

Chapter 2

Here we learned about norms on spaces of matrices. Two cases are here important:

i) Induced matrix norms

Definition 0.0.1. Let $\mathbf{A} \in \mathbb{K}^{m \times n}$ and we equip \mathbb{K}^n with $\|\cdot\|_{(n)}$ and \mathbb{K}^m with $\|\cdot\|_{(m)}$. The induced matrix norm is then

$$\|\mathbf{A}\|_{(m,n)} = \sup_{\substack{\mathbf{x} \in \mathbb{K}^n \\ \mathbf{x} \neq 0}} \frac{\|\mathbf{A}\mathbf{x}\|_{(m)}}{\|\mathbf{x}\|_{(n)}} = \sup_{\substack{\mathbf{x} \in \mathbb{K}^n \\ \|\mathbf{x}\|_{(n)}=1}} \|\mathbf{A}\mathbf{x}\|_{(m)} \quad (1)$$

ii) Matrix norm on the vector space of matrices

Definition 0.0.2. A function $\|\cdot\| : \mathbb{K}^m \rightarrow \mathbb{R}$ is called a norm if

1. $\|\mathbf{x}\| \geq 0$ for all $\mathbf{x} \in \mathbb{K}^m$ and $\|\mathbf{x}\| = 0 \Leftrightarrow \mathbf{x} = \mathbf{0}$
2. $\|\mathbf{x} + \mathbf{y}\| \leq \|\mathbf{x}\| + \|\mathbf{y}\|$ for all $\mathbf{x}, \mathbf{y} \in \mathbb{K}^m$
3. $\|\alpha\mathbf{x}\| = |\alpha|\|\mathbf{x}\|$ for all $\mathbf{x} \in \mathbb{K}^m$ and for all $\alpha \in \mathbb{K}$.

We have also seen central statements like

- Young's product inequality:

Lemma 0.0.1 (Young's product inequality). Let $a, b \in \mathbb{R}_{\geq 0}$. Then

$$ab \leq \frac{1}{p}a^p + \frac{1}{q}b^q \quad (2)$$

for $1 \leq p, q \leq \infty$ and $\frac{1}{p} + \frac{1}{q} = 1$.

Proof. Let $a, b \in \mathbb{R}_{\geq 0}$, and $t = \frac{1}{p}$ and $1 - t = \frac{1}{q}$. Then

$$\ln(ta^p + (1-t)b^q) \underset{(*)}{\geq} t \ln(a^p) + (1-t) \ln(b^q) = \ln(a) + \ln(b) = \ln(ab) \quad (3)$$

where we used that \ln is concave in $(*)$. □

- Hölder inequality:

Theorem 0.1 (Hölder inequality). Let $\mathbf{x}, \mathbf{y} \in \mathbb{K}^m$. Then

$$|\langle \mathbf{x}, \mathbf{y} \rangle| \leq \|\mathbf{x}\|_p \|\mathbf{y}\|_q, \quad (4)$$

where $1 \leq p, q \leq \infty$ and $\frac{1}{p} + \frac{1}{q} = 1$.

- Cauchy-Schwarz inequality:

Corollary 0.1.1 (Cauchy-Schwarz inequality). For $p, q = 2$ the Hölder inequality yields

$$|\langle \mathbf{x}, \mathbf{y} \rangle| \leq \|\mathbf{x}\|_2 \|\mathbf{y}\|_2. \quad (5)$$

- Minkowski inequality

Theorem 0.2 (Minkowski inequality). *Let $\mathbf{x}, \mathbf{y} \in \mathbb{K}^m$ and $p \geq 1$. Then*

$$\|\mathbf{x} + \mathbf{y}\|_p \leq \|\mathbf{x}\|_p + \|\mathbf{y}\|_p \quad (6)$$

- The standard inner product on matrix spaces:

