Colum Stores & BigTable

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Traditional Row Stores

- Review of the logical data model
 - A database has a collection of tables
 - Each table has a collection of columns, defined in the schema
 - Each tuple has a value or NULL for each column
- Review of the storage model: A Row Store
 - File: a table is implemented as a file
 - Pages: a file contains a chain of pages
 - Records: multiple records are packed in a page

Why row stores?

- Natural: storage model similar to the logical model
- Good performance for writes (inserts, updates, deletes)



Column Stores

- Change the storage model w.o. changing the logical model
 - Store a table in a column-based fashion
- Why column stores?
 - IO performance: save I/O if queries only access a few columns
 - Storage efficiency: compression is easier.
 - Flexible for schema expansion: adding a new column is easy
 - Great for sparse tables or denomalized schema
 - Easy replication of a column with diff. ordering properties
- Potential problems?
 - Expensive updates: suitable for read-most workloads
 - Potentially many (expensive) joins

Overall Architecture

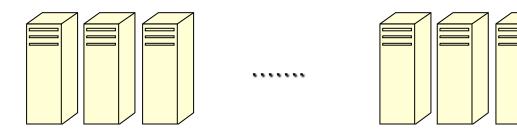
- Shared-nothing architecture of thousands of nodes!
 - A node is an off-the-shelf, commodity PC.

SQL (Hive) or Scripting (Yahoo! Pig Latin)

Map/Reduce Programming / Execution

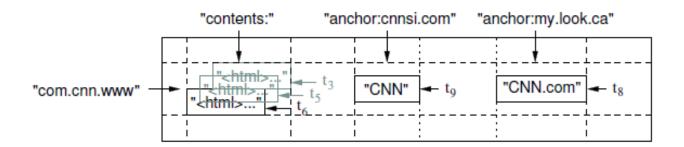
Key-Value Stores (Bigtable, HBase)

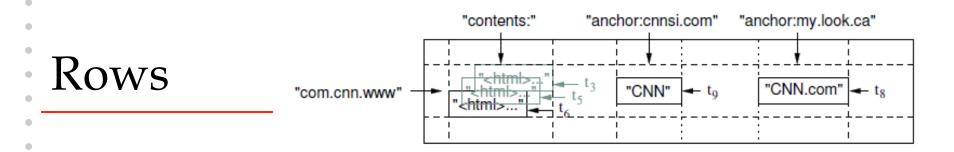
Distributed File System (GFS, HDFS)



Google Bigtable

- A data model (a schema).
- A sparse, distributed persistent multi-dimensional sorted map.
- Data is <u>partitioned</u> across the nodes seamlessly.
- The map is <u>indexed</u> by a row key, column key, and a timestamp.
 - Output value in the map is an un-interpreted array of bytes.
 - (row: byte[], column: byte[], time: int64) → byte[]

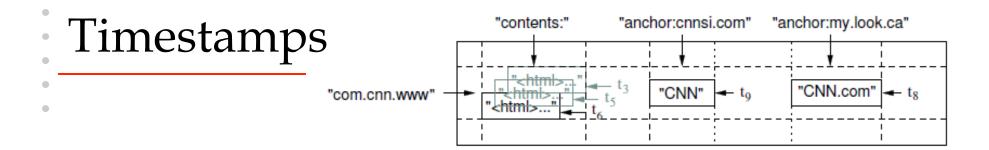




- A row key is an arbitrary string.
 - Typically 10-100 bytes in size, up to 64 KB.
 - Every read or write of data under a single row is atomic.
- Data is maintained in lexicographic order by row key.
- The row range for a table is dynamically partitioned.
 - Each partition (row range) is named a *tablet*.
 - Unit of distribution and load-balancing.
- Objective: make read operations single-sited!
 - E.g., In Webtable, pages in the same domain are grouped together by reversing the hostname components of the URLs: com.google.maps instead of maps.google.com.

Column Families "com.cnn.www" "com.cnn.www" "contents:" "anchor:cnnsi.com" "anchor:my.look.ca" "anchor:my.look.ca" "anchor:my.look.ca" "anchor:my.look.ca" "contents:" "anchor:cnnsi.com" "anchor:my.look.ca" "contents:" "anchor:cnnsi.com" "anchor:my.look.ca" "contents:" "anchor:my.look.ca" "contents:" "anchor:my.look.ca" "contents:" "anchor:my.look.ca" "contents:" "anchor:my.look.ca" "contents:" "anchor:my.look.ca"

- Column keys are grouped into sets called column families.
- A column family must be created before data can be stored in a column key.
- Hundreds of static column families.
- Syntax is family:key
 - e.g., Language:English, Language:German, etc.



- Store different versions of data in a cell in decreasing timestamp order.
 - New writes default to current time, but timestamps for writes can also be set explicitly by clients
- Lookup options:
 - "Return most recent K values"
 - "Return all values in timestamp range (or all values)"
- Column families can be marked w/ attributes:
 - "Only retain most recent K values in a cell"
 - "Keep values until they are older than K seconds"

Bigtable Usage

• Used in different applications supported by Google.

| Project | Table size | Compression | # Cells | # Column | # Locality | % in | Latency- |
|---------------------|------------|-------------|------------|----------|------------|--------|------------|
| name | (TB) | ratio | (billions) | Families | Groups | memory | sensitive? |
| Crawl | 800 | 11% | 1000 | 16 | 8 | 0% | No |
| Crawl | 50 | 33% | 200 | 2 | 2 | 0% | No |
| Google Analytics | 20 | 29% | 10 | 1 | 1 | 0% | Yes |
| Google Analytics | 200 | 14% | 80 | 1 | 1 | 0% | Yes |
| Google Base | 2 | 31% | 10 | 29 | 3 | 15% | Yes |
| Google Earth | 0.5 | 64% | 8 | 7 | 2 | 33% | Yes |
| Google Earth | 70 | _ | 9 | 8 | 3 | 0% | No |
| Orkut | 9 | _ | 0.9 | 8 | 5 | 1% | Yes |
| Personalized Search | 4 | 47% | 6 | 93 | 11 | 5% | Yes |

App 1: Google Analytics

- Enables webmasters to analyze traffic pattern at their web sites. Statistics such as:
 - number of unique visitors per day and the page views per URL per day,
 - percentage of users that made a purchase given that they earlier viewed a specific page.

