R Textbook Companion for Operations Research: An Introduction by Hamdy A Taha¹

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Book Description

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R numbering policy used in this document and the relation to the above book.

Exa Example (Solved example)

Eqn Equation (Particular equation of the above book)

For example, Exa 3.51 means solved example 3.51 of this book. Sec 2.3 means an R code whose theory is explained in Section 2.3 of the book.

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Chapter 2

Modelling with Linear Programming

R code Exa 2.2.1 Graphical Solution to LP Problems

```
1 ##Chapter 2: Modelling with Linear Programming
2 \# Example 2-1 : Page 16
4 #To plot the line, we have to consider them as
      equation instead of inequality and express/
5 #them in terms of x2 :
7 #Constraint 1 : 6 * x1 + 4 * x2 \le 24
                  x^2 = (24 - 6 * x^1)/4
9 \text{ con1} \leftarrow \text{function}(x1) (24 - 6 * x1)/4
10 plot (con1, xlab = "x1", ylab = "x2", xlim = c(0,7),
       ylim = c(0,7), col = "red",
         main = "Example 2-1", yaxs= "i", xaxs = "i")
11
12 #xlab & ylab : x and y label respectively
13 #xlim
                 : limits of x value on the plot
14 #ylim
                 : limits of y value on the plot
15 #col
                 : color of the line
                : Title of the plot
16 #main
17 #yaxs & xaxs : the style of axis interval
```

```
calculation to be used by R. The default
18 #value is a 4% gap at each end of axis
20 \# Constraint 2 : x1 + 2 * x2 <=6
21 \# \text{Con} 2 : x2 = (6 - x1)/2
22 \operatorname{con2} \leftarrow \operatorname{function}(x1) (6 - x1)/2
23 plot (con2, add=T, xlim = c(0,7), ylim = c(0,7), col
       = "blue")
24 #add
                  : adds to an existing plot
25
26
27 \# Constraint 3 : -x1 + x2 <= 1
28 #Con3
            : x2 = (1 + x1)
29 \text{ con} 3 \leftarrow \text{function}(x1) (1 + x1)
30 plot (con3, add=T, xlim = c(0,7), ylim = c(0,7), col
       = "green")
31
32 \# Constraint 4 : x2 <= 2
33 #Con4
                  : x2 = 2
34 \text{ con} 4 \leftarrow \text{function}(x1) (2 + 0*x1)
35 plot (con4, add=T, xlim = c(0,7), ylim = c(0,7), col
       = "green")
                   : horizontal line at y=2
36 #h
37
38 #Points of intersections of constraints : (0,1)
      (1,2),(2,2),(2,1.5),(4,0)
39 points(c(0,1,2,3,4),c(1,2,2,1.5,0))
40
41 #Add a shaded area
42 polygon(c(0,1,2,3,4,0),c(1,2,2,1.5,0,0), col = rgb
      (0.48, 0.46, 0.46, 0.5),
43
            border=NA)
                   : option to add border to the shaded
44 #border
      area
45
46 #Adding "solution space" text to the shaded area at
      (2,1)
47 text(2,1, "Solution \nSpace")
```

R code Exa 2.2.2 Graphical solution of LP for Diet Problem

```
1 ##Chapter 2 : Modelling with Linear Programming
2 \# Example 2-2 : Page 24
4 #To plot the line, we have to consider them as
      equation instead of inequality and express/
5 #them in terms of x2 :
7 \# Constraint 1 : x1 + x2 <= 800
8 #Con1
                : x2 = (800 - x1)
9 \text{ con1} \leftarrow \text{function}(x1) (800 - x1)
10 plot (con1, xlab = "x1", ylab = "x2", xlim = c
      (0,1500), ylim = c(0,1500), col = "red",
         main = "Example 2-2", yaxs= "i", xaxs = "i")
12 #xlab & ylab : x and y label respectively
13 #xlim
                 : limits of x value on the plot
14 #ylim
                 : limits of y value on the plot
15 #col
                 : color of the line
16 #main
                 : Title of the plot
17 #yaxs & xaxs : the style of axis interval
      calculation to be used by R. The default
```

```
18 #value is a 4% gap at each end of axis
19
20 \# Constraint 2 : 0.21 * x1 - 0.3 * x2 <= 0
21 \# \text{Con} 2 : x2 = (0.21 * x1) / 0.3
22 \text{ con2} \leftarrow \text{function}(x1) (0.21 * x1)/0.3
23 plot (con2, add=T, xlim = c(0,1500), ylim = c
      (0,1500), col = "blue")
24 #add
                  : adds to an existing plot
25
26
27 \# Constraint 3 : 0.03 * x1 - 0.01 * x2 >= 0
28 \# \text{Con} 3 : x2 = (0.03 * x1) / 0.01
29 con3 \leftarrow function(x1) (0.03 * x1)/0.01
30 plot (con3, add=T, xlim = c(0,1500), ylim = c
      (0,1500), col = "green")
31
32 #Points of intersections of constraints : (0,1)
      (1,2),(2,2),(2,1.5),(4,0)
33 points(c(470.6,200),c(329.4,600))
34
35 #Add a shaded area
36 polygon(c(470.6,200,500,15000/7),c
      (329.4,600,1500,1500), col = rgb(0.48, 0.46,
      0.46, 0.5),
           border=NA)
37
38 #border : option to add border to the shaded area
39
40 #Adding "solution space" text to the shaded area at
      (750, 1000)
41 text (750, 1000, "Solution \nSpace")
42
43 ##Objective function : Min 0.3 *x1 + 0.9 *x2
44 # Minimum objective is 437.64, Therefore 0.3 *x1 +
      0.9 * x2 = 437.64
45 # Obj : x2 = (437.64 - 0.3 * x1)/0.9
46 Obj \leftarrow function(x1) (437.64 - 0.3 * x1)/0.9
47 plot (Obj, add=T, xlim = c(0,1500), ylim = c(0,1500)
      , lty =2 )
```

```
48 #lty : option to set the type of line,2 for dashed line
49
50 text(1000,250,"x1 = 470.6 \nx2 = 329.4 \n z=437.64")
```

R code Exa 2.4.1 LP solution for Bank Loan Problem

```
1 ##Chapter 2: Modelling with Linear Programming
2 ##Example 4-1 : Page 35
4 # Objective function : Max (0.126 * x1 + 0.1209 * x2 +
       0.1164 * x3 + 0.11875 * x4 + 0.098 * x5
                         - (0.1 * x1 + 0.07 * x2 + 0.03)
5 #
      * x3 + 0.05 * x4 + 0.02 * x5
6 a <- c
      (0.126-0.1, 0.1209-0.07, 0.1164-0.03, 0.11875-0.05, 0.098-0.02)
8 \# Constraint 1 : x1 + x2 + x3 + x4 + x5 \le 12
9 C1 \leftarrow c(1,1,1,1,1)
10 bc1<- 12
11
12 \# Constraint 2 : 0.4*x1 + 0.4*x2 + 0.4*x3 - 0.6*x4 -
       0.6 * x5 <=0
13 C2 \leftarrow c(0.4,0.4,0.4,-0.6,-0.6)
14 bc2<-0
15
16 # Constraint 3 : 0.5 * x1 + 0.5 * x2 - 0.5 * x3 <=0
17 C3 \leftarrow c(0.5,0.5,-0.5,0,0)
18 bc3<-0
19
20 \# Constraint 4 : 0.06 * x1 + 0.03 * x2 - 0.01 * x3
      +0.01 * x4 -0.02 * x5 <=0
21 \quad C4 \leftarrow c(0.06, 0.03, -0.01, 0.01, -0.02)
22 bc4<-0
```

R code Exa 2.4.2 LP solution for Single period Production Model

```
1 ##Chapter 2: Modelling with Linear Programming
2 \# \text{Example } 4-2 : \text{Page } 41
4 # Objective function : Max (30 * x1 + 40 * x2 + 20 *
      x3 + 10 * x4
5 #
                          - (15 * s1 + 20 * s2 + 10 * s3)
     + 8 * s4
6 a \leftarrow c(30,40,20,10,-15,-20,-10,-8)
8 # Constraint 1 : 0.3 * x1 + 0.3 * x2 + 0.25 * x3 +
      0.15 * x4 <= 1000
9 C1 \leftarrow c(0.3,0.3,0.25,0.15,0,0,0,0)
10 bc1<- 1000
11
12 # Constraint 2 : 0.25 * x1 + 0.35 * x2 + 0.3 * x3 +
      0.1 * x4 <= 1000
13 C2 \leftarrow c(0.25,0.35,0.3,0.1,0,0,0,0)
14 bc2<-1000
15
16 # Constraint 3 : 0.45 * x1 + 0.5 * x2 + 0.4 * x3 +
      0.22 * x4 <= 1000
17 C3 \leftarrow c(0.45,0.5,0.4,0.22,0,0,0,0)
18 bc3<-1000
19
20 # Constraint 4 : 0.15 * x1 + 0.15 * x2 + 0.1 * x3 +
      0.05 * x4 <=0
```

```
21 C4 \leftarrow c(0.15,0.15,0.1,0.05,0,0,0,0)
22 bc4<-1000
23
24 \# Constraint 5 : x1 + s1 = 800
25 C5 \leftarrow c(1,0,0,0,1,0,0,0)
26 bc5<-800
27
28 \# Constraint 6 : x2 + s2 = 750
29 \quad C6 \leftarrow c(0,1,0,0,0,1,0,0)
30 bc6<-750
31
32 \# Constraint 7: x3 + s3 = 600
33 C7 \leftarrow c(0,0,1,0,0,0,1,0)
34 bc7<-600
35
36 \# Constraint 8 : x4 + s4 = 500
37 C8 \leftarrow c(0,0,0,1,0,0,0,1)
38 bc8<-500
39
40 library ("boot")
41
42 \text{ simplex(a, A1 = rbind(C1, C2, C3, C4), b1 = c(bc1, C2, C3, C4)}
      bc2, bc3, bc4),
             A2=NULL, b2=NULL, A3 = rbind(C5, C6, C7, C8),
43
                b3 = c(bc5, bc6, bc7, bc8), maxi = TRUE)
```

R code Exa 2.4.3 LP solution for Multi period Production Inventory Model

```
7
8 \# Constraint 1 : x1 - I1 = 100
9 C1 <- c(1,0,0,0,0,0,-1,0,0,0,0,0)
10 bc1<- 100
11
12 \# Constraint 2 : I1 + x2 - I2 = 250
13 C2 \leftarrow c(0,1,0,0,0,0,1,-1,0,0,0,0)
14 bc2<-250
15
16 \# Constraint 3 : I2 + x3 - I3 = 190
17 C3 \leftarrow c(0,0,1,0,0,0,0,1,-1,0,0,0)
18 bc3<-190
19
20 \# Constraint 4 : I3 + x4 - I4 = 140
21 C4 \leftarrow c(0,0,0,1,0,0,0,0,1,-1,0,0)
22 \text{ bc4} < -140
23
24 \# Constraint 5 : I4 + x5 - I5 = 220
25 C5 \leftarrow c(0,0,0,0,1,0,0,0,0,1,-1,0)
26 bc5<-220
27
28 \# Constraint 6 : I5 + x6 = 110
29 \quad C6 \leftarrow c(0,0,0,0,0,1,0,0,0,0,1,0)
30 bc6<-110
31
32 library("boot")
33
34 simplex(a, A1=NULL, b1=NULL, A2=NULL, b2=NULL, A3 =
      rbind(C1, C2, C3, C4,C5,C6),
            b3 = c(bc1, bc2, bc3, bc4, bc5, bc6), maxi = F
35
               )
```

R code Exa 2.4.4 LP solution for Multi period Production smoothing Model

```
1 ##Chapter 2: Modelling with Linear Programming
```

```
2 \# Example 4-4 : Page 43
3
4 # Objective function : Min 50(I1 + I2 + I3 + I4) +
      200(Sm1 + Sm2 + Sm3 + Sm4) + 400(Sp1 + Sp2 + Sp3 + Sp4)
5 a <- c
      (0,0,0,0,50,50,50,50,200,200,200,200,400,400,400,400)
7 \# Constraint 1 : 10*x1 - I1 = 400
8 C1 \leftarrow c(10,0,0,0,-1,0,0,0,0,0,0,0,0,0,0,0)
9 bc1<- 400
10
11 # Constraint 2 : I1 + 10 * x^2 - I^2 = 600
12 C2 \leftarrow c(0,10,0,0,1,-1,0,0,0,0,0,0,0,0,0,0)
13 bc2<-600
14
15 # Constraint 3 : I2 + 10 * x3 - I3 = 400
16 C3 \leftarrow c(0,0,10,0,0,0,1,-1,0,0,0,0,0,0,0,0)
17 bc3<-400
18
19 # Constraint 4 : I3 + 10 * x4 = 500
20 C4 \leftarrow c(0,0,0,10,0,0,0,1,0,0,0,0,0,0,0)
21 bc4<-500
22
23 \# Constraint 5 : x1 - Sm1 + Sp1 = 0
24 C5 \leftarrow c(1,0,0,0,0,0,0,0,-1,0,0,0,1,0,0,0)
25 bc5<-0
26
27 \# Constraint 6 : x2 -x1 -Sm2 +Sp2 = 0
28 C6 \leftarrow c(-1,1,0,0,0,0,0,0,0,-1,0,0,0,1,0,0)
29 bc6<-0
30
31 # Constraint 7 : x3 - x2 - Sm3 + Sp3 = 0
32 C7 \leftarrow c(0,-1,1,0,0,0,0,0,0,0,-1,0,0,0,1,0)
33 bc7<-0
34
35 # Constraint 8 : x4 - x3 - Sm4 + Sp4 = 0
36 \quad C8 \leftarrow c(0,0,-1,1,0,0,0,0,0,0,0,-1,0,0,0,1)
```

R code Exa 2.4.5 LP solution for Bus Scheduling Problem

```
1 ##Chapter 2: Modelling with Linear Programming
2 \# \text{Example } 4-5 : \text{Page } 50
3
4 # Objective function : Min x1 + x2 + x3 + x4 + x5 +
      x6
5 \text{ a} \leftarrow c(1,1,1,1,1,1)
7 \# Constraint 1 : x1 + x6 >= 4
8 C1 \leftarrow c(1,0,0,0,0,1)
9 bc1<- 4
10
11 # Constraint 2 : x1 + x2 >= 8
12 C2 \leftarrow c(1,1,0,0,0,0)
13 bc2<-8
14
15 \# Constraint 3 : x2 + x3 >= 10
16 \quad C3 \leftarrow c(0,1,1,0,0,0)
17 bc3<-10
18
19 # Constraint 4 : x3 + x4 >= 7
20 \quad C4 \leftarrow c(0,0,1,1,0,0)
21 \text{ bc4} < -7
22
23 \# Constraint 5 : x4 + x5 >= 12
```