Definition 0.2.1. *Let $\mathbf{A}, \mathbf{B} \in \mathbb{K}^{m \times n}$. The standard inner product of matrices is defined as*

$$\langle \mathbf{A}, \mathbf{B} \rangle = \text{Tr}(\mathbf{A}^* \mathbf{B}), \quad (7)$$

- The Frobenius norm is induced by the standard inner product:

Proposition 0.2.1. *The Frobenius norm is induced by the standard inner product of matrices, i.e., for $\mathbf{A} \in \mathbb{K}^{m \times n}$*

$$\|\mathbf{A}\|_F = \sqrt{\langle \mathbf{A}, \mathbf{A} \rangle}. \quad (8)$$

- Unitarily invariant norms:

Theorem 0.3. *For any $\mathbf{A} \in \mathbb{K}^{m \times n}$ and unitary $\mathbf{U} \in \mathbb{K}^{m \times m}$, we have*

$$\|\mathbf{U}\mathbf{A}\|_2 = \|\mathbf{A}\|_2 \quad \text{and} \quad \|\mathbf{U}\mathbf{A}\|_F = \|\mathbf{A}\|_F \quad (9)$$

Proof. Let $\mathbf{A} \in \mathbb{K}^{m \times n}$ and $\mathbf{U} \in \mathbb{K}^{m \times m}$ be unitary. Then

$$\|\mathbf{U}\mathbf{A}\mathbf{x}\|_2 = \sqrt{\mathbf{x}^* \mathbf{A}^* \mathbf{U}^* \mathbf{U} \mathbf{A} \mathbf{x}} = \|\mathbf{A}\mathbf{x}\|_2 \Rightarrow \|\mathbf{U}\mathbf{A}\|_2 = \|\mathbf{A}\|_2 \quad (10)$$

and

$$\|\mathbf{U}\mathbf{A}\|_F = \sqrt{\text{Tr}((\mathbf{U}\mathbf{A})^* \mathbf{U}\mathbf{A})} = \sqrt{\text{Tr}(\mathbf{A}^* \mathbf{A})} = \|\mathbf{A}\|_F \quad (11)$$

□

In the homework assignments you have seen central statements like:

- Hermitian matrices have real-valued eigenvalues.
- Skew hermitian matrices have purely imaginary eigenvalues
- And matrix inequalities: $\|x\|_2 \leq \sqrt{m} \|x\|_\infty$

Proof. Starting from the definition of the 2-norm, we find

$$\|x\|_2 = \left(\sum_{i=1}^m |x_i|^2 \right)^{1/2} \leq \left(\sum_{i=1}^m \max_{i \in [m]} |x_i|^2 \right)^{1/2} = \sqrt{m} \left(\max_{i \in [m]} |x_i|^2 \right)^{1/2} = \sqrt{m} \left(\max_{i \in [m]} |x_i| \right),$$

hence $\sqrt{m} \|x\|_\infty$. The inequality is sharp for $x = (1, 1, \dots, 1)^\top$, i.e., the vector with all entries equal to one, since $\|x\|_2 = \sqrt{m}$ and $\|x\|_\infty = 1$. □

Chapter 3

The most central object of this course and large parts of numerical linear algebra; the singular value decomposition:

Definition 0.3.1. Let $\mathbf{A} \in \mathbb{K}^{m \times n}$. We call the factorization

$$\mathbf{A} = \mathbf{U}\mathbf{\Sigma}\mathbf{V}^*, \quad (12)$$

where $\mathbf{U} \in \mathbb{K}^{m \times m}$ and $\mathbf{V} \in \mathbb{K}^{n \times n}$ are unitary, and $\mathbf{\Sigma} \in \mathbb{K}^{m \times n}$ is diagonal, singular value decomposition of \mathbf{A} .

