4. Basic building blocks of concurrency

Nelma Moreira & José Proença

Concurrent programming (CC3040) 2023/2024

CISTER - U.Porto, Porto, Portugal

https://fm-dcc.github.io/pc2324







Overview

We are here



Blocks of sequential code running concurrently and sharing memory:

- What is Scala?
- Concurrency in Java and its memory model
- Basic concurrency blocks and libraries
- Actor model

What we will see



- Tread pools: Executor and ExecutionContext
- Non-blocking synchronisation compare-and-set (CAS)
- Lazy (concurrent) values
- Concurrent collections
- Running OS processes

Existing thread pools in Scala

Executor interface



```
import scala.concurrent._
import java.util.concurrent.ForkJoinPool
object ExecutorsCreate extends App {
 val executor = new ForkJoinPool
  executor.execute(new Runnable {
   def run() = log("Thisutaskuisurunu
        asynchronously.")
 7)
 Thread.sleep(500) // not needed with
     fork:=false in SBT
```

- Executor: can start a new thread, an existing one, or the current one
- Abstracts from the management of threads
- ExecutorService: API that extends
 Executor with shutdown
 - executor.shutdown → executes all tasks and then stops working threads
 - executor.awaitTermination(...) →
 force termination if, after a given
 time, the tasks are not completed

Scala's ExecutionContext



```
import scala.concurrent._
object ExecutionContextGlobal extends App {
  val ectx = ExecutionContext.global
  ectx.execute(new Runnable {
    def run() = log("Running_on_the_execution_context.")
  })
  Thread.sleep(500)
}
```

```
object ExecutionContextCreate extends App {
  val pool = new forkjoin.ForkJoinPool(2)
  val ectx = ExecutionContext.fromExecutorService(pool)
  ectx.execute(new Runnable {
    def run() = log("Running_on_the_execution_context_again.")
  })
  Thread.sleep(500)
```

- scala.concurrent: has ExecutionContext
- Similar to Executor but more Scala specific
- often the ExecutionContext is used as implicit parameter
- global: default execution context (internally uses a ForkJoinPool)
- fromExecutorService: creates ExecutionContext
 from ExecutorService

Simplifying the execution



Similar to threads:

We now define execute

```
def execute(body: =>Unit) =
    ExecutionContext.global.execute(
     new Runnable { def run() = body }
// For example:
object ExecutionContextSleep extends App {
  for (i<- 0 until 32) execute {</pre>
    Thread. sleep (2000)
    log(s"Tasku$iucompleted.")
  Thread. sleep (10000)
```



```
object ExecutionContextSleep
     extends App {
  for (i<- 0 until 32) execute {
    Thread.sleep(2000)
    log(s"Tasku$iucompleted.")
  }
  Thread.sleep(10000)
}</pre>
```

- Expected: all executions terminate after 2s
- Result: only some execute after 2s



- Expected: all executions terminate after 2s
- Result: only some execute after 2s
- Using quad-core CPU with hyper threading
- global has 8 threads in the thread pool



```
object ExecutionContextSleep
        extends App {
  for (i<- 0 until 32) execute {
    Thread.sleep(2000)
    log(s"Tasku$iucompleted.")
  }
  Thread.sleep(10000)
}</pre>
```

- Expected: all executions terminate after 2s
- Result: only some execute after 2s
- Using quad-core CPU with hyper threading
- global has 8 threads in the thread pool
- executes tasks in batches of 8
- after 2s, 8 tasks print "completed"
- after 2s more, 8 more print "completed"
- sleep: all enter a timed waiting state



```
object ExecutionContextSleep
        extends App {
  for (i<- 0 until 32) execute {
    Thread.sleep(2000)
    log(s"Tasku$iucompleted.")
  }
  Thread.sleep(10000)
}</pre>
```

- Expected: all executions terminate after 2s
- Result: only some execute after 2s
- Using quad-core CPU with hyper threading
- global has 8 threads in the thread pool
- executes tasks in batches of 8
- after 2s, 8 tasks print "completed"
- after 2s more, 8 more print "completed"
- sleep: all enter a timed waiting state
- if T1 waits for T10 to notify: blocks indefinitely

