

Project SMU

Slimme Meter Uitlezer

Docenten:

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Wat gaan we doen?

Week 1: Kennismaken met de Slimme Meter en de SMU

Week 2: Het meten van spanning en stroom

Week 3: Digitale signalen en schakelaars

Week 4: Booleaanse algebra en het 7-segmentsdisplay

Week 5: De microcontroller programmeren

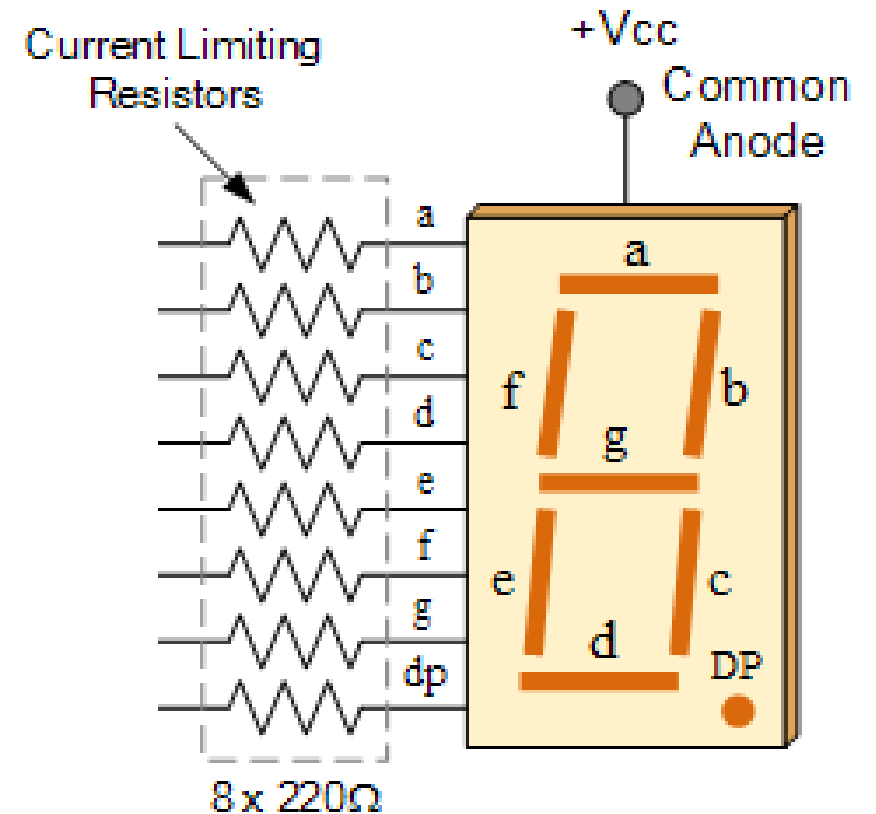
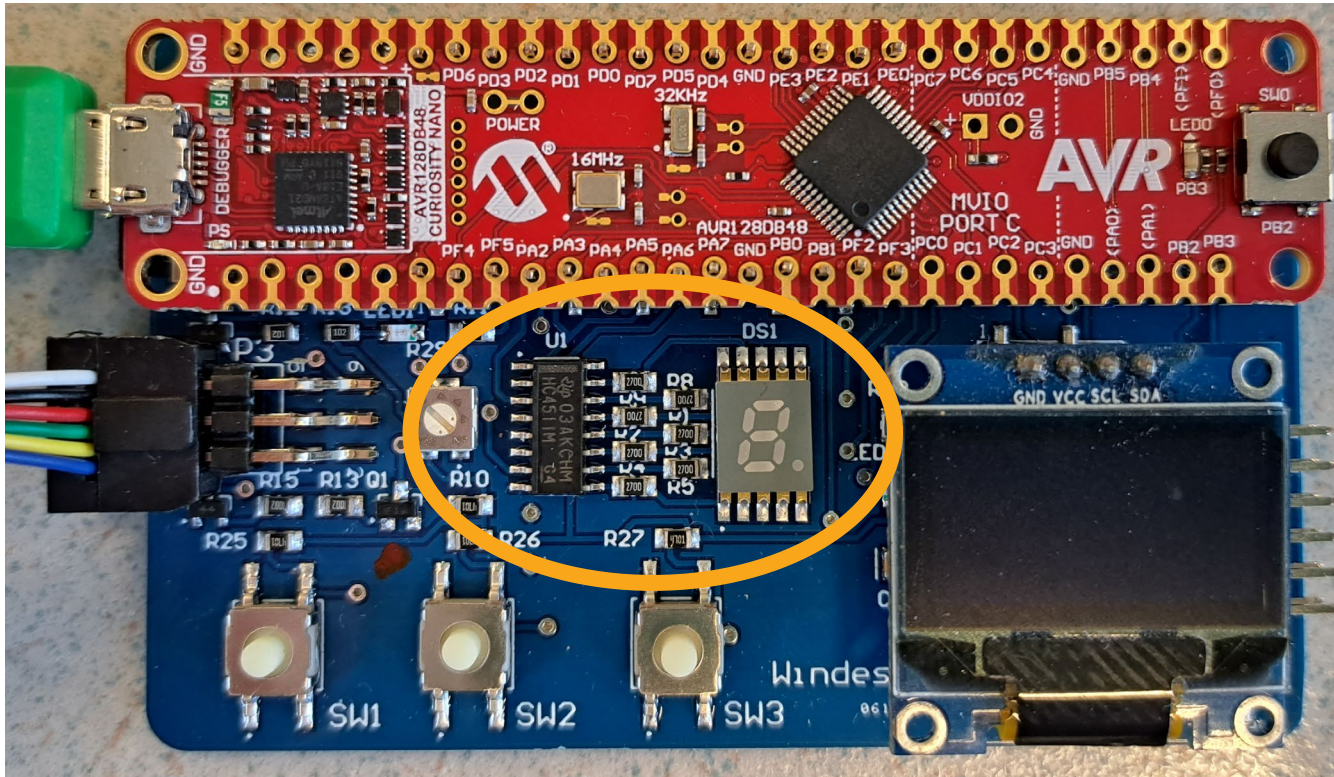
Week 6: Seriële communicatie, het OLED display en de SD-kaart