R code Exa 2.4.6 LP solution for Urban Renewal Model Problem

```
1 ##Chapter 2: Modelling with Linear Programming
   2 ##Example 4-6 : Page 52
   3
   4 # Objective function : Max 1000 * x1 + 1900 * x2 +
                               2700 * x3 + 3400 * x4
   5 \text{ a} \leftarrow c(1000, 1900, 2700, 3400, 0)
   7 # Constraint 1 : 0.18 * x1 + 0.28 * x2 + 0.4 * x3 +
                                0.5 * x4 - 0.2125 * x5 <= 0
   8 C1 \leftarrow c(0.18, 0.28, 0.4, 0.5, -0.2125)
   9 bc1<- 0
10
11 \# Constraint 2 : x5 <= 300
12 \quad C2 \leftarrow c(0,0,0,0,1)
13 bc2<-300
14
15 # Constraint 3 : -0.8 * x1 + 0.2 * x2 + 0.2 * x3 +
                               0.2 * x4 <= 0
```

```
16 \quad C3 \leftarrow c(-0.8, 0.2, 0.2, 0.2, 0)
17 bc3<-0
18
19 # Constraint 4 : 0.1 * x1 - 0.9 * x2 + 0.1 * x3 +
      0.1 * x4 <= 0
20 \text{ C4} \leftarrow c(0.1, -0.9, 0.1, 0.1, 0)
21 \text{ bc4} < -7
22
23 # Constraint 5 : 0.25 * x1 + 0.25 * x2 - 0.75 * x3 -
       0.75 * x4 <= 0
24 C5 \leftarrow c(0.25,0.25,-0.75,-0.75,0)
25 bc5<-0
26
27 \# Constraint 6 : 50 * x1 + 70 * x2 + 130 * x3 + 160
      * x4 + 2* x5 <= 15000
28 	ext{ C6 } \leftarrow c(50, 70, 130, 160, 2)
29 bc6<-15000
30
31 #to install the lpSolve package, run the following
      command
32 #install.packages("lpSolve")
33 library ("lpSolve")
34
35 solution = lp("max", a, rbind(C1, C2, C3, C4, C5, C6
      ), rep("<=",6),
36
                   c(bc1, bc2, bc3, bc4, bc5, bc6))
37 solution
38 solution $ solution
```

R code Exa 2.4.7 LP solution for Crude Oil Refining and Gasoline Blending Problem

```
1 ##Chapter 2 : Modelling with Linear Programming
2 ##Example 4-7 : Page 57
```

```
4 # Objective function : Max 6.7 * (x11 + x21) + 7.2 * (
      x12 + x22 + 8.1 * (x13 + x23)
5 \text{ a} \leftarrow c(6.7, 7.2, 8.1, 6.7, 7.2, 8.1)
7 # Constraint 1 : 5 *(x11 + x12 + x13) + 10 * (x21 +
      x22 + x23) <= 1500000
8 C1 \leftarrow c(5, 5, 5, 10, 10, 10)
9 bc1<- 1500000
10
11 # Constraint 2 : 2 * (x21 + x22 + x23) \le 200000
12 C2 \leftarrow c(0, 0, 0, 2, 2, 2)
13 bc2<-200000
14
15 \# Constraint 3 : (x11 + x21) \le 50000
16 \quad C3 \leftarrow c(1,0,0,1,0,0)
17 bc3<-50000
18
19 # Constraint 4 : (x12 + x22) \le 30000
20 \text{ C4} \leftarrow c(0, 1, 0, 0, 1, 0)
21 bc4<-30000
22
23 \# Constraint 5 : (x13 + x23) <= 40000
24 \text{ C5} \leftarrow c(0, 0, 1, 0, 0, 1)
25 bc5<-40000
26
27 \# Constraint 6 : 5 * x11 - 11 * x21 <= 0
28 \quad C6 \quad \leftarrow \quad c(5,0,0,-11,0,0)
29 bc6<-0
30
31 \# Constraint 7 : 7* x12 - 9 * x22 <= 0
32 \quad C7 \leftarrow c(0, 7, 0, 0, -9, 0)
33 bc7<-0
34 \# Constraint 8 : 10 * x12 - 6 * x22 <= 0
35 \text{ C8} \leftarrow c(0, 0, 10, 0, 0, -6)
36 bc8<-0
37
38 #to install the lpSolve package, run the following
      command
```

Chapter 3

The Simplex Method and Sensitivity Analysis

R code Exa 3.2.1 Graphical solution of LP with 2 variables

```
1 ##Chapter 3: The Simplex Method and Sensitivity
     Analysis
2 ##Example 2-1 : Page 83
4 #To plot the line, we have to consider them as
     equation instead of inequality and express/
5 #them in terms of x2 :
7 #Constraint 1 : 2 * x1 + x2 \le 4
           : x2 = (4 - 2 * x1)
8 #Con1
9 \text{ con1} \leftarrow \text{function}(x1) (4 - 2 * x1)
10 plot (con1, xlab = "x1", ylab = "x2", xlim = c(0,6),
      ylim = c(0,5), col = "red",
         main = "Example 2-1", yaxs= "i", xaxs = "i")
11
12 #xlab & ylab : x and y label respectively
13 #xlim
                : limits of x value on the plot
                : limits of y value on the plot
14 #ylim
                 : color of the line
15 #col
16 #main
                 : Title of the plot
```

```
17 #yaxs & xaxs : the style of axis interval
      calculation to be used by R. The default
18 #value is a 4% gap at each end of axis
19
20 \# Constraint 2 : x1 + 2 * x2 <= 5
21 \# \text{Con} 2 : x2 = (5 - x1)/2
22 \text{ con2} \leftarrow \text{function}(x1) (5 - x1)/2
23 plot (con2, add=T, xlim = c(0,6), ylim = c(0,5), col
       = "blue")
24 #add
                  : adds to an existing plot
25
26 #Points of intersections of constraints and the axis
       (0,0),(0,2.5),(1,2),(2,0),(5,0),(0,4)
27 points(c(0,0,1,2,5,0),c(0,2.5,2,0,0,4))
28
29 #Inserting texts to name the points
30 \text{ text}(c(0,0,1,2,5,0)+0.1,c(0,2.5,2,0,0,4)+0.2,LETTERS)
      [1:6], cex = 0.7)
31 #First two arguments are x(+0.1) and y(+0.2)
      coordinates
                 : a stored array variable of all
32 #LETTERS
      alphabets
                  : ratio of modification to font size
33 #cex
34
35 #Add a shaded area
36 polygon(c(0,0,1,2),c(0,2.5,2,0), col = rgb(0.48,
      0.46, 0.46, 0.5),
           border=NA)
37
38 #border
                  : option to add border to the shaded
      area
39
40 #Adding "solution space" text to the shaded area at
      (0.7,1)
41 text (0.7,1," Solution \ \ nSpace", cex = 0.9)
43 ##Objective function : Max 2 *x1 + 3 * x2
44 # maximum objective is 8, Therefore 2 *x1 + 3 *x2 =
       8
```

R code Exa 3.3.1 LP solution for Reddy Mikks Problem using lpSolve

```
1 ##Chapter 3: The Simplex Method and Sensitivity
       Analysis
2 ##Example 3-1 : Page 89
4 # Objective function : Max 5*x1 + 4*x2 + 0*s1 + 0*s2
      + 0*s3 + 0*s4
5 \text{ a} \leftarrow c(5,4,0,0,0,0)
7 \# Constraint 1 : 6*x1 + 4*x2 +s1 = 24
8 C1 \leftarrow c(6,4,1,0,0,0)
9 bc1<- 24
10
11 # Constraint 2 : x1 + 2*x2 + s2 = 6
12 C2 \leftarrow c(1,2,0,1,0,0)
13 bc2<-6
14
15 \# Constraint 3 : -x1 + x2 + s3 = 1
16 \quad C3 \quad \leftarrow \quad c(-1,1,0,0,1,0)
17 bc3<-1
18
19 # Constraint 4 : 1*x^2 + s^4 = 2
20 \quad C4 \leftarrow c(0,1,0,0,0,1)
21 \text{ bc4} \leftarrow 2
```

R code Exa 3.4.1 LP solution using M Method

```
1 ##Chapter 3: The Simplex Method and Sensitivity
      Analysis
2 ##Example 4-1 : Page 100
4 # Objective function : Min 4 * x1 + x2 + M * R1 + M *
       R2
5 BigM <- 1000
6 \text{ a} \leftarrow c(4,1,0,0,BigM,BigM)
8 \# Constraint 1 : 3*x1 + x2 + R1 = 3
9 C1 \leftarrow c(3,1,0,0,1,0)
10 bc1<- 3
11
12 \# Constraint 2 : 4*x1 + 3*x2 -x3 + R2 >= 6
13 C2 \leftarrow c(4,3,-1,0,0,1)
14 bc2<-6
15
16 \# Constraint 3 : x1 + 2*x2 +x4 <=4
17 C3 \leftarrow c(1,2,0,1,0,0)
18 bc3<-4
19
20 library ("boot")
21
22 simplex(a, NULL, NULL, NULL, NULL, rbind(C1, C2, C3), c(bc1
```

R code Exa 3.4.2 LP solution using Two Phase Method

```
1 ##Chapter 3: The Simplex Method and Sensitivity
      Analysis
2 \# \text{Example } 4-2 : \text{Page } 105
4 # Objective function : Min 4 * x1 + x2
5 \text{ a} < -c(4,1)
7 \# Constraint 1 : 3*x1 + x2 = 3
8 \text{ C1} \leftarrow c(3,1)
9 bc1<- 3
10
11 # Constraint 2 : 4*x1 + 3*x2 >= 6
12 \ C2 \leftarrow c(4,3)
13 bc2<-6
14
15 \# Constraint 3 : x1 + 2*x2 <=4
16 \quad C3 < -c(1,2)
17 bc3<-4
18
19 library ("boot")
20
21 \text{ simplex}(a,C3,bc3,C2,bc2, C1,bc1,maxi = F)
22
23 #The method employed by this function is the two
      phase tableau simplex method.
24 #If there are >= or equality constraints an initial
      feasible solution is not
25 #easy to find. To find a feasible solution an
      artificial variable is introduced
26 #into each >= or equality constraint and an
      auxiliary objective function is defined
```

R code Exa 3.5.1 Degenerate Optimal Solution

```
1 ##Chapter 3: The Simplex Method and Sensitivity
      Analysis
2 ##Example 5-1 : Page 108
4 # Objective function : Max 3 * x1 + 9 * x2
5 \text{ a} \leftarrow c(3,9)
7 \# Constraint 1 : x1 + 4 * x2 <= 8
8 \text{ C1} \leftarrow c(1,4)
9 bc1<- 8
10
11 # Constraint 2 : x1 + 2*x2 \le 4
12 \quad C2 < - c(1,2)
13 bc2<-4
14
15
16 library ("boot")
17
18 solution=simplex(a, rbind(C1,C2),c(bc1,bc2),maxi = T
```

19 ##The simplex function arbitrarily breaks the tie in minimum ratio for the leaving variable.

R code Exa 3.5.2 Infinite Solutions to an LP

```
1 ##Chapter 3: The Simplex Method and Sensitivity
      Analysis
2 \# \text{Example } 5-2 : \text{Page } 112
4 # Objective function : Max 2 * x1 + 4 * x2
5 \text{ a} < -c(2,4)
7 # Constraint 1 : 1 * x1 + 2 * x2 \le 5
8 \text{ C1} \leftarrow c(1,2)
9 bc1<- 5
10
11 # Constraint 2 : x1 + x2 <= 4
12 C2 < -c(1,1)
13 bc2<-4
14
15 library ("boot")
16
17 simplex(a, rbind(C1, C2), c(bc1, bc2), maxi = T)
18 #The simplex function as well as lpsolve stops as
      soon as a feasible optima is obtained and doesn't
       evaluate
19 #alternate optima
```

R code Exa 3.5.3 Unbounded objective value

```
1 ##Chapter 3 : The Simplex Method and Sensitivity
        Analysis
2 ##Example 5-3 : Page 115
```

```
3
4 # Objective function : Max 2 * x1 + x2
5 \text{ a} < -c(2,1)
6
7 \# Constraint 1 : x1 - x2 <= 10
8 C1 < c(1,-1)
9 bc1<- 10
10
11 # Constraint 2 : 2 * x1 \le 40
12 C2 < -c(2,0)
13 bc2<-40
14
15 library ("lpSolve")
16
17 solution = lp("max", a, rbind(C1, C2), rep("<=",2),
     c(bc1, bc2))
18 #Error: status 3 implies that the model is unbounded
```

R code Exa 3.5.4 Infeasible Solution Space

R code Exa 3.6.1 Changing RHS for Sensitivity analysis

```
1 ##Chapter 3: The Simplex Method and Sensitivity
      Analysis
2 ##Example 6-1 : Page 118
4 # Objective function : Max 30 * x1 + 20 * x2
5 \text{ a} < -c(30,20)
7 \# Constraint 1 : 2 * x1 + x2 <= 8
8 \text{ C1} \leftarrow c(2,1)
9 bc1<- 8
10
11 \# Constraint 2 : x1 + 3 * x2 <= 8
12 C2 < -c(1,3)
13 bc2<-8
14
15 library ("lpSolve")
16
17 solution=lp("\max", a, rbind(C1, C2), rep("<=",2), c(
      bc1, bc2),compute.sens = 1)
18
19 ##Unit worth of resources in $/hr
20 solution $duals [1:2]
21
22 ##Lower limit of hours resource for respective
      machines for which dual prices are valid
23 solution $duals.from [1:2]
24
```

```
25 ##Upper limit of hours resource for respective machines for which dual prices are valid 26 solution$duals.to[1:2]
```

R code Exa 3.6.2 Changing objective coefficient for Sensitivity analysis

```
1 ##Chapter 3: The Simplex Method and Sensitivity
      Analysis
2 \# \text{Example } 6-2 : \text{Page } 121
4 # Objective function : Max 30 * x1 + 20 * x2
5 \text{ a} < -c(30,20)
7 \# Constraint 1 : 2 * x1 + x2 <= 8
8 \text{ C1} \leftarrow c(2,1)
9 bc1<- 8
10
11 # Constraint 2 : x1 + 3 * x2 \le 8
12 C2 < -c(1,3)
13 bc2<-8
14
15 library ("lpSolve")
16
17 solution=lp("\max", a, rbind(C1, C2), rep("<=",2), c(
      bc1, bc2),compute.sens = 1)
18
19 ##Lower limit of respective Objective coefficient
      for which the objective value will not change
20 solution $ sens. coef.from
21 ##Upper limit of respective Objective coefficient
      for which the objective value will not change
22 solution $ sens.coef.to
```

R code Exa 3.6.3 Sensitivity Analysis using TOYCO Model

```
1 ##Chapter 3: The Simplex Method and Sensitivity
      Analysis
2 \# \text{Example } 6-3 : \text{Page } 124
4 # Objective function : Max 3 * x1 + 2 * x2 + 5 * x3
5 \text{ a} \leftarrow c(3,2,5)
7 \# Constraint 1 : x1 + 2 * x2 + x3 <= 430
8 \text{ C1} \leftarrow c(1,2,1)
9 bc1<- 430
10
11 # Constraint 2 : 3 * x1 + 2 * x3 <= 460
12 \quad C2 \leftarrow c(3,0,2)
13 bc2<-460
14
15 # Constraint 3 : 1 * x1 + 4 * x2 \le 420
16 \quad C3 < -c(1,4,0)
17 bc3<-420
18
19 library ("lpSolve")
20 solution=lp("\max", a, rbind(C1, C2, C3), rep("\leq=",3)
      , c(bc1, bc2, bc3), compute.sens = 1)
21
22 ##Lower limit of hours resource for respective
      machines for which dual prices are valid
23 solution $duals.from
24
25 ##Upper limit of hours resource for respective
      machines for which dual prices are valid
26 solution $duals.to
27
28 ##Refer Footnote 10, Page 127
```

Chapter 4

Duality and Post optimality Analysis

R code Exa 4.2.1 Solving primal and dual problem using linprog pacakge

```
1 ##Chapter 4: Duality and Post-optimality Analysis
2 ##Example 2-1 : Page 160
4 # Objective function : Max 5 * x1 + 12 * x2 + 4 * x3
5 \text{ a} < -c(5,12,4)
7 \# Constraint 1 : x1 + 2 * x2 + x3 <= 10
8 \text{ C1} \leftarrow c(1,2,1)
9 bc1<- 10
10
11 # Constraint 2 : 2 * x1 - x2 + 3 * x3 = 8
12 C2 \leftarrow c(2,-1,3)
13 bc2<-8
14
15 library ("linprog")
16
17 solveLP(a,c(bc1, bc2), rbind(C1, C2), c("<=","="),
      maximum = T, lpSolve = T)
18
```

```
19 ## At the moment the dual problem can not be solved
     with equality constraints in the function solveLP
20 ##Nevertheless we change the equality constrait to
     two inequality constraints. i.e.
21 ## f1(x)=b ==> f1(x)<=b , f1(x)>=b
22 solveLP(a,c(bc1, bc2, bc2), rbind(C1, C2, C2), c("<=
     ","<=",">="), maximum = T,lpSolve = T,solve.dual
     = T)
```