Theorem 0.4. Every matrix $\mathbf{A} \in \mathbb{K}^{m \times n}$ has a singular value decomposition and the singular values $\{\sigma_i\}$ are uniquely determined. Moreover, if \mathbf{A} is square and σ_i distinct, the left and right singular vectors $\{\mathbf{u}_j\}$ and $\{\mathbf{v}_j\}$ are uniquely determined up to complex signs, i.e., complex scaling factors of length one.

The proof is long but parts can be asked:

- Then $\mathbf{A}^* \mathbf{A}$ is positive semi-definite, indeed,

$$\mathbf{x}^* \mathbf{A}^* \mathbf{A} \mathbf{x} = (\mathbf{A} \mathbf{x})^* (\mathbf{A} \mathbf{x}) = \|\mathbf{A} \mathbf{x}\|_2^2 \geq 0. \quad (13)$$

Proposition 0.4.1. Let $\mathbf{A} \in \mathbb{K}^{m \times n}$. Then $\|\mathbf{A}\|_2 = \sigma_{\max}(\mathbf{A})$, i.e., the largest singular value.

Proof. Let $\mathbf{A} \in \mathbb{K}^{m \times n}$, with singular value decomposition $\mathbf{A} = \mathbf{U}\mathbf{\Sigma}\mathbf{V}^*$ and σ_1 being the largest singular value. Then for $\|\mathbf{x}\|_2 = 1$

$$\|\mathbf{A} \mathbf{x}\|_2^2 = \langle \mathbf{x}, \mathbf{A}^* \mathbf{A} \mathbf{x} \rangle = \sum_{i=1}^n \sigma_i^2 \langle \mathbf{x}, \mathbf{v}_i \mathbf{v}_i^* \mathbf{x} \rangle \leq \sigma_1^2 \sum_{i=1}^n |\mathbf{v}_i^* \mathbf{x}|^2 = \sigma_1^2 \|\mathbf{V}^* \mathbf{x}\|^2 \leq \sigma_1^2 \|\mathbf{V}^*\|^2 = \sigma_1^2 \quad (14)$$

which is tight for $\mathbf{x} = \mathbf{v}_1$. □

We learned two key applications:

- Low-rank approximation
- Moore-Penrose inverse:

Definition 0.4.1. Let $\mathbf{A} \in \mathbb{K}^{m \times n}$. The matrix $\mathbf{A}^+ \in \mathbb{K}^{n \times m}$ is called the pseudo inverse (Moore-Penrose) inverse of \mathbf{A} if

$$\begin{array}{ll} \text{i)} & \mathbf{A} \mathbf{A}^+ \mathbf{A} = \mathbf{A} \\ \text{ii)} & \mathbf{A}^+ \mathbf{A} \mathbf{A}^+ = \mathbf{A}^+ \end{array} \quad \begin{array}{ll} \text{iii)} & (\mathbf{A} \mathbf{A}^+)^* = \mathbf{A} \mathbf{A}^+ \\ \text{iv)} & (\mathbf{A}^+ \mathbf{A})^* = \mathbf{A}^+ \mathbf{A} \end{array}$$

That have different properties:

Low-rank

Theorem 0.5 (Eckart-Young-Mirsky – spectral norm). *Let $\mathbf{A} \in \mathbb{K}^{m \times m}$ with $\text{rank}(\mathbf{A}) = r$. For any k with $1 \leq k < r$, define*

$$\mathbf{A}_k = \sum_{j=1}^k \sigma_j \mathbf{u}_j \mathbf{v}_j^*. \quad (15)$$

Then

$$\|\mathbf{A} - \mathbf{A}_k\|_2 = \inf_{\substack{\mathbf{B} \in \mathbb{K}^{m \times m} \\ \text{rank}(\mathbf{B}) \leq k}} \|\mathbf{A} - \mathbf{B}\|_2 = \sigma_{k+1} \quad (16)$$

