Homework 1 solutions

COMS W3251, Summer Session 2

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1 Matrix and vector computations (20 points)

For each of the following expressions, compute the result. You can use properties of basic operations to simplify the expression:

$$\left(\begin{bmatrix} 2 & 3 & -1 \\ 0 & 1 & 1 \end{bmatrix}^T \begin{bmatrix} 2 & 3 & -1 \\ 0 & 1 & 1 \end{bmatrix} + 4I \right) e_2, \tag{1.1}$$

$$\left(\begin{bmatrix}2&1\\0&-1\end{bmatrix}^T + \begin{bmatrix}0&-1\\4&1\end{bmatrix}^T\right)^T \begin{bmatrix}2&3&-1\\0&1&1\end{bmatrix},\tag{1.2}$$

$$\left\langle \left(6 \begin{bmatrix} 2 & 3 & -1 \\ 0 & 1 & 1 \\ 1 & 0 & 0 \end{bmatrix} - 3I \right) (2e_1 + e_3), \begin{bmatrix} 1 \\ -1 \\ 1 \end{bmatrix} \right\rangle,$$
 (1.3)

$$\langle 2e_1 + 3e_2 + 4e_3, 8e_2 - 6e_3 \rangle.$$
 (1.4)

Solution:

1.

$$\begin{pmatrix}
\begin{bmatrix} 2 & 3 & -1 \\ 0 & 1 & 1 \end{bmatrix}^T \begin{bmatrix} 2 & 3 & -1 \\ 0 & 1 & 1 \end{bmatrix} + 4I \end{pmatrix} e_2 =
\begin{pmatrix}
\begin{bmatrix} 2 & 0 \\ 3 & 1 \\ -1 & 1 \end{bmatrix} \begin{bmatrix} 2 & 3 & -1 \\ 0 & 1 & 1 \end{bmatrix} + 4I \end{pmatrix} e_2 =
\begin{pmatrix}
\begin{bmatrix} 4 & 6 & -2 \\ 6 & 10 & -2 \\ -2 & -2 & 2 \end{bmatrix} + 4I \end{pmatrix} e_2 =
\begin{pmatrix}
\begin{bmatrix} 8 & 6 & -2 \\ 6 & 14 & -2 \\ -2 & -2 & 6 \end{bmatrix} e_2 =
\begin{pmatrix}
\begin{bmatrix} 6 \\ 14 \\ -2 \end{bmatrix}
\end{pmatrix}$$

2.

$$\begin{pmatrix}
\begin{bmatrix} 2 & 1 \\ 0 & -1 \end{bmatrix}^T + \begin{bmatrix} 0 & -1 \\ 4 & 1 \end{bmatrix}^T \end{pmatrix}^T \begin{bmatrix} 2 & 3 & -1 \\ 0 & 1 & 1 \end{bmatrix} = \\
\begin{pmatrix}
\begin{bmatrix} 2 & 1 \\ 0 & -1 \end{bmatrix} + \begin{bmatrix} 0 & -1 \\ 4 & 1 \end{bmatrix} \end{pmatrix} \begin{bmatrix} 2 & 3 & -1 \\ 0 & 1 & 1 \end{bmatrix} = \\
\begin{bmatrix} 2 & 0 \\ 4 & 0 \end{bmatrix} \begin{bmatrix} 2 & 3 & -1 \\ 0 & 1 & 1 \end{bmatrix} = \\
\begin{bmatrix} 4 & 6 & -2 \\ 8 & 12 & -4 \end{bmatrix}$$

3.