R code Exa 4.2.2 Associated objective values for arbitrary feasible primal and dual solutions

```
1 ##Chapter 4: Duality and Post-optimality Analysis
2 ##Example 2-2 : Page 162
4 ##Returns primal objective value
5 PrimalObj <-function(x1,x2,x3){</pre>
       return (5*x1+12*x2+4*x3)
6
7 }
8
9 #Returns dual objective value
10 DualObj <-function(y1,y2){
       return (10*y1+8*y2)
11
12 }
13
14 ##Calling primal and dual functions with the
      arbitrary feasible solutions
15 PrimalObj(0,0,8/3)
16 DualObj(6,0)
```

R code Exa 4.2.3 Getting the simplex tableau at any iteration from the original data and inverse of the iteration

```
1 ##Chapter 4: Duality and Post-optimality Analysis
2 \text{ ##Example } 2-3 : Page 166
4 #Optimal Inverse
5 OptimalInv=matrix(c(2/5,-1/5,1/5,2/5),nrow=2,byrow=T
6
7 #Original X1 column
8 OrigX1=c(1,2)
10 ##Optimal X1 column using formaula 1
11 X1Optimal=OptimalInv%*%OrigX1
12 X1Optimal
13
14 ##Functions to calculate Z coefficients of X1 and R
15 ZCoefX1<-function(y1,y2) {return(y1+2*y2-5)}
16 ZCoefR<-function(y1,y2) {return(paste(y2,"+ M"))}
17
18 ## Calling the function with optimal dual values
19 ZCoefX1(29/5,-2/5)
20 \text{ ZCoefR}(29/5, -2/5)
```

R code Exa 4.3.1 Economic interpretation of Dual variables of Reddy Mikks Problem

```
1 ##Chapter 4 : Duality and Post-optimality Analysis
2 ##Example 2-1 : Page 171
3
4 # Objective function :Max 5 * x1 + 4 * x2
5 a <- c(5,4)
6
7 # Constraint 1 : 6 * x1 + 4 * x2 <= 24
8 C1 <- c(6,4)
9 bc1<- 24</pre>
```

```
11 \# Constraint 2 : x1 + 2 * x2 <= 6
12 C2 < -c(1,2)
13 bc2<-6
14
15 \# Constraint 3 : -x1 + x2 <= 1
16 \quad C3 < -c(-1,1)
17 bc3<-1
18
19 # Constraint 4 : x^2 <= 2
20 \text{ C4} \leftarrow c(0,1)
21 bc4<-2
22
23 library("linprog")
24
25 solveLP(a,c(bc1, bc2, bc3, bc4), rbind(C1, C2, C3,
      C4), rep("<=",4"), maximum = T,lpSolve = T)
26
27 ##solve.dual arguement can be passed to solveLP
      function to solve the dual of the problem
28 solveLP(a,c(bc1, bc2, bc3, bc4), rbind(C1, C2, C3,
      C4), rep("<=",4), maximum = T, lpSolve = T, solve.
      dual = T)
```

R code Exa 4.3.2 Economic interpretation of Dual constraints of Reddy Mikks Problem

```
1 ##Chapter 4 : Duality and Post-optimality Analysis
2 ##Example 3-2 : Page 173
3
4 # Objective function :Max 3 * x1 + 2 * x2 + 5 * x3
5 a <- c(3,2,5)
6
7 # Constraint 1 : x1 + 2 * x2 + x3 <= 430
8 C1 <- c(1,2,1)
9 bc1<- 430</pre>
```

```
10
11 # Constraint 2 : 3 * x1 + 2 * x3 \le 460
12 \quad C2 \leftarrow c(3,0,2)
13 bc2<-460
14
15 # Constraint 3 : 1 * x1 + 4 * x2 \le 420
16 \quad C3 < -c(1,4,0)
17 bc3<-420
18
19 library ("linprog")
20 solveLP(a,c(bc1, bc2, bc3), rbind(C1, C2, C3), rep("
      <=",3), maximum = T,1pSolve = T)
21
22 ##solve.dual arguement can be passed to solveLP
      function to solve the dual of the problem
23 solveLP(a,c(bc1, bc2, bc3), \frac{\text{rbind}}{\text{c1}}, C2, C3), \frac{\text{rep}}{\text{c2}}
      <=",3), maximum = T,lpSolve = T,solve.dual = T)
```

R code Exa 4.4.1 Dual Simplex using glpk library

```
1 ##Chapter 4 : Duality and Post-optimal Analysis
2 ##Example 2-1 : Page 175
3
4 ##Please note that the constraints are not given in the example(printing error) and
5 ##it has been deduced from simplex tableau
6
7 ##We will be using glpkAPI library. If you are using Debian system, run the following command
8 ##before running the R script :: sudo apt-get install libglpk-dev
9 ##For other linux and windows systems, install GLPK which is available at https://www.gnu.org/software/glpk/
10 ##or from their respective repositories
```

```
11 library(glpkAPI)
12
13 # Objective function : Min 3 * x1 + 2 * x2 + 1 * x3
14 a <-c(3,2,1)
15
16 # Constraint 1 : 3 * x1 + 1 * x2 + x3 >= 3
17 # Standardizing constraint by multiplying by -1
18 C1 \leftarrow c(-3,-1,-1)
19 bc1<- -3
20
21 \# Constraint 2 : -3 * x1 + 3 * x2 + x3 >= 6
22 \# Standardizing constraint by multiplying by <math>-1
23 C2 < c(3,-3,-1)
24 bc2<--6
25
26 \# Constraint 3 : x1 + x2 + x3 \le 3
27 \quad C3 \leftarrow c(1,1,1)
28 bc3<-3
29
30 #upper bound vector
31 bc <- c(bc1,bc2,bc3)
32
33 #Constraint matrix
34 ConstraintMatrix <- rbind(a,C1,C2,C3)
35
36 #Initiating row and coloumn index variable as well
      as constraint coefficient value variable
37 rowindex <- numeric()
38 colindex <- numeric()
39 value <- numeric()</pre>
40
41 #initiate GLPK object and name
42 dualSimplex <- initProbGLPK()
43 setProbNameGLPK(dualSimplex, "Example 4-1")
44 setObjNameGLPK(dualSimplex, "Minimize using Dual
      Simplex")
45
46 #Setting objective direction and number of coloumns
```

```
47 setObjDirGLPK(dualSimplex, GLP_MIN)
48 addColsGLPK(dualSimplex, 3)
49
50 #setting decision variable names, bounds and
      coefficients
51 for (i in 1:3) {
     setColsNamesGLPK(dualSimplex,i,toString(c("x",i)))
52
     setColBndGLPK(dualSimplex, i, GLP_LO, 0.0, 0.0)
53
     setObjCoefsGLPK(dualSimplex, i, a[i])
54
55 }
56
57 #add 4 rows (including the objective)
58 addRowsGLPK(dualSimplex, 4)
59
60 #set row name as objective name itself
61 setRowsNamesGLPK(dualSimplex, 1, getObjNameGLPK(
      dualSimplex))
62
63 #set row names and bounds for constraint
64 for (i in 1:3) {
     setRowsNamesGLPK(dualSimplex, i+1, toString(c("
65
        Constraint", i)))
     setRowBndGLPK(dualSimplex, i+1, GLP_UP, 0, bc[i])
66
67 }
68
69 #initiating row and coloumn index and the values
70 \text{ counter=1}
71 for (i in 1:4) {
    for (j in 1:3) {
       rowindex[counter] <- i</pre>
73
       colindex[counter] <- j</pre>
74
       value[counter] <- ConstraintMatrix[i,j]</pre>
75
76
       counter=counter+1
     }
77
78 }
79
80 #change the soving algorithm to dual simplex
81 setSimplexParmGLPK(METH,GLP_DUAL)
```

```
82
83 #shows the current solver parameters
84 getSimplexParmGLPK()
85
86 #load and initiate all the data
87 loadMatrixGLPK(dualSimplex, 12, rowindex, colindex,
      value)
88
89 #Solve
90 solveSimplexGLPK(dualSimplex)
91
92 #Prints the status, optimal objective value and
      decisision variable value
93 getSolStatGLPK(dualSimplex)
94 getObjValGLPK(dualSimplex)
95 getColsPrimGLPK(dualSimplex)
96
97 #prints the summary of the optimization to your
      working directory
98 printSolGLPK(dualSimplex, 'textfile.txt')
99
100 #deletes the glpk object
101 delProbGLPK(dualSimplex)
```

R code Exa 4.5.1 Changes in RHS

```
8 TableauNewRightHandSide
9
10 ##Situation 2
11 ConstraintNNewRightHandSide <-rbind(450,460,400)
12 TableauNewRightHandSide <- Inverse %*%
ConstraintNNewRightHandSide
13 TableauNewRightHandSide
```

R code Exa 4.5.2 Addition of a new constraint

```
1 ##Chapter 4: Duality and Post-optimal Analysis
2 \# \text{Example } 5-2 : \text{Page } 185
4 ##Situation 1 is a theoretical explanation
6 ##Situation 2
7 tableau=matrix(c
     , nrow=4, ncol =8, byrow=T )
8
9 OldX7Row <- tableau[4,]
10 CoeffX2 <-OldX7Row[2]
11 CoeffX3 <-OldX7Row[3]
12 X2Row \leftarrow c(-0.25, 1, 0, 0.5, -0.25, 0, 0, 100)
13 X3Row \leftarrow c(1.5,0,1,0,1.5,0,0,230)
14 tableau[4,] <- OldX7Row - (CoeffX2 %*% tableau[1,] +
      CoeffX3 %*% tableau[2,])
15 tableau
```

R code Exa 4.5.3 Changes in objective coefficient

```
1 ##Chapter 4 : Duality and Post-optimal Analysis
2 ##Example 5-3 : Page 187
```

```
3
4 ##Situation 1
5 NewObjCoeffBasic <- c(3,4,0)
6 Inverse \leftarrow rbind(c(0.5,-0.25,0),c(0,0.5,0),c(-2,1,1)
      )
8 NewDualVariables <-NewObjCoeffBasic %*%
                                              Inverse
9 NewDualVariables
10
11 ##Element-wise multiplication to get Reduced Costs
12 ReducedCostX1<-sum(c(1,3,1)*NewDualVariables) -2
13 ReducedCostX4<-sum(c(1,0,0)*NewDualVariables) -0
14 ReducedCostX5<-sum(c(0,1,0)*NewDualVariables) -0
15 ReducedCostX1
16 ReducedCostX4
17 ReducedCostX5
18
19
20 CurrentOptimal \leftarrow c(0,100,230)
21 NewObjCoeff <-c(2,3,4)
22
23 #Optimal Objective value
24 NewRevenue <- sum(NewObjCoeff*CurrentOptimal)
25 NewRevenue
26
27 ##Situation 2
28 NewObjCoeffBasic <- c(3,4,0)
29 Inverse \leftarrow rbind(c(0.5,-0.25,0),c(0,0.5,0),c(-2,1,1)
30
31 NewDualVariables <-NewObjCoeffBasic %*% Inverse
32 NewDualVariables
33
34 ##Element-wise multiplication to get Reduced Costs
35 ReducedCostX1<-sum(c(1,3,1)*NewDualVariables) -6
36 \text{ ReducedCostX4} < -\text{sum}(c(1,0,0) * \text{NewDualVariables}) -0
37 ReducedCostX5<-sum(c(0,1,0)*NewDualVariables) -0
38 ReducedCostX1
```

```
39 ReducedCostX4
40 ReducedCostX5
41
42 NewObjCoeff<-c(-0.75,0,0,1.5,1.25,0)
43 C1<-c(-0.25,1,0,0.5,-0.25,0)
44 bc1<-100
45 C2<-c(1.5,0,1,0,0.5,0)
46 bc2<-230
47 C3<-c(2,0,0,-2,1,1)
48 bc3<-20
49 library("lpSolve")
50 solution=lp("max", NewObjCoeff, rbind(C1, C2, C3), rep("=",3), c(bc1, bc2, bc3))
51 solution</pre>
```

R code Exa 4.5.4 Addition of a new activity

```
1 ##Chapter 4 : Duality and Post-optimal Analysis
2 ##Example 5-4 : Page 189
3 OptimalDualValues <-c(1,2,0)
4
5 ReducedCostX7 <-sum(c(1,1,2)*OptimalDualValues)-4
6 ReducedCostX7
7
8 Inverse <- rbind(c(0.5,-0.25,0),c(0,0.5,0),c(-2,1,1)
)
9 OldX7Col <-rbind(1,1,2)
10 NewX7Col <-Inverse%*%OldX7Col
11 NewX7Col</pre>
```

Chapter 5

Transportation Model and its Variant

R code Exa 5.1.1 A general transportation model

```
1 ##Chapter 5: Transportation Model and its Variant
2 \# Example 1-1 : Page 195
3 ##Refer to the transportation tableau in Page 197
4 costs <- matrix (c(80,215,100,108,102,68), 3, 2,
     byrow = T)
6 #Constraints 1,2 & 3 are row constraints as they
     corresponds to rows of transportation tableau
7 row.signs <-rep("=",3)
8 \text{ row.rhs} < -c (1000, 1500, 1200)
10 #Constraints 4 & 5 are coloumn constraints as they
     corresponds to coloumns of transportation tableau
11 col.signs <-rep("=",2)
12 \text{ col.rhs} < -c(2300, 1400)
13
14 library(lpSolve)
15 ##lpSolve library has lp.transport function
      specially for problems which can be formulated as
```

R code Exa 5.1.2 Balancing transportation models

```
1 ##Chapter 5 : Transportation Model and its Variant
2 \# \text{Example } 1-2 : \text{Page } 197
4 ##Part1- Adding a dummy origin
5 costs <- matrix (c(80,215,100,108,102,68,0,0), 4, 2,
      byrow = T)
7 #Constraints 1 to 4 are row constraints as they
      corresponds to rows of transportation tableau
8 row.signs <-rep("=",4)</pre>
9 \text{ row.rhs} < -c (1000, 1300, 1200, 200)
10
11 #Constraints 5 & 6 are coloumn constraints as they
      corresponds to coloumns of transportation tableau
12 col.signs <-rep("=",2)
13 \text{ col.rhs} < -c(2300, 1400)
14
15 library(lpSolve)
16 solution <- lp.transport (costs, "min", row.signs,
      row.rhs, col.signs, col.rhs)
17 solution $ solution
18
19
20 ##Part2 - Adding a dummy destination
21 costs <- matrix (c(80,215,0,100,108,0,102,68,0), 3,
      3, byrow = T
22
```

```
#Constraints 1,2 & 3 are row constraints as they
corresponds to rows of transportation tableau

24 row.signs <-rep("=",3)
25 row.rhs <-c(1000,1500,1200)

26

27 #Constraints 4 & 5 are coloumn constraints as they
corresponds to coloumns of transportation tableau

28 col.signs <-rep("=",3)
29 col.rhs <-c(1900,1400,400)

30

31 library(lpSolve)
32 solution <- lp.transport (costs, "min", row.signs,
row.rhs, col.signs, col.rhs)
33 solution$solution
```

