Formal Simulation and Visualisation of Hybrid Programs in <u>LINCE</u>

Slides based on a presentation given in FMAS @ iFM 2024

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https://jose.proenca.org/publication/fmas-lince-2024/

José Proença

System Verification (CC4084) 2024/2025

CISTER - U.Porto, Porto, Portugal

https://fm-dcc.github.io/sv2425

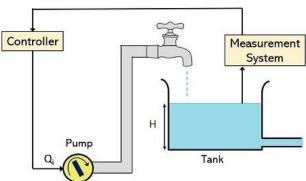
Hybrid systems





Computational devices that interact with physical environment

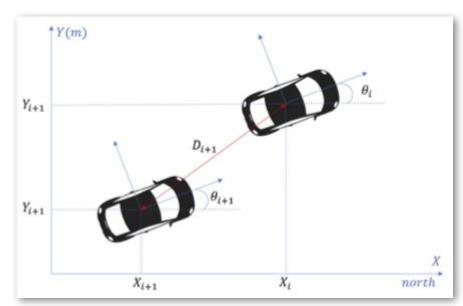




Another hybrid system

Platooning

- Acceleration (1D)
- Steering (2D)
- Failures



By Ênio Filho, Anis Koubâ, Ricardo Severino, Eduardo Tovar @ CISTER

Discrete behaviour

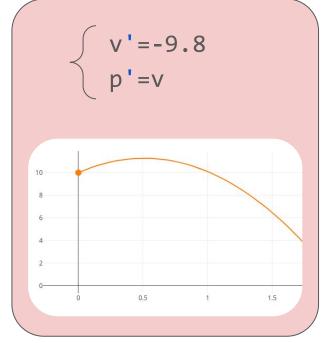


Continuous behaviour

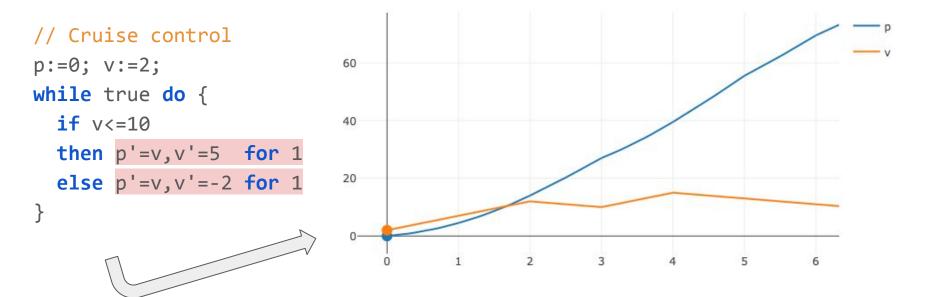
```
v := 10;
while v <= 10 do {</pre>
```





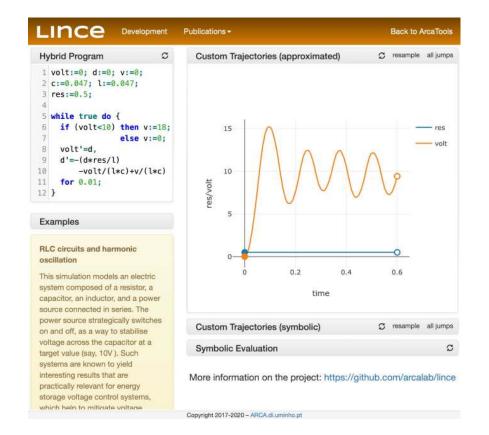


A cruise controller in LINCE



The expected output

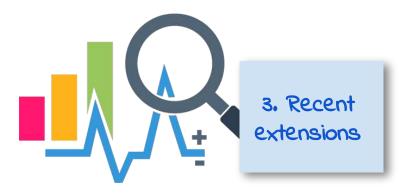
Why Lince?



- No installation
 - just a website (+ server)
- **Easy** to experiment
- **Simple** language
 - No need for complex frameworks
- **Precise** semantics
- Low effort to **extend**
 - new extensions
 - involve students
 - involve partners

What we will see

```
// Cruise control
p:=0; v:=2;
while true do {
  if v<=10
  then p'=v,v'=5 for 1
  else p'=v,v'=-2 for 1
}</pre>
```







Syntax

Discrete control

Continuous control

deq
$$\Rightarrow$$
 $x_1' = t_1$, (systems of differential equations) for t

$$t,s \Rightarrow real \mid real * x \mid t + s \mid x$$

(linear terms)