Week 7: Alles afmaken

Week 8: Toetsing



Het 7-segmentsdisplay + decoder



Om decoder aan te sturen: binaire getallen

| Binair getal | Macht | Decimaal getal |
|--------------|-------|----------------|
| 00001 | 2^0 | 1 |
| 00010 | 2^1 | 2 |
| 00100 | 2^2 | 4 |
| 01000 | 2^3 | 8 |
| 10000 | 2^4 | 16 |

- Elke één in een binair getal komt overeen met een macht van twee.
- Je begint aan de rechterkant met 2^0 en elke plek die je naar links schuift wordt de macht 1 hoger.

Bij een binair getal met meerdere enen tel je de bijbehorende machten bij elkaar op om de decimale waarde te vinden:

$$1011 = 2^3 + 2^1 + 2^0 = 8 + 2 + 1 = 11$$



Van decimaal naar binair

Wil je een decimaal getal omzetten naar een binair getal, dan kijk je welke machten van twee erin passen, te beginnen bij de grootste macht.

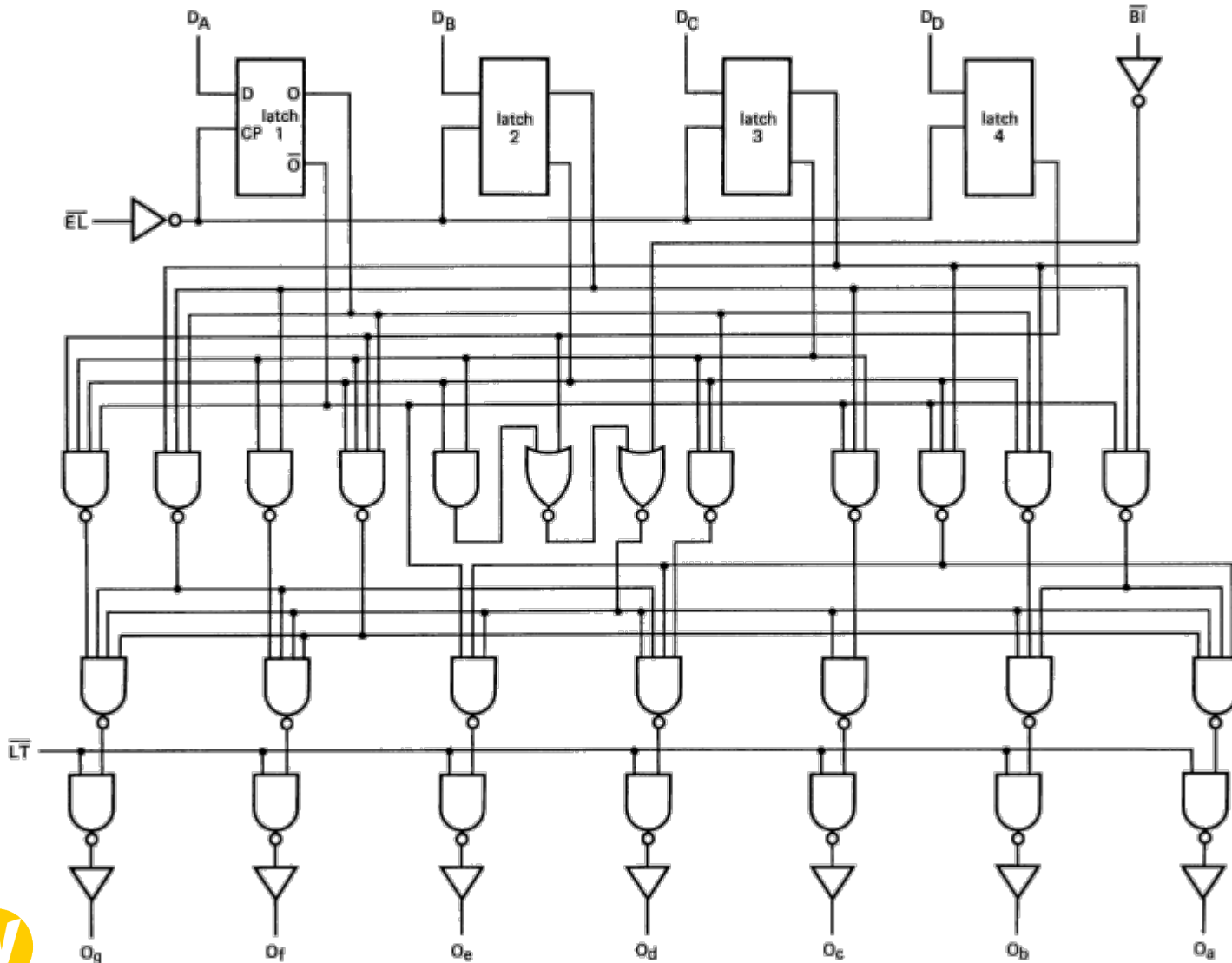
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| 01000 | 2^3 | 8 |
| 10000 | 2^4 | 16 |

Het decimale getal 23 wordt op deze manier:

- De grootste macht die in 23 past is $2^4 = 16$
- $23 - 16 = 7$, de grootste macht die hierin past is $2^2 = 4$
- $7 - 4 = 3$, hier past $2^1 = 2$ in
- $3 - 2 = 1$, hier past $2^0 = 1$ in.
- $2^4 + 2^2 + 2^1 + 2^0 = 10111$











Binnen in de decoder: combinatorische logica




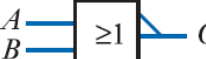

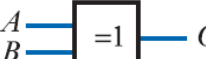

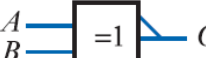


Een combinatie van de logische poorten die je vorige week geleerd hebt!