Lock-free programming

Avoiding syncrhonized with atomic variables



- atomic variable: memory location that supports complex linearizable operations
- ... i.e., appears to occur atomically
- write of a volatile operation: simple linearizable operation
- at least two reads and/or writes: complex linearizable operation

Avoiding syncrhonized with atomic variables



- atomic variable: memory location that supports complex linearizable operations
- ... i.e., appears to occur atomically
- write of a volatile operation: simple linearizable operation
- at least two reads and/or writes: complex linearizable operation
- java.util.concurrent.atomic supports some complex ones:
 - AtomicBoolean
 - AtomicInteger
 - AtomicLong
 - AtomicReference

Variation of Example 1 (getUniqueId)

```
import
    java.util.concurrent.atomic._
object AtomicUid extends App {
  private val uid =
    new AtomicLong(OL)
  def getUniqueId(): Long =
    uid.incrementAndGet()
  execute {
    log(s"Uid asynchronously: u$
        {getUniqueId()}")
  log(s"Got | a | unique | id: | $
      {getUniqueId()}")
```

Compare-And-Set (CAS) – the ♥ of complex linearizable operations



- CAS can be used to implement others:
 - getAndSet
 - decrementAndGet
 - addAndGet
- available in all atomic variables
- including AtomicReference[T]

Long-CAS conceptually equivalent to:

Ref-CAS conceptually equivalent to:

Using CAS



- Back to Example 1 (getUniqueld)
- Need to keep-on-trying
- Looks like busy-waiting, but it is much better
- Here: using (cheap) recursion instead of a loop

Lock-free programming – really?



- Lock-free programs: without locks (with synchronized)
- Achieved using atomic variables (and some re-trying)
- No locks, no deadlocks...

Lock-free programming – really?



- Lock-free programs: without locks (with synchronized)
- Achieved using atomic variables (and some re-trying)
- No locks, no deadlocks...
- (almost):
 - lock-free ⇒ use atomic variables (for atomicity)
 - use atomic variables *>>* lock-free

Lock-free programming - really?



- Lock-free programs: without locks (with synchronized)
- Achieved using atomic variables (and some re-trying)
- No locks, no deadlocks...
- (almost):
 - lock-free ⇒ use atomic variables (for atomicity)
 - use atomic variables *>>* lock-free

```
object AtomicLock extends App {
  private val lock = new
      AtomicBoolean(false)
  def mySynchronized(body: =>Unit):
      Unit = \{
    while (!lock.compareAndSet(false,
        true)) {}
    try body finally lock.set(false)
  var count = 0
 for (i<- 0 until 10) execute {</pre>
      mySynchronized { count += 1 } }
 Thread. sleep (1000)
  log(s"Count_is:_\$count")
```

Lock-freedom definition



Lock-freedom

Given a set of threads and an operation OP.

OP is lock-free if at least one thread always completes the operation after a finite number of steps, regardless of the speed at which different threads progress.

One more example: Concurrent filesystem



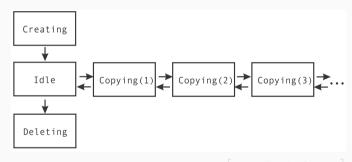
- Example 1: getUniqueId()
- Example 2: Logging Bank Transfers
- Example 3: Thread pool
- Example 4: Batman
- **Example 5:** Concurrent filesystem

Concurrent filesystem



Filesystem API

- T1 is creating F:
- T2 cannot copy or delete F
- T1 & T2 are copying F:
- T3 cannot delete F
- T1 is deleting F:
- T2 cannot copy nor delete F



in "Learning Concurrent Programming in Scala", pg. 79

Concurrent filesystem - Scala data types



Filesystem API

```
T1 is creating F:
```

T2 cannot copy or delete F

T1 & T2 are copying F:

T3 cannot delete F

T1 is deleting F:

