



On the Syntactic Expressiveness of Variability Languages

FOSD 2025 | Benjamin Moosherr, Paul Bittner, Thomas Thüm | 2025-03-26

Variability Languages

i.e. variability in the solution space

preprocessor

```
#ifdef Feature1  
    Line1  
#else  
    Line2  
#endif
```

binary choice calculus (2CC)

```
file < Feature1 ⟨ Line1 <>, Line2 <> ⟩ >
```

Motivation

variability languages

- binary choice calculus (2CC)
- n-ary choice calculus (NCC)
- core choice calculus (CCC)
- algebraic decision trees (ADT)
- feature structure trees (FST)
- option calculus (OC)
- clown-and-own (CaO)

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question: which one should you to choose for your project?

Previous Work

```
choice-idempotency :  $\forall \{i : \text{Size}\} \{A : \text{Domain}\} \{D : \text{Dimension}\} \{e : \text{BCC } i \ A\}$ 
-----
 $\rightarrow \text{BCC } [\_] \vdash D \ (e \ , \ e) \approx e$ 
choice-idempotency  $\{i\} \{A\} \{D\} \{e\} = \text{extensionality } (\lambda \ c \rightarrow$ 
 $[ \ D \ (e \ , \ e) \ ] \ c \quad \quad \quad \equiv ()$ 
 $[ \ \text{if } (c \ D) \ \text{then } e \ \text{else } e \ ] \ c \quad \quad \quad \equiv \{ \text{Eq.cong } (\text{flip } [\_] \ c) \ (\text{if-idemp } (c \ D)) \}$ 
 $[ \ e \ ] \ c \quad \quad \quad \equiv ()$ 
-----
- 8.9k BCC.lagda.md Agda

completeness-by-expressiveness :  $\forall \{L_1 \ L_2 : \text{VarLang}\} \{C_1 \ C_2 : \text{Conflang}\} \{S_1 : \text{Set}\}$ 
 $\rightarrow \text{Complete } L_1 \ C_1 \ S_1$ 
 $\rightarrow L_2 \ , \ S_2 \text{ is-as-expressive-as } L_1 \ , \ S_1$ 
-----
 $\rightarrow \text{Complete } L_2 \ C_2 \ S_2$ 
completeness-by-expressiveness  $\{L_1\} \{L_2\} \{_\_ \} \{_\_ \} \{S_1\} \{S_2\} \text{ encode-in-} L_1 \ L_1\text{-to-} L_2$ 
let { }
-----
- 11k Completeness.lagda.md Agda
```

Formal Languages for Solution-Space Variability

Paul Bittner, Jeffrey Young, Parisa Ataei, Alexander Schultheiß, Eric Walkingshaw, Leopoldo Teixeira, Thomas Thüm | FOSD 2023



FOSD 2023 (Paul Bittner)

¹ Bittner et al. "On the Expressive Power of Languages for Static Variability". In: *Proc. ACM Program. Lang.* 8.OOPSLA2 (Oct. 2024)



On the Expressive Power of Languages for Static Variability

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Variability permeates software development to satisfy ever-changing requirements and mass-customization needs. A prime example is the Linux kernel, which employs the C preprocessor to specify a set of related but distinct kernel variants. To study, analyze, and verify variational software, several formal languages have been proposed. For example, the choice calculus has been successfully applied for type checking and symbolic execution of configurable software, while other formalisms have been used for variational model checking, change impact analysis, among other use cases. Yet, these languages have not been formally compared, hence, little is known about their relationships. Crucially, it is unclear to what extent one language subsumes another, how research results from one language can be applied to other languages, and which language is suitable for which purpose or domain. In this paper, we propose a formal framework to compare the expressive power of languages for static (i.e. compile-time) variability. By establishing a common semantic domain to capture a widely used notion of explicit variability, we can formulate the basic, yet to date neglected, properties of soundness, completeness, and expressiveness for variability languages. We then prove the (in)soundness and (in)completeness of a range of existing languages, and relate their ability to express the same variational systems. We implement our framework as an extensible open source Agda library in which proofs act as correct comparators between languages or differentiating algorithms. We find different levels of expressiveness as well as complete and incomplete languages w.r.t. our unified semantic domain, with the choice calculus being among the most expressive languages.