Proof. First note that

$$\|\mathbf{A} - \mathbf{A}_k\|_2 = \sigma_{k+1}. \quad (17)$$

It remains to show that \mathbf{A}_k is the infimum. To that end, assume the exist $\mathbf{B}_k = \mathbf{X}\mathbf{Y}^*$ where \mathbf{X}, \mathbf{Y} have k -columns and that

$$\|\mathbf{A} - \mathbf{B}_k\|_2 < \|\mathbf{A} - \mathbf{A}_k\|_2 = \sigma_{k+1}. \quad (18)$$

However, since

$$\text{rank}(\mathbf{Y}) = k < k + 1 = \text{rank}([\mathbf{v}_1 | \dots | \mathbf{v}_{k+1}]) \quad (19)$$

there exists a linear combination of right singular vectors of

$$\mathbf{w} = c_1 \mathbf{v}_1 + \dots + c_{k+1} \mathbf{v}_{k+1} \quad (20)$$

with

$$\mathbf{Y}^* \mathbf{w} = \mathbf{0}. \quad (21)$$

W.l.o.g. we assume \mathbf{w} is normalized, otherwise we normalize \mathbf{w} . Then,

$$\|\mathbf{A} - \mathbf{B}_k\|_2^2 \geq \|(\mathbf{A} - \mathbf{B}_k) \mathbf{w}\|_2^2 = \|\mathbf{A} \mathbf{w}\|_2^2 = c_1^2 \sigma_1^2 + \dots + c_{k+1}^2 \sigma_{k+1}^2 \geq \sigma_{k+1}^2 \quad (22)$$

□

Theorem 0.6 (Courant-Fisher min-max – singular values). *For $\mathbf{A} \in \mathbb{K}^{m \times n}$, we have*

$$\sigma_k = \max_{\substack{V \subset \mathbb{K}^n \\ \dim(V)=k}} \min_{\substack{\|\mathbf{v}\|=1 \\ \mathbf{v} \in V}} \|\mathbf{A} \mathbf{v}\|_2 \quad (23)$$

and

$$\sigma_{k+1} = \min_{\substack{V \subset \mathbb{K}^n \\ \dim(V)=n-k}} \max_{\substack{\|\mathbf{v}\|=1 \\ \mathbf{v} \in V}} \|\mathbf{A} \mathbf{v}\|_2. \quad (24)$$

Theorem 0.7 (Weyl's inequality). *Let $\mathbf{A}, \mathbf{B} \in \mathbb{K}^{m \times n}$ and denote its singular values by $\sigma_i(\mathbf{A})$ and $\sigma_i(\mathbf{B})$, respectively. We then have*

$$\sigma_{i+j-1}(\mathbf{A} + \mathbf{B}) \leq \sigma_i(\mathbf{A}) + \sigma_j(\mathbf{B}). \quad (25)$$

Theorem 0.8 (Eckart-Young-Mirsky for Frobenius norm). *Let $\mathbf{A} \in \mathbb{K}^{m \times m}$ with $\text{rank}(\mathbf{A}) = r$. For any k with $1 \leq k < r$, define*

$$\mathbf{A}_k = \sum_{j=1}^k \sigma_j \mathbf{u}_j \mathbf{v}_j^*. \quad (26)$$

Then

$$\|\mathbf{A} - \mathbf{A}_k\|_F = \inf_{\substack{\mathbf{B} \in \mathbb{K}^{m \times m} \\ \text{rank}(\mathbf{B}) \leq k}} \|\mathbf{A} - \mathbf{B}\|_F = \sqrt{\sigma_{k+1}^2 + \dots + \sigma_r^2}. \quad (27)$$

Moore Penrose inverse

Proposition 0.8.1. *Let $\mathbf{A} \in \mathbb{K}^{m \times n}$ with $m \leq n$ and \mathbf{A}^+ its Moore-Penrose inverse. Then*

$$\text{range}(\mathbf{A}^+) \perp \ker(\mathbf{A}). \quad (28)$$

Proof. Let $\mathbf{A} \in \mathbb{K}^{m \times n}$ with \mathbf{A}^+ its Moore-Penrose inverse. Recall that

$$\mathbf{A}\mathbf{A}^+\mathbf{A} = \mathbf{A} \quad \text{and} \quad (\mathbf{A}^+\mathbf{A})^* = \mathbf{A}^+\mathbf{A}. \quad (29)$$