Syntax

Discrete control

Continuous control

```
t,s \Rightarrow real | real * x | t + s | x (linear terms)

e \Rightarrow e | f(e<sub>1</sub>,...,e<sub>1</sub>) (non-linear terms)
```

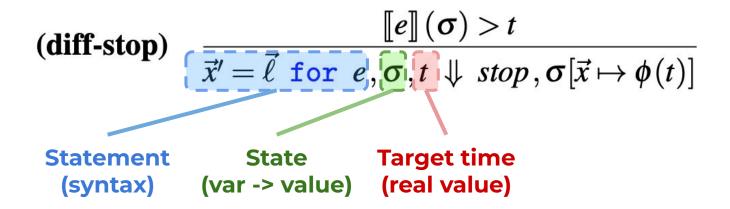
(diff-skip)
$$\frac{\llbracket e \rrbracket(\sigma) = t}{\vec{x}' = \vec{\ell} \text{ for } e, \sigma, t \Downarrow skip, \sigma[\vec{x} \mapsto \phi(t)]}$$

(diff-stop)
$$\frac{\llbracket e \rrbracket(\sigma) > t}{\vec{x}' = \vec{\ell} \text{ for } e, \sigma, t \Downarrow stop, \sigma[\vec{x} \mapsto \phi(t)]}$$
 (diff-err)
$$\frac{\llbracket e \rrbracket(\sigma) \text{ undefined}}{\vec{x}' = \vec{\ell} \text{ for } e, \sigma, t \Downarrow err}$$

(asg-skip)
$$\frac{\llbracket e \rrbracket(\sigma) \text{ defined}}{x := e, \sigma, 0 \Downarrow skip, \sigma[x \mapsto \llbracket e \rrbracket(\sigma)]}$$
 (asg-err)
$$\frac{\llbracket e \rrbracket(\sigma) \text{ undefined}}{x := e, \sigma, t \Downarrow err}$$

(seq-skip)
$$\frac{p, \sigma, t \Downarrow skip, \tau \qquad q, \tau, u \Downarrow v}{p; q, \sigma, t + u \Downarrow v}$$

(if-rules) ... (while-rules) ...



(diff-err)
$$\frac{\llbracket e \rrbracket(\sigma) \text{ undefined}}{\vec{x}' = \vec{\ell} \text{ for } e, \sigma, t \Downarrow err}$$

Partial functions (can throw errors)

(Small-Step) Semantics

```
(\mathbf{asg}^{\rightarrow})
                                                                                       x := t, \sigma, t \rightarrow skip, \sigma \nabla [t\sigma/x], t
                                                                  \bar{\mathtt{x}}' = \bar{\mathtt{u}} \, \mathsf{for} \, \mathtt{t} \,, \sigma, \mathtt{t} \, 	o \, stop \,, \sigma \triangledown [\phi_{\sigma}(\mathtt{t})/\bar{\mathtt{x}}] \,, \mathsf{0}
(diff-stop^{\rightarrow})
                                                                                                                                                                                                                           (if t < t\sigma)
(diff-skip^{\rightarrow})
                                                     \bar{\mathbf{x}}' = \bar{\mathbf{u}} \mathbf{for} \mathbf{t}, \sigma, \mathbf{t} \rightarrow skip, \sigma \nabla [\phi_{\sigma}(\mathbf{t}\sigma)/\bar{\mathbf{x}}], \mathbf{t} - (\mathbf{t}\sigma)
                                                                                                                                                                                                                           (if \ t \geqslant t\sigma)
(if-true^{\rightarrow})
                                                                                                                                                                                                                         (if b\sigma = \top)
                                                                          if b then p else q, \sigma, t \rightarrow p, \sigma, t
(if-false^{\rightarrow})
                                                                          if b then p else q, \sigma, t \rightarrow q, \sigma, t
                                                                                                                                                                                                                         (if \ b\sigma = \bot)
(\text{wh-true}^{\rightarrow})
                                                       while b do \{p\}, \sigma, t \rightarrow p; while b do \{p\}, \sigma, t
                                                                                                                                                                                                                        (if \ \mathbf{b}\sigma = \top)
(\text{wh-false}^{\rightarrow})
                                                                            while b do { p }, \sigma, t \rightarrow skip, \sigma, t
                                                                                                                                                                                                                         (if b\sigma = \bot)
                                                                                                                                           (\mathbf{seq\text{-}skip}^{\rightarrow}) \quad \frac{\mathsf{p}, \sigma, \mathsf{t} \rightarrow \mathit{skip}, \sigma', \mathsf{t}'}{\mathsf{p}; \mathsf{q}, \sigma, \mathsf{t} \rightarrow \mathsf{q}, \sigma', \mathsf{t}'}
(\mathbf{seq\text{-}stop}^{\rightarrow}) \quad \frac{\mathsf{p}, \sigma, \mathsf{t} \rightarrow \mathit{stop}, \sigma', \mathsf{t}'}{\mathsf{p}; \mathsf{q}, \sigma, \mathsf{t} \rightarrow \mathit{stop}, \sigma', \mathsf{t}'}
                          (\mathbf{seq}^{\rightarrow}) \quad \frac{\mathbf{p}, \sigma, \mathbf{t} \rightarrow \mathbf{p}', \sigma', \mathbf{t}'}{\mathbf{p}; \mathbf{q}, \sigma, \mathbf{t} \rightarrow \mathbf{p}'; \mathbf{q}, \sigma', \mathbf{t}'}
                                                                                                                                           (if p' \neq stop \ and p' \neq skip)
```

