Compilers Parse Generators and ASTs

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FEUP-FCUP

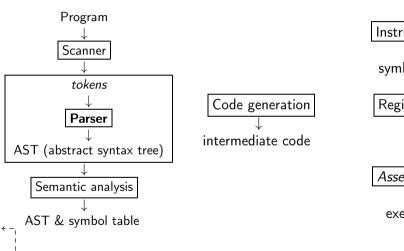
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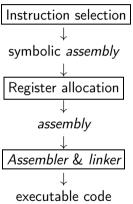
Parser generators JavaCC

Abstract Syntax

Abstract Syntax
Functional Languages
C
Java and OO language
Other languages

Compiler





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Parser generators

► Parsers generators:

```
JavaCC LL parser generator for Java
ANTLR LL parser generator for Java
StableCC LALR parser generator for Java, C++, C, Python and OCaml
Yacc Unix LALR parser generator for C
Bison a GNU version of Yacc for C and C++
Happy LALR parser generator for Haskell
```

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JavaCC

- ► LL Parser generator
- Grammar Rules are of the form:

```
void Assignment() : {} { Identifier() "=" Expression() ";" }
```

```
PARSER_BEGIN(MyParser)
    public class MyParser {}
PARSER_END(MyParser)
SKIP :
{ " " | "\t" | "\n" }
TOKEN:
{ < WHILE: "while" >
| < BEGIN: "begin" >
| < END: "end" > | < DO: "do" > | < IF: "if" >
| < THEN: "then" >
| < ELSE: "else" > | < SEMI: ";" >
| < ASSIGN: "=" >
| < ID: ["a"-"z"](["a"-"z"] | ["0"-"9"])* >
```

```
void Prog() :
{}
{ StmList() <EOF> }
void StmList() :
{}
{ Stm() StmListPrime() }
void StmListPrime() :
{}
{ ( ": " Stm() StmListPrime() )? }
```

```
void Stm() :
{}
{ <ID> "=" <ID>
| "while" <ID> "do" Stm()
| "begin" StmList() "end"
| LOOKAHEAD(5) /* we need to lookahead until the else */
    "if" <ID> "then" Stm()
| "if" <ID> "then" Stm()
}
```

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Abstract Syntax

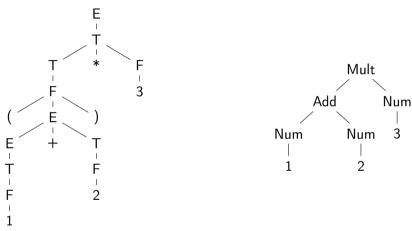
Consider the following grammar:

$$E \rightarrow E + T$$
 $T \rightarrow T * F$ $F \rightarrow \text{num}$ $E \rightarrow T$ $T \rightarrow F$ $F \rightarrow (E)$

and the derivation:

$$E \Rightarrow T \Rightarrow T * F \Rightarrow F * F \Rightarrow (E) * F \Rightarrow \cdots \Rightarrow (1+2)*3$$

Abstract Syntax (cont.)



The parse tree on the left has redundant information.

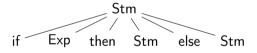
The abstract syntax tree (AST) on the right has one node per operation.

Abstract Syntax (cont.)

Example: commands.

$$Stm \rightarrow if \ Exp \ then \ Stm \ else \ Stm \ Stm \rightarrow etc.$$

In the parse tree each keyword is a different node.



In the AST we only need one node for the if/then/else:

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Functional Languages

In ML, OCaml, Haskell or F# the AST is an algebraic type:

- Alternatives identified by the constructors
- ► Datatypes may be recursive
- Processing AST uses pattern matching

```
Example of an AST:
example :: Stm
example =
 CompoundStm
  (AssignStm "a"
    (OpExp (NumExp 5) Plus (NumExp 3))
  (AssignStm "b"
    (OpExp (IdExp "a") Minus (NumExp 2))
```

```
Processing using pattern matching:
process :: Stm -> ...
process (IncrStm id) = ...
process (AssignStm id exp) = ...
process (CompoundStm s1 s2) = ...
Or cases:
process stm = case stm of
   IncrStm id -> ...
   AssignStm id exp -> ...
   CompoundStm s1 s2 -> ...
```