Moreover let $\mathbf{y} \in \text{range}(\mathbf{A}^+)$, i.e., $\mathbf{y} = \mathbf{A}^+\mathbf{b}$ for some $\mathbf{b} \in \mathbb{K}^m$, and $\mathbf{x} \in \ker(\mathbf{A})$. Then

$$\langle \mathbf{y}, \mathbf{x} \rangle = \langle \mathbf{A}^+\mathbf{b}, \mathbf{x} \rangle = \langle \mathbf{A}^+\mathbf{A}\mathbf{A}^+\mathbf{b}, \mathbf{x} \rangle = \langle \mathbf{A}^+\mathbf{b}, (\mathbf{A}^+\mathbf{A})^*\mathbf{x} \rangle = \langle \mathbf{A}^+\mathbf{b}, \mathbf{A}^+\mathbf{A}\mathbf{x} \rangle = 0 \quad (30)$$

□

Theorem 0.9. *Let $\mathbf{A} \in \mathbb{K}^{m \times n}$, the Moore-penrose inverse \mathbf{A}^+ is unique.*

Proposition 0.9.1. *Let $\mathbf{A} \in \mathbb{K}^{m \times n}$ with $m > n$ and $\text{rank}(\mathbf{A}) = n$. Then*

$$\mathbf{A}^+ = (\mathbf{A}^*\mathbf{A})^{-1}\mathbf{A}^*. \quad (31)$$

Theorem 0.10. *If $\mathbf{A} \in \mathbb{K}^{m \times m}$ attains an inverse, then $\mathbf{A}^{-1} = \mathbf{A}^+$.*

Proof. Note that

$$\mathbf{I} = \mathbf{A}\mathbf{A}^{-1} = \mathbf{A}\mathbf{A}^+\mathbf{A}\mathbf{A}^{-1} = \mathbf{A}\mathbf{A}^+ \quad (32)$$

hence

$$\mathbf{A}^{-1} = \mathbf{A}^{-1}\mathbf{I} = \mathbf{A}^{-1}\mathbf{A}\mathbf{A}^+ = \mathbf{A}^+ \quad (33)$$

□

Theorem 0.11. *Let $\mathbf{A} \in \mathbb{K}^{m \times n}$ with $\mathbf{A}^+ \in \mathbb{K}^{n \times m}$ its pseudo inverse, then*

$$(\mathbf{A}^+)^+ = \mathbf{A}. \quad (34)$$

Application of MP inverse:

The MP inverse solves the over-determined least squares problem, i.e., minimize

$$\|\mathbf{A}\mathbf{x} - \mathbf{b}\|_2. \quad (35)$$

where $\mathbf{A} \in \mathbb{K}^{m \times n}$ and $m \geq n$ – we say “ \mathbf{A} is tall and skinny”. We have more equations than variables and consequently zero solutions to the system. We therefore seek $\mathbf{x} \in \mathbb{K}^n$ that minimizes the above residual, i.e.,

$$\min_{\mathbf{x} \in \mathbb{K}^n} \|\mathbf{A}\mathbf{x} - \mathbf{b}\|_2^2. \quad (36)$$

To that end, we compute the gradient of with respect to \mathbf{x} :

$$\frac{\partial}{\partial \mathbf{x}} \|\mathbf{A}\mathbf{x} - \mathbf{b}\|_2^2 = 2\mathbf{A}^*(\mathbf{A}\mathbf{x} - \mathbf{b}). \quad (37)$$

