



Chapter 2

Instructions: Language of the Computer

Instruction Set

- The repertoire of instructions of a computer
- Different computers have different instruction sets
 - But with many aspects in common
- Early computers had very simple instruction sets
 - Simplified implementation
- Many modern computers also have simple instruction sets

The RISC-V Instruction Set

- Used as the example throughout the book
- Developed at UC Berkeley as open ISA
- Now managed by the RISC-V Foundation
(riscv.org)
- Typical of many modern ISAs
 - See RISC-V Reference Data tear-out card
- Similar ISAs have a large share of embedded core market
 - Applications in consumer electronics, network/storage equipment, cameras, printers, ...



Arithmetic Operations

- Add and subtract, three operands
 - Two sources and one destination

```
add a, b, c // a gets b + c
```
- All arithmetic operations have this form
- *Design Principle 1: Simplicity favours regularity*
 - Regularity makes implementation simpler
 - Simplicity enables higher performance at lower cost

Arithmetic Example

- C code:

```
f = (g + h) - (i + j);
```

- Compiled RISC-V code:

```
add t0, g, h    // temp t0 = g + h  
add t1, i, j    // temp t1 = i + j  
add f, t0, t1   // f = t0 - t1
```

Register Operands

- Arithmetic instructions use register operands
- RISC-V has a 32×64 -bit register file
 - Use for frequently accessed data
 - 64-bit data is called a “doubleword”
 - 32×64 -bit general purpose registers x0 to x30
 - 32-bit data is called a “word”
- *Design Principle 2: Smaller is faster*
 - c.f. main memory: millions of locations

RISC-V Registers

- x0: the constant value 0
- x1: return address
- x2: stack pointer
- x3: global pointer
- x4: thread pointer
- x5 – x7, x28 – x31: temporaries
- x8: frame pointer
- x9, x18 – x27: saved registers
- x10 – x11: function arguments/results
- x12 – x17: function arguments

Register Operand Example

- C code:

$f = (g + h) - (i + j);$

- f, ..., j in x19, x20, ..., x23

- Compiled RISC-V code:

add x5, x20, x21

add x6, x22, x23

sub x19, x5, x6

Memory Operands

- Main memory used for composite data
 - Arrays, structures, dynamic data
- To apply arithmetic operations
 - Load values from memory into registers
 - Store result from register to memory
- Memory is byte addressed
 - Each address identifies an 8-bit byte
- RISC-V is Little Endian
 - Least-significant byte at least address of a word
 - *c.f.* Big Endian: most-significant byte at least address
- RISC-V does not require words to be aligned in memory
 - Unlike some other ISAs



Memory Operand Example

- C code:

$A[12] = h + A[8];$

- h in $x21$, base address of A in $x22$

- Compiled RISC-V code:

- Index 8 requires offset of 64
 - 8 bytes per doubleword

ld x9, 64(x22)

add x9, x21, x9

sd x9, 96(x22)

Registers vs. Memory

- Registers are faster to access than memory
- Operating on memory data requires loads and stores
 - More instructions to be executed
- Compiler must use registers for variables as much as possible
 - Only spill to memory for less frequently used variables
 - Register optimization is important!



Immediate Operands

- Constant data specified in an instruction

```
addi x22, x22, 4
```

- Make the common case fast
 - Small constants are common
 - Immediate operand avoids a load instruction

Unsigned Binary Integers

- Given an n-bit number

$$x = x_{n-1}2^{n-1} + x_{n-2}2^{n-2} + \dots + x_12^1 + x_02^0$$

- Range: 0 to $+2^n - 1$
- Example
 - 0000 0000 ... 0000 1011₂
= 0 + ... + 1×2³ + 0×2² + 1×2¹ + 1×2⁰
= 0 + ... + 8 + 0 + 2 + 1 = 11₁₀
- Using 64 bits: 0 to +18,446,774,073,709,551,615

