

# CS146 - Homework

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## 1.5

```
public static int NumOfOnes(int n)
{
    if(n == 0) //Base Case
        return 0;

    if(n%2 == 0) //n is an even number
        return NumOfOnes(n/2);
    else
        return 1 + NumOfOnes(n/2);
}
```

## 1.7a.

To Prove:  $\log x < x$  for all  $x > 0$

Consider the function  $g(x) = x - \log(x)$  for any real value of  $x > 0$

$g(x) = x - \frac{\ln(x)}{\ln(2)}$  re-written using change of base

$g'(x) = 1 - \frac{1}{\ln(2)*x}$  by taking the first derivative

Critical point is at  $x = \frac{1}{\ln(2)}$ . The function  $g(x)$  has a negative slope before this point and a positive slope after this point. Thus  $x = \frac{1}{\ln(2)}$  gives an absolute minimum for the function  $g(x)$  on  $x > 0$ . Since  $g(\frac{1}{\ln(2)}) > 0$  the value of  $g(x)$  before and after  $x = \frac{1}{\ln(2)}$  must be positive.

Therefore,  $\log(x)$  must be less than  $x$  for  $x - \log(x)$  to be  $> 0$  which is what was to be proved.

## 1.7b.

To Prove:  $\log(A^B) = (B)\log(A)$  Note: Using base 2 for logarithm

Let  $\log A = C$  equation 1\*

Which is  $2^C = A$  by definition of logarithm

Raise each side to the  $B$  power for some integer  $B$ .

$$(2^C)^B = A^B$$

Re-write L.H.S using rule of exponent multiplication.

$$2^{CB} = A^B$$

Take the  $\log$  of both sides.

$$\log 2^{CB} = \log A^B$$

Using logarithm rule change the L.H.S.

$$CB = \log A^B$$

Substitute the value of  $C$  from equation 1\* on the L.H.S

$$\log AB = \log A^B$$

Re-arrange the product on the L.H.S

$$(B)\log A = \log A^B$$

Which is what was to be proved.

## 1.8a.

$$S_{\infty} = \sum_{i=0}^{\infty} \frac{1}{4^i} = \frac{1}{4^0} + \frac{1}{4^1} + \frac{1}{4^2} + \dots + \frac{1}{4^{\infty}}$$

$$S_{\infty} = \sum_{i=0}^{\infty} \frac{1}{4^i} = 1 + \frac{1}{4} + \frac{1}{16} + \dots + \frac{1}{4^{\infty}}$$

Formula for the sum of a convergent geometric series:  $\frac{\text{first term}}{1 - \text{common ratio}}$  or  $\frac{a_1}{1-r}$

$$S_{\infty} = \frac{1}{1 - \frac{1}{4}} = \frac{1}{\frac{3}{4}} = \frac{4}{3}$$

## 1.8b.

$$S_{\infty} = \sum_{i=0}^{\infty} \frac{i}{4^i}$$

$$S_1 = 0 + \frac{1}{4} = \frac{1}{4}$$

$$S_5 = 0 + \frac{1}{4} + \frac{1}{8} + \frac{3}{64} + \frac{1}{64} + \frac{5}{1024} = 0.44238$$

$$S_8 = 0 + \frac{1}{4} + \frac{1}{8} + \frac{3}{64} + \frac{1}{64} + \frac{5}{1024} + \frac{6}{4096} + \frac{7}{16384} + \frac{1}{8192} = 0.44439$$

By considering what number the partial sums are approaching:

$$S_{\infty} = \sum_{i=0}^{\infty} \frac{i}{4^i} = 0.444444444444 = \frac{4}{9}$$

## 1.9

Estimate  $\sum_{i=\frac{n}{2}}^n \frac{1}{i}$

$$S_1 = \frac{1}{\frac{1}{2}} = 2$$

$$S_2 = \frac{1}{\frac{2}{2}} + \frac{1}{\frac{2}{2}} = 1 + 1 = 2$$

$$S_3 = \frac{1}{\frac{3}{2}} + \frac{1}{\frac{3}{2}} + \frac{1}{\frac{3}{2}} = \frac{2}{3} + \frac{2}{3} + \frac{2}{3} = 2$$

$$S_4 = \frac{1}{\frac{4}{2}} + \frac{1}{\frac{4}{2}} + \frac{1}{\frac{4}{2}} + \frac{1}{\frac{4}{2}} = \frac{1}{2} + \frac{1}{2} + \frac{1}{2} + \frac{1}{2} = 2$$

$$S_n = 2$$

## 1.11a.

$\sum_{i=1}^{n-2} F_i = F_n - 2$  where  $F_i$  is defined as Fibonacci number.

Proof will be by induction.

Base Case:  $n = 3$   $\sum_{i=1}^{3-2} F_i = 1 = F_3 - 2 = 3 - 2$  base case is good

Assume true for all integers from base case up to  $k$ .  $\sum_{i=1}^{k-2} F_i = F_k - 2$  (Inductive Hypothesis)

Prove also true for  $k + 1$   $\sum_{i=1}^{(k+1)-2} F_i = F_{k+1} - 2$

By definition of Fibonacci number is  $F_{k+1} = F_k + F_{k-1}$

Change R.H.S. to  $= F_k + F_{k-1} - 2$

Rearrange R.H.S  $= F_k - 2 + F_{k-1}$

On R.H.S. substitute for  $F_k - 2$  from the inductive hypothesis  $= \sum_{i=1}^{k-2} F_i + F_{k-1}$

$= \sum_{i=1}^{k-2} F_i + \sum_{i=1}^{k+1} F_{k-1} = \sum_{i=1}^{(k+1)-2} F_i$  which is the same as the L.H.S. End Proof.

## 1.11b.

Prove  $F_n < \phi^n$  with  $\phi = \frac{1+\sqrt{5}}{2}$

In mathematics  $\frac{1 + \sqrt{5}}{2}$  is defined as the golden ratio.

The Fibonacci numbers are related to the golden ratio and its conjugate by the equation  $F_i = \frac{\phi^i - \phi^i}{\sqrt{5}}$  (this is noted in the book Introduction To Algorithms, by Cormen, pg.59)

Using this fact, the inequality can be re-written as  $\frac{\phi^i - \phi^i}{\sqrt{5}} < \phi^i$

Or more simply  $\frac{\phi^i}{\sqrt{5}} < \phi^i$

Which makes it obvious that the L.H.S. is less than the right since it is being divided by  $\sqrt{5}$