Enforcing first-order optimality yields the normal equation

$$\mathbf{A}^*\mathbf{A}\mathbf{x} = \mathbf{A}^*\mathbf{b}. \quad (38)$$

Assuming $\mathbf{A}^*\mathbf{A}$ is invertible, which holds if \mathbf{A} has full rank, we can solve the normal equation, i.e.,

$$\mathbf{x} = (\mathbf{A}^*\mathbf{A})^{-1}\mathbf{A}^*\mathbf{b} = \mathbf{A}^+\mathbf{b}. \quad (39)$$

Chapter 4

This Chapter was all about QR factorization. We learned

Definition 0.11.1. Let $\mathbf{P} \in \mathbb{K}^{m \times m}$. We call \mathbf{P} a projector if and only if

$$\mathbf{P}^2 = \mathbf{P}, \quad (40)$$

i.e., \mathbf{P} is idempotent.

Remark 0.11.1. This definition includes both, orthogonal and non-orthogonal projectors. To avoid confusion, we call non-orthogonal projectors oblique projectors.

Proposition 0.11.1. If $\mathbf{P} \in \mathbb{K}^{m \times m}$ is a projector, then $\mathbf{I} - \mathbf{P}$ is also a projector.

Proof. Note that

$$(\mathbf{I} - \mathbf{P})^2 = (\mathbf{I} - \mathbf{P})(\mathbf{I} - \mathbf{P}) = \mathbf{I} - 2\mathbf{P} + \mathbf{P}^2 = \mathbf{I} - \mathbf{P} \quad (41)$$

which shows the claim. \square

Definition 0.11.2. Let $\mathbf{P} \in \mathbb{K}^{m \times m}$ be a projector. We call \mathbf{P} an orthogonal projector if and only if

$$\langle \mathbf{P}\mathbf{x}, \mathbf{y} \rangle = \langle \mathbf{x}, \mathbf{P}\mathbf{y} \rangle \quad \forall \mathbf{x}, \mathbf{y} \in \mathbb{K}^m, \quad (42)$$

i.e., $\mathbf{P} \in \mathbb{H}_m(\mathbb{K})$.

Definition 0.11.3. Let $\mathbf{A} \in \mathbb{K}^{m \times n}$. We call the factorization

$$\mathbf{A} = \mathbf{Q}\mathbf{R} \quad (43)$$

where $\mathbf{Q} \in \mathbb{K}^{m \times m}$ unitary, and $\mathbf{R} \in \mathbb{K}^{m \times n}$ is an upper triangular matrix, a QR-factorization of \mathbf{A} .

Remark 0.11.2. We shall now take a closer look at the QR-factorization:

Consider a reduced QR-factorization of $\mathbf{A} \in \mathbb{K}^{m \times n}$ with $n \leq m$, i.e.,

$$[\mathbf{a}_1 | \dots | \mathbf{a}_n] = [\mathbf{q}_1 | \dots | \mathbf{q}_n] \begin{bmatrix} r_{11} & r_{12} & \dots & r_{1n} \\ 0 & r_{22} & \dots & r_{2n} \\ \vdots & \ddots & \ddots & \vdots \\ 0 & \dots & 0 & r_{nn} \end{bmatrix} \quad (44)$$