$$\left\langle \left(6 \begin{bmatrix} 2 & 3 & -1 \\ 0 & 1 & 1 \\ 1 & 0 & 0 \end{bmatrix} - 3I \right) (2e_1 + e_3), \begin{bmatrix} 1 \\ -1 \\ 1 \end{bmatrix} \right\rangle = \left\langle \left(\begin{bmatrix} 12 & 18 & -6 \\ 0 & 6 & 6 \\ 6 & 0 & 0 \end{bmatrix} - 3I \right) \begin{bmatrix} 2 \\ 0 \\ 1 \end{bmatrix}, \begin{bmatrix} 1 \\ -1 \\ 1 \end{bmatrix} \right\rangle = \left\langle \begin{bmatrix} 9 & 18 & -6 \\ 0 & 3 & 6 \\ 6 & 0 & -3 \end{bmatrix} \begin{bmatrix} 2 \\ 0 \\ 1 \end{bmatrix}, \begin{bmatrix} 1 \\ -1 \\ 1 \end{bmatrix} \right\rangle = \left\langle \begin{bmatrix} 12 \\ 6 \\ 9 \end{bmatrix}, \begin{bmatrix} 1 \\ -1 \\ 1 \end{bmatrix} \right\rangle = 15$$

4.

$$\langle 2e_1 + 3e_2 + 4e_3, 8e_2 - 6e_3 \rangle = \langle \begin{bmatrix} 2 \\ 3 \\ 4 \end{bmatrix}, \begin{bmatrix} 0 \\ 8 \\ -6 \end{bmatrix} \rangle = 0$$

2 Vector spaces (20 points)

Let $P_3(\mathbb{R})$ be the set of all polynomials of **degree 3 or less** with real coefficients. Prove that $P_3(\mathbb{R})$ is a vector space when considering + as point-wise sum and \cdot as point-wise scalar multiplication (that is, if $p_1, p_2 \in P_3(\mathbb{R})$, then $p_1 + p_2$ is the polynomial such that $p_1 + p_2(x) = p_1(x) + p_2(x)$, while $\alpha \cdot p_1$ is the polynomial such that $\alpha \cdot p_1(x) = \alpha p_1(x)$). Can you find a finite subset of $P_3(\mathbb{R})$ that spans $P_3(\mathbb{R})$?

Solution:

Properties of the sum:

1. Commutativity:

$$(p_1 + p_2)(x) = p_1(x) + p_2(x) = p_2(x) + p_1(x) = (p_2 + p_1)(x)$$

2. Associativity:

$$(p_1 + (p_2 + p_3))(x) = p_1(x) + p_2(x) + p_3(x) = ((p_1 + p_2) + p_3)(x)$$

3. **Identity element:** Let p_0 be the polynomial such that $p_0(x) = 0$ for every $x \in \mathbb{R}$. Then $p_0 \in P_3(\mathbb{R})$ and for every $p \in P_3(\mathbb{R})$ we have

$$(p + p_0)(x) = p(x) + p_0(x) = p(x)$$

4. Existence of an inverse For every $p \in P_3(\mathbb{R})$, let -p be the polynomial such that (-p)(x) = -p(x). Then

$$p + (-p)(x) = p(x) - p(x) = 0 = p_0(x)$$

Properties of the scalar multiplication:

1. Compatibility with field multiplication: Let $\alpha, \beta \in \mathbb{R}$. Then for every p

$$\alpha \cdot (\beta \cdot p)(x) = \alpha \cdot (\beta p(x)) = \alpha \beta p(x) = (\alpha \beta) \cdot p(x)$$

2. Existence of identity element:

$$1 \cdot p(x) = p(x)$$

3. Distributivity with respect to +: for every $\alpha \in \mathbb{R}$

$$\alpha \cdot (p_1 + p_2)(x) = \alpha p_1(x) + \alpha p_2(x) = (\alpha p_1 + \alpha p_2)(x)$$

4. Distributivity with respect to scalar addition: for every $\alpha, \beta \in \mathbb{R}$:

$$((\alpha + \beta) \cdot p)(x) = \alpha p(x) + \beta p(x) = (\alpha \cdot p)(x) + (\beta \cdot p)(x)$$

Finally, we need to prove that $P_3(\mathbb{R})$ is closed under linear combinations. That is, for every $\alpha, \beta \in \mathbb{R}$ and $p_1, p_2 \in P_3(\mathbb{R})$, then $\alpha p_1 + \beta p_2 \in P_3(\mathbb{R})$. But this is true because a linear combination of polynomials of degree less than or equal to 3 is still a polynomial of degree less than or equal to 3.