Logische poorten: notatie en functie

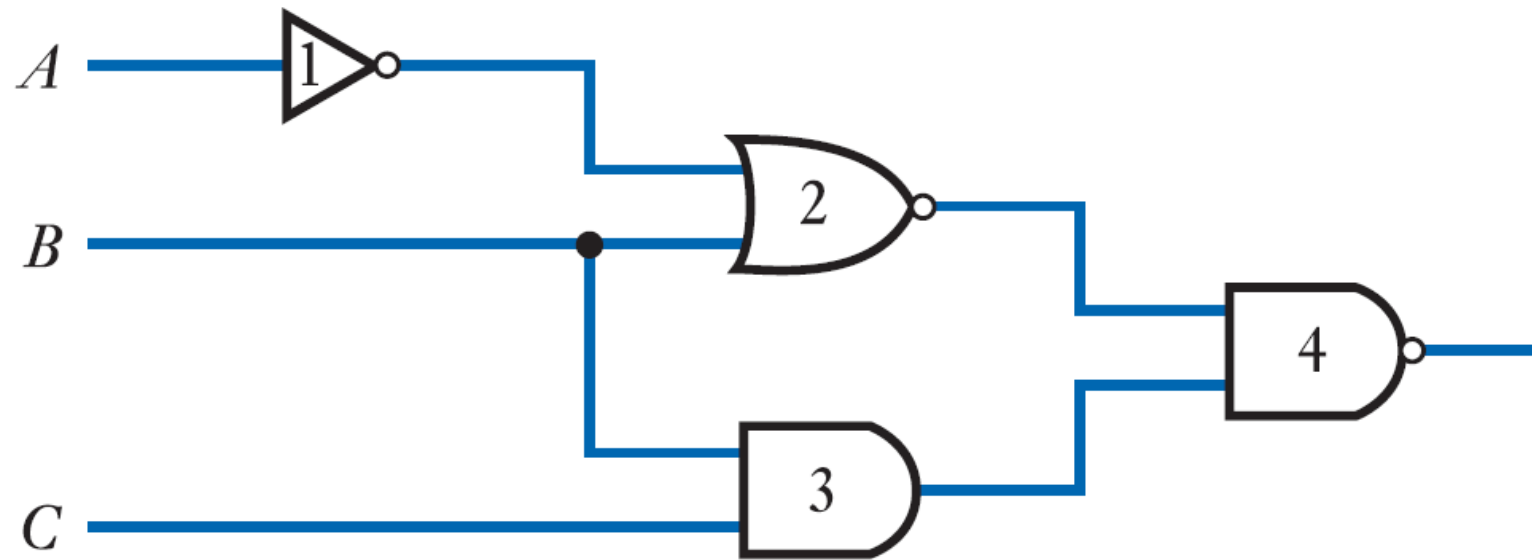
| Function | Symbol | Alternative symbol | Boolean expression | Truth table | | | | | | | | | | | | | | | |
|----------|---|---|--------------------|--|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| Buffer |  |  | $B = A$ | <table><tr><th>A</th><th>B</th></tr><tr><td>0</td><td>0</td></tr><tr><td>1</td><td>1</td></tr></table> | A | B | 0 | 0 | 1 | 1 | | | | | | | | | |
| A | B | | | | | | | | | | | | | | | | | | |
| 0 | 0 | | | | | | | | | | | | | | | | | | |
| 1 | 1 | | | | | | | | | | | | | | | | | | |
| NOT |  |  | $B = \bar{A}$ | <table><tr><th>A</th><th>B</th></tr><tr><td>0</td><td>1</td></tr><tr><td>1</td><td>0</td></tr></table> | A | B | 0 | 1 | 1 | 0 | | | | | | | | | |
| A | B | | | | | | | | | | | | | | | | | | |
| 0 | 1 | | | | | | | | | | | | | | | | | | |
| 1 | 0 | | | | | | | | | | | | | | | | | | |
| AND |  |  | $C = A \cdot B$ | <table><tr><th>A</th><th>B</th><th>C</th></tr><tr><td>0</td><td>0</td><td>0</td></tr><tr><td>0</td><td>1</td><td>0</td></tr><tr><td>1</td><td>0</td><td>0</td></tr><tr><td>1</td><td>1</td><td>1</td></tr></table> | A | B | C | 0 | 0 | 0 | 0 | 1 | 0 | 1 | 0 | 0 | 1 | 1 | 1 |
| A | B | C | | | | | | | | | | | | | | | | | |
| 0 | 0 | 0 | | | | | | | | | | | | | | | | | |
| 0 | 1 | 0 | | | | | | | | | | | | | | | | | |
| 1 | 0 | 0 | | | | | | | | | | | | | | | | | |
| 1 | 1 | 1 | | | | | | | | | | | | | | | | | |
| OR |  |  | $C = A + B$ | <table><tr><th>A</th><th>B</th><th>C</th></tr><tr><td>0</td><td>0</td><td>0</td></tr><tr><td>0</td><td>1</td><td>1</td></tr><tr><td>1</td><td>0</td><td>1</td></tr><tr><td>1</td><td>1</td><td>1</td></tr></table> | A | B | C | 0 | 0 | 0 | 0 | 1 | 1 | 1 | 0 | 1 | 1 | 1 | 1 |
| A | B | C | | | | | | | | | | | | | | | | | |
| 0 | 0 | 0 | | | | | | | | | | | | | | | | | |
| 0 | 1 | 1 | | | | | | | | | | | | | | | | | |
| 1 | 0 | 1 | | | | | | | | | | | | | | | | | |
| 1 | 1 | 1 | | | | | | | | | | | | | | | | | |