Exercises

1. Verify if the transitions exist, and justify.

```
x:=3, [], 4 \rightarrow stop, [x=3], 4
x:=5, [], 0 \rightarrow skip, [x=5], 0
x:=5;y:=6, [], 0 \rightarrow skip, [x=5,y=6], 0
```

3. Run the program

```
x:=3;x'=1 for 1; x'=-2 for 2, [], 2
```

2. Evolve twice each of these states

```
x:=3;x:=5;y:=7, [], 4

if (x>2)
then y:=0
else y:=1, [x=25], 0

while (x>2) do
  {x:=x-1} , [x=25], 0
```

Hybrid programs in LINCe

- Demo -

Run in our server

https://arcatools.org/lince

Internet browser

Download and run **your server**

https://github.com/arcalab/lince

- SageMath (http://www.sagemath.org/)
- SBT (https://www.scala-sbt.org)
- Java runtime
- Internet browser

Examples

```
// Simple composition
p:=0; v:=0;
p'=v,v'=-2 for 1;
p'=v,v'=-2 for 1;
```

```
// Initial values of the water tank
level := 5; drain := -1/2;
while true do {
  // keep level between 3..10
  if level<=3 then drain:= 1/2;</pre>
  else if level>=10 then drain:=-1/2;
  else skip;
 level'= drain, drain'=0 for 0.1;
```

Exercise

4. Re-implement the water tank, but calculating the optimal duration at each step (when to reach the maximum or minimum level)

```
// Initial values of the water tank
level := 5; drain := -1/2;
while true do {
  // keep level between 3..10
  if level<=3 then drain:= 1/2;</pre>
  else if level>=10 then drain:=-1/2;
 else
          skip;
 level'= drain, drain'=0 for 0.1;
```

More Examples

```
// Adaptive Cruise control
p:=0; v:=0; a:=5; b:=-2; // follower
pl:=50; vl:=10; al:=0; // leader
period:=1;
while true do {
   ???
}
```





More Examples



```
// Adaptive Cruise control
p:=0; v:=0; a:=5; b:=-2; // follower
pl:=50; vl:=10; al:=0;
                       // leader
period:=1;
while true do {
  if ((p+v*period+ a/2*period^2 <</pre>
       pl+vl*period+al/2*period^2) &&
      (((v-v)+(a-a))*period)^2 -
       4*(p-pl+(v-vl)*period +
        (a-a1)/2*period^2*(b-a1)/2 < 0)
  then p'=v,v'=a,pl'=vl,vl'=al for period;
  else p'=v,v'=b,pl'=vl,vl'=al for period;
```

More Examples

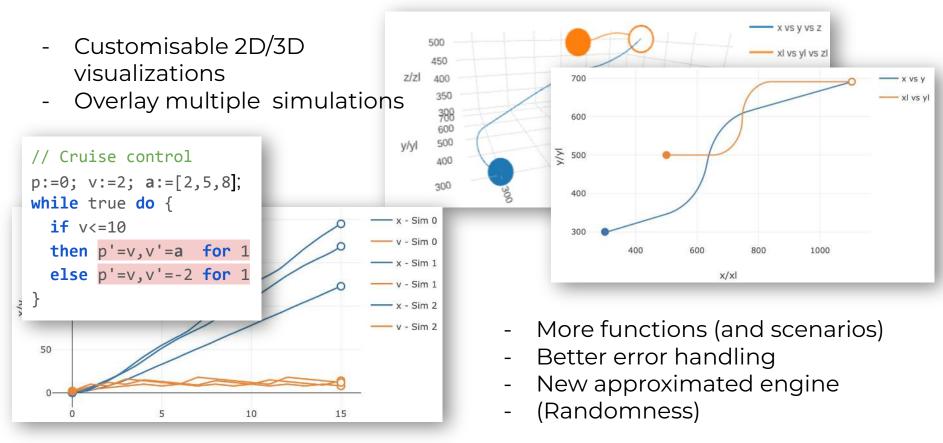
```
while true do {
    if v<=10
    then p'=v,v'=5 for 1
    else p'=v,v'=-2 for 1
// MORE:
// Automated braking system
// Pursuit games
// (2D, 3D views)
// Electric RLC circuit
// (many simulations, approx)
```

