0, new Copying(c.n+1))) true
      else acquire(entry)
```

Copying in our Filesystem



Same CAS retry-approach for releasing.

```
0tailrec
private def release(entry: Entry): Unit = {
 val s0 = entry.state.get
 s0 match {
    case c: Creating =>
      if (!entry.state.compareAndSet(s0, new Idle)) release(entry)
    case c: Copving =>
      val nstate = if (c.n == 1) new Idle else new Copying(c.n-1)
      if (!entry.state.compareAndSet(s0, nstate)) release(entry)
```

Copying in our Filesystem



Finally: wrapper for copying a file.

```
def copyFile(src: String, dest: String): Unit = {
 files.get(src) match {
    case Some(srcEntry) if !srcEntry.isDir => execute {
      if (acquire(srcEntry)) try {
        val destEntry = new Entry(isDir = false)
        destEntry.state.set(new Creating)
        if (files.putIfAbsent(dest, destEntry) == None) try {
          FileUtils.copyFile(new File(src), new File(dest))
        } finally release(destEntry)
      } finally release(srcEntry)
```

Creating and handling processes



- So far: run in a single JVM
- Now: run processes outside JVM
- Why:



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 - Some programs do not exist in Scala/Java



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- Now: run processes outside JVM
- Why:
 - Some programs do not exist in Scala/Java
 - Want to sandbox untrusted code
 - Performance (running independent code)
- Using the scala.sys.process package

Using processes - examples



```
import scala.sys.process._
object ProcessRun extends App {
  val command = "ls"
  val exitcode = command.! // run process (with side effects)
  log(s"command_exited_with_status_$exitcode") }
```

```
def lineCount(filename: String): Int = {
  val output = s"wc_u$filename".!! // run and retreive stdout
  output.trim.split("_u").head.toInt }
```

```
object ProcessAsync extends App {
  val lsProcess = "lsu-Ru/".run() // run and returns a Process object
  Thread.sleep(1000)
  log("Timeoutu-ukillinguls!")
  lsProcess.destroy() } // kill a slow process
```

https://www.scala-lang.org/api/2.13.x/scala/sys/process/ProcessBuilder.html

Wrapping up "concurrency blocks"



- executor.execute(...)
- lock-free programming with atomic variables
- av.compareAndSet(...)
- ABA problem
- Lazy values & "lazy" objects
- java.util.concurrent.BlockingQueue
- scala.collection.concurrent.Map
- weakly consistent iterators
- custom concurrent data structures

- Filesystem example
- Processes outside JVM

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Next

- Futures and Promises
- Data-Parallel Collections
- Reactive Programming (Concurrently)
- Software Transactional Memory
- Actors

Extra: Quick guide to the Future

Futures and Promises



- High-level asynchronous constructs
- Used very often by libraries that use concurrent tasks
- Future: read-only placeholder for the result of an ongoing computation
- Promise: writable, single-assignment container that completes a Future

Creating a Feature



```
object LazyValsCreate extends App {
  import scala.concurrent.
 import Future
  import ExecutionContext.Implicits.global
 val futFive = Future {
    log("Computation | started.")
    val result = {
      Thread. sleep (5000)
      5
    log("Computation__finally__finished.")
    result // : Future[Int]
```

Small notes

- Using our ExecutionContext.global
- import of Implicits.global ⇔ implicit val ec =
 ExecutionContext.global

Creating a Feature



```
object LazyValsCreate extends App {
  import scala.concurrent.
 import Future
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 val futFive = Future {
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    result // : Future[Int]
```

Small notes

- Using our ExecutionContext.global
- import of Implicits.global ⇔ implicit val ec = ExecutionContext.global

```
Future.apply[T]
  (body: => T)
  (implicit ec: ExecutionContext)
: Future[T]
```

Getting the values 1: onComplete