| Function | Symbol | Alternative symbol | Boolean expression | Truth table | | | | | | | | | | | | |
|---------------|---|---|-----------------------------|---|---|---|---|---|---|---|---|---|---|---|---|---|
| NAND |  |  | $C = \overline{A \cdot B}$ | <table><tr><td>0</td><td>0</td><td>1</td></tr><tr><td>0</td><td>1</td><td>1</td></tr><tr><td>1</td><td>0</td><td>1</td></tr><tr><td>1</td><td>1</td><td>0</td></tr></table> | 0 | 0 | 1 | 0 | 1 | 1 | 1 | 0 | 1 | 1 | 1 | 0 |
| 0 | 0 | 1 | | | | | | | | | | | | | | |
| 0 | 1 | 1 | | | | | | | | | | | | | | |
| 1 | 0 | 1 | | | | | | | | | | | | | | |
| 1 | 1 | 0 | | | | | | | | | | | | | | |
| NOR |  |  | $C = \overline{A + B}$ | <table><tr><td>0</td><td>0</td><td>1</td></tr><tr><td>0</td><td>1</td><td>0</td></tr><tr><td>1</td><td>0</td><td>0</td></tr><tr><td>1</td><td>1</td><td>0</td></tr></table> | 0 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |
| 0 | 0 | 1 | | | | | | | | | | | | | | |
| 0 | 1 | 0 | | | | | | | | | | | | | | |
| 1 | 0 | 0 | | | | | | | | | | | | | | |
| 1 | 1 | 0 | | | | | | | | | | | | | | |
| Exclusive OR |  |  | $C = A \oplus B$ | <table><tr><td>0</td><td>0</td><td>0</td></tr><tr><td>0</td><td>1</td><td>1</td></tr><tr><td>1</td><td>0</td><td>1</td></tr><tr><td>1</td><td>1</td><td>0</td></tr></table> | 0 | 0 | 0 | 0 | 1 | 1 | 1 | 0 | 1 | 1 | 1 | 0 |
| 0 | 0 | 0 | | | | | | | | | | | | | | |
| 0 | 1 | 1 | | | | | | | | | | | | | | |
| 1 | 0 | 1 | | | | | | | | | | | | | | |
| 1 | 1 | 0 | | | | | | | | | | | | | | |
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| 0 | 0 | 1 | | | | | | | | | | | | | | |
| 0 | 1 | 0 | | | | | | | | | | | | | | |
| 1 | 0 | 0 | | | | | | | | | | | | | | |
| 1 | 1 | 1 | | | | | | | | | | | | | | |