R code Exa 5.2.1 Production inventory control

R code Exa 5.2.2 Tool sharpening

```
1 ##Chapter 5: Transportation Model and its Variant
2 \# \text{Example } 2-2 : \text{Page } 203
3
4 #BigM is Big
5 BigM = 1000
7 #Initializing the matrix with all values as M
8 costs<-matrix(BigM,8,8,byrow = T)</pre>
9 #All values of first row are 12
10 costs[1,]=c(rep(12,8))
11 #All values of 8th column is 0
12 costs[,8]=rbind(rep(0,8))
13
14 ##Adding the rest of the values
15 for (i in 2:7) {
     costs[i,]=c(rep(BigM,i-1),c(6,5,rep(3,4))[1:(8-i)
16
        ],0)
17 }
18
19 #Constraints 1 to 8 are row constraints as they
      corresponds to rows of transportation tableau
20 row.signs <-rep("=",8)
21 row.rhs <-c(124,24,12,14,20,18,14,22)
22
23 #Constraints 8 to 16 are coloumn constraints as they
       corresponds to coloumns of transportation
      tableau
24 col.signs <-rep("=",8)
```

```
25 col.rhs <-c(24,12,14,20,18,14,22,124)
26
27 library(lpSolve)
28 solution <- lp.transport (costs, "min", row.signs,
         row.rhs, col.signs, col.rhs)
29 solution$solution</pre>
```

R code Exa 5.3.1 Sunray Transport

```
1 ##Chapter 5: Transportation Model and its Variant
2 \# \text{Example } 3-1 : \text{Page } 207
3 costs <- matrix (c(10,2,20,11,12,7,9,20,4,14,16,18),
       3, 4, byrow = T
5 #Constraints 1,2 & 3 are row constraints as they
      corresponds to rows of transportation tableau
6 row.signs <-rep("=",3)
7 \text{ row.rhs} < -c(15,25,10)
9 #Constraints 4 & 5 are coloumn constraints as they
      corresponds to coloumns of transportation tableau
10 col.signs <-rep("=",4)
11 col.rhs < -c(5,15,15,15)
12
13 library(lpSolve)
14 solution <- lp.transport (costs, "min", row.signs,
      row.rhs, col.signs, col.rhs)
15 solution $ objval
16 solution $ solution
```

R code Exa 5.3.2 Northwest corner method

1 ##Chapter 5: Transportation Model and its Variant

```
2 ##Example 3-2 : Page 208
3 costs <- matrix (c(10,2,20,11,12,7,9,20,4,14,16,18),
       3, 4, byrow = T
4 \text{ row.rhs} < -c(15, 25, 10)
5 \text{ col.rhs} < -c(5,15,15,15)
6 i=1; j=1
7 allotment \leftarrow matrix (rep(0,12), 3, 4, byrow = T)
9 while (sum(row.rhs) & sum(col.rhs)) {
     ##Till we reach the last row and coloumn
10
     while (i<=3 & j<=4) {</pre>
11
12
       ##if demand is >= supply
13
       if (row.rhs[i] >= col.rhs[j]) {
         ##assign the demand to that cell
14
         allotment[i,j]=col.rhs[j]
15
         ##deduct supply from demand
16
17
         row.rhs[i]=row.rhs[i]-col.rhs[j]
18
         ##assign zero to supply
19
         col.rhs[j]=0
       }else {
20
21
         ##assign the supply to that cell
22
         allotment[i,j]=row.rhs[i]
         ##deduct demand from supply
23
24
         col.rhs[j] = col.rhs[j] - row.rhs[i]
         ##assign zero to demand
25
26
         row.rhs[i]=0
27
28
29
       #if demand=0, go to the next demand, else go to
          next supply
       ifelse(row.rhs[i] == 0,(i=i+1),(j=j+1))
30
31
     }
32 }
33 allotment
```

R code Exa 5.3.3 Least cost method

```
1 ##Chapter 5 : Transportation Model and its Variant
2 ##Example 3-3 : Page 209
3 costs <- matrix (c(10,2,20,11,12,7,9,20,4,14,16,18),
       3, 4, byrow = T
4 \text{ row.rhs} < -c(15, 25, 10)
5 \text{ col.rhs} \leftarrow (5, 15, 15, 15)
6 costsdup <- costs
7 allotment \leftarrow matrix (rep(0,12), 3, 4, byrow = T)
9 #until there are supply and demand
10 while (sum(row.rhs) & sum(col.rhs)) {
     #index of min cost
11
     index=which.min(costsdup)
12
     #get the row index
13
     rowindex=index %% length(row.rhs)
14
     #get the coloumn index
15
     colindex=ceiling(index/length(row.rhs))
16
     #if row index=0 ,assign 3(since we are takinf
17
        modulus)
     if (rowindex == 0) {rowindex = 3}
18
19
20
     #if demand > supply
     if (row.rhs[rowindex]>=col.rhs[colindex]){
21
22
       #allocate supply to that cell
       allotment[index] <- col.rhs[colindex]</pre>
23
       ##deduct supply from demand
24
       row.rhs[rowindex] <- row.rhs[rowindex]-col.rhs[</pre>
25
          colindex
       ##assign zero to supply
26
       col.rhs[colindex]=0
27
28
     }else{
       #allocate demand to that cell
29
       allotment[index] <- row.rhs[rowindex]</pre>
30
       ##deduct demand from supply
31
32
       col.rhs[colindex] <- col.rhs[colindex]-row.rhs[</pre>
          rowindexl
```

```
33  ##assign zero to demand
34    row.rhs[rowindex]=0
35    }
36    costsdup[index]=1000
37
38  }
39
40
41  allotment
```

R code Exa 5.3.5 Method of multipliers

```
1 ##Chapter 5: Transportation Model and its Variant
2 ##Example 3-5 : Page 212
3 #u-v incidence matrix with columns - u1, v1, v2, u2, v3,
     v4, u3 and rows as uv equations
4 a=matrix(c
     ,6,7,byrow = T)
5
6 #cost matrix
7 costs <- matrix (c(10,2,20,11,12,7,9,20,4,14,16,18),
      3, 4, byrow = T
8 #rhs of uv equations
9 n=c(10,2,7,9,20,18)
10
11 library(limSolve)
12 ##Least Squares with Equality and Inequality
     Constraints
13 Sol=lsei(E=a, F=n, A = diag(7), B = rep(0, 7), G=diag(7)
     (7), H=rep(0, 7), verbose = FALSE)$X
14
15 c(Sol[1]+Sol[5]-20,Sol[1]+Sol[6]-11,Sol[4]+Sol
     [2]-12, Sol [7]+Sol [2]-4, Sol [7]+Sol [3]-14, Sol [7]+
     Sol [5] -16)
```

R code Exa 5.4.1 Hungarian method 1

```
1 ##Chapter 5: Transportation Model and its Variant
2 \# Example 4-1 : Page 221
3 \text{ costs} \leftarrow \text{matrix} (c(15,10,9,9,15,10,10,12,8), 3, 3,
      byrow = T)
4 rowmin<-numeric()
5 colmin<-numeric()</pre>
7 ##Subtracting minimum cost element of row from all
      the elements of rows
8 for(i in 1:nrow(costs)){
     costs[i,] <-costs[i,]-min(costs[i,])</pre>
10 }
11
12 ##Subtracting minimum cost element of column from
      all the elements of column
13 for(i in 1:ncol(costs)){
     costs[,i] <-costs[,i]-min(costs[,i])</pre>
15 }
16
17 ##logic1 is a boolean matrix which contains true if
      cost matrix after the above operations is 0 and 0
       otherwise
18 \log ic1 < - costs == 0
19 eqnrow <-numeric()
20
21 ##We formulate it as a transportation tableau such
      that only one zero is selected from every row
      and every column
22 row.signs <-rep("=",nrow(costs))
23 row.rhs <-rep(1,nrow(costs))
24
25 col.signs <-rep("=",ncol(costs))
```

R code Exa 5.4.2 Hungarian method 2

```
1 ##Chapter 5: Transportation Model and its Variant
2 \# \text{Example } 4-2 : \text{Page } 222
3 costs <- matrix (c</pre>
      (1,4,6,3,9,7,10,9,4,5,11,7,8,7,8,5), 4, 4, byrow =
4 rowmin<-numeric()
5 colmin<-numeric()</pre>
6 costsdup <- costs
8 ##Subtracting minimum cost element of row from all
      the elements of rows
9 for(i in 1:nrow(costsdup)){
10
     costsdup[i,] <-costsdup[i,]-min(costsdup[i,])</pre>
11 }
12
13 ##Subtracting minimum cost element of column from
      all the elements of column
14 for(i in 1:ncol(costsdup)){
     costsdup[,i] <- costsdup[,i] -min(costsdup[,i])</pre>
15
16 }
17
18
19 ##We formulate it as a transportation tableau such
      that only one zero is selected from every row
```

```
and every column
20 ##maketable function returns the constraint matrix
21 eqnrow<-numeric()
22 maketable <-function (costsdup) {
23
     logic1<-costsdup==0</pre>
24
     return(logic1)
25 }
26
27 matr <-maketable (costsdup)
29 row.signs <-rep("=",nrow(costsdup))
30 row.rhs <-rep(1,nrow(costsdup))
31
32 col.signs <-rep("=",ncol(costsdup))
33 col.rhs <-rep(1,ncol(costsdup))
34 library(lpSolve)
35
36 Solution <-lp. transport (matr, "max", row.signs, row.
      rhs, col.signs, col.rhs)
37 while (Solution $objval!=nrow(costsdup)) {
     costsdup <- costsdup -min (costsdup [costsdup > 0])
38
     for (i in 1:length(costsdup)){
39
       ifelse(costsdup[i]<0, (costsdup[i]=0),0)</pre>
40
41
     matr <-maketable(costsdup)</pre>
42
43
     Solution <- lp. transport (matr, "max", row.signs,
        row.rhs, col.signs, col.rhs)
44 }
45 Solution $ solution
```

Chapter 6

Network Model

R code Exa 6.2.1 Minimal spanning tree algorithm

```
1 ##Chapter 6 : Network Model
2 ##Example 2-1 : Page 239
4 #If you have trouble installing the package/library,
      please reinstall R form the following link: https
     ://cran.r-project.org/bin/
5 library(igraph)
6 #creating the undirected graph with 6 nodes
7 A=graph(edges=c
     (1,2,1,5,1,3,1,4,2,5,2,3,2,4,3,4,3,6,4,5,4,6), n
     =6, directed=F)
8 #mst function generates the minimum spanning tree
9 MST \leftarrow mst(A, weights = c(1,9,5,7,3,6,4,5,10,8,3))
11 #The Fruchterman-Reingold Algorithm is a force-
     directed layout algorithm. The idea of a force
     directed layout algorithm is to consider
12 #a force between any two nodes. In this algorithm,
     the nodes are represented by steel rings and the
     edges are springs between them.
13 #The attractive force is analogous to the spring
```

```
force and the repulsive force is analogous to the
       electrical force. The basic idea is
14 #to minimize the energy of the system by moving the
      nodes and changing the forces between them.
15 #Here we create the cordinates for our graph using
      layout.fruchterman.reingold function
16 lay <- layout.fruchterman.reingold(A)
17
18 #Assigning the coordinates to the nodes of A
19 V(A)$x <- lay[, 1]
20 \ V(A) \$ y \leftarrow lay[, 2]
21
22 #assigning range of x and y
23 xlim <- range(lay[,1])
24 ylim <- range(lay[,2])
25
26 #plot graph A
27 plot.igraph(A, layout = lay, vertex.size=20,
               xlim = xlim, ylim = ylim, rescale =
28
                  FALSE)
29 #plot MST with red edges and nodes over the previous
       graph
30 plot.igraph(MST, layout = lay, vertex.color="red",
      edge.color="red", vertex.size=20,
               add = TRUE, rescale = FALSE)
31
```

R code Exa 6.3.1 Equipment replacement as shortest route problem

```
6
7 #creating the directed graph with 5 nodes
8 A=graph(edges=c(1,2,1,3,1,4,2,3,2,4,2,5,3,4,3,5,4,5)
      , n=5, directed=T)
9 #creating weights vector for each edges of the graph
10 weightsg<-c
      (4000,5400,9800,4300,6200,8700,4800,7100,4900)
11 #get shortest path from 1 to 5
12 sP<-get.shortest.paths(A, 1, to=5, weights = weightsg)
     $vpath
13 #get shortest path cost from 1 to 5
14 sPCost <- shortest.paths(A, 1, to=5, weights = weightsg)
15 sP
16 sPCost
17
18 #creating coordinates for layout
19 1 < -cbind(seq(0,9,2),0)
20
21 #plot the graph with straight edges (edge 1,4,7 & 9
      are straight edges. The order of edges is taken
     from graph initialization in line 6)
22 plot(delete.edges(A,c(2,3,5,6,8)),layout=1*2,vertex.
     size=15, edge.arrow.width = 0.2, asp = 0.5, edge.
     label=weightsg[c(1,4,7,9)])
23 #plot the graph with straight edges over the
      previous graph
24 plot(delete.edges(A,c(1,4,7,9)),edge.curved=.8,
     layout=1*2, vertex.size=15, edge.arrow.width = 0.2,
     asp = 0.5, edge.label=weightsg[c(2,3,5,6,8)],add=T
     )
25 A[] <- 0
26 #assign color red to each node of the graph in the
      shortest path
27 for (ed in 1:(length(sP[[1]])-1)){
     A<-A+edge(c(sP[[1]][ed],sP[[1]][ed+1]),color="red"
    V(A)[sP[[1]][ed]]$color <-" red"
29
30 }
```

R code Exa 6.3.2 Most reliable route

```
1 ##Chapter 6 : Network Model
2 ##Example 3-2 : Page 244
4 #If you have trouble installing the package/library,
       please reinstall R form the following link: https
      ://cran.r-project.org/bin/
5 library(igraph)
6 #creating the directed graph with 7 nodes
7 A=make_directed_graph(edges=c
      (1,2,1,3,2,4,2,3,3,4,3,5,4,5,4,6,5,7,6,7), n=7)
8 #creating probability vector
9 prob < -c(.2,.9,.8,.6,.1,.3,.4,.35,.25,.5)
10 #creating weight vector from probability vector
11 weightsg<--round(log10(prob),digits = 5)
12 ##calculating shortest path and its cost
13 sP<-get.shortest.paths(A, 1, to=7, weights = weightsg)
     $vpath
14 sPCost <- shortest.paths(A, 1, to=7, weights = weightsg,
     mode = "out")
15 \text{ sP}
16 10^(-sPCost)
17 #plotting the graph A
18 l <-layout.auto(A)
19 plot(A, vertex.size=15, layout=1, edge.arrow.width =
     0.2, asp = 0.5, edge.label=weightsg)
20
```

R code Exa 6.3.3 Three jub puzzle

```
1 ##Chapter 6 : Network Model
2 \# \text{Example } 3-3 : \text{Page } 245
3 library(igraph)
4
5 #creating the directed graph with 15 nodes
6 A=make_directed_graph(edges=c
      (1,2,2,3,3,4,4,5,5,6,6,7,7,8,8,9,1,10,10,11,11,12,12,13,13,14,14,
      , n=15)
7 #weights vector
8 \text{ weightsg} \leftarrow \text{rep}(1,21)
9 #shortest path
10 sP<-get.shortest.paths(A, 1, to=9, weights =weightsg)
      $vpath
11 #shortest path cost
12 sPCost <- shortest.paths(A, 1, to=9, weights = weightsg,
      mode = "out")
13 sP
14 sPCost
15 #circle layout of graph
16 l<-layout.circle (A)
17 #plot graph A
18 plot(A, vertex.size=15, layout=1, edge.arrow.size =
```

```
0.2, asp = 0.5)

19
20 #make an empty graph and add edges of the shortest
    path to the graph with red color nodes
21 A[] <- 0
22 for (ed in 1:(length(sP[[1]])-1)){
        A<-A+edge(c(sP[[1]][ed],sP[[1]][ed+1]),color="red"
        )
        V(A)[sP[[1]][ed]]$color<-"red"
25 }
26 V(A)[9]$color<-"red"
27 #plot the shortest path over the previous graph
28 plot(A,vertex.size=15,layout=1, edge.arrow.size = 0.2,edge.color="orange", asp = 0.5,add=T)</pre>
```

