

# Applied Cryptography

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To perform well in this course

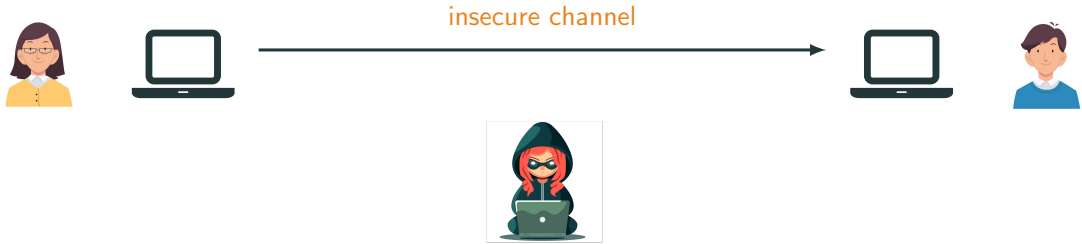
- Follow the topics from lectures: solve the given problems
- Ask (questions) if you can not solve them
- Go through the suggested reading: ask if you face difficulties
- Take notes during lecture
- I encourage discussing with your colleagues (on a topic or problems)

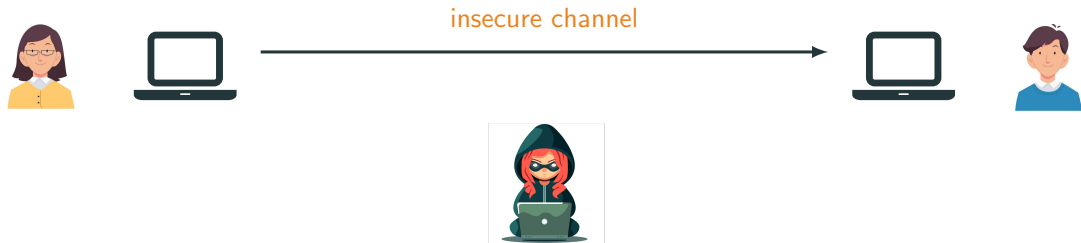
**Recommended textbook:** C. Paar, J. Pelzl: *Understanding Cryptography*