CCS Concepts • Software and its engineering → Formal language definitions; Software configuration management and version control systems; Software product lines
Additional Key Words and Phrases: variation, software product lines, configuration, language semantics

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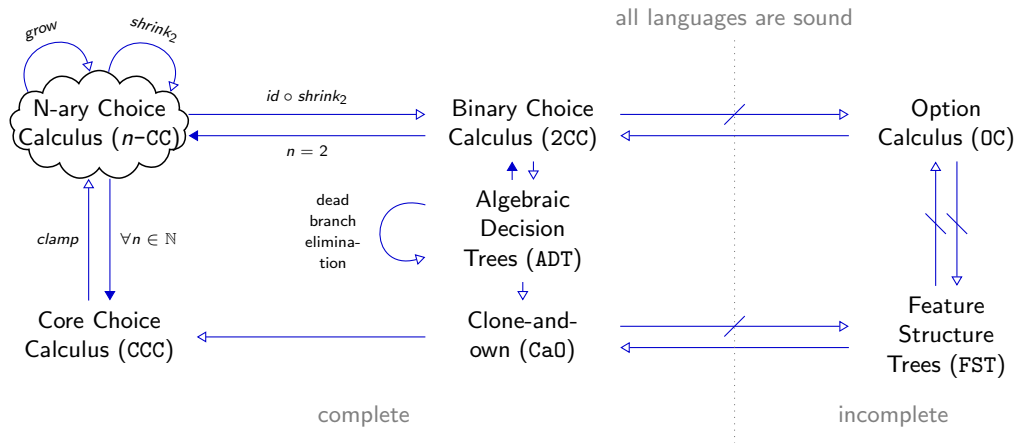
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OOPSLA 2024¹

Previous Work¹



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Motivation

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- option calculus (OC)
- clown-and-own (CaO)

question: which one should you to choose for your project?

goal for this thesis:

compare these variability languages regarding their expressiveness/compression

Example of Size Differences

#ifdef Feature1

Line1

#else

Line2

#endif

Example of Size Differences

#ifdef Feature1

Line1

#else

Line2

#endif

#ifdef Feature2

Line3

#else

Line4

#endif

Example of Size Differences

#ifdef Feature1

Line1

#else

Line2

#endif

#ifdef Feature2

Line3

#else

Line4

#endif

#ifdef Feature3

Line5

#else

Line6

#endif

Example of Size Differences

#ifdef Feature1

Line1

#else

Line2

#endif

#ifdef Feature2

Line3

#else

Line4

#endif

#ifdef Feature3

Line5

#else

Line6

#endif

#ifdef Feature4

Line7

#else

Line8

#endif

Example of Size Differences

2CC

```
file <
  Feature1 < Line1 <>, Line2 <> >
>
```

ADT

```
Feature1 < file < Line1 <> >, file
< Line2 <> > >
```

Example of Size Differences

2CC

```
file <
  Feature1 < Line1 <>, Line2 <> >,
  Feature2 < Line3 <>, Line4 <> >
>
```

ADT

```
Feature1 < Feature2 < file < Line1
<>, Line3 <> >, file < Line1 <>,
Line4 <> > >, Feature2 < file <
Line2 <>, Line3 <> >, file <
Line2 <>, Line4 <> > > >
```

Example of Size Differences

2CC

```
file <
  Feature1 < Line1 <>, Line2 <> >,
  Feature2 < Line3 <>, Line4 <> >,
  Feature3 < Line5 <>, Line6 <> >
>
```

ADT

```
Feature1 < Feature2 < Feature3 < file <
  Line1 <>, Line3 <>, Line5 <> >, file <
  Line1 <>, Line3 <>, Line6 <> > >,
  Feature3 < file < Line1 <>, Line4 <>,
  Line5 <> >, file < Line1 <>, Line4 <>,
  Line6 <> > > >, Feature2 < Feature3 < file
  < Line2 <>, Line3 <>, Line5 <> >, file <
  Line2 <>, Line3 <>, Line6 <> > >,
  Feature3 < file < Line2 <>, Line4 <>,
  Line5 <> >, file < Line2 <>, Line4 <>,
  Line6 <> > > > > >
```

Example of Size Differences

2CC

```
file <
  Feature1 < Line1 <>, Line2 <> >,
  Feature2 < Line3 <>, Line4 <> >,
  Feature3 < Line5 <>, Line6 <> >,
  Feature4 < Line7 <>, Line8 <> >
>
```