T2 cannot copy nor delete F

```
class Entry(val isDir: Boolean) {
  val state = new AtomicReference[State](new Idle)
}
sealed trait State
class Idle extends State
class Creating extends State
class Copying(val n: Int) extends State
class Deleting extends State
```

Deleting and Copying



Deleting: prepare (checks for permission) then delete (perform delete)

Copying: aquire (get permission); copy (perform action); then release (give permission)

Prepare for deleting



```
@tailrec
private def prepareForDelete(entry: Entry): Boolean = {
  val s0 = entry.state.get
  s0 match {
    case i: Idle =>
      if (entry.state.compareAndSet(s0, new Deleting)) true
      else prepareForDelete(entry)
    case c: Creating =>
      logMessage("File | currently | created , | cannot | delete."); false
    case c: Copving =>
      logMessage("File | currently | copied , | cannot | delete."); false
    case d: Deleting =>
      false
```

logMessage: presented later - similar to log, but stores the log message

Bad copy - the ABA problem



"ABA" problem: two readings of the same value A lead to believe that B was never present (type of race condition)

Illustrated by a bad acquire-release for Copying, using a mutable n in:

Copying(var n: Int)

Bad code – acquire/release Copying



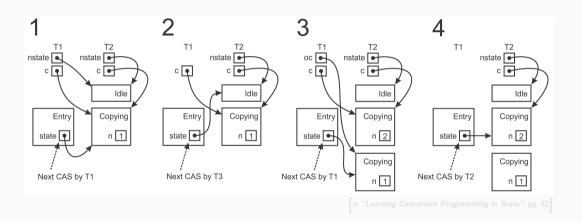
```
def releaseCopy(e: Entry): Copying = e.state.get match {
  case c: Copying =>
    val nstate = if (c.n == 1) new Idle else new Copying(c.n - 1)
    if (e.state.compareAndSet(c, nstate)) c
    else releaseCopy(e)
}
```

```
def acquireCopy(e: Entry, c: Copying) = e.state.get match {
  case i: Idle =>
    c.n = 1
    if (!e.state.compareAndSet(i, c)) acquireCopy(e, c)
  case oc: Copying =>
    c.n = oc.n + 1
    if (!e.state.compareAndSet(oc, c)) acquireCopy(e, c)
}
```

Optimization: reusing previous Copying if possible

Bad trace – T1&T2 release; T3&T1 aquire – (T2 is slow)





Some guidelines to avoid the ABA problem



- use fresh objects in AtomicReference
- use immutable objects in AtomicReference
- avoid re-assigning the same value to an atomic variable
- only increment or decrement values of numeric atomic variables (if possible)

Lazy values

Lazy values can cause deadlocks



- lazy values: initialized when read for the first time
- these should not depend-on/modify state (non-determinism)
- code in singleton objects: lazy execution
- under the hood: first write uses a lock – to ensure at most a thread initialises a lazy value
- stack overflow (sequential code)
 can become
 deadlock (concurrent code)

```
object LazyValsCreate extends App {
  var x = 5
  lazy val y = x+2
  execute {log(s"Wrk: "y" = "$y")}
  x = 10
  log(s"Main: "y" = "$y")
  // y = 7 or 12 in both cases
  Thread.sleep(500)
}
```

```
object LazyValsDeadlock extends App {
  object A { lazy val x: Int = B.y }
  object B { lazy val y: Int = A.x }
  execute { B.y }
  A.x
}
```

Concurrent (mutable) collections

Default mutable collections are not concurrent



- Naive code: arbitrarily returns different results and exceptions
- Corruption of the internal state
- Possible fixes:
 - immutable collections + atomic variables
 - mutable collections + synchronized
 - dedicated libraries

```
import scala.collection.
object CollectionsBad extends App {
 val buffer =
     mutable.ArrayBuffer[Int]()
 def asyncAdd(numbers: Seq[Int]) =
     execute {
   buffer ++= numbers
   asyncAdd(0 until 10)
 asyncAdd(10 until 20)
 Thread.sleep (500)
```