Logisch schema en Booleaanse expressies

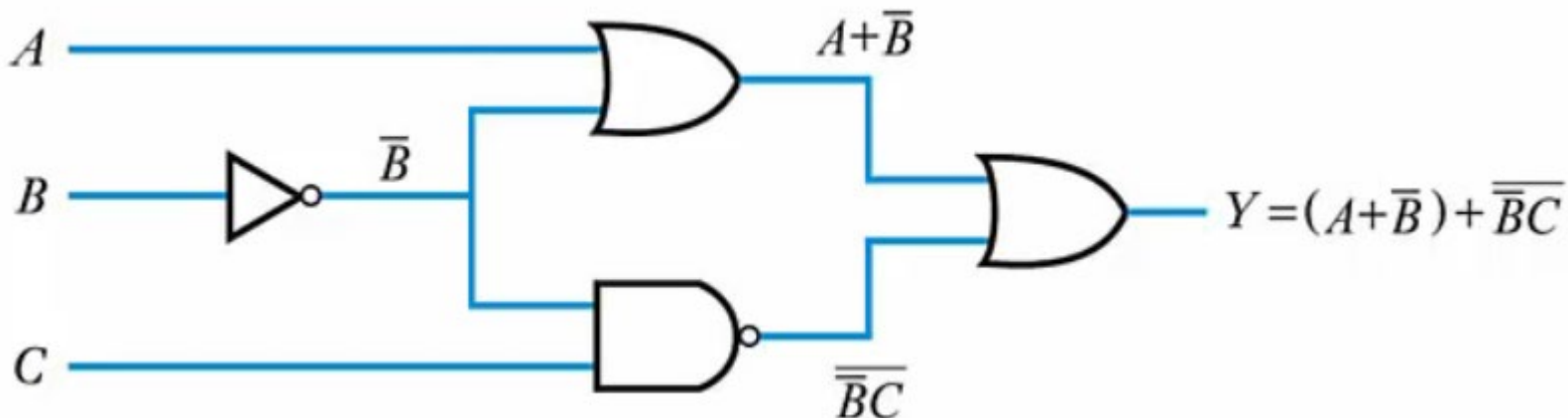
Van schema naar expressie



Logisch schema en Booleaanse expressies

Van expressie naar schema:

- Let op de volgorde van de Booleaanse algebra!
(haakjes \rightarrow AND-poorten \rightarrow OR-poorten)
- Let op de lengte van de streepjes!



Logisch schema vanuit een waarheidstabel

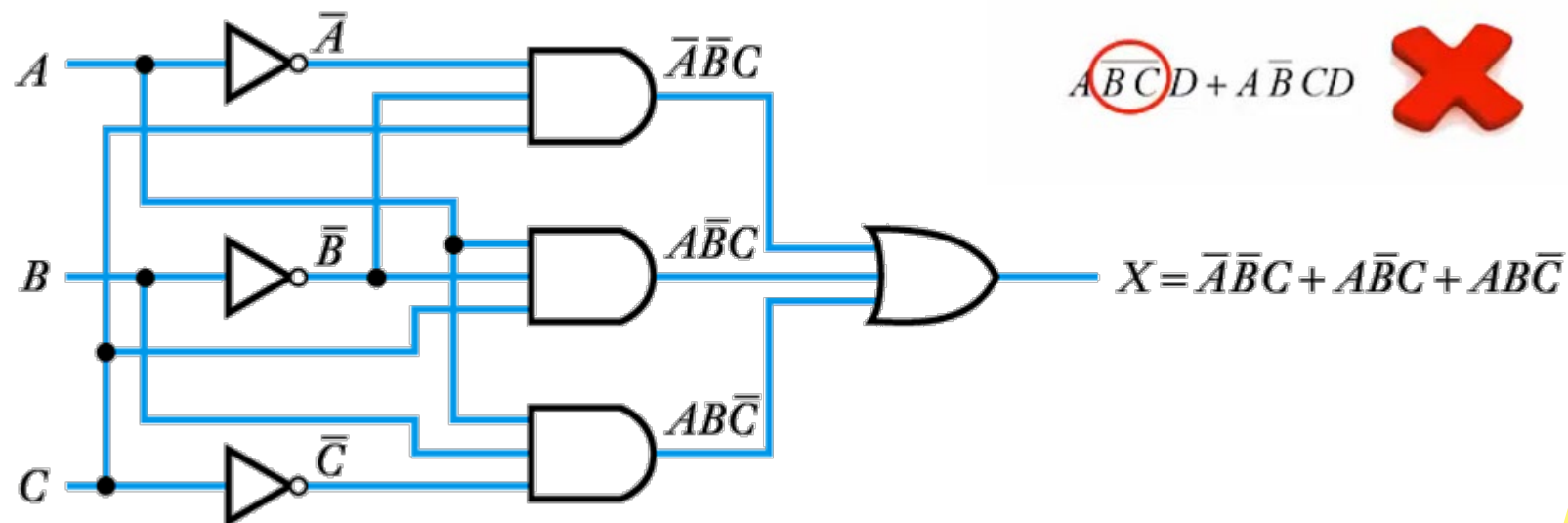
| A | B | C | X |
|---|---|---|---|
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 1 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 0 |