ADT

```
Feature1 < Feature2 < Feature3 < Feature4 < file < Line1 <>,
Line3 <>, Line5 <>, Line7 <> >, file < Line1 <>, Line3
<>, Line5 <>, Line8 <> > >, Feature4 < file < Line1 <>,
Line3 <>, Line6 <>, Line7 <> >, file < Line1 <>, Line3
<>, Line6 <>, Line8 <> > > >, Feature3 < Feature4 < file <
Line1 <>, Line4 <>, Line5 <>, Line7 <> >, file < Line1
<>, Line4 <>, Line5 <>, Line8 <> > >, Feature4 < file <
Line1 <>, Line4 <>, Line6 <>, Line7 <> >, file < Line1
<>, Line4 <>, Line6 <>, Line8 <> > > > >, Feature2 <
Feature3 < Feature4 < file < Line2 <>, Line3 <>, Line5 <>,
Line7 <> >, file < Line2 <>, Line3 <>, Line5 <>, Line8 <>
> >, Feature4 < file < Line2 <>, Line3 <>, Line6 <>, Line7
<> >, file < Line2 <>, Line3 <>, Line6 <>, Line8 <> > >
>, Feature3 < Feature4 < file < Line2 <>, Line4 <>, Line5
<>, Line7 <> >, file < Line2 <>, Line4 <>, Line5 <>,
Line8 <> > >, Feature4 < file < Line2 <>, Line4 <>, Line6
<>, Line7 <> >, file < Line2 <>, Line4 <>, Line6 <>,
Line8 <> > > > > > >
```

Definition of Size

- derived from the syntax
- count productions/constructors

Definition of Size

- derived from the syntax
- count productions/constructors
- example: 2CC

- grammar:

$$\begin{aligned} 2CC &:= a \prec 2CC, 2CC, \dots, 2CC \succ && \text{(artifact)} \\ &| \textcolor{brown}{D} \langle 2CC, 2CC \rangle && \text{(choice)} \end{aligned}$$

Definition of Size

- derived from the syntax
- count productions/constructors
- example: 2CC

- grammar:

$2CC := a \prec 2CC, 2CC, \dots, 2CC \succ$ (artifact)
| $\mathbf{D} \langle 2CC, 2CC \rangle$ (choice)

- expression:

$2CC$

size: 0

Definition of Size

- derived from the syntax
- count productions/constructors
- example: 2CC

- grammar:

$2CC := a \prec 2CC, 2CC, \dots, 2CC \succ$ (artifact)
| $\mathbf{D} \langle 2CC, 2CC \rangle$ (choice)

- expression:

file $\prec 2CC \succ$
size: 1

Definition of Size

- derived from the syntax
- count productions/constructors
- example: 2CC
 - grammar:
$$2CC := a \prec 2CC, 2CC, \dots, 2CC \succ \quad (\text{artifact})$$
$$| \text{D} \langle 2CC, 2CC \rangle \quad (\text{choice})$$
 - expression:
$$\text{file} \prec \text{Feature1} \langle 2CC, 2CC \rangle \succ$$
size: 2

Definition of Size

- derived from the syntax
- count productions/constructors
- example: 2CC
 - grammar:
$$2CC := a \prec 2CC, 2CC, \dots, 2CC \succ \quad (\text{artifact})$$
$$| \text{D} \langle 2CC, 2CC \rangle \quad (\text{choice})$$
 - expression:
$$\text{file} \prec \text{Feature1} \langle \text{Line1} \prec \succ, 2CC \rangle \succ$$
size: 3

Definition of Size

- derived from the syntax
- count productions/constructors
- example: 2CC
 - grammar:
$$2CC := a \prec 2CC, 2CC, \dots, 2CC \succ \quad (\text{artifact})$$
$$| \text{D} \langle 2CC, 2CC \rangle \quad (\text{choice})$$
 - expression:
$$\text{file} \prec \text{Feature1} \langle \text{Line1} \prec \succ, \text{Line2} \prec \succ \rangle \succ$$
size: 4

Size comparison

- should be independent of small differences
- inspiration: Big-O notation
- used for comparing the growth of functions
- $f = O(g) \quad := \quad \exists c, n_0 \in \mathbb{N} : \forall n \geq n_0 : f(n) \leq c \cdot g(n)$
- this requires enumeration

Formalization

$$f = O(g) \quad := \quad \exists c, n_0 \in \mathbb{N} : \forall n \geq n_0 : f(n) \leq c \cdot g(n)$$

Formalization

$$\begin{aligned} f = O(g) &:= \exists c, n_0 \in \mathbb{N} : \forall n \geq n_0 : f(n) \leq c \cdot g(n) \\ &\Leftrightarrow \exists c \in \mathbb{N} : \forall n \in \mathbb{N} : f(n) \leq c \cdot g(n) \end{aligned} \quad (\text{if } g > 0)$$