# Introduction to Cryptography

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# Communication security

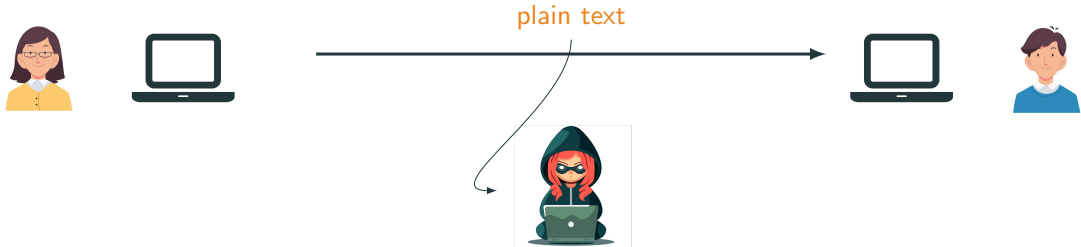


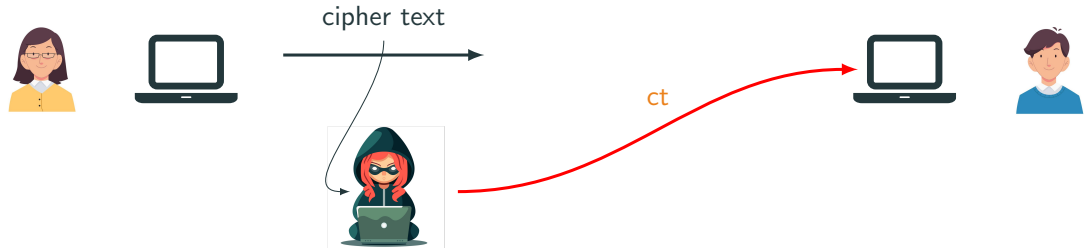


## Security aims (data in communication or under storage)

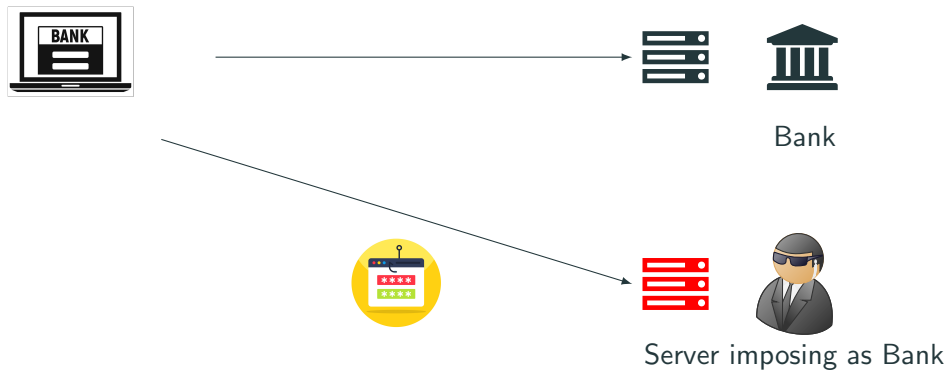
- **Confidentiality:** Adversary must not be able to **read** message  $m$
- **Integrity:** Adversary must not be able to **modify** message  $m$
- **Authenticity:** Receiver should be able to ensure that  $m$  originated from sender

# Confidentiality or secrecy





# Authentication





# Adversary

Adversarial powers

Adversarial goals

## Adversarial powers

- With finite computing power: computational security
- Access to physical device and implementations (observable or subject to manipulation): side channel security or implementation security (not covered in this course)
- Unlimited computational power: information theoretic security or perfect secrecy (not covered in this course)
- Access to input and output: choosing or obtaining plaintext(s) and/or ciphertext(s)

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## Adversarial goals

- Recovering secret key: allows to read all messages encrypted with that key
- Plaintext recovery: recovering plaintext(s) from observable or chosen ciphertexts

## General adversarial (attack) models

- Ciphertext only: adversary gains access to ciphertexts
- Known plaintext: access to  $(x, \text{Enc}(x))$  but no control over plaintext
- Chosen plaintext: chooses  $x$  and obtains corresponding  $\text{Enc}(x)$
- Chosen plaintext (adaptive): choose  $x_1$  and obtains  $\text{Enc}(x_1)$ ; then chooses  $x_2$  and obtains  $\text{Enc}(x_2)$
- Known ciphertext: access to  $(y, \text{Dec}(y))$  but no control over ciphertext
- Chosen ciphertext (adaptive)

## Security against generic attacks

- Adversary only have access to input and/or outputs from Enc
- No knowledge of Enc algorithm or its description

**Opening the box:** Adversary knows the algorithm or function description of Enc (cryptanalysis)

**Side-channel:** Adversary knows the algorithm and additional information from the execution of the algorithm [NOT covered in this course]

# Symmetric Cryptography

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# Syntax of symmetric encryption

Symmetric encryption  $\mathcal{E} = (\text{Gen}, \text{Enc}, \text{Dec})$

- $\text{Gen}(\kappa) \xrightarrow{\$} sk$
- $\text{Enc} : \mathcal{M} \times \mathcal{K} \rightarrow \mathcal{C}$  s.t  $\text{Enc}(x, sk) = y \in \mathcal{C}$
- $\text{Dec} : \mathcal{C} \times \mathcal{K} \rightarrow \mathcal{M}$  s.t  $\text{Dec}(y, sk) = x \in \mathcal{M}$
- $\text{Dec}(\text{Enc}(x, sk), sk) = x$  (correctness)

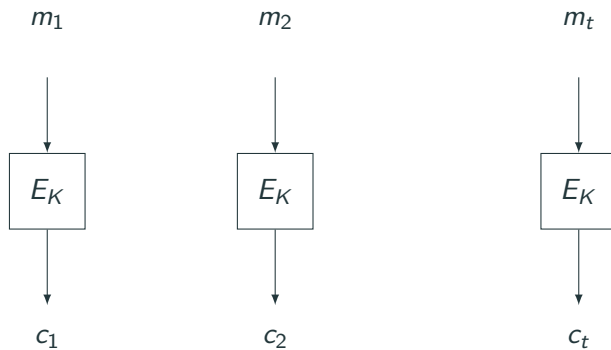
**Note:** Given a fixed  $k \in \mathcal{K}$ , the  $\text{Enc}(\cdot, k)$  is a permutation ( $|\mathcal{M}| = |\mathcal{C}|$ )

- **Qn:** Let  $\mathcal{M} = \mathcal{C} = \{0, 1\}^n$  how many different permutations  $f : \mathcal{M} \rightarrow \mathcal{C}$  are there?
- **Qn:** If  $\mathcal{K} = \{0, 1\}^m$  how many permutations are possible to have with  $\text{Enc}$ ?