Default mutable collections are not concurrent



- Naive code: arbitrarily returns different results and exceptions
- Corruption of the internal state
- Possible fixes:
 - immutable collections + atomic variables (does not scale)
 - mutable collections + synchronized (assuming collections do not block; may not scale)
 - dedicated libraries
 (far better performance and scalability)

```
import scala.collection.
object CollectionsBad extends App {
 val buffer =
     mutable.ArrayBuffer[Int]()
 def asyncAdd(numbers: Seq[Int]) =
     execute {
   buffer ++= numbers
   asyncAdd(0 until 10)
 asyncAdd(10 until 20)
 Thread.sleep (500)
```

Some concurrent collections



- Concurrent queues
 - java.util.concurrent.BlockingQueue interface
 - ...ArrayBlockingQueue class (bounded)
 - ...LinkedBlockingQueue class (unbounded)
- Concurrent Sets and Maps
 - scala.collection.concurrent.Map trait
 - java.util.concurrent.ConcurrentHashMap class

BlockingQueue API



Operation	Exception	Special value	Timed	Blocking
Dequeue	remove(): T	poll(): T	poll(t: Long, u: TimeUnit): T	take(): T
Enqueue	add(x: T)	offer(x: T): Boolean	offer(x: T, t: Long, u: TimeUnit)	put(x: T)
Inspect	element: T	peek: T	N/A	N/A

in "Learning Concurrent Programming in Scala", pg. 90

Back to Example 5: logging in our concurrent filesystem



We will compile a queue of messages when logging messages in our file system

```
class FileSystem(...) {
    ...
    private val messages = new LinkedBlockingQueue[String]
    val logger = new Thread {
        setDaemon(true)
        override def run() = while (true) log(messages.take())
    }
    logger.start()
    def logMessage(msg: String): Unit = messages.offer(msg)
}
```

```
...
val fileSystem = new FileSystem(".") // to be defined later
fileSystem.logMessage("Testingulog!")
```

Note on iterators



- concurrentQueue.iterator
- can produce inconsistent results
- while traversing and modifying, the iterator can be updated
- (heavier) exceptions create a copy when producing an iterator

Example 5: files as a Map in our FileSystem



```
import scala.collection.convert.decorateAsScala._
import java.io.File
import org.apache.commons.io.FileUtils // needs "commons-io" in build.sbt
class FileSystem(val root: String) {
 val rootDir = new File(root)
 val files: concurrent.Map[String, Entry] =
    new ConcurrentHashMap().asScala
 for (f <- FileUtils.iterateFiles(rootDir, null, false).asScala)</pre>
    files.put(f.getName, new Entry(false))
  . . .
```

Deleting a file



Recall the prepareForDelete(entry)

```
def deleteFile(filename: String): Unit = {
  files.get(filename) match {
    case None =>
      logMessage(s"Pathu', filename, udoes not exist!")
    case Some(entry) if entry.isDir =>
      logMessage(s"Pathu', filename, uisuaudirectory!")
    case Some(entry) => execute {
      if (prepareForDelete(entry))
        if (FileUtils.deleteQuietly(new File(filename)))
          files . remove (filename)
```

Some complex linearizable methods of concurrent Map



Signature	Description		
putIfAbsent (k: K, v: V): Option[V]	This atomically assigns the value v to the key k if k is not in the map. Otherwise, it returns the value associated with k .		
remove (k: K, v: V): Boolean	This atomically removes the key k if it is associated to the value equal to v and returns true if successful.		
replace (k: K, v: V): Option[V]	This atomically assigns the value ν to the key k and returns the value previously associated with k .		
replace (k: K, ov: V, nv: V): Boolean	This atomically assigns the key k to the value nv if k was previously associated with ov and returns true if successful.		