Formalization

$$\begin{aligned} f = O(g) &:= \exists c, n_0 \in \mathbb{N} : \forall n \geq n_0 : f(n) \leq c \cdot g(n) \\ &\Leftrightarrow \exists c \in \mathbb{N} : \forall n \in \mathbb{N} : f(n) \leq c \cdot g(n) \quad (\text{if } g > 0) \end{aligned}$$

$$L_1 \leq_s L_2 \quad := \quad \exists c \in \mathbb{N} : \forall e_2 \in L_2 : \exists e_1 \in L_1 : e_1 \cong e_2 \wedge \text{size}(e_1) \leq c \cdot \text{size}(e_2)$$

Formalization

$$\begin{aligned} f = O(g) &:= \exists c, n_0 \in \mathbb{N} : \forall n \geq n_0 : f(n) \leq c \cdot g(n) \\ &\Leftrightarrow \exists c \in \mathbb{N} : \forall n \in \mathbb{N} : f(n) \leq c \cdot g(n) \quad (\text{if } g > 0) \end{aligned}$$

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$$\begin{aligned} L_1 \not\leq_s L_2 &:= \neg(L_1 \leq_s L_2) \\ &\Leftrightarrow \forall c \in \mathbb{N} : \exists e_2 \in L_2 : \forall e_1 \in L_1 : e_1 \cong e_2 \rightarrow \text{size}(e_1) > c \cdot \text{size}(e_2) \end{aligned}$$

Example of Size Differences

$$ADT \not\leq_s 2CC \Leftrightarrow \forall c \in \mathbb{N} : \exists e_2 \in 2CC : \forall e_1 \in ADT : e_1 \cong e_2 \rightarrow \text{size}(e_1) > c \cdot \text{size}(e_2)$$

2CC
size = 4

ADT
size = 5

```
file <
  Feature1 < Line1 <>, Line2 <> >
>
Feature1 < file < Line1 <> >, file
< Line2 <> > >
```

Example of Size Differences

$$ADT \not\leq_s 2CC \Leftrightarrow \forall c \in \mathbb{N} : \exists e_2 \in 2CC : \forall e_1 \in ADT : e_1 \cong e_2 \rightarrow size(e_1) > c \cdot size(e_2)$$

2CC
size = 7

ADT
size = 27

```
file <
  Feature1 < Line1 <>, Line2 <> >,
  Feature2 < Line3 <>, Line4 <> >
>
```

```
Feature1 < Feature2 < file < Line1
<>, Line3 <> >, file < Line1 <>,
Line4 <> > >, Feature2 < file <
Line2 <>, Line3 <> >, file <
Line2 <>, Line4 <> > > >
```

Example of Size Differences

$$ADT \not\leq_s 2CC \Leftrightarrow \forall c \in \mathbb{N} : \exists e_2 \in 2CC : \forall e_1 \in ADT : e_1 \cong e_2 \rightarrow size(e_1) > c \cdot size(e_2)$$

2CC
size = 10

```
file <
  Feature1 < Line1 <>, Line2 <> >,
  Feature2 < Line3 <>, Line4 <> >,
  Feature3 < Line5 <>, Line6 <> >
>
```

ADT
size = 39

```
Feature1 < Feature2 < Feature3 < file <
  Line1 <>, Line3 <>, Line5 <> >, file <
  Line1 <>, Line3 <>, Line6 <> > >,
  Feature3 < file < Line1 <>, Line4 <>,
  Line5 <> >, file < Line1 <>, Line4 <>,
  Line6 <> > > >, Feature2 < Feature3 < file
  < Line2 <>, Line3 <>, Line5 <> >, file <
  Line2 <>, Line3 <>, Line6 <> > >,
  Feature3 < file < Line2 <>, Line4 <>,
  Line5 <> >, file < Line2 <>, Line4 <>,
  Line6 <> > > > > >
```

Example of Size Differences

$$ADT \not\leq_s 2CC \Leftrightarrow \forall c \in \mathbb{N} : \exists e_2 \in 2CC : \forall e_1 \in ADT : e_1 \cong e_2 \rightarrow size(e_1) > c \cdot size(e_2)$$

2CC
size = 13

ADT
size = 95

file <

```
Feature1 < Line1 <>, Line2 <> >,
Feature2 < Line3 <>, Line4 <> >,
Feature3 < Line5 <>, Line6 <> >,
Feature4 < Line7 <>, Line8 <> >
```