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- ► A struct for each syntactic element:
 - ► a (tag) for the kind;
 - a union of alternatives;
- Constructors
- ▶ Use the *tag* to process the alternatives

```
struct stm {
 enum
                               typedef enum {PLUS, MINUS, TIMES, DIV} binop;
  {COMPOUND, ASSIGN, INCR} tag;
 union {
                                struct _exp {
                                 enum {ID, NUM, OP} tag;
   struct { // for COMPOUND
     struct _stm *fst, *snd; union {
   } compound;
                                   int val; // for NUM
                                   char *id; // for ID
   struct { // for ASSIGN
     char *ident;
                                   struct { // for OP
     struct _exp *expr;
                                     binop op;
   } assign;
                                     struct _exp *left, *right;
   char *ident: // for INCR
                                   } binop:
 };
};
                               typedef struct _exp *Exp;
typedef struct _stm *Stm;
```

```
Exp mk_num(int v) {
  Exp e = (Exp) malloc(sizeof(struct _exp));
  e->tag = NUM;
  e->val = v;
  return e:
Exp mk_ident(char *txt) {
  Exp e = (Exp)malloc(sizeof(struct _exp));
  char *str = malloc(strlen(txt)+1);
  strcpy(str, txt);
  e->tag = ID;
  e->id = str;
  return e;
```

```
Exp mk_op(binop op, Exp e1, Exp e2) {
   Exp e = (Exp) malloc(sizeof(struct _exp));
   e->tag = OP;
   e->binop.op = op;
   e->binop.left = e1;
   e->binop.right = e2;
   return e;
}
(The other constructors are similar)
```

```
Example of an AST:
Stm example =
    mk_compound
    (mk_assign("a", mk_op(PLUS, mk_num(5), mk_num(3))),
        mk_assign("b", mk_op(TIMES, mk_ident("a"), mk_num(2))));
```

Building the AST in Yacc/Bison:

```
Processing the AST using tags:
void process(Stm stm) {
   switch(stm->tag) {
     case COMPOUND: // use stm->compound.fst and stm->compound.snd
     break;
     case ASSIGN: // use stm->assign.ident and stm->assign.expr
     break:
     case INCR: // etc.
     break;
```

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Em OO languages:

- An abstract class for each syntactic category;
- ► A subclass for each alternative
- ► To identify the alternatives use instanceof and downcasts

```
public abstract class Stm {}
                                       public class AssignStm extends Stm
public class CompoundStm extends Stm
                                         public String id;
 public Stm fst, snd;
                                         public Exp exp;
  public CompoundStm(Stm s1, Stm s2)
                                         public AssignStm(String i, Exp e)
    fst = s1; snd = s2;
                                           id=i; exp=e;
```

(Other subclasses for the other alternatives)

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```
public abstract class Exp {}
public class IdExp extends Exp {
                                       public class NumExp extends Exp {
 public String id;
                                         public int num;
 public IdExp(String i) {
                                         public NumExp(int n) {
   id=i:
                                           num=n;
public class PlusExp extends Exp {
 private Exp left, right;
  public PlusExp(Exp e1, Exp e2) {
   left=e1; right=e2;
Example of an AST:
```

Building the AST in *JavaCC*:

```
Exp Exp():
{ Exp e1,e2; }
{ e1=Term()
    ( "+" e2=Term() { e1=new PlusExp(e1,e2); }
    | "-" e2=Term() { e1=new MinusExp(e1,e2); } )*
  { return e1: }
Exp Term() :
{ Exp e1,e2; }
{ e1=Factor()
    ( "*" e2=Factor() { e1=new TimesExp(e1,e2); }
    | "/" e2=Factor() { e1=new DivideExp(e1,e2); })*
  { return e1; }
```

```
Processing using cases:
public void process(Stm stm) {
  if(stm instanceof CompoundStm) {
      CompoundStm cstm = (CompoundStm)stm;
  else if(stm instanceof AssignStm) {
       AssignStm astm = (AssignStm)stm;
  else ...
```

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Other languages

- ► Modern multi-paradigm languages (OO+Funcional) use pattern matching (as in Haskell):
 - "case classes" em Scala
 - "enumerations" em Swift e Rust