2s-Complement Signed Integers

- Given an n-bit number

$$x = -x_{n-1}2^{n-1} + x_{n-2}2^{n-2} + \dots + x_12^1 + x_02^0$$

- Range: -2^{n-1} to $+2^{n-1} - 1$
- Example
 - $1111\ 1111\ \dots\ 1111\ 1100_2$
 $= -1 \times 2^{31} + 1 \times 2^{30} + \dots + 1 \times 2^2 + 0 \times 2^1 + 0 \times 2^0$
 $= -2,147,483,648 + 2,147,483,644 = -4_{10}$
 - Using 64 bits: $-9,223,372,036,854,775,808$ to $9,223,372,036,854,775,807$

2s-Complement Signed Integers

- Bit 63 is sign bit
 - 1 for negative numbers
 - 0 for non-negative numbers
- $-(-2^n - 1)$ can't be represented
- Non-negative numbers have the same unsigned and 2s-complement representation
- Some specific numbers
 - 0: 0000 0000 ... 0000
 - -1: 1111 1111 ... 1111
 - Most-negative: 1000 0000 ... 0000
 - Most-positive: 0111 1111 ... 1111

Signed Negation

- Complement and add 1
 - Complement means $1 \rightarrow 0, 0 \rightarrow 1$

$$x + \bar{x} = 1111\dots111_2 = -1$$

$$\bar{x} + 1 = -x$$

- Example: negate +2
 - $+2 = 0000\ 0000\dots0010_{\text{two}}$
 - $-2 = 1111\ 1111\dots1101_{\text{two}} + 1$
 $= 1111\ 1111\dots1110_{\text{two}}$

Sign Extension

- Representing a number using more bits
 - Preserve the numeric value
- Replicate the sign bit to the left
 - c.f. unsigned values: extend with 0s
- Examples: 8-bit to 16-bit
 - +2: 0000 0010 => 0000 0000 0000 0010
 - -2: 1111 1110 => 1111 1111 1111 1110
- In RISC-V instruction set
 - 1b: sign-extend loaded byte
 - 1bu: zero-extend loaded byte

Representing Instructions

- Instructions are encoded in binary
 - Called machine code
- RISC-V instructions
 - Encoded as 32-bit instruction words
 - Small number of formats encoding operation code (opcode), register numbers, ...
 - Regularity!

Hexadecimal

■ Base 16

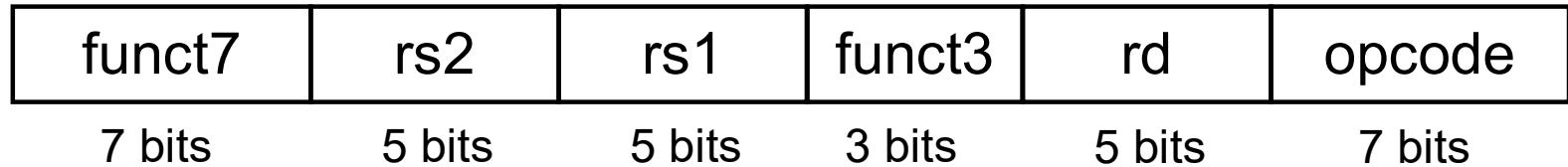
- Compact representation of bit strings
- 4 bits per hex digit

0	0000	4	0100	8	1000	c	1100
1	0001	5	0101	9	1001	d	1101
2	0010	6	0110	a	1010	e	1110
3	0011	7	0111	b	1011	f	1111

■ Example: eca8 6420

- 1110 1100 1010 1000 0110 0100 0010 0000

RISC-V R-format Instructions



Instruction fields

- opcode: operation code
- rd: destination register number
- funct3: 3-bit function code (additional opcode)
- rs1: the first source register number
- rs2: the second source register number
- funct7: 7-bit function code (additional opcode)

