Q1. Numerical Linear Algebra: (Dr. Peterson) Summer 2016

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a)

Stationary Iterative Method:

$$x^{k+1} = px^k + c$$

These are methods where the data in the equation remains fixed. Solve a linear system with an operator approximating the original one and based on a measurement of the error in the result. Jacobi, Gauss-Seidel, and Successive Over Relaxation (SOR) are three examples.

Non-Stationary Iterative Method:

$$x^{k+1} = \sigma_k x^k + d_k$$

These are methods where the data changes at each iteration. σ_k and d_k change for each iteration. d_k called search direction and σ_k is the step length. Krylov space method, Conjugate Gradient (CD), Generalized Minimal Residual (GM-RES), BICG, and BICGSTAB are some examples of this method.

b)

$$A = T + L + U$$

For each numerical method, we have two ways: Stable and Cheap.

As we recall, for Gauss-Seidel we keep a lower triangular matrix on the left hand side and move the upper triangular to the right hand side, so we could use forward substitution.

But, here we are dealing with a full lower triangular matrix which cost the order of n^2 . Solving a tridiagonal matrix is the order of n.

$$Ax = b$$

$$(T+L+U)x = b \Rightarrow Tx = b - (L+U)x$$

$$X^{k+1} = -T^{-1}(L+U)x^k + T^{-1}b$$

$$\Rightarrow p = -T^{-1}(L+U)x^k, c = T^{-1}b$$

 $\mathbf{c})$

Recall SOR method:

Given a square system of n linear equations with unknown x:

$$Ax = b$$

A can be decomposed into a diagonal component D, and strictly lower and upper triangular components L and U:

$$A = D + L + U$$

The system of linear equations may be rewritten as:

$$(D + \omega L)x = \omega b - (\omega U + (\omega - 1)D)x$$

$$x_i^{k+1} = (1 - \omega)x_i^k + \frac{\omega}{a_{ii}}(b_i - \sum_{j < i} a_{ij}x_j^{k+1} - \sum_{j > i} a_{ij}x_j^k), with : i = 1, 2, ..., n$$

For finding the similarity with SOR we have:

$$Tx^{k+1} = b - (L+U)x^k \Rightarrow x^{k+1} - x^k = T^{-1}[b - (L+U+T)x^k]$$

We also have from SOR: $x^{k+1} = x^k + \omega(x^{k+1} - x^k)$

$$x^{k+1} = x^k + \omega(T^{-1}[b - (L+U+T)x^k])$$

$$x^{k+1} = T^{-1}[(1-\omega)T - \omega(L+U)]x^k + \omega b T^{-1}$$

It is obvious with $\omega = 1$ we have the same formula like part b.

 \mathbf{d}

Pseudo Code:

input: Matrix A, Matrix b for RHS, initial guess x_0 We have to decompose matrix A to Matrices L, U, T

Set the counter = 1

While $Counter < P = Max_{iteration}$

Calculate Right-Hand-Side: $b - (L + U)x_0$

Call TDMA Solver to solve: Tx = RHS

Then Substitute x: $x_0 \leftarrow x$

Increase the counter: $counter \leftarrow counter + 1$

End While

We should also know the Thomas Algorithm to solve Tridiagonal Matrix T:

$$Tx = d \Rightarrow \begin{bmatrix} b_1 & c_1 & 0 & \dots & 0 \\ a_2 & b_2 & c_2 & \dots & 0 \\ 0 & a_3 & b_3 & \dots & 0 \\ 0 & \dots & \dots & c_{n-1} \\ 0 & 0 & \dots & a_n & b_n \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} d_1 \\ d_2 \\ \vdots \\ d_n \end{bmatrix}$$

TDMA Pseudo Code:

Forward Elimination Phase:

for
$$k=2$$
 until n do:

$$m = \frac{a_k}{b_k - 1}$$

$$b_k = b_k - mc_{k-1}$$

$$d_k = d_k - md_{k-1}$$

$$d_i = d_i = md_i$$

End for Loop(k)

Backward Substitution Phase:

for k = n-1 stepdown until 1 do:
$$x_k = \frac{d_n}{b_n}$$

$$x_k = \frac{d_k - c_k x_{k+1}}{b_k}$$
 End Loop(k)

For the Operation Count for each iteration we have:

- 1) LU: 2N multiplications, N additions
- 2) Forward: N multiplications, N additions
- 3) Backward: 2N multiplications, N additions

Total = 8N For P iteration we will have: Total = $8PN \sim O(N)$

e)

I have tested out three different examples, Symmetric Positive Definite, Diagonally dominant matrix and a general matrix.

Theorem:

Thomas Algorithm will work if:

$$|b_i| > |a_i| + |c_i|$$
 with $i = 2, 3, ..., N - 1$
 $|a_1| > |c_1| and |a_N| > |c_N|$