// Cruise control

p:=0; v:=2;

```
// Adaptive Cruise control
p:=0; v:=0; a:=5; b:=-2; // follower
pl:=50; vl:=10; al:=0; // leader
period:=1;
while true do {
 if ((p+v*period+ a/2*period^2 
      pl+vl*period+al/2*period^2) &&
     (((v-v)+(a-a))*period)^2 -
       4*(p-pl+(v-vl)*period +
       (a-a1)/2*period^2*(b-a1)/2 < 0)
  then p'=v,v'=a,pl'=vl,vl'=al for period;
 else p'=v,v'=b,pl'=vl,vl'=al for period;
```

Recent improvements to LINCE



Probabilities in LINCE

```
// Cruise control
p:=0; v:=unif(10,20); // initial value random in [10,20]
sim:=[0,1,2,3,4]; // 5 random simulations
while true do {
   if v<=10
    then p'=v,v'=5 for expn(0.5);
   else p'=v,v'=-2 for expn(0.5);
}</pre>
```

5. Implement the 3 examples in Uppaal with probabilities

https://fm-dcc.github.io/sv2425 /slides/7-mchains.pdf 6. Implement a probabilistic version of our Lamp example

https://fm-dcc.github.io/sv2425 /slides/5-TA-modelling.pdf slide 12

Conclusions and challenges



What to do with hybrid programs?

Generate software (e.g., controllers)

Predict physical behaviours

Simulate scenarios

Model checking (verify properties of scenarios)

Theorem proving (deductive reasoning to prove generic properties)





What to do with hybrid programs?

Generate software (e.g., controllers)

Predict physical behaviours





Simulate scenarios



Model checking (verify properties of scenarios)

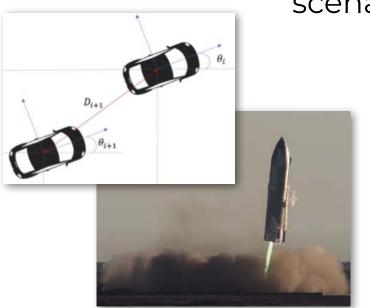




Theorem proving (deductive reasoning to prove generic properties)



Model more concrete scenario in LINCE



```
p:=0; v:=2;
while x<=50 do {
    // extend the language if needed
    x(t) = 5/2*t^2+10*t+x0 for 1
}
// Find the need for new extensions</pre>
```

Export program to another tool

```
// Cruise control
p:=0; v:=2;
while true do {
  if v<=10
    then p'=v,v'=5 for 1
  else p'=v,v'=-2 for 1
}

KeymaeraX.org
```

```
// Cruise control
p:=0; v:=2;
while true do {...}
check [v>=10?...] p>5
```

Show some property?

Import program from another tool

```
// Cruise control
p:=0; v:=2;
while true do {
   if v<=10
    then p'=v,v'=5 for 1
   else p'=v,v'=-2 for 1
}

KeymaeraX.org
```

```
// Cruise control
p:=0; v:=2;
while true do {...}
check [v>=10?...] p>5
```

Useful subset?

Checking properties: using logics for runtime verification

```
// Cruise control
p:=0; v:=2;
while true do {
   if v<=10
    then p'=v,v'=5 for 1
   else p'=v,v'=-2 for 1

assert <formula>
}
monitor <formula>
```

Metric temporal logic

RMTL-J3 (three-valued restricted metric temporal logic with durations)

(Interval) duration logic

Running many simulations: Statistical Analysis

```
// Cruise control
p:=0; v:=2; a:=Unif[2..8];
while true do {
   if v<=10
    then p'=v,v'=5 for Exp[0..1]
   else p'=v,v'=-2 for Exp[0..1]
}
check Prob(conf=95%, v<=10)</pre>
```

Running multiple times

Aim at statistical relevance

Control number of runs, range of a run, confidence, error margin, etc.

Improve framework (technology)