$$\bar{A}\bar{B}C$$

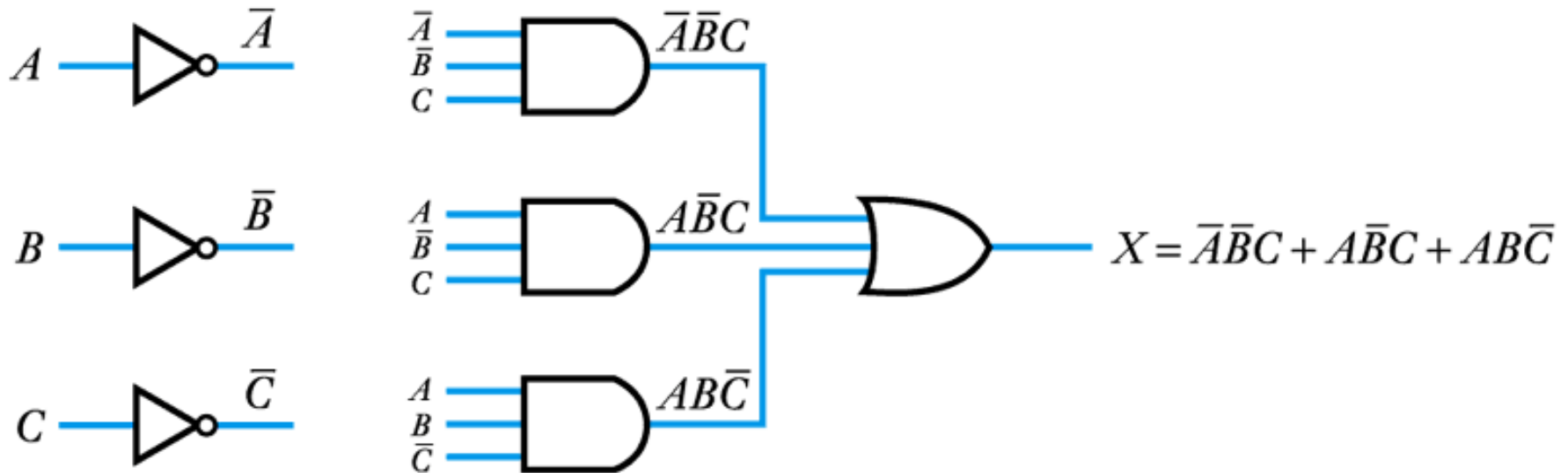
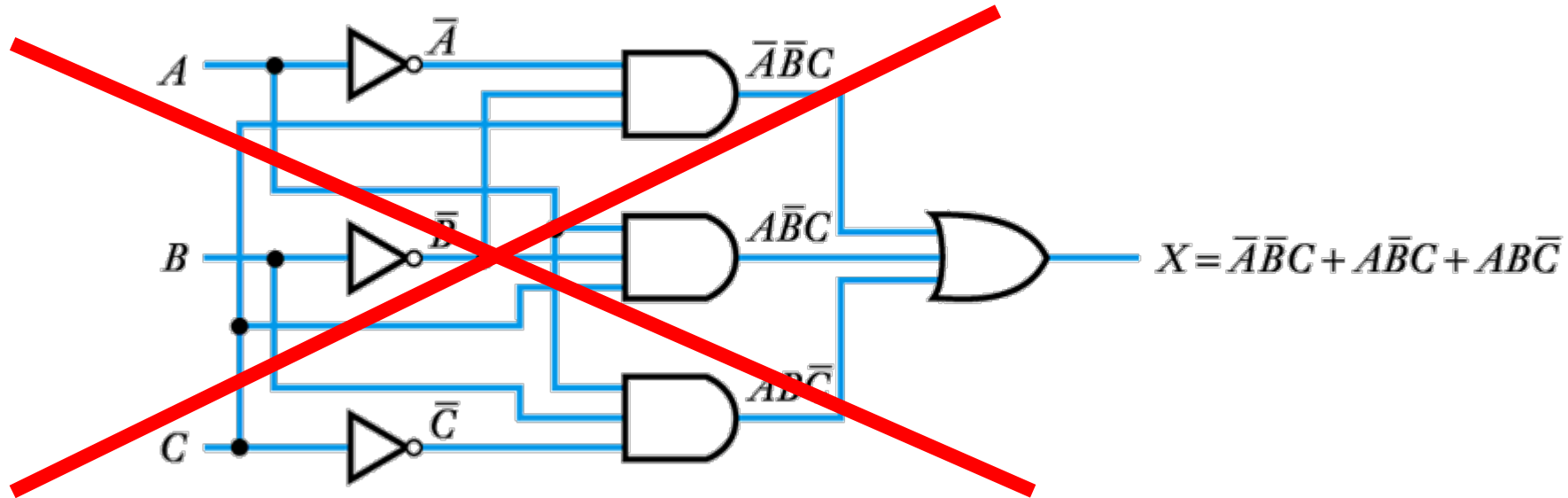
$$A\bar{B}C$$