R-format Example

funct7	rs2	rs1	funct3	rd	opcode
7 bits	5 bits	5 bits	3 bits	5 bits	7 bits

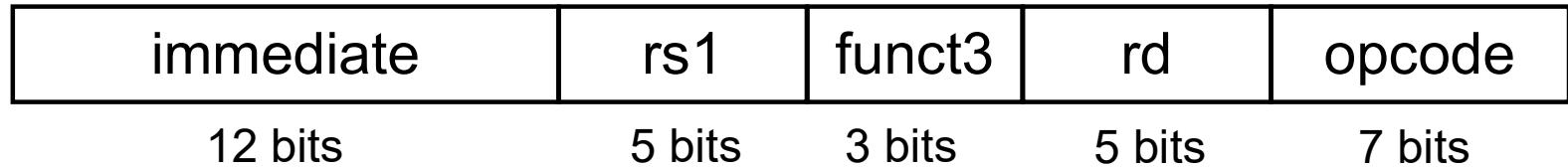
add x9, x20, x21

0	21	20	0	9	51
---	----	----	---	---	----

0000000	10101	10100	000	01001	0110011
---------	-------	-------	-----	-------	---------

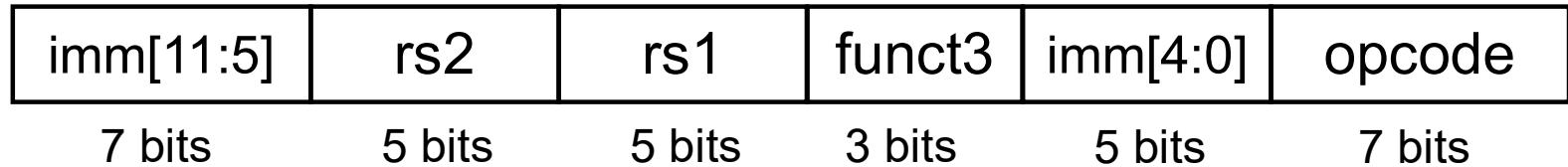
0000 0001 0101 1010 0000 0100 1011 0011_{two} =
015A04B3₁₆

RISC-V I-format Instructions



- Immediate arithmetic and load instructions
 - rs1: source or base address register number
 - immediate: constant operand, or offset added to base address
 - 2s-complement, sign extended
- *Design Principle 3: Good design demands good compromises*
 - Different formats complicate decoding, but allow 32-bit instructions uniformly
 - Keep formats as similar as possible

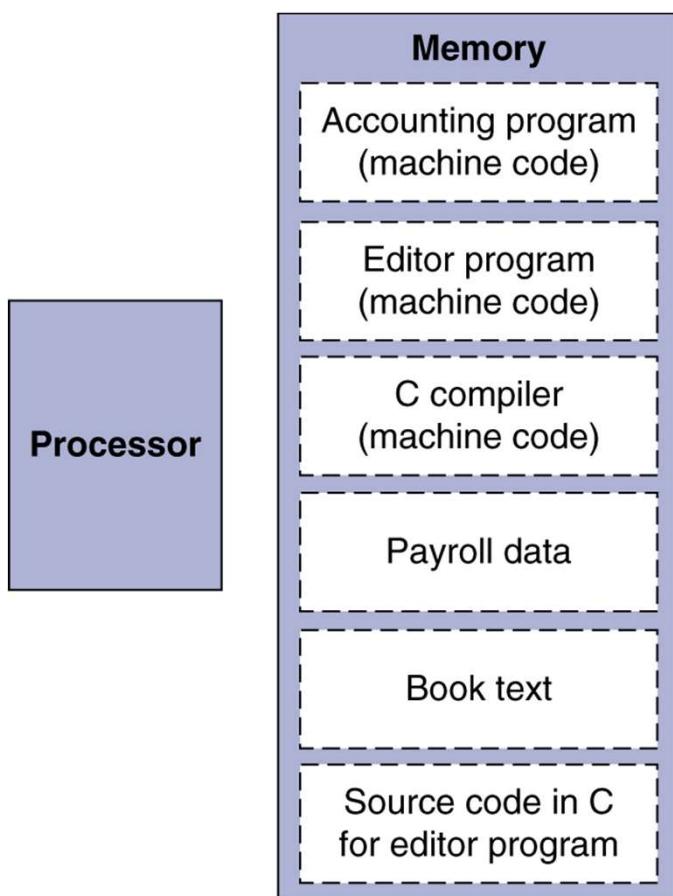
RISC-V S-format Instructions



- Different immediate format for store instructions
 - rs1: base address register number
 - rs2: source operand register number
 - immediate: offset added to base address
 - Split so that rs1 and rs2 fields always in the same place