R code Exa 6.3.4 Djiktras algorithm

```
1 ##Chapter 6 : Network Model
2 ##Example 3-4 : Page 248
3
4 # Initializing no of nodes and the distance matrix
5 n = 5
6 d <- matrix (Inf, nrow=n, ncol = n)
7 d[1,2]=100;d[1,3]=30;d[2,3]=20;d[3,4]=10;d[3,5]=50;d
      [4,2]=15;d[4,5]=50
9 #initalizing the djiktra's algorithm table as shown
      in the book
10 djiktable < -matrix(c(1:n, rep(0, 3*n)), nrow=n)
11
12 \text{ now=1}
13 djiktable[1,4]=1
14
15 #if any of the node doesn't have any outgoing edge,
     make it permanent. Ex: node 5
```

```
16 for (i in 1:5){
17
     if (length(which((d[i,]!=Inf) %in% TRUE))==0){
       djiktable[i,4]=1
18
     }
19
20 }
21 #while there are nodes with temporary status
22 while (sum (djiktable [,4])!=n) {
     #find all the possibile nodes form current node
23
     possibles <-which ((d[now,]!=Inf) %in% TRUE)</pre>
24
     #for each node in possible nodes
25
26
     for (i in possibles){
27
       #assign current node to temp
28
       temp<-djiktable[i,2]</pre>
       #if current node is not assigned
29
       if (djiktable[i,2]!=0){
30
         #the minimum of distance is added
31
         djiktable[i,2]=min(djiktable[i,2],djiktable[
32
            now,2]+ d[now,i])
33
       }else{
         djiktable[i,2]=djiktable[now,2]+ d[now,i]
34
35
       #if there is no change in the next node
36
       if (djiktable[i,2]!=temp){
37
         #backtrack
38
         djiktable[i,3]=now
39
40
       }
41
     }
42
     #assign permanent status to the minimum index
43
     min.indx<-which.min(djiktable[possibles,2])</pre>
44
     djiktable[possibles[min.indx],4]=1
45
     now<- possibles[min.indx]</pre>
46
47 }
48
49 #prints out the shortest route
50 djiktable [now,2]
51 path <-character()
52 while (now!=0) {
```

```
53     path <-paste("->", now, path)
54     now=djiktable[now,3]
55 }
56 path
```

R code Exa 6.3.5 Floyds algorithm

```
1 ##Chapter 6 : Network Model
2 \# \text{Example } 3-5 : \text{Page } 252
3 #Initializing the nodes and floyd's D matrix and S
      matrix
4 n = 5
5 floydD<-array(Inf,dim= c(5,5))</pre>
6 diag(floydD) <- 0
7 floydD[1,2]=3;floydD[1,3]=10;floydD[2,4]=5;floydD
      [3,4]=6; floydD [4,5]=4;
9 #symmetric matrix
10 for(i in 1:5){
11
     j = 1
12
     while(j<i){</pre>
13
        floydD[i,j]=floydD[j,i]
14
       j = j + 1
     }
15
16 }
17 floydseq <-matrix (1:n, nrow=n, ncol=n, byrow = T)
18 diag(floydseq) <- 0
19
20 ##Floyd's Algorithm
21 for(k in 1:n){
     for(i in 1:n){
22
       for(j in 1:n){
23
          if(i!=k & j!=k & i!=j){
24
25
            if(floydD[i,k]+ floydD[k,j]< floydD[i,j]){</pre>
26
              floydD[i,j] <- floydD[i,k]+ floydD[k,j]</pre>
```

```
27
              floydseq[i,j] <- k</pre>
            }
28
         }
29
       }
30
31
     }
32 }
33 floydD
34 floydseq
35
36 ##Printing the shortest route
37 path <-character()
38 i=1; j=5
39 while(floydseq[i,j]!=j){
     path <-paste("->", j, path)
     j=floydseq[i,j]
41
42 }
43 path <-paste(i,"->",j,path)
44 path
```

R code Exa 6.3.6 LP formulation and solution of shortest route problem

```
14
15 # Constraint 3 : x13 + x23 - x34 - x35 = 0
16 C3 <- c(0,1,1,-1,-1,0,0)
17 bc3<-0
18
19 # Constraint 4 : x34 - x42 - x45 = 0
20 \quad C4 \leftarrow c(0,0,0,1,0,-1,-1)
21 bc4<-0
22
23 \# Constraint 5 : x35 + x45 = 0
24 \text{ C5} \leftarrow c(0,0,0,0,1,0,1)
25 bc5<-0
26
27 library ("lpSolve")
28
29 solution \leftarrow lp("min", a, rbind(C1, C2,C3,C4,C5), rep("
      =",5), c(bc1, bc2,bc3,bc4,bc5))
30 solution $ objval
31 solution $ solution
```

R code Exa 6.4.2 Maximal flow algorithm

```
1 ##Chapter 6 : Network Model
2 ##Example 4-2 : Page 263
4 #If you have trouble installing the package/library,
      please reinstall R form the following link: https
      ://cran.r-project.org/bin/
5 library(optrees)
6 #edge matrix with weights
7 arcs<-matrix(c(1,2,20,</pre>
8
       1,3,30,
9
       1,4,10,
       1,3,30,
10
11
       1,4,10,
```

R code Exa 6.4.3 LP formulation and solution of Maximal flow Mode

```
1 ##Chapter 6 : Network Model
2 ##Example 4-3 : Page 271
3
4 # Objective function 1 : Max x12 + x13 + x14
5 a \leftarrow c(1,1,1,0,0,0,0,0,0)
7 # Objective function 2 : Max x25 + x35 + x45
8 b \leftarrow c(0,0,0,0,1,0,1,0,1)
9
10 # Constraint 1 : x12 - x23 - x25 = 0
11 C1 \leftarrow c(1,0,0,-1,-1,0,0,0,0)
12 bc1<- 0
13
14 # Constraint 2 : x13 + x23 - x34 - x35 + x43 = 0
15 C2 \leftarrow c(0,1,0,1,0,-1,-1,1,0)
16 bc2<-0
17
18 # Constraint 3 : x14 + x34 - x43 - x45 = 0
19 C3 \leftarrow c(0,0,1,0,0,1,0,-1,-1)
20 bc3<-0
21
22 \# Constraint 4 : x12 <= 20
```

```
23 C4 \leftarrow c(1,0,0,0,0,0,0,0,0)
24 bc4<-20
25
26 \# Constraint 5 : x13 <= 30
27 \text{ C5} \leftarrow c(0,1,0,0,0,0,0,0,0)
28 bc5<-30
29
30 \# Constraint 6 : x14 <= 10
31 C6 \leftarrow c(0,0,1,0,0,0,0,0,0)
32 bc6<-10
33
34 \# Constraint 7 : x23 <= 40
35 \quad C7 \leftarrow c(0,0,0,1,0,0,0,0,0)
36 bc7<-40
37
38 \# Constraint 8 : x25 <= 30
39 C8 \leftarrow c(0,0,0,0,1,0,0,0,0)
40 bc8<-30
41
42 # Constraint 9 : x34 <= 10
43 C9 \leftarrow c(0,0,0,0,0,1,0,0,0)
44 bc9<-10
45
46 # Constraint 10 : x35 <= 20
47 \text{ C10} \leftarrow c(0,0,0,0,0,0,1,0,0)
48 bc10<-20
49
50 # Constraint 11 : x43 <= 5
51 \text{ C11} \leftarrow c(0,0,0,0,0,0,0,1,0)
52 bc11<-5
53
54 \# Constraint 12 : x45 <= 20
55 C12 <- c(0,0,0,0,0,0,0,0,1)
56 bc12<-20
57
58 library ("lpSolve")
59
60 solution <-lp("max", a, rbind(C1, C2,C3,C4,C5,C6,C7,
```

R code Exa 6.5.1 Network representation for PERT and CPM

```
1 ##Chapter 6 : Network Model
2 \# \text{Example } 4-3 : \text{Page } 275
3 library(igraph)
4 #create a directed graph
5 A=make_directed_graph(edges=c
      (1,2,3,4,4,6,6,7,7,8,8,9,1,3,1,5,5,7,1,8), n=9)
6
7 #creating layout for plot
8 \ 1 < -cbind(c(1,1.5,3,5,4.5,7,9,11,13),c)
      (1,3,3,3,2,3,2,1,1))
10 #creating plot
11 plot(A, vertex.size=15, layout=1, edge.arrow.size =
      0.5, asp = 0.2, edge.label=c("A - 3", "E - 2", "F - 2")
      ", "G - 2", "I - 2", "J - 4", "B - 2", "D - 3", "H - 1"
      "C - 4"
12 plot(graph(c(2,3),n=9),layout=1,edge.arrow.size =
      0.5, asp = 0.2, edge.lty=2, add=T)
```

R code Exa 6.5.2 Critical path method

```
1 ##Chapter 6 : Network Model
2 \# \text{Example } 5-2 : \text{Page } 280
4 #create distance matrix with -1 meaning no edge and
      0 meaning no cost edge
5 D<-matrix(c
      (-1,5,6,-1,-1,-1,-1,-1,3,8,-1,-1,-1,-1,-1,-1,2,11,-1,-1,-1,-1,0,1
      , nrow=6, ncol = 6, byrow = T)
6
7 ##Forward pass
8 ECT<-numeric()</pre>
9 ECT[1]<-0
10 for(i in 2:6){
     index <-which((D[,i]>=0) %in% TRUE)
11
     ECT[i] <-max(ECT[index]+D[index,i])</pre>
12
13 }
14
15 ##Backward pass
16 BP<-numeric()</pre>
17 BP [6] <-ECT [6]
18 for(i in 5:1){
     index <-which((D[i,]>=0) %in% TRUE)
19
20
     BP[i] <-min(BP[index]-D[i,index])</pre>
21 }
22
23 ##Finding critical nodes
24 critical <-character()
25 criti <-numeric()
26 critj <-numeric()
27 for(i in 1:6){
28
     for(j in 1:6){
       if(BP[i] == ECT[i] & BP[j] == ECT[j] & BP[j] - ECT[i
```

```
]==D[i,j]){
30
          critical <-cbind (critical, paste(i, "-", j))</pre>
          criti=cbind(criti,i)
31
32
          critj=cbind(critj,j)
33
        }
34
     }
35 }
36 critical
37
38 ##Duration of the project
39 TotalDays <- 0
40 for (i in 1:length(criti)){
     TotalDays <- TotalDays +D [criti[i], critj[i]]</pre>
41
42 }
43 TotalDays
```

R code Exa 6.5.3 Preliminary schedule

```
1 ##Chapter 6 : Network Model
2 \# \text{Example } 5-3 : \text{Page } 282
3 #Add an empty plot
4 plot(1, type="n", axes=T, xlab="Days", ylab="", yaxt
      = 'n', xlim=c(0,25), ylim=c(0,10), yaxs= "i", xaxs =
       " i ")
6 #Add line segment for each segment
7 \text{ segments}(0, 9, x1 = 5, y1 = 9)
8 \text{ segments}(5, 8, x1 = 13, y1 = 8)
9 \text{ segments} (13, 7, x1 = 25, y1 = 7)
10 segments (0, 6, x1 = 11, y1 = 6, 1ty = 2)
11 segments (5, 5, x1 = 11, y1 = 5, 1ty = 2)
12 \text{ segments}(8, 4, x1 = 13, y1 = 4, 1ty = 2)
13 segments (8, 3, x1 = 25, y1 = 3, 1ty = 2)
14 \text{ segments}(13, 2, x1 = 25, y1 = 2, 1ty = 2)
15 #Add text at specific coordinates
```

R code Exa 6.5.4 Determination of floats and red flag rule

```
1 ##Chapter 6 : Network Model
2 ##Example 5-4 : Page 284
4 #create distance matrix with -1 meaning no edge and
      0 meaning no cost edge
5 D<-matrix(c
      (-1,5,6,-1,-1,-1,-1,-1,3,8,-1,-1,-1,-1,-1,-1,2,11,-1,-1,-1,-1,0,1
      , nrow=6, ncol = 6, byrow = T)
7 ##Forward pass
8 ECT<-numeric()</pre>
9 ECT[1]<-0
10 for(i in 2:6){
     index <- which ((D[,i]>=0) %in% TRUE)
11
12
     ECT[i] <-max(ECT[index]+D[index,i])</pre>
13 }
14
15 ##Backward pass
16 BP<-numeric()
17 BP [6] <-ECT [6]
18 for(i in 5:1){
     index <- which ((D[i,]>=0) %in% TRUE)
19
     BP[i] <-min(BP[index]-D[i,index])</pre>
20
21 }
22
23 ##Finding critical nodes
```

```
24 critical <-character()
25 criti<-numeric()
26 critj <-numeric()
27 for(i in 1:6){
28
     for(j in 1:6){
29
          if (BP[i] == ECT[i] & BP[j] == ECT[j] & BP[j] - ECT[i
             ]==D[i,j]){
            critical <-cbind (critical, paste(i, "-", j))</pre>
30
            criti=cbind(criti,i)
31
            critj=cbind(critj,j)
32
       }
33
     }
34
35 }
36 critical
37
38 ##Duration of the project
39 TotalDays<-0
40 for (i in 1:length(criti)){
     TotalDays <- TotalDays +D [criti[i], critj[i]]
41
42 }
43 TotalDays
44
45 #Calculating total float and free float for non-
      critical activities
46 NonCritical \leftarrow matrix (c(1,3,2,3,3,5,3,6,4,6), ncol=2,
      byrow = T)
47 NCA <- character()
48 duration <- numeric()
49 TotalF<-numeric()
50 FreeF<-numeric()
51 for(i in 1:length(NonCritical[,1])){
     j<-NonCritical[i,1]</pre>
52
     k <- NonCritical[i,2]
53
     NCA[i] <-paste(j,"-->",k)
54
     duration[i] <- D[j,k]</pre>
55
     TotalF[i] <- BP[k]-ECT[j]-D[j,k]
56
     FreeF[i] \leftarrow ECT[k]-BP[j]-D[j,k]
57
58 }
```

R code Exa 6.5.5 LP formulation and solution of project scheduling problem

```
1 ##Chapter 6 : Network Model
2 ##Example 5-5 : Page 288
3
4 # Objective function : Max 6*x12 + 6*x13 + 3*x23 + 8*
      x24 + 2*x35 + 11*x36 + 1 *x46 + 12*x56
5 a \leftarrow c(6,6,3,8,2,11,0,1,12)
6
7 \# Constraint 1 : -x12 - x13 = -1
8 C1 \leftarrow c(-1,-1,0,0,0,0,0,0,0)
9 bc1<- -1
10
11 # Constraint 2 : x12 - x23 - x24 = 0
12 C2 \leftarrow c(1,0,-1,-1,0,0,0,0,0)
13 bc2<-0
14
15 # Constraint 3 : x13 + x23 - x35 - x36 = 0
16 C3 \leftarrow c(0,1,1,0,-1,-1,0,0,0)
17 bc3<-0
18
19 # Constraint 4 : x24 - x45 - x46 = 0
20 \quad C4 \leftarrow c(0,0,0,1,0,0,-1,-1,0)
21 bc4<-0
22
23 # Constraint 5 : x35 + x45 - x56 = 0
24 C5 \leftarrow c(0,0,0,0,1,0,1,0,-1)
```