$$AB\bar{C}$$

1. Schrijf de mintermen op bij de regels waarvan de output gelijk is aan '1'.
2. De bijbehorende Booleaanse expressie is de som van deze mintermen: $X = \bar{A}\bar{B}C + A\bar{B}C + AB\bar{C}$



Notatie zonder spaghetti:



Schakelalgebra

Wanneer we vanuit een waarheidstabel een Booleaanse vergelijking opstellen, wordt dit altijd een lange som van mintermen:

$$E = \overline{B}\overline{C}\overline{D} + \overline{A}BD + ABD + BC\overline{D} + \overline{B}CD + \overline{A}\overline{B}\overline{C}D + A\overline{B}\overline{C}D$$

Het is mogelijk deze vergelijking korter en overzichtelijker te maken met behulp van Booleaanse wetten en regels.

Bijkomend voordeel: hoe korter de vergelijking, hoe minder componenten er nodig zijn om de schakeling te bouwen!



Boolean identities

AND function

$$0 \cdot 0 = 0$$

$$0 \cdot 1 = 0$$

$$1 \cdot 0 = 0$$

$$1 \cdot 1 = 1$$

$$A \cdot 0 = 0$$

$$0 \cdot A = 0$$

$$A \cdot 1 = A$$

$$1 \cdot A = A$$

$$A \cdot A = A$$

$$A \cdot \bar{A} = 0$$

OR function

$$0 + 0 = 0$$

$$0 + 1 = 1$$

$$1 + 0 = 1$$

$$1 + 1 = 1$$

$$A + 0 = A$$

$$0 + A = A$$

$$A + 1 = 1$$

$$1 + A = 1$$

$$A + A = A$$

$$A + \bar{A} = 1$$

NOT function

$$\bar{0} = 1$$

$$\bar{1} = 0$$

$$\bar{\bar{A}} = A$$

Boolean laws

Commutative law

$$AB = BA$$

$$A + B = B + A$$

Absorption law

$$A + AB = A$$

$$A(A + B) = A$$

Distributive law

$$A(B + C) = AB + AC$$

$$A + BC = (A + B)(A + C)$$

De Morgan's law

$$\overline{A + B} = \bar{A} \cdot \bar{B}$$

$$\overline{A \cdot B} = \bar{A} + \bar{B}$$

Associative law

$$A(BC) = (AB)C$$

$$A + (B + C) = (A + B) + C$$

Note also

$$A + \bar{A}B = A + B$$

$$A(\bar{A} + B) = AB$$



Schakelalgebra

We gaan de volgende vergelijking vereenvoudigen:

$$E = B\bar{C}\bar{D} + \bar{A}BD + ABD + B\bar{C}\bar{D} + \bar{B}CD + \bar{A}\bar{B}\bar{C}D + A\bar{B}\bar{C}D$$

1. Neem termen die maar één variabele verschillen samen:

$$E = B\bar{D}(\bar{C} + C) + BD(\bar{A} + A) + \bar{B}\bar{C}D(\bar{A} + A) + \bar{B}CD$$

2. Gebruik de wet: $\bar{A} + A = 1$

$$E = B\bar{D} + BD + \bar{B}\bar{C}D + \bar{B}CD$$

3. Herhaling van stappen:

$$E = B(\bar{D} + D) + \bar{B}D(\bar{C} + C) = B + \bar{B}D$$

4. Gebruik de wet: $A + \bar{A}B = A + B$

$$E = B + D$$



Aan de slag!

- Ga naar leren.windesheim.nl
(zoek de cursus EDPD.22, project SMU)
- Voer de opdrachten van week 4 uit.
- Ben je klaar? Ga vast verder met de voorbereiding van week 5!