hence

$$\begin{aligned} \mathbf{a}_1 = r_{11}\mathbf{q}_1 & \Leftrightarrow \mathbf{q}_1 = \frac{\mathbf{a}_1}{r_{11}} = \frac{\mathbf{a}_1}{\|\mathbf{a}_1\|} \\ \mathbf{a}_2 = r_{12}\mathbf{q}_1 + r_{22}\mathbf{q}_2 & \Leftrightarrow \mathbf{q}_2 = \frac{\mathbf{a}_2 - r_{12}\mathbf{q}_1}{r_{22}} = \frac{\mathbf{a}_2 - \langle \mathbf{q}_1, \mathbf{a}_2 \rangle \mathbf{q}_1}{r_{22}} = \frac{(\mathbf{I} - \mathbf{q}_1\mathbf{q}_1^*)\mathbf{a}_2}{\|(\mathbf{I} - \mathbf{q}_1\mathbf{q}_1^*)\mathbf{a}_2\|} \\ \mathbf{a}_3 = r_{13}\mathbf{q}_1 + r_{23}\mathbf{q}_2 + r_{33}\mathbf{q}_3 & \Leftrightarrow \mathbf{q}_3 = \frac{\mathbf{a}_3 - r_{13}\mathbf{q}_1 - r_{23}\mathbf{q}_2}{r_{33}} = \frac{(\mathbf{I} - \mathbf{q}_1\mathbf{q}_1^* - \mathbf{q}_2\mathbf{q}_2^*)\mathbf{a}_3}{\|(\mathbf{I} - \mathbf{q}_1\mathbf{q}_1^* - \mathbf{q}_2\mathbf{q}_2^*)\mathbf{a}_3\|} \\ & \vdots \\ \mathbf{a}_i = \sum_{j=1}^i r_{ji}\mathbf{q}_j & \Leftrightarrow \mathbf{q}_i = \frac{\mathbf{a}_i - \sum_{j=1}^{i-1} r_{ji}\mathbf{q}_j}{r_{ii}} = \frac{(\mathbf{I} - \sum_{j=1}^{i-1} \mathbf{q}_j\mathbf{q}_j^*)\mathbf{a}_i}{\|(\mathbf{I} - \sum_{j=1}^{i-1} \mathbf{q}_j\mathbf{q}_j^*)\mathbf{a}_i\|} \end{aligned} \quad (45)$$

This gave rise to three algorithms

- Classical Gram-Schmidt
- Modified Gram-Schmidt
- Iterative Gram-Schmidt

Together with their operational count.

Definition 0.11.4. Let $\mathbf{v} \in \mathbb{K}^n$ be a normal vector defining a hyperplane. The transformation

$$f_H : \mathbb{K}^n \rightarrow \mathbb{K}^n : x \mapsto x - 2\langle x, v \rangle v$$

is the Householder transformation about the hyperplane defined by the normal vector $\mathbf{v} \in \mathbb{K}^n$.

Proposition 0.11.2. Let $\mathbf{v} \in \mathbb{K}^n$ be a normal vector defining a hyperplane and f_H be the Householder transformation about the hyperplane defined by the normal vector $\mathbf{v} \in \mathbb{K}^n$. Then f_H is a linear map and its matrix representation is

$$\mathbf{P}_\mathbf{v} = \mathbf{I} - 2\mathbf{v}\mathbf{v}^*$$

Proposition 0.11.3. Let $\mathbf{v} \in \mathbb{K}^n$ be a normal vector defining a hyperplane and f_H be the Householder transformation about the hyperplane defined by the normal vector $\mathbf{v} \in \mathbb{K}^n$. The householder matrix $\mathbf{P}_\mathbf{v}$ fulfills:

- | | | |
|-----------------|--|---|
| i) Hermitian | ($\mathbf{P}_\mathbf{v} = \mathbf{P}_\mathbf{v}^*$) | iv) $\mathbf{P}_\mathbf{v}$ has eigenvalues ± 1 |
| ii) Unitary | ($\mathbf{P}_\mathbf{v}^{-1} = \mathbf{P}_\mathbf{v}^*$) | v) $\det(\mathbf{P}_\mathbf{v}) = -1$ |
| iii) Involutory | ($\mathbf{P}_\mathbf{v}^{-1} = \mathbf{P}_\mathbf{v}$) | |