Advanced Linear Programming

R code Exa 7.1.2 All basic feasible and infeasible solutions of an equation

```
1 ##Chapter 7: Advanced Linear Programming
2 \# \text{Example } 1-2 : \text{Page } 300
4 #Creating the A and b matrix
5 A=matrix(c(1,3,-1,2,-2,-2),nrow=2,byrow=T)
6 \text{ b=matrix}(c(4,2),nrow=2)
7
8
9 s<-character()</pre>
10 basic <- character()
11 Type <- character()</pre>
12
13 for(i in 3:1){
14
     for(j in i:3){
       ##if i!=j and ith and jth column are not
15
           linearly dependent
16
        if(i!=j & !all(abs(A[,i]/A[,j])==c(1,1))){
17
          ##Solve for X with the ith and jth column
          a < - solve (A[,c(i,j)]) % * % b
18
          s<-rbind(s,toString(a))</pre>
19
          basic <-rbind(basic,c(paste(i,"&",j)))
20
```

R code Exa 7.1.3 Simmplex tableau in matrix form

```
1 ##Chapter 7: Advanced Linear Programming
2 ##Example 1-3 : Page 303
4 # Objective function : Max x1 + 4*x2 + 2*x3 + 5*x4
5 \text{ a} \leftarrow c(1,4,7,5)
7 # Constraint 1 : 2*x1 + x2 + 2*x3 + 4*x4 = 10
8 \text{ C1} \leftarrow c(2,1,2,4)
9 bc1<- 10
10
11 # Constraint 2 : 3*x1 - x2 - 2*x3 + 6*x4 = 5
12 C2 \leftarrow c(3,-1,-2,6)
13 bc2<-5
14
15 A <- rbind (C1, C2)
16 b <- rbind (bc1, bc2)
17
18 Binv<-solve(A[,c(1,2)])
19 X < - Binv % * % b
20
21 Binv
22 X
```

R code Exa 7.2.1 Revised simplex algoithmr

```
1 ##Chapter 7: Advanced Linear Programming
2 \# \text{Example } 2-1 : \text{Page } 309
3
4 # Objective function
5 \text{ a} \leftarrow c(5,4,0,0,0,0)
6
7 #Constraints
8 C \leftarrow rbind(c(6,4,1,0,0,0),c(1,2,0,1,0,0),c
      (-1,1,0,0,1,0),c(0,1,0,0,0,1))
9 b<-rbind(24,6,1,2)
10 library ("lpSolve")
11
12 solution <-lp("max", a, C, rep("=",4), b)
13 solution $ objval
14 solution $ solution
15
16 ##This solver is based on the revised simplex method
```

R code Exa 7.3.1 Bounded variables algorithm

```
10
11 solution <-lp("max", a, C, rep("<=",6), b)
12 solution $ objval
13 solution $ solution</pre>
```

R code Exa 7.4.1 Dual simplex algorithm

```
1 ##Chapter 7 : Advanced Linear Programming
2 \# \text{Example } 4-1 : \text{Page } 322
3 # Objective function
4 \text{ w} \leftarrow c(3,5,0,0)
6 #Constraints
7 C \leftarrow rbind(c(1,2,1,0),c(-1,3,0,1))
8 b<-rbind(5,2)</pre>
9
10 B < -C[, c(1,4)]
11 Binv <-solve(B)
12
13 #Associated primal and dual variables are
       assocprimal and assocdual
14 assocprimal <-Binv%*%b
15 assocdual \leftarrow w[c(1,4)] \% *\%Binv
16
17 ##the objective values are
18 primalobj <-c(3,0) %*% assocprimal
19 dualobj -assocdual \% *\% c (5,0)
20
21 primalobj
22 dualobj
```

Integer Linear Programming

R code Exa 8.1.1 Project selection

```
1 ##Chapter 8 : Integer Linear Programming
2 ##Example 1-1 : Page 336
3
4 # Objective function
5 \text{ a} \leftarrow c(20,40,20,15,30)
7 # Constraint 1 :
8 C1 \leftarrow c(5,4,3,7,8)
9 bc1<- 25
10
11 # Constraint 2
12 \quad C2 \leftarrow c(1,7,9,4,6)
13 bc2<-25
14
15 # Constraint 3
16 \quad C3 \leftarrow c(8,10,2,1,10)
17 bc3<-25
18
19 library ("lpSolve")
20 solution \leftarrow lp ("max", a, rbind (C1, C2, C3), rep ("<=",3)
       , c(bc1, bc2,bc3),all.bin=T)
```

```
21 solution $objval
22 solution $solution
```

R code Exa 8.1.2 Installing security phones

```
1 ##Chapter 8: Integer Linear Programming
 2 \# \text{Example } 1-2 : \text{Page } 340
4 # Objective function
 5 \text{ a} \leftarrow \text{rep}(1,8)
 7 \# Constraint 1 :
8 C1 \leftarrow c(1,1,rep(0,6))
9 bc1<- 1
10
11 # Constraint 2
12 C2 \leftarrow c(0,1,1,rep(0,5))
13 bc2<-1
14
15 # Constraint 3
16 \quad C3 \leftarrow c(0,0,0,1,1,0,0,0)
17 bc3<-1
18
19 # Constraint 4
20 \text{ C4} \leftarrow c(rep(0,6),1,1)
21 bc4<-1
22
23 # Constraint 5
24 \text{ C5} \leftarrow c(rep(0,5),1,1,0)
25 bc5<-1
26
27 # Constraint 6
28 \quad C6 \leftarrow c(0,1,0,0,0,1,0,0)
29 bc6<-1
30
```

```
31 # Constraint 7
32 \quad C7 \leftarrow c(1,0,0,0,0,1,0,0)
33 bc7<-1
34
35 # Constraint 8
36 \quad C8 \leftarrow c(0,0,0,1,0,0,1,0)
37 bc8<-1
38
39 # Constraint 9
40 \quad C9 \quad \leftarrow \quad c(0,1,0,1,0,0,0,0)
41 bc9<-1
42
43 \# Constraint 10
44 C10 \leftarrow c(0,0,0,0,1,0,0,1)
45 bc10<-1
46
47 # Constraint 11
48 C11 \leftarrow c(0,0,1,0,1,0,0,0)
49 bc11<-1
50
51 library("lpSolve")
52 solution <-lp("min", a, rbind(C1, C2,C3,C4,C5,C6,C7,
       C8, C9, C10, C11), rep(">=",11), c(bc1, bc2, bc3, bc4,
       bc5, bc6, bc7, bc8, bc9, bc10, bc11), all.bin=T)
53 solution $ objval
54 solution $ solution
```

R code Exa 8.1.3 Choosing a telephone company

```
1 ##Chapter 8 : Integer Linear Programming
2 ##Example 1-3 : Page 346
3
4 # Objective function
5 a <- c(0.25,0.21,0.22,16,25,18)</pre>
```

```
7 # Constraint 1 :
8 C1 \leftarrow c(1,1,1,0,0,0)
9 bc1<- 200
10
11 # Constraint 2
12 \quad C2 \leftarrow c(1,0,0,-200,0,0)
13 bc2<-0
14
15 # Constraint 3
16 \quad C3 \leftarrow c(0,1,0,0,-200,0)
17 bc3<-0
18
19 # Constraint 4
20 \quad C4 \leftarrow c(0,0,1,0,0,-200)
21 bc4<-0
22
23 library ("lpSolve")
24 solution <-lp("min", a, rbind(C1, C2,C3,C4), c("=","
      <= ","<= ","<= "), {\color{red}c} (bc1, bc2,bc3,bc4),int.vec={\color{red}c}
       (4,5,6))
25 solution $ objval
26 solution $ solution
```

R code Exa 8.1.4 Job scheeduling model

```
1 ##Chapter 8 : Integer Linear Programming
2 ##Example 1-4 : Page 350
3
4 M<-1000
5 # Objective function
6 a <- c(0,0,0,0,0,0,19,0,12,0,34,0)
7
8 # Constraint 1 :
9 C1 <- c(1,-1,0,M,0,0,0,0,0,0,0)
10 bc1<- 20</pre>
```

```
11
12 # Constraint 2
13 C2 \langle -c(-1,1,0,-M,0,0,0,0,0,0,0,0,0)
14 bc2<-5-M
15
16 # Constraint 3
17 C3 \leftarrow c(1,0,-1,0,M,0,0,0,0,0,0,0)
18 bc3<-15
19
20 # Constraint 4
21 C4 \leftarrow c(-1,0,1,0,-M,0,0,0,0,0,0,0)
22 \text{ bc4} \leftarrow 5 - M
23
24 # Constraint 5
25 \quad C5 \quad \leftarrow \quad c(0,1,-1,0,M,0,0,0,0,0,0,0,0)
26 bc5<-15
27
28 # Constraint 6
29 C6 \leftarrow c(0,-1,1,0,-M,0,0,0,0,0,0,0)
30 \text{ bc6} < -20 - M
31
32
33 # Constraint 7
34 \quad C7 \leftarrow c(1,0,0,0,0,0,-1,1,0,0,0,0)
35 \text{ bc7} < -20
36
37 # Constraint 8
38 C8 \leftarrow c(0,1,0,0,0,0,0,0,-1,1,0,0)
39 bc8<-2
40
41 # Constraint 9
42 C9 \leftarrow c(0,0,1,0,0,0,0,0,0,0,-1,1)
43 bc9<-20
44
45 library ("lpSolve")
46 solution <-lp("min", a, rbind(C1, C2,C3,C4,C5,C6,C7,
       C8,C9), c(rep(">=",6),rep("=",3)), c(bc1, bc2,bc3)
       , bc4, bc5, bc6, bc7, bc8, bc9), int.vec=c(4,5,6))
```

```
47 solution sobjval 48 solution solution
```

R code Exa 8.2.1 Branch and bound algorithm

```
1 ##Chapter 8: Integer Linear Programming
2 ##Example 2-1 : Page 356
4 # Objective function
5 \text{ a} < -c(5,4)
7 # Constraint 1 :
8 \text{ C1} \leftarrow c(1,1)
9 bc1<- 5
10
11 # Constraint 2
12 C2 < -c(10,6)
13 bc2<-45
14
15 library ("lpSolve")
16 solution \leftarrow 1p ("max", a, rbind (C1, C2), rep ("<=",2), c
      (bc1, bc2), int.vec=c(1,2)
17 solution $ objval
18 solution $ solution
19
20 ## lpSolve solver is based on the revised simplex
      method and a branch-and-bound (B&B) approach.
```

R code Exa 8.2.2 Cutting plane algorithm

```
1 ##Chapter 8 : Integer Linear Programming
2 ##Example 2-2 : Page 364
3
```

```
4 # Objective function
5 \text{ a} < -c(7,10)
7 # Constraint 1 :
8 \text{ C1} \leftarrow c(-1,3)
9 bc1<- 6
10
11 # Constraint 2
12 C2 < -c(7,1)
13 bc2<-35
14
15 library("lpSolve")
16 solution \leftarrow lp ("max", a, rbind (C1, C2), rep ("<=",2), c
      (bc1, bc2), int.vec=c(1,2))
17 solution $ objval
18 solution $ solution
19
20 ## lpSolve solver is based on the revised simplex
      method and a branch-and-bound (B\&B) approach.
```

Heuristic Programming

R code Exa 9.2.1 Discrete variable heuristics

```
1 ##Chapter 8 : Heuristic Programming
2 ##Example 2-1 : Page 382
4 ##Initializing function F, Random number R and index
5 \text{ F} = \text{c} (90,60,50,80,100,40,20,70)
6 \quad R = 0.1002
7 index=ceiling(R*8)
8 #while any of neighbouring values are lower, keep
      looping
9 while (F[index]>=F[index+1] | F[index]>=F[ifelse(
      index == 1, 1, index - 1) ]) {
10
     #if the next index is less than the current index
     if (F[index+1] < F[index]) {</pre>
11
12
        index=index+1
13
     #if the previous index is less than the current
14
15
     if (F[index -1] < F[index]) {</pre>
       index=index+1
16
17
18 }
```

```
19 index20 F[index]
```

R code Exa 9.2.2 Random walk heuristic

```
1 ##Chapter 8 : Heuristic Programming
2 ##Example 2-2 : Page 383
4 ##Initializing function F, Random number R, current
      index and index
5 F = c(90,60,50,80,100,40,20,70)
6 \quad \mathbf{R} = 0.1002
7 \quad count = 0
8 current=ceiling(R*8)
9 index<-ceiling(runif(1,min = 0,max = 8))</pre>
10 #while count < 3, keep looping
11 while (count <3) {</pre>
12
13
     count = count +1
14
     if (F[index] < F[current]) {</pre>
       #set current to index and a new random value is
15
           alloted to index
        current=index
16
17
        index<-ceiling(runif(1,min = 0,max = 8))</pre>
       #count is set to 0 when another index lesser
18
           than current is found
19
        count=0
20
     }
21 }
22 current
23 F[current]
```

R code Exa 9.2.3 Random walk heuristic for continuous variables

```
1 ##Chapter 8 : Heuristic Programming
2 \text{ ##Example } 2-3 : Page 385
4 #Function F returns the value of F for a given x
5 F \leftarrow function(x) \{return(x**5-10*x**4+35*x**3-50*x**
      2+24*x)
6
7 #Intial values
8 x = 0.5
9 \text{ newx=5}
10 count = 0
11 #while count <3
12 while(count <3){</pre>
     #generate news random uniform sampling
13
     while (\text{newx} > 4 \mid \text{newx} < 0) \{ \text{newx} < -x + (\text{runif}(1) - 0.5) *
14
         (4-0)}
     #if the new value is better than the old
15
16
     if (F(newx) < F(x))
        #assign news as x, news as 5 (so that it enters
17
           news while loop in next iteration) and count
           to 0
18
        x = newx
19
        newx=5
20
        count=0
     }else{
21
22
       #else increment count
23
        count = count +1
     }
24
25 }
26
27 print(paste("Optimal using uniform distribution = ",
      x))
28
29 x = 0.5
30 \text{ newx}=5
31 count=0
32 #while count<3
33 while(count <3){
```

```
#generate news random uniform sampling
34
      while (\text{newx} > 4 \mid \text{newx} < 0) \{ \text{newx} < -x + (4 * \text{rnorm} (1) / 6) \}
35
      if (F(newx) < F(x))
36
        #assign news as x, news as 5 (so that it enters
37
            newx while loop in next iteration) and count
            to 0
38
        x = n e w x
39
        newx=5
        count=0
40
      }else{
41
        #else increment count
42
43
        count = count +1
44
        newx=5
45
      }
46 }
47
48 print(paste("Optimal using normal distribution = ",x
      ))
```

 ${f R}$ code ${f Exa}$ 9.3.1 Minimization of single varibale function using tabu search algorithm

```
13
         }
       }
14
15
     }
     return(A)
16
17 }
18
19 #for 100 repetitions
20 for(m in 1:100){
     #take a random uniform index
21
22
     index=ceiling(runif(1)*8)
23
     Tenure=3
24
     L=c(0)
25
     #a random index from the neighbourhood not in tabu
         list
     i=ceiling(runif(1)*sum(sum(!(N(index) %in% L))))
26
27
28
     #while new index is lesser than old
29
     while(F[index]>F[N(index)[!(N(index) %in% L)]][i])
       #add current node to tabu list
30
31
       L<-c(L, index)</pre>
32
       #pick new node
       index=N(index)[!(N(index) %in% L)][i]
33
34
       #remove the node after their tenure
       if(i>2){
35
36
         L<-L[-1]
37
       }
       ##pick the next index
38
39
       i=ceiling(runif(1)*sum(sum(!(N(index) %in% L))))
40
     #counter for bar plot
41
     G[index]=G[index]+1
42
43 }
44
45 barplot(G/sum(G), names.arg = 1:8)
```