• How?

- A small JavaScript program that the webmaster embeds in their web pages.
- Every time the page is visited, the program is executed.
- Program records the following information about each request:
 - User identifier
 - The page being fetched...

App 1: Google Analytics (Cont...)

- Two of the Bigtables
 - Raw click table (~ 200 TB)
 - A row for each end-user session.
 - Row name includes the website's name and the time at which the session was created.
 - Clustering of sessions that visit the same web site. And a sorted chronological order.
 - Compression factor of 6-7.
 - Summary table (~ 20 TB)
 - Stores predefined summaries for each web site.
 - Generated from the raw click table by periodically scheduled MapReduce jobs.
 - Each MapReduce job extracts recent session data from the raw click table.
 - Row name includes the website's name and the column family is the aggregate summaries.
 - Compression factor is 2-3.

App 2: Google Earth & Maps

- Functionality: Pan, view, and annotate satellite imagery at different resolution levels.
- One Bigtable stores raw imagery (~ 70 TB):
 - Row name is a geographic segment. Names are chosen to ensure adjacent geographic segments are clustered together.
 - Column family maintains sources of data for each segment.
- There are different sets of tables for serving client data, e.g., index table.

App 3: Personalized Search

- Records user queries and clicks across Google properties.
- Users browse their search histories and request for personalized search results based on their historical usage patterns.
- One Bigtable:
 - Row name is userid
 - A column family is reserved for each action type, e.g., web queries, clicks.
 - User profiles are generated using MapReduce.
 - These profiles personalize live search results.
 - Replicated geographically to reduce latency and increase availability.

Bigtable API

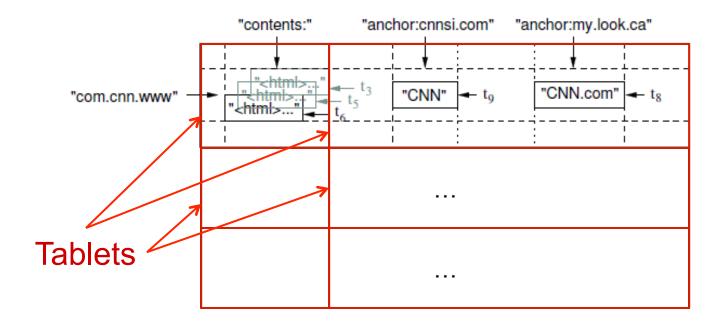
- Implements interfaces to
 - create and delete tables and column families,
 - modify cluster, table, and column family metadata such as access control rights,
 - Write or delete values in Bigtable,
 - Look up values from individual rows,
 - Iterate over a subset of the data in a table,
 - Atomic R-M-W sequences on data stored in a single row key (no support for Xacts across multiple rows).

Building Blocks

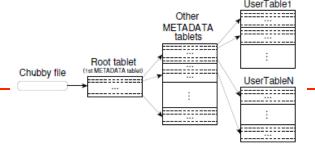
- Google File System:
 - for storing log and data files
 - high availability
- SSTable
 - a key/value database
- Chubby
 - name space

Tablets

- Large tables broken into tablets at row boundaries
 - Tablet holds contiguous range of rows
 - Clients can often choose row keys to achieve locality
 - Aim for ~100MB to 200MB of data per tablet



Locating Tablets (Ranges)



- A 3-level hierarchy:
- 1st Level: A file stored in chubby contains location of the root tablet, i.e., a directory of ranges (tablets) and associated meta-data.
 - The root tablet never splits.
- 2nd Level: Each meta-data tablet contains the location of a set of user tablets.
- 3rd Level: A set of SSTable identifiers for each tablet.
- Analysis:
 - Each meta-data row stores ~ 1KB of data,
 - With 128 MB tablets, the three level store addresses 2³⁴ tablets (2⁶¹ bytes in 128 MB tablets).
 - Approaches a Zetabyte (million Petabytes).

Write & Read Operations

- Write operation arrives at a tablet server:
 - Server ensures the client has sufficient privileges for the write operation (Chubby),
 - A log record is generated to the commit log file,
 - Once the write commits, its contents are inserted into the memtable.
- Read operation arrives at a tablet server:
 - Server ensures client has sufficient privileges for the read operation (Chubby),
 - Read is performed on a merged view of (a) the SSTables that constitute the tablet, and (b) the memtable.

Write Operations

- As writes execute, size of memtable increases.
- Once memtable reaches a threshold:
 - Memtable is frozen,
 - A new memtable is created,
 - Frozen metable is converted to an SSTable and written to GFS.
- This *minor compaction* minimizes memory usage of tablet server, and reduces recovery time in the presence of crashes (checkpoints).
- *Merging compaction* (in the background) reads a few SSTables and memtable to produce one SSTable. (Input SSTables and memtable are discareded.)
- *Major compaction* rewrites all SSTables into exactly one SSTable (containing no deletion entries).

Questions