```
import scala.util.{Try,Success,Failure}
val futFive = Future { ... }
def getFive = futFive.onComplete {
  case Success(n) =>
    log(s"Got_number: ...$n")
  case Failure(exception) =>
    log(s"Failed..to..get..number:...
        $exception")
getFive
Thread.sleep(6000)
getFive // What to expect?
```

A future has 2 stages

- Not completed (ongoing computation)
- Completed (finished computation)
 - Success(res) (no errors)
 - Failure(err) (got errors)

Method 1: onComplete

```
futFive.onComplete[U]
  (callback: Try[Int] => U)
  (implicit exec: ...)
: Unit.
```

Getting the values 2: Await



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```
import scala.concurrent.duration.
. . .
val futFive = Future { ... }
def getFive = // : Int
  Await.result(futFive, Duration.inf)
def doneFive = // : Future[Int]
 Await.ready (futFive, 6.seconds)
val r1 = getFive // doneFive
val r2 = getFive // doneFive
Thread.sleep(6000)
val r3 = getFive // doneFive
```

A future has 2 stages

- Not completed (ongoing computation)
- Completed (finished computation)
 - Success(res) (no errors)
 - Failure(err) (got errors)

Method 2: Await.result/ready

Can throw TimeoutException

Polling a future



```
import scala.concurrent.duration.
. . .
val futFive = Future { ... }
def checkFive = // : Boolean
  futFive.isCompleted
def maybeFive = // : Option[Try[Int]]
  futFive.value
val r1 = checkFive // maybeFive
val r2 = checkFive // maybeFive
Thread. sleep (6000)
val r3 = checkFive // maybeFive
```

A future has 2 stages

- Not completed (ongoing computation)
- Completed (finished computation)
 - Success(res) (no errors)
 - Failure(err) (got errors)

Polling: isCompleted/value

How to check the status without blocking

Nelma Moreira & José Proença Extra: Quick guide to the Future 41 / 44

Composing futures



```
import scala.concurrent.duration._
. . .
val futTwo = Future { ... } // takes 2s
val futFive = Future { ... } // takes 5s
val r1 = futTwo.map(x => x*3)
val r2 = for (
    x <- futFive
    v \leftarrow r1
 ) vield
    x + y
log(s"a)_${r1.value}/${r2.value}")
Thread. sleep (3000)
log(s"b)_\${r1.value}/${r2.value}")
Thread. sleep (6000)
log(s"c)_{\square} {r1.value}/${r2.value}")
```

- read and update futures "in advance"
- map, flatMap, for-loops

Composing futures



```
import scala.concurrent.duration._
. . .
val futTwo = Future { ... } // takes 2s
val futFive = Future { ... } // takes 5s
val r1 = futTwo.map(x => x*3)
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log(s"a)_${r1.value}/${r2.value}")
Thread.sleep(3000)
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Thread. sleep (6000)
log(s"c)_{\square} {r1.value}/${r2.value}")
```

- read and update futures "in advance"
- map, flatMap, for-loops
- very common mechanism to run tasks in parallel
- used in many libraries

Promises: the other side of the Future



A promise p can...

- p.success(res)
- p.failure(exception)
- p.future

```
import scala.concurrent.{Future, Promise}
import scala.concurrent...Implicits.global
val prom = Promise[Int]()
val alice = Future { // execute/thread
  val res = {Thread.sleep(2000); 2}
  prom. success (res)
  log(s"gave,,'$res'")
  Thread.sleep(2000) // more work
val fut = prom.future
val bob = Future { // execute/thread
  Thread. sleep (3000)
  fut.foreach { r =>
    log("gotu'$r'")
```

Wrapping up "concurrency blocks" (again)



- executor.execute(...)
- lock-free programming with atomic variables
- av.compareAndSet(...)
- ABA problem
- Lazy values & "lazy" objects
- java.util.concurrent.BlockingQueue
- scala.collection.concurrent.Map
- weakly consistent iterators
- custom concurrent data structures

- Filesystem example
- Processes outside JVM
- Futures and Promises (overview)

Next

- Data-Parallel Collections
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