Proof. First note that

$$\mathbf{P}_\mathbf{v}^* = (\mathbf{I} - 2\mathbf{v}\mathbf{v}^*)^* = \mathbf{I} - 2\mathbf{v}\mathbf{v}^* = \mathbf{P}_\mathbf{v} \quad (46)$$

which shows i). Next, we consider

$$\mathbf{P}_\mathbf{v}^2 = (\mathbf{I} - 2\mathbf{v}\mathbf{v}^*)(\mathbf{I} - 2\mathbf{v}\mathbf{v}^*) = \mathbf{I} - 4\mathbf{v}\mathbf{v}^* + 4\mathbf{v}\mathbf{v}^* = \mathbf{I} \quad (47)$$

showing that $\mathbf{P}_\mathbf{v}$ is involutory. This in turn yields that $\mathbf{P}_\mathbf{v}$ is unitary, since

$$\mathbf{P}_\mathbf{v}^{-1} = \mathbf{P}_\mathbf{v} = \mathbf{P}_\mathbf{v}^*. \quad (48)$$

Note that for $\mathbf{u} \perp \mathbf{v}$ we have $\mathbf{P}_\mathbf{v}\mathbf{u} = \mathbf{u}$. Since there are $n - 1$ linearly independent vectors $\mathbf{u} \in \mathbb{K}^n$ fulfilling $\mathbf{u} \perp \mathbf{v}$, the eigenspace of $\mathbf{P}_\mathbf{v}$ corresponding to the eigenvalue $\lambda = 1$ is $n - 1$ dimensional. Moreover $\mathbf{P}_\mathbf{v}\mathbf{v} = -\mathbf{v}$, showing iv). By iv), we know that $\mathbf{P}_\mathbf{v}$ is diagonalizable with $n - 1$ eigenvalues $\lambda_1 = 1$ and one eigenvalue $\lambda_1 = -1$. Applying the determinant multiplication Theorem we have

$$\det(\mathbf{P}_\mathbf{v}) = \det \begin{bmatrix} 1 & 0 & \cdots & 0 \\ 0 & \ddots & \ddots & \vdots \\ \vdots & \ddots & 1 & 0 \\ 0 & \cdots & 0 & -1 \end{bmatrix} = (-1) 1^{n-1} = -1 \quad (49)$$

□

The most important application:

Householder QR

We also did its operational count and argued how to keep it low:

Work with Householder vectors

Definition 0.11.5. Let $i, j \in \llbracket m \rrbracket$ and $\theta \in [0, 2\pi)$. A matrix $\mathbf{G}(i, j, \theta) \in \mathbb{K}^{m \times m}$ defined through

$$[\mathbf{G}(i, j, \theta)]_{l,m} = \begin{cases} 1 & , \text{if } l = m, \text{ and } l \neq i, j \\ \cos(\theta) & , \text{if } l = m = i, j \\ \sin(\theta) & , \text{if } l = i, \text{ and } m = j \\ -\sin(\theta) & , \text{if } l = j, \text{ and } m = i \\ 0 & , \text{else.} \end{cases} \quad (50)$$

is called Givens rotation around θ in the i - j -plane.

Proposition 0.11.4. Givens rotations are orthogonal matrices, i.e., $\mathbf{G}^\top = \mathbf{G}^{-1}$.

Remark 0.11.3. Givens rotations indeed rotate in the i - j -plane. Consider

$$\mathbf{G}(i, j, \theta)\mathbf{x} = \begin{bmatrix} x_1 \\ \vdots \\ x_{i-1} \\ cx_i - sx_j \\ x_{i+1} \\ \vdots \\ cx_j + sx_i \\ x_{j+1} \\ \vdots \\ x_m \end{bmatrix} = \begin{bmatrix} c & -s \\ s & c \end{bmatrix} \begin{bmatrix} x_i \\ x_j \end{bmatrix} \quad (51)$$

substituting c and s with $\cos(\theta)$ and $\sin(\theta)$, respectively, we see that this corresponds to a (counter-clockwise) rotation through an angle θ in the i - j -plane.

We designed an algorithm that uses Givens rotations to compute a QR factorization and discussed the operational count, and how to keep it low:

Track only the Givens angles.