Stored Program Computers

The BIG Picture



- Instructions represented in binary, just like data
- Instructions and data stored in memory
- Programs can operate on programs
 - e.g., compilers, linkers, ...
- Binary compatibility allows compiled programs to work on different computers
 - Standardized ISAs

Logical Operations

Instructions for bitwise manipulation

Operation	C	Java	RISC-V
Shift left	<<	<<	slli
Shift right	>>	>>>	srli
Bit-by-bit AND	&	&	and, andi
Bit-by-bit OR			or, ori
Bit-by-bit XOR	^	^	xor, xorri
Bit-by-bit NOT	~	~	

Useful for extracting and inserting groups of bits in a word

Shift Operations



- immed: how many positions to shift
- Shift left logical
 - Shift left and fill with 0 bits
 - sll_i by i bits multiplies by 2^i
- Shift right logical
 - Shift right and fill with 0 bits
 - srl_i by i bits divides by 2^i (unsigned only)

AND Operations

- Useful to mask bits in a word
 - Select some bits, clear others to 0

and x_9, x_{10}, x_{11}

x_{10} 00000000 00000000 00000000 00000000 00000000 00000000 00001101 11000000

x_{11} 00000000 00000000 00000000 00000000 00000000 00000000 00111100 00000000

x_9 00000000 00000000 00000000 00000000 00000000 00000000 00001100 00000000

OR Operations

- Useful to include bits in a word
 - Set some bits to 1, leave others unchanged

or $x9, x10, x11$

x10	00000000 00000000 00000000 00000000 00000000 00000000 00001101 11000000
x11	00000000 00000000 00000000 00000000 00000000 00000000 00111100 00000000
x9	00000000 00000000 00000000 00000000 00000000 00000000 00111101 11000000

XOR Operations

- Differencing operation
 - Set some bits to 1, leave others unchanged

```
xor x9,x10,x12 // NOT operation
```

x10	00000000 00000000 00000000 00000000 00000000 00000000 00001101 11000000
x12	11111111 11111111 11111111 11111111 11111111 11111111 11111111 11111111
x9	11111111 11111111 11111111 11111111 11111111 11111111 11110010 00111111

Conditional Operations

- Branch to a labeled instruction if a condition is true
 - Otherwise, continue sequentially
- **beq rs1, rs2, L1**
 - if ($rs1 == rs2$) branch to instruction labeled L1
- **bne rs1, rs2, L1**
 - if ($rs1 != rs2$) branch to instruction labeled L1

Compiling If Statements

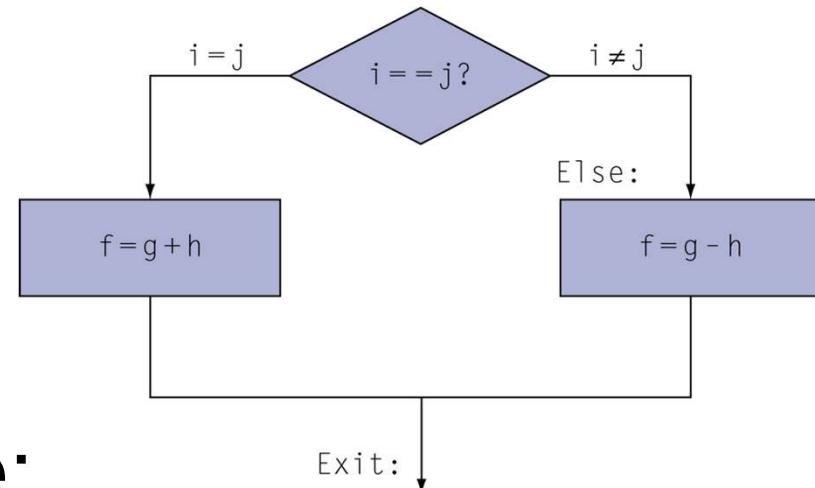
- C code:

```
if (i==j) f = g+h;  
else f = g-h;
```

- f, g, ... in x19, x20, ...