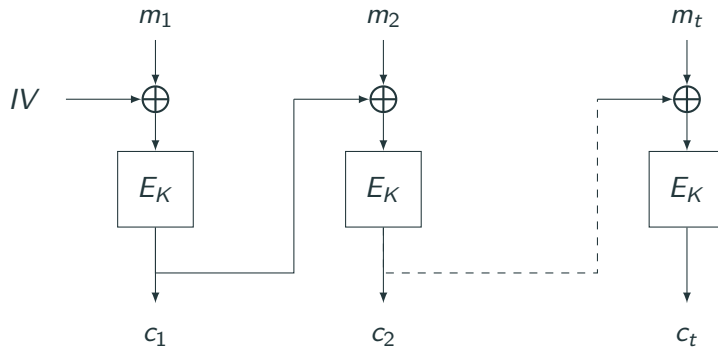
- Block ciphers are fixed length primitives
  - AES (Advanced Encryption Standard) can process 128 bits of input, key size = 128, 192 and 256 bits
  - DES (Data Encryption Standard): input size = 64 bits, key size = 56 bits
- ❓ How to encrypt data with arbitrary length? Secure mode of operation
  - Examples: ECB, OFB, CBC, CTR, GCM



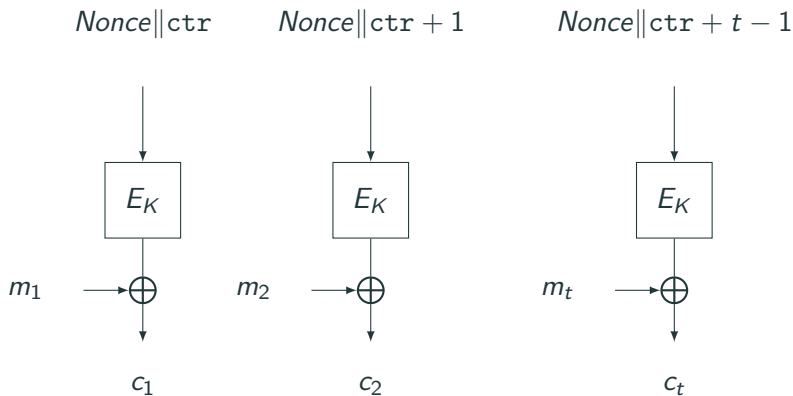
# Electronic Code Book (ECB) Mode



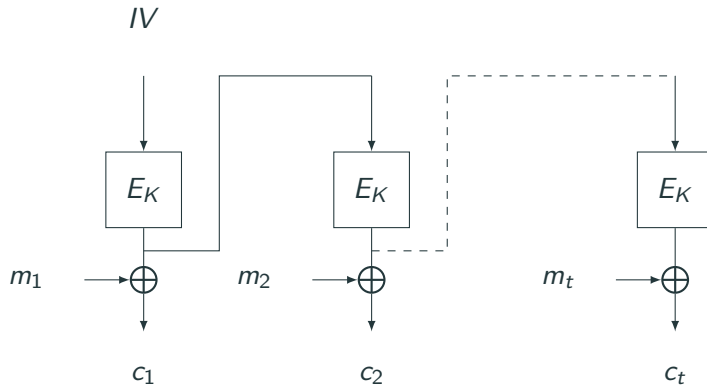
# Cipher Block Chaining (CBC) Mode





## Counter (CTR) Mode



## Output Feedback (OFB) Mode



Symmetric-key crypto (SKC) is everywhere


TLS: on our Web-browsers  

Cards: Payment cards  

Wireless communications: Wifi , Bluetooth 

Mobile communications: 

Encrypted data storage: 

Crypto-currencies: 

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Thank you!