- These use "equals" instead of the reference (which CAS does)
- Avoid null as key or valye (often used as special values)

in "Learning Concurrent Programming in Scala", pg. 95

Methods +=, -=, put, update, get, apply, remove are (non-complex) linearizable

Wrapping up our Filesystem

(Example 5)

Copying in our Filesystem



Recall our broken aquireCopy/releaseCopy methods (ABA problem) - slide19

```
0tailrec
private def acquire(entry: Entry): Boolean = {
  val s0 = entry.state.get
  s0 match {
    case _: Creating | _: Deleting =>
      logMessage("File inaccessible, cannot copy."); false
    case i: Idle =>
      if (entry.state.compareAndSet(s0, new Copying(1))) true
      else acquire(entry)
    case c: Copying =>
      if (entry.state.compareAndSet(s0, new Copying(c.n+1))) true
      else acquire(entry)
```

Copying in our Filesystem



Same CAS retry-approach for releasing.

```
0tailrec
private def release(entry: Entry): Unit = {
 val s0 = entry.state.get
 s0 match {
    case c: Creating =>
      if (!entry.state.compareAndSet(s0, new Idle)) release(entry)
    case c: Copving =>
      val nstate = if (c.n == 1) new Idle else new Copying(c.n-1)
      if (!entry.state.compareAndSet(s0, nstate)) release(entry)
```

Copying in our Filesystem



Finally: wrapper for copying a file.

```
def copyFile(src: String, dest: String): Unit = {
 files.get(src) match {
    case Some(srcEntry) if !srcEntry.isDir => execute {
      if (acquire(srcEntry)) try {
        val destEntry = new Entry(isDir = false)
        destEntry.state.set(new Creating)
        if (files.putIfAbsent(dest, destEntry) == None) try {
          FileUtils.copyFile(new File(src), new File(dest))
       } finally release(destEntry)
      } finally release(srcEntry)
```

Creating and handling processes



- So far: run in a single JVM
- Now: run processes outside JVM
- Why:



- So far: run in a single JVM
- Now: run processes outside JVM
- Why:
 - Some programs do not exist in Scala/Java



- So far: run in a single JVM
- Now: run processes outside JVM
- Why:
 - Some programs do not exist in Scala/Java
 - Want to sandbox untrusted code



- So far: run in a single JVM
- Now: run processes outside JVM
- Why:
 - Some programs do not exist in Scala/Java
 - Want to sandbox untrusted code
 - Performance (running independent code)



- So far: run in a single JVM
- Now: run processes outside JVM
- Why:
 - Some programs do not exist in Scala/Java
 - Want to sandbox untrusted code
 - Performance (running independent code)
- Using the scala.sys.process package

Using processes - examples



```
import scala.sys.process._
object ProcessRun extends App {
  val command = "ls"
  val exitcode = command.! // run process (with side effects)
  log(s"command_exited_with_status_$exitcode") }
```

```
def lineCount(filename: String): Int = {
  val output = s"wc_u$filename".!! // run and retreive stdout
  output.trim.split("_u").head.toInt }
```

```
object ProcessAsync extends App {
  val lsProcess = "lsu-Ru/".run() // run and returns a Process object
  Thread.sleep(1000)
  log("Timeoutu-ukillinguls!")
  lsProcess.destroy() } // kill a slow process
```

https://www.scala-lang.org/api/2.13.x/scala/sys/process/ProcessBuilder.html

Wrapping up "concurrency blocks"



- executor.execute(...)
- lock-free programming with atomic variables
- av.compareAndSet(...)
- ABA problem
- Lazy values & "lazy" objects
- java.util.concurrent.BlockingQueue
- scala.collection.concurrent.Map
- weakly consistent iterators
- custom concurrent data structures

- Filesystem example
- Processes outside JVM

Wrapping up "concurrency blocks"



- executor.execute(...)
- lock-free programming with atomic variables
- av.compareAndSet(...)
- ABA problem
- Lazy values & "lazy" objects
- java.util.concurrent.BlockingQueue
- scala.collection.concurrent.Map
- weakly consistent iterators
- custom concurrent data structures

- Filesystem example
- Processes outside JVM

Next

- Futures and Promises
- Data-Parallel Collections
- Reactive Programming (Concurrently)
- Software Transactional Memory
- Actors