- Compiled RISC-V code:

```
bne x22, x23, Else  
add x19, x20, x21  
beq x0,x0,Exit // unconditional  
Else: sub x19, x20, x21  
Exit: ...
```



Assembler calculates addresses

Compiling Loop Statements

- C code:

```
while (save[i] == k) i += 1;
```

- i in x22, k in x24, address of save in x25

- Compiled RISC-V code:

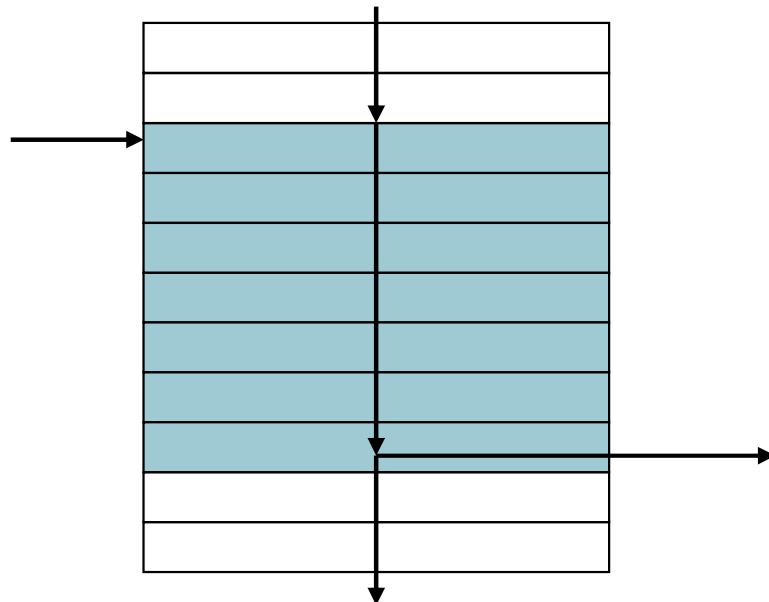
```
Loop: slli x10, x22, 3
      add x10, x10, x25
      ld x9, 0(x10)
      bne x9, x24, Exit
      addi x22, x22, 1
      beq x0, x0, Loop
```

```
Exit: ...
```



Basic Blocks

- A basic block is a sequence of instructions with
 - No embedded branches (except at end)
 - No branch targets (except at beginning)



- A compiler identifies basic blocks for optimization
- An advanced processor can accelerate execution of basic blocks

More Conditional Operations

- `blt rs1, rs2, L1`
 - if ($rs1 < rs2$) branch to instruction labeled L1
- `bge rs1, rs2, L1`
 - if ($rs1 \geq rs2$) branch to instruction labeled L1
- Example
 - if ($a > b$) $a += 1$;
 - a in x22, b in x23

`bge x23, x22, Exit // branch if $b \geq a$`

`addi x22, x22, 1`

Exit:



Signed vs. Unsigned

- Signed comparison: blt, bge
- Unsigned comparison: bltu, bgeu
- Example
 - $x_{22} = 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111$
 - $x_{23} = 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0001$
 - $x_{22} < x_{23}$ // signed
 - $-1 < +1$
 - $x_{22} > x_{23}$ // unsigned
 - $+4,294,967,295 > +1$

Procedure Calling

■ Steps required

1. Place parameters in registers x10 to x17
2. Transfer control to procedure
3. Acquire storage for procedure
4. Perform procedure's operations
5. Place result in register for caller
6. Return to place of call (address in x1)

Procedure Call Instructions

- Procedure call: jump and link
`jal x1, ProcedureLabel`
 - Address