R code Exa 9.3.3 Minimization of single varibale function using simulated annealing algorithm

```
1 ##Chapter 8 : Heuristic Programming
2 \# Example 3-3 : Page 396
3 library(GenSA)
4 #function that returns the function value
5 Fx<-function(x){</pre>
     F = c(90,60,50,80,100,40,20,70)
7
     return(F[ceiling(x)])
8 }
9 #GenSA is a continues function, But as you see in the
       return value in function Fx, we return the value
       of the ceiling of the input
10 Solution=GenSA(fn=Fx,lower=c(1),upper=c(8),control =
      list(smooth =F))
11 Solution $ value
12 ceiling(Solution$par)
```

R code Exa 9.3.4 Job scheduling using simulated annealing algorithm

```
1 ##Chapter 8 : Heuristic Programming
2 ##Example 3-1 : Page 388
3
4 #function to return the cost of a sequence
5 Fx<-function(index1){
6    cost<-0
7    date<-0
8    ctable=matrix(c
        (10,15,3,10,8,20,2,22,6,10,5,10,7,30,4,8),nrow
= 4,byrow=T)</pre>
```

```
for(i in 1:4){
10
11
        date<-date+ctable[index1[i],1]</pre>
12
        if (date > ctable[index1[i],2]){
13
14
          cost<-cost+(date-ctable[index1[i],2])*ctable[</pre>
             index1[i],4]
15
        }else{
          cost<-cost+(ctable[index1[i],2]-date)*ctable[</pre>
16
             index1[i],3]
        }
17
18
19
     return(cost)
20 }
21
22 #function to get the neighbourhoods of a sequence
23 N<-function(index1){
24
     swap \leftarrow function(x,i) {x[c(i,i+1)] \leftarrow x[c(i+1,i)];
          x}
     A<-array(numeric(),c(0,4))
25
     for(i in 1:3){
26
        A<-rbind(A,swap(index1,i))</pre>
27
28
     }
     return(A)
29
30 }
31
32 #Initializations
33 index1=sample(c(1,2,3,4))
34 i=1; Ti \leftarrow Fx (index1)
35
36 #for 60 repetitions
37 while(i<=50){
     Ti = 0.5 * Ti
38
39
     p = 0
     ##Step 2 of the algorithm
40
     if(p<3){
41
        Neighbours <-N(index1)</pre>
42
        RandomN<-Neighbours[ceiling(runif(1,0,3)),]</pre>
43
        if (Fx(index1) < Fx(RandomN)) {</pre>
44
```

```
if (runif(1) <exp(-abs(Fx(index1)-Fx(RandomN))/</pre>
45
             Ti)){
            index1 <- RandomN
46
47
            p=p+1
48
            i = i + 1
          }else{
49
50
            i=i+1
          }
51
52
        }else{
53
          index1 <- RandomN
54
          p=p+1
55
          i=i+1
        }
56
     }else{
57
        i=i+1
58
     }
59
60 }
61 Fx(index1)
```

Travelling Salesman Problem

 ${f R}$ code Exa 10.3.1 Travelling Salesman Problem using branch and bound algorithm

```
##Chapter 10 : Travelling Salesman Problem
##Example 3-1 : Page 438

#TSP library
library("TSP")
#Distance matrix
dij<-matrix(c(Inf,10,3,6,9,5,Inf,5,4,2,4,9,Inf,7,8,7,1,3,Inf,4,3,2,6,5,Inf),nrow=5,byrow = T)
#making a asymetric TSP instance
stsp<-ATSP(dij)
#Solve TSP
d=solve_TSP(atsp,method = "nearest_insertion", control = list(start,"1"))
d
d d[1:5]</pre>
```

 ${f R}$ code Exa 10.3.2 Travelling Salesman Problem using cutting plane algorithm

```
##Chapter 10 : Travelling Salesman Problem
##Example 3-2 : Page 441

#TSP library
library("TSP")
#Distance matrix
dij<-matrix(c(Inf,13,21,26,10,Inf,29,20,30,20,Inf,5,12,30,7,Inf),nrow=4,byrow = T)

#making a asymetric TSP instance
stsp<-ATSP(dij)
#Solve TSP
d=solve_TSP(atsp,method = "nearest_insertion", control = list(start,"1"))

d
d d[1:4]</pre>
```

 $\bf R$ code Exa $\bf 10.4.1$ Travelling Salesman Problem using nearest neighbour heuristic

```
1 ##Chapter 10 : Travelling Salesman Problem
2 ##Example 3-1: Page 443
4 #Distance matrix
5 dij <-matrix (c(Inf, 120, 220, 150, 210, 120, Inf
      ,100,110,130,220,80, Inf,160,185,150, Inf,160, Inf
      ,190,210,130,185,Inf,Inf),nrow=5,byrow = T)
6 #Source node
7 source=3
8 #All edges leading to 3 has infinite length
9 dij[,3]=Inf
10 tour <- source
11 ##Nearest neighbour heuristic
12 for(i in 1:4){
13
       #choose nearest neighbour to the node
14
     mini <- which . min (dij [source,])</pre>
```

```
#add it to the tour
tour<-c(tour,mini)
#make it as the source for next iteration
source=mini
#set all edges leading to current node as Inf
dij[,mini]=Inf
}</pre>
```

Deterministic Dynamic Programming

R code Exa 11.1.1 Shortest route problem using dynamic programming

```
1 ##Chapter 11 : Deterministic Dynamic Programming
2 ##Example 1-1: Page 461
4 #Create a matrix with all Inf
5 d <- matrix (Inf, 7, 7)
6 #Add edge weights/lengths
7 d[8]=7; d[15]=8; d[22]=5; d[30:32]=c(12,8,7); d[38:39]=c
      (9,13); d[47:48] = c(9,6)
8 #Dynamic algorithm for Shortest path algorithm
9 ShortestDistance <-function (node, stage) {
10
     index <-which (d[1:7, node]!=Inf)</pre>
11
     #if it is node 1 at stage 2, return 1
12
13
     if(c(1) %in% index & stage==2){
14
       return(d[1,node])
     #else return the minimum distance of all possible
15
        nodes from current node
16
     }else{
17
       dist<-numeric()</pre>
```

Deterministic Inventory Modelling

R code Exa 12.2.1 Analyzing the nature of demand

```
1 ##Chapter 12 : Deterministic Inventory Modelling
2 \# Example 2-1 : Page 493
4 ##Natural gas consumption details
5 d<-matrix(c
      (100,110,90,70,65,50,40,42,56,68,88,95,110,125,98,80,60,53,44,45,
                88,79,56,57,38,39,60,70,82,90,121,130,95,90,70,58,41,44,
6
                68,55,43,41,65,79,88,94,130,122,100,85,73,58,42,43,64,75
7
                55, 45, 40, 67, 78, 98, 97, 130, 115, 100, 95, 80, 60, 49, 48, 64, 85, 96
8
9
                39,69,90,100,110,87,80,78,75,69,48,39,41,50,70,88,93)
                   ,byrow=T, ncol=12)
10 #row and coloumn names
11 rownames (d) <-1990:1999
12 colnames (d) <-month.abb
13
```

R code Exa 12.3.1 Classic EOQ model problem

```
1 ##Chapter 12 : Deterministic Inventory Modelling
2 ##Example 3-1 : Page 495
4 #input parametres
5 D = 100
6 K = 100
7 h = 0.02
8 L = 12
9 #optimal quantity
10 y = sqrt(2*K*D/h)
11
12 print(paste("EOQ =",y,"units"))
13 #cycle length
14 t = y/D
15 print(paste("Cycle length =",t,"days"))
16 #number of cycles
17 L0=L-floor(L/t)*t
18
19 print(paste("Reorder point =",L0*D,"units"))
20 print(paste("Daily inventory cost = \$", (K*D/y)+h*y/
     2))
```

R code Exa 12.3.2 Optimal order policy

```
1 ##Chapter 12 : Deterministic Inventory Modelling
2 \# \text{Example } 3-2 : \text{Page } 501
4 #input parameters
5 \quad D = 187.5
6 h = 0.02
7 K = 20
8 L=2
9 c1 = 3
10 c2=2.5
11 \quad q = 1000
12 #optimal quantity
13 ym = sqrt(2*K*D/h)
14 #function to calculate the least cost
15 if (q<ym) {
     Y = ym
16
17 }else{
18
     x = polyroot(c(2*K*D/h, (2*(c2*D-(c1*D+(K*D/ym)+(h*ym)))))
         /2) ))/h),1))
     if(Re(x[2])>q){
19
20
        Y = q
21
     }else{
22
        Y = ym
     }
23
24 }
25 Y
```

R code Exa 12.4.1 Dynamic EOQ models with no setup

```
1 ##Chapter 12 : Deterministic Inventory Modelling
2 ##Example 4-1 : Page 508
3
4 #Setting up initial cost matrix and other values
5 cost<-matrix(Inf,8,5)
6 for(i in 1:4){</pre>
```

```
for (j in i:4){
7
       cost[2*i-1,j]=6 + (j-i)*0.1
8
       cost[2*i,j]=9 + (j-i)*0.1
9
     }
10
11 }
12 cost[,5]=0
13 totalcost=0
14 Regular=c(90,100,120,110)
15 Overtime=c(50,60,80,70)
16 Demand=c(100,190,210,160,20)
17 Supply = c(rbind(Regular, Overtime))
18
19 #Setting up the resultant table
20 rown=numeric()
21 for(i in 1:4) {rown=c(rown, paste("R", i, sep = ""),
      paste("O",i,sep = ""))}
22 allocation=matrix(0,8,5,dimnames = list(rown,c(1:4,"
      Surplus")))
23
24
25 for(i in 1:5){
26
     while (Demand[i]>0) {
27
       minindex=which.min(cost[,i])
28
       #if it can still meet the demand
29
30
       if (Supply[minindex]>0){
31
         #allocate min of demand or supply
32
         allocation[minindex,i]=min(Demand[i],Supply[
            minindex])
33
         #calculate cost
         totalcost=totalcost+allocation[minindex,i]*
34
            cost[minindex,i]
         ##Subtract the allocated form supply and
35
            demand
         tmp=Demand[i]
36
         Demand[i] = Demand[i] - min (Demand[i], Supply[
37
            minindex])
         Supply [minindex] = Supply [minindex] -min(tmp,
38
```

```
Supply[minindex])
         cost[minindex,i]=Inf
39
       #else set all cost of that week to Inf so that
40
          it doesnt get chosen again
41
       }else{
         cost[minindex,]=Inf
42
43
       }
44
45
46
     }
47 }
48 allocation
49 totalcost
```

Decision Analysis and Games

R code Exa 13.1.1 Overall Idea of AHP

```
1 ##Chapter 13: Decision Analysis and Games
2 ##Example 1-1 : Page 527
4 #Setting up input values
5 LocationC=c(12.9, 27.2, 59.4)
6 ReputationC=c(54.5, 27.3, 18.2)
7 LocationP=0.17
8 ReputationP=0.83
9 Criteria=c(LocationP, ReputationP)
10 #composite weights
11 Choice=Criteria%*%rbind(LocationC, ReputationC)
12
13 colnames (Choice) <-c("UofA", "UofB", "UofC")
14
15
16 #Setting up the data tree for the particular problem
17 #If you have trouble installing the package/library,
       please reinstall R form the following link: https
      ://cran.r-project.org/bin/
18 library(data.tree)
19 SelectAUniversity <- Node$new("Select a University")
```

```
Location <- Select AUniversity $ AddChild (paste ("
20
        Location", LocationP))
21
       for(child in 1:3){
          Location $ AddChild (paste ("Uof", LETTERS [child],
22
             LocationC[child]))
23
     Reputation <- Select AUniversity $ AddChild (paste("
24
        Reputation", ReputationP))
25
       for(child in 1:3){
          Reputation $AddChild (paste ("Uof", LETTERS [child
26
             ],LocationC[child]))
27
       }
28 print(SelectAUniversity)
29 print (Choice)
```

R code Exa 13.1.2 Weights and consistency for AHP

```
1 ##Chapter 13 : Decision Analysis and Games
2 ##Example 1-2 : Page 530
3
4 A=matrix(c(1,0.2,5,1),nrow = 2,byrow = T)
5 rowMeans(sweep(A,2,colSums(A),'/'))
6 AL=matrix(c(1,0.5,0.2,2,1,0.5,5,2,1),nrow = 3,byrow = T)
7 rowMeans(sweep(AL,2,colSums(AL),'/'))
8 AR=matrix(c(1,2,3,0.5,1,1.5,1/3,2/3,1),nrow = 3,byrow = T)
9 rowMeans(sweep(AR,2,colSums(AR),'/'))
```

R code Exa 13.1.3 Consistency ratio

```
1 ##Chapter 13 : Decision Analysis and Games
2 ##Example 1-3 : Page 533
```

R code Exa 13.2.1 Decision tree

```
##Chapter 13 : Decision Analysis and Games
##Example 2-1 : Page 538

CompanyA=c(5000,-2000)
CompanyB=c(1500,500)
#Probability of occurance
POfOccurance=c(0.6,0.4)
#Espected return of stock A and B
ExpectedStockA=POfOccurance%*%CompanyA
ExpectedStockB=POfOccurance%*%CompanyB
ExpectedStockA
ExpectedStockA
ExpectedStockB
```

R code Exa 13.2.2 Expected value criterion

```
1 ##Chapter 13 : Decision Analysis and Games
2 ##Example 2-2 : Page 544
3
```

```
4 #Probability matrix
5 \text{ CondP=matrix}(c(0.9,0.1,0.5,0.5),nrow=2,byrow = T)
7 PriorP=c(0.6,0.4)
8 JointP=CondP*PriorP
9 AbsP=colSums(JointP)
10 BayesP=sweep(JointP,2,AbsP,'/')
11
12 CompanyA = c(5000, -2000)
13 CompanyB = c (1500, 500)
14
15 ExpectedStockAat4=BayesP[,1]%*%CompanyA
16 ExpectedStockBat5=BayesP[,1]%*%CompanyB
17 ExpectedStockAat6=BayesP[,2]%*%CompanyA
18 ExpectedStockBat7=BayesP[,2]%*%CompanyB
19
20 ExpectedStockAat4
21 ExpectedStockBat5
22 ExpectedStockAat6
23 ExpectedStockBat7
```

R code Exa 13.3.1 Decision under uncertainity

```
min(E)*1000

#Minimax criterion
print("Minimax")

E=apply(costmatrix,1,max)
min(E)*1000

#Savage criterion
print("Savage")

r=sweep(costmatrix,2,apply(costmatrix,2,min))

E=apply(r,1,max)
min(E)
```

R code Exa 13.4.1 Two person zero sum game

R code Exa 13.4.3 Mixed strategy games

```
1 ##Chapter 13: Decision Analysis and Games
```

```
2 \# Example 4-3 : Page 559
4 #plot an empty graph
5 \text{ plot}(1, \text{ type="n"}, \text{ axes=F,xlab = "x1", ylab = "x2",}
      xlim = c(-0.5, 1.5), ylim = c(-2,7), col = "red",
         yaxs = "i", xaxs = "i")
6
7
8 #Add custom axis on sides 1 and 2
9 \text{ axis}(\text{side} = 1, \text{pos} = 0, \text{at} = \text{seq}(0, 1, 0.25))
10 axis(side = 2, pos = 1, at = seq(-2, 9, 2), padj = 3.5)
11 axis(side = 2, pos = 0, at = seq(-2, 9, 2))
12
13 #Add line segments for the constriants
14 segments (0,6,1,-1,col = "red")
15 segments (0,4,1,2,col = "yellow")
16 segments (0,3,1,2,col = "green")
17 segments (0,2,1,3,col = "blue")
18 #adding a line segment to indicate the maximum
19 segments (0.5,0,0.5,2.5,col = "black",lwd = 1.5)
20 #Adding the name of the constriants on the graph
21 \text{ text} (0,4,"B1",pos=4,cex=0.5)
22 \text{ text}(0,3,"B2",pos=4,cex=0.5)
23 text(0,2,"B3",pos=4,cex=0.5)
24 \text{ text} (0,6,"B4",pos=4,cex=0.5)
```