Finally, we know that all polynomials of degree less than or equal to 3 can be written as

$$p(x) = \alpha_0 + \alpha_1 x + \alpha_2 x^2 + \alpha_3 x^3$$

for some $\alpha_0, \alpha_1, \alpha_2, \alpha_3$ in \mathbb{R} . This means that every polynomial is a linear combination of 4 monomials $\{1, x, x^2, x^3\}$, so that $P_3(\mathbb{R}) = \text{span}(1, x, x^2, x^3)$.

3 Linear functions (25 points)

For each of the following function T, prove if they are linear or not. If they are, find the matrix A such that T(x) = Ax. If they are not, find a counterexample to the linearity condition.

$$T(x_1, x_2) = (x_1 + 4x_2, |x_2|) (3.1)$$

$$T(x_1, x_2, x_3) = (x_1 - 8x_3, 0, 3x_2)$$
(3.2)

$$T(x_1, x_2, x_3) = (x_1 + 3, x_2, x_3 - x_1)$$
(3.3)

$$T(x_1, x_2) = x_1 x_2 + x_2^2 (3.4)$$

$$T(x_1) = (0, x_1, x_1^2, x_1^4) (3.5)$$

$$T(x_1) = (x_1, x_1, 0, x_1) (3.6)$$

Solution:

1. Not linear: absolute value in the second component is not linear. Consider $v_1 = \begin{bmatrix} 0 \\ -1 \end{bmatrix}, v_2 = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$. Then

$$T(v_1 + v_2) = T(\begin{bmatrix} 0 \\ 0 \end{bmatrix}) = (0, 0)$$

while

$$T(v_1) + T(v_2) = (-4, 1) + (4, 1) = (0, 2) \neq T(v_1 + v_2)$$

2. Linear. In fact we have

$$A = \begin{bmatrix} T(e_1) & T(e_2) & T(e_3) \end{bmatrix} = \begin{bmatrix} 1 & 0 & -8 \\ 0 & 0 & 0 \\ 0 & 3 & 0 \end{bmatrix}$$

We can double-check by considering

$$\begin{bmatrix} 1 & 0 & -8 \\ 0 & 0 & 0 \\ 0 & 3 & 0 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} x_1 - 8x_3 \\ 0 \\ 3x_2 \end{bmatrix}$$

3. Non-linear: $T(0,0,0) = (3,0,0) \neq (0,0,0)$. Therefore

$$T(2(0,0,0)) = T(0,0,0) = (3,0,0)$$

while

$$2T(0,0,0) = (6,0,0) \neq T(2(0,0,0))$$

4. Non-linear: the second component is quadratic. For $v_1 = (0,1)$, we have

$$T(3v_1) = T(0,3) = 9$$

while

$$3T(v_1) = 3 \neq T(3v_1)$$

5. Non-linear: the third and fourth component are respectively quadratic and a fourth power. For $v_1 = 1$, we have

$$T(2v_1) = T(2) = (0, 2, 4, 16)$$

while

$$2T(v_1) = 2T(1) = (0, 2, 2, 2)$$

6. Linear. In this case the matrix is given by

$$A = [T(e_1)] = \begin{bmatrix} 1\\1\\0\\1 \end{bmatrix}$$

4 Transformations of the plane (15 points)

Find the linear or affine transformations that satisfy the desired properties and write it in the form T(x) = Ax + b:

- 1. The transformation $T: \mathbb{R}^2 \to \mathbb{R}^2$ sending the origin to itself and a triangle of vertices (0,0), (1,0), (0,1) to a triangle of vertices $(0,0), (\frac{\sqrt{2}}{2}, \frac{\sqrt{2}}{2}), (-\frac{\sqrt{2}}{2}, \frac{\sqrt{2}}{2})$.
- 2. The transformation $T: \mathbb{R}^2 \to \mathbb{R}^2$ sending the unit square $[0,1] \times [0,1]$ to the parallelogram with vertices (4,1), (4,2), (7,2), (7,3).