>

```
Feature1 < Feature2 < Feature3 < Feature4 < file < Line1 <>,
Line3 <>, Line5 <>, Line7 <> >, file < Line1 <>, Line3
<>, Line5 <>, Line8 <> > >, Feature4 < file < Line1 <>,
Line3 <>, Line6 <>, Line7 <> >, file < Line1 <>, Line3
<>, Line6 <>, Line8 <> > > >, Feature3 < Feature4 < file <
Line1 <>, Line4 <>, Line5 <>, Line7 <> >, file < Line1
<>, Line4 <>, Line5 <>, Line8 <> > >, Feature4 < file <
Line1 <>, Line4 <>, Line6 <>, Line7 <> >, file < Line1
<>, Line4 <>, Line6 <>, Line8 <> > > > >, Feature2 <
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> >, Feature3 < Feature4 < file < Line2 <>, Line4 <>, Line5
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Line8 <> > >, Feature4 < file < Line2 <>, Line4 <>, Line6
<>, Line7 <> >, file < Line2 <>, Line4 <>, Line6 <>,
Line8 <> > > > > > >
```

Formalization

$$\begin{aligned} f = O(g) &:= \exists c, n_0 \in \mathbb{N} : \forall n \geq n_0 : f(n) \leq c \cdot g(n) \\ &\Leftrightarrow \exists c \in \mathbb{N} : \forall n \in \mathbb{N} : f(n) \leq c \cdot g(n) \quad (\text{if } g > 0) \end{aligned}$$

$$L_1 \leq_s L_2 \quad := \quad \exists c \in \mathbb{N} : \forall e_2 \in L_2 : \exists e_1 \in L_1 : e_1 \cong e_2 \wedge \text{size}(e_1) \leq c \cdot \text{size}(e_2)$$

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$$L_1 =_s L_2 \quad := \quad L_1 \leq_s L_2 \wedge L_2 \leq_s L_1$$

$$L_1 <_s L_2 \quad := \quad L_1 \leq_s L_2 \wedge L_2 \not\leq_s L_1$$

Properties

- \leq_s is a partial order
- $<_s$ is a strict partial order
- $=_s$ is an equivalence relation
- \leq_s and $<_s$ are not total (i.e., there are uncomparable languages)

All of these are already verified using Agda (in our Vatrás library)

Current Progress

- $2CC =_s CCC \leq_s NCC$
- $2CC <_s ADT$
- $2CC \not\leq_s FST$
- $FST \not\leq_s 2CC$ (because FST is incomplete)

legend:

2CC	binary choice calculus
NCC	n-ary choice calculus
CCC	core choice calculus
ADT	algebraic decision trees
FST	feature structure trees

All of these are already verified using Agda (in our Vatrás library)

Current Progress

legend:

2CC	binary choice calculus
NCC	n-ary choice calculus
CCC	core choice calculus
ADT	algebraic decision trees
FST	feature structure trees

- $2CC =_s CCC \leq_s NCC$
- $2CC <_s ADT$
- $2CC \not\leq_s FST$
- $FST \not\leq_s 2CC$ (because FST is incomplete)
 $\forall c \in \mathbb{N} : \exists e_2 \in 2CC : \forall e_1 \in FST : e_1 \cong e_2 \rightarrow size(e_1) > c \cdot size(e_2)$

All of these are already verified using Agda (in our Vatrás library)

Future Work

open questions:

- How to handle incomplete languages?

$$\forall c \in \mathbb{N} : \exists e_2 \in L_2 : (\forall e_1 \in L_1 : e_1 \cong e_2 \rightarrow \text{size}(e_1) > c \cdot \text{size}(e_2)) \\ \wedge (\exists e_1 \in L_1 : e_1 \cong e_2)$$

breaks transitivity

- Do I need to consider feature model? (i.e., do feature models make a difference?)

future work:

- explore relationships between variability languages using \leq_s
- explore the influence of sharing primitives (i.e., subexpression reuse)

Questions?

open questions:

- How to handle incomplete languages?

$$\forall c \in \mathbb{N} : \exists e_2 \in L_2 : (\forall e_1 \in L_1 : e_1 \cong e_2 \rightarrow \text{size}(e_1) > c \cdot \text{size}(e_2))$$

$$\wedge (\exists e_1 \in L_1 : e_1 \cong e_2)$$

breaks transitivity

- Do I need to consider feature model? (i.e., do feature models make a difference?)

recap of the main definition:

$$L_1 \leq_s L_2 \quad := \quad \exists c \in \mathbb{N} : \forall e_2 \in L_2 : \exists e_1 \in L_1 : e_1 \cong e_2 \wedge \text{size}(e_1) \leq c \cdot \text{size}(e_2)$$

$$L_1 \not\leq_s L_2 \quad := \quad \forall c \in \mathbb{N} : \exists e_2 \in L_2 : \forall e_1 \in L_1 : e_1 \cong e_2 \rightarrow \text{size}(e_1) > c \cdot \text{size}(e_2)$$