Probabilistic Inventory Models

R code Exa 14.1.1 Probabilistic Inventory Models with normal distribution of demand

```
1 ##Chapter 14: Probabilistic Inventory Models
2 ##Example 1-1 : Page 575
4 #Daily demand
5 D = 100
6 #standard deviation
7 \text{ dev} = 10
8 #probability of running out of stock
9 \text{ alpha=0.05}
10 #lead time
11 L=2
12 #average demand during lead time
13 \text{ muL} = D * L
14 #standard deviation of demand during lead time
15 \text{ devL=} \text{sqrt}((\text{dev}**2)*L)
16 #optimal inventory level for reordering
17 \text{ x=qnorm}(0.05,100,10,lower.tail = F)
18 z = (x-D)/dev
19 B = devL*z
20 B
```

R code Exa 14.2.1 Newsvendor problem

```
1 ##Chapter 14: Probabilistic Inventory Models
2 \# \text{Example } 2-1 : \text{Page } 582
4 #holding cost
5 h = 25
6 #penalty cost
7 p = 45
8 #critical ratio
9 CR=p/(p+h)
10
11 D=matrix(c(200,220,300,320,340,0.1,0.2,0.4,0.2,0.1),
      nrow = 2, byrow = T)
12
13 #The demand is a normal distribuion
14 print("Case A")
15 ystar=qnorm(0.643,300,20)
16
17 #The demand is a discrete PDF
18 print("Case B")
19 CDF=numeric()
20 for(i in 1:5){CDF[i]=sum(D[2,1:i])}
21 rbind(D,CDF)
```

Markov Chain

R code Exa 15.2.1 Absolute probabilities after transitions

```
1 ##Chapter 15 : Markov Chain
2 ##Example 2-1 : Page 596
4 #Library expm is needed to get element-wise power of
      a matrix
5 library(expm)
6 P=matrix(c(0.3,0.6,0.1,0.1,0.6,0.3,0.05,0.4,0.55),
     nrow = 3, byrow = T)
7 P8=P%^%8
8 P16=P%^%16
10 #Steady-state probabilities
11 a1=c(1,0,0)%*%P
12 a8=c(1,0,0)%*%P8
13 a16=c(1,0,0)%*%P16
14
15 a1
16 a8
17 a16
```

R code Exa 15.3.1 Absorbing and transient states

```
1 ##Chapter 15 : Markov Chain
2 ##Example 3-1 : Page 598
3
4 ##Editing mistake in textbook
5 ##The last element of P should be one as (sum over j
      )(p)=1 for all i
6 P=matrix(c(0.2,0.5,0.3,0,0.5,0.5,0,0,1),nrow = 3,
      byrow = T)
7 library(expm)
8 P%~%100
```

R code Exa 15.3.2 Periodic states

```
1 ##Chapter 15 : Markov Chain
2 ##Example 3-2 : Page 599
3
4 #transition matrix
5 P=matrix(c(0,0.6,0.4,0,1,0,0.6,0.4,0),nrow = 3,byrow = T)
6 library(expm)
7 #n=2
8 P%^%2
9 #n=3
10 P%^%3
11 #n=4
12 P%^%4
13 #n=5
14 P%^%5
```

R code Exa 15.4.1 Steady state probabilities

R code Exa 15.4.2 Cost model

```
##Chapter 15 : Markov Chain
##Example 4-2 : Page 602

##transition matrix

A=matrix(c(0.6,-0.4,0.4,0.1,0.3,-0.45,1,1,1),nrow = 3,byrow = T)

b=c(rep(0,2),1)

#solve the linear equations
pi=solve(A,b)
#mean recurrance time
mu=1/pi
#expected abbual cost of fertilizer
```

```
bagsOfFertilizer=c(2,2*1.25,2*1.6)
sum(bagsOfFertilizer*50*pi)
```

R code Exa 15.5.1 Mean first passage time

R code Exa 15.6.1 Analysis of absorbing states

13 POfAbsorption

Queuing Systems

R code Exa 16.4.1 Pure birth model

R code Exa 16.4.2 Pure death model

```
1 ##Chapter 16 : Queuing Systems
2 \# \text{Example } 4-1 : \text{Page } 632
4 #probability of placing an order in any one day of
      the week
5 Pfunction <- function(mu,t,n,Slimit){</pre>
     #empty matrix for the table
     A \leftarrow matrix(0, nrow = 3, ncol = 7)
7
8
     p0=0
9
     for(t in 1:7){
10
       P = p0
11
       for(n in 1:5){
12
          P=P+((mu*t)^(Slimit-n) *exp(-mu*t))/factorial(
             Slimit-n)
       }
13
       #appending to the table
14
       A[1,t]=t
15
16
       A[2,t]=mu*t
17
       A[3,t] = round(P, digits = 4)
18
     }
19
     return(A)
20 }
21 Pfunction (3,7,5,18)
22
23 #average number of dozen roses discarded at the end
      of the week
24 discardedRoses <-function(N, mu, n, t){
     P = 0
25
     for(i in 1:n){
26
       P=P+i*((mu*t)^(N-i) * exp(-mu*t))/factorial(N-i)
27
     }
28
29
     return(P)
30 }
31 discardedRoses (18,3,18,7)
```

R code Exa 16.6.1 Measures of performance

```
1 ##Chapter 16 : Queuing Systems
2 ##Example 6-1 : Page 642
4 #Function to calculate p0
5 PO<-function(n, servers, temp){
     x = 1
     for(i in 1:n){
7
       if (i < 6) {</pre>
        x=x+(temp^i)/factorial(i)
9
10
11
          x=x+(temp^i)/(factorial(servers)*servers^(i-
             servers))
       }
12
     }
13
14
     return(1/x)
15 }
16 p0=P0(8,5,3)
17 p0
18
19 #calculate Pn
20 \text{ temp=3}
21 servers=5
22 A=matrix(0,2,8)
23 for(i in 1:8){
24
     if (i < 6) {</pre>
       A[1,i]=i
25
       A[2,i]=p0*(3^i)/factorial(i)
26
27
     }else{
       A[1,i]=i
28
       A[2,i]=p0*(temp^i)/(factorial(servers)*servers^(
29
          i-servers))
     }
30
31 }
32 A
33 #arrival rate
34 \quad lambda=6
```

```
35 #arrivals lost
36 lambdalost=lambda*A[2,8]
37 lambdalost
38 #effective arrivals
39 lamdaeff=lambda-lambdalost
40 lamdaeff
41 #average lengths in the systems
42 Ls=0*p0+sum(1:8*A[2,])
43 Ls
44 #waiting time in the systems
45 Ws=Ls/lamdaeff
47 #average lengths in the queue
48 \quad \text{Wq=Ws}-1/2
49 Wq
50 #average number of occupied spaces
51 cbar=lamdaeff/2
52 cbar
53 #parking lot utilization
54 utilization=cbar/servers
55 utilization
```

R code Exa 16.6.2 MM1 GDInfInf model

```
1 ##Chapter 16 : Queuing Systems
2 ##Example 6-2 : Page 645
3
4 #Queueing library to process different queueing models
5 library(queueing)
6 #creating a MM1 instance with the following parameters
7 x=NewInput.MM1(lambda=4, mu=6,n=25)
8 #Solving the model which returns a list
9 y=QueueingModel(x)
```

R code Exa 16.6.4 MM1 GDNInf model

```
##Chapter 16 : Queuing Systems
##Example 6-4 : Page 649

#Queueing library to process different queueing models

library(queueing)

#creating a MMK instance with the following parameters

x=NewInput.MM1K(lambda=4, mu=6,k=5)

#Solving the model which returns a list

y=QueueingModel(x)

summary(y)
```

R code Exa 16.6.5 MMc GDInfInf model

```
1 ##Chapter 16 : Queuing Systems
2 ##Example 6-5 : Page 653
3
4 #Queueing library to process different queueing models
5 library(queueing)
```

R code Exa 16.6.6 MMc GDNInf model

R code Exa 16.6.7 MMInf GDInfNEInf or Slef service models

```
1 ##Chapter 16 : Queuing Systems
2 \# \text{Example } 6-7 : \text{Page } 660
4 #Queueing library to process different queueing
     models
5 library (queueing)
6 #creating a MMc instance with the following
      parameters
7 x=NewInput.MMInf(lambda=12, mu=0.333)
8 #Solving the model which returns a list
9 y=QueueingModel(x)
10 summary(y)
11 estimate = (0.25*round(y$L)*1000)*(1-0.2)+(0.75*round(
     v$L)*1000)*(1+0.12)
12 estimate
13 ##The answer for the estimate is given wrong in the
     book
```

R code Exa 16.6.8 MMR GDKK or Machine servicing model

```
##Chapter 16 : Queuing Systems
##Example 6-8 : Page 662

#Queueing library to process different queueing models
library(queueing)

#Empty matrix table for Machine productivity
McProductivity=matrix(0,2,4)
#Making a matrix for different number of repairpersons
for(i in 1:4){
#creating a MMc instance with the following parameters
x=NewInput.MMCKK(lambda=0.5, mu=5, c=i, k=22)
```

R code Exa 16.7.1 MG1 GDInfInf model or Pollaczek Khintchine formula

```
1 ##Chapter 16 : Queuing Systems
2 \# \text{Example } 7-1 : \text{Page } 664
4 #arrival rate
5 \quad lambda=4
6 #Expectation of arrival
7 \quad EOfT=1/6
8 #variance of arrival
9 VarOfT=0
10 #Length of the system
11 Ls=lambda*E0fT+((lambda^2 *(E0fT^2 + VarOfT))/(2*(1-
      lambda*EOfT)))
12 #Length of the queue
13 Lq=Ls-lambda*EOfT
14 #Waiting time in the system
15 Ws=Ls/lambda
16 #Waiting time in the queue
17 Wq=Lq/lambda
18 Ls
19 Lq
20 Ws
21 \text{ Wq}
```

R code Exa 16.9.1 Cost models

```
1 ##Chapter 16 : Queuing Systems
2 \# \text{Example } 9-1 : \text{Page } 667
4 #Queueing library to process different queueing
      models
5 library(queueing)
6 #inputs
7 Speed=c(30,36,50,66)
8 OperatingCost=c(15,20,24,27)
10 #service rate for each model
11 ServiceRate=numeric()
12 for(i in 1:4){
     ServiceRate[i] = 24/(10000/(Speed[i]*60))
13
14 }
15 ServiceRate
16 Ls=numeric()
17
18 #length of system for each model
19 for(i in 1:4){
20
    #creating a MMc instance with the following
        parameters
21
     y=NewInput.MM1(lambda=4, mu=ServiceRate[i])
     #Solving the model which returns a list
22
23
     z=QueueingModel(y)
24
     Ls[i]=zL
25 }
26 Ls
27 #cost for the 4 models
28 \quad cost=matrix(0,4,3)
29 colnames(cost)=c("EOC", "EWC", "ETC")
30 for(i in 1:4){
```

```
31    cost[i,1]=24*OperatingCost[i]
32    cost[i,2]=80*Ls[i]
33    cost[i,3]=cost[i,1]+cost[i,2]
34  }
35  cost
```

R code Exa 16.9.2 Practical problem for MMc GDInfInf model

```
1 ##Chapter 16 : Queuing Systems
2 \# \text{Example } 9-2 : \text{Page } 670
4 #Queueing library to process different queueing
      models
  library(queueing)
7 Ls=numeric()
8 ETC=numeric()
9 for (c in 2:6){
    #creating a MMc instance with the following
10
        parameters
     y=NewInput.MMC(lambda=17.5, mu=10,c)
11
     #Solving the model which returns a list
12
     z=QueueingModel(y)
13
14
     Ls[c]=zL
     ETC[c] = 12*c + 50*Ls[c]
15
16 }
17 ETC[2:6]
```

R code Exa 16.9.3 Aspiration level model

```
1 ##Chapter 16 : Queuing Systems
2 ##Example 9-3 : Page 672
3
```

```
4 #Queueing library to process different queueing
     models
5 library(queueing)
6 Ws=numeric()
7 RO=numeric()
9
10 for (c in 2:8){
    #creating a MMc instance with the following
11
        parameters
     y=NewInput.MMC(lambda=17.5, mu=10,c)
12
    #Solving the model which returns a list
13
    z=QueueingModel(y)
14
    #waiting time in the system
15
    Ws[c] = z$W*60
16
    #idleness percentage of the system
17
     RO[c] = (1-z\$R0)*100
18
19 }
20 rbind(Ws,RO)[,2:8]
```

Simulation Modelling

R code Exa 17.1.1 Monte carlo sampling

```
1 ##Chapter 17 : Simulation Modelling
2 ##Example 1-1 : Page 681
4 #number of trials
5 n = 1000
6 #Setting up parameters for the plot
7 par(pty="s")
8 #plot empty graph
9 plot(1, axes=T, asp = 1, xlim = c(-4, 6), ylim = c(-3, 6)
       7))
10 #getting a sample of 1000 random uniform numbers
      within the given range
11 x = -4 + (10) * runif(n)
12 y = -3 + 10 * runif(n)
13 #counting the number of true values for the give
      expression
14 m = sum((x-1)^2+(y-2)^2<25)
15 #plotting the points on the graph
16 \text{ points}(x, y)
17 #plotting the circle
18 symbols(x=1,y=2,circles = 5,add = T,inches = F)
```

```
19 #plotting the square
20 symbols(x=1,y=2,squares =10,add = T,inches = F)
21 #Area of the circle
22 ApproxArea=m*100/n
23 ApproxArea
```

R code Exa 17.3.3 Erlang distribution

```
##Chapter 17 : Simulation Modelling
##Example 3-3 : Page 691

#Function which returns an erlang distribution
erlang=function(m,lambda){
    R=c(0.0589,0.6733,0.4799)
    y=-log(prod(R))/lambda
    return(y)
}
erlang(3,4)
```

R code Exa 17.4.1 Multiplicative congruential method

```
1 ##Chapter 17 : Simulation Modelling
2 ##Example 3-3 : Page 691
3
4 #inputs for the random number generator
5 b=9
6 c=5
7 m=12
8
9
10 u=numeric()
11 #initial random number
12 u[1]=11
```

```
13
14 R=numeric()
15 for(i in 2:4){
16  u[i]=(b*u[i-1]+c)%%m
17  R[i-1]=u[i]/m
18 }
19 R
```

Classical Optimization theory

R code Exa 18.1.1 Necessary and sufficient conditions

```
1 ##Chapter 18 : Classical Optimization theory
2 ##Example 1-1 : Page 714
4 #Function to be given as input to optim function
5 minimize <-function(x){</pre>
6
     x1 = x[1]
     x2=x[2]
     x3 = x[3]
     return(-(x1+2*x2+x2*x3-x1^2-x2^2-x3^2))
10 }
11
12 #Calculates the optimal value
13 \operatorname{optim}(c(0,0,0), \min i mize, hessian = T)
14 #the Jacobian for each
15 f = expression(x1+2*x2+x2*x3-x1^2-x2^2-x3^2)
16 for(i in 1:3){
     x=D(f,paste("x",i,sep = ""))
17
18
     print(x)
19 }
```

R code Exa 18.1.3 Newton Raphson method