Solutions:

1. This is a rotation of $\frac{\pi}{4}$ in the counter-clockwise direction. Therefore it is given by the matrix"

$$A = \begin{bmatrix} \cos(\frac{\pi}{4}) & -\sin(\frac{\pi}{4}) \\ \sin(\frac{\pi}{4}) & \cos(\frac{\pi}{4}) \end{bmatrix} = \begin{bmatrix} \frac{\sqrt{2}}{2} & -\frac{\sqrt{2}}{2} \\ \frac{\sqrt{2}}{2} & \frac{\sqrt{2}}{2} \end{bmatrix}$$

Notice that $A\begin{bmatrix}0\\0\end{bmatrix}=\begin{bmatrix}0\\0\end{bmatrix}$ since it's linear. Moreover

$$A \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \begin{bmatrix} \frac{\sqrt{2}}{2} \\ \frac{\sqrt{2}}{2} \end{bmatrix}$$
$$A \begin{bmatrix} 0 \\ 1 \end{bmatrix} = \begin{bmatrix} \frac{-\sqrt{2}}{2} \\ \frac{\sqrt{2}}{2} \end{bmatrix}$$

2. This is an affine transformation, since $T(0,0)=(4,1)\neq (0,0)$. Therefore we know that $T(x)=Ax+\begin{bmatrix} 4\\1 \end{bmatrix}$. In order to find the residual matrix, we consider

a linear transformation $T_l(x) = T(x) - b$. Then $T_l(0,0) = (0,0)$, the origin is sent to the origin, $T_l(1,0) = T(1,0) - (4,1) = (7,2) - (4,1) = (3,1)$ and $T_l(0,1) = T(0,1) - (4,1) = (0,1)$. Then $T_l(x) = Ax$ for

$$A = \begin{bmatrix} 3 & 0 \\ 1 & 1 \end{bmatrix}$$

Therefore we have that

$$T(x) = T_l(x) + b = \begin{bmatrix} 3 & 0 \\ 1 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} + \begin{bmatrix} 4 \\ 1 \end{bmatrix}$$

5 Flop count (20 points)

Let $A, B \in \mathbb{R}^{n \times n}$ and $x, y \in \mathbb{R}^n$. Following the order given by parenthesis, calculate the complexity of the following operations in term of flops.

$$\langle (A(Bx)), (\alpha y) \rangle \tag{5.1}$$

$$\alpha(\langle (AB)x, y \rangle) \tag{5.2}$$

Which one has the lowest complexity? Can you reduce the complexity even more?

Solutions:

1, non-sparse case: (Bx) is a matrix-vector product, so it has n inner products, each with 2n-1 flops.

 (αy) is a scalar multiplication of a vector, so it has n flops.

(A(Bx)) is another matrix-vector product, so it has the same complexity as the previous one, with n(2n-1) flops.

Finally, we have the scalar product between two $n \times 1$ vectors, so there are 2n-1 flops.

$$tot = n(2n-1) + n + n(2n-1) + 2n - 1 = 4n^{2} + n - 1 \sim 4n^{2}$$

2, non-sparse case: (AB) is a matrix-matrix product, so it has n^2 operations, each with 2n-1 flops.

(AB)x is a matrix-vector product, so it has n operations, each with 2n-1 flops.

The inner product $\langle (AB)x, y \rangle$ has 2n-1 operations.

Finally, the multiplication by scalar α has 1 flop. Therefore

$$tot = n^{2}(2n - 1) + n(2n - 1) + 2n - 1 + 1 = 2n^{3} + n^{2} + n \sim 2n^{3}$$

Therefore the first method has a smaller computational complexity. Since there is always a matrix-vector product, the complexity cannot be lower than n(2n-1). However, a marginal improvement can be obtained by multipliying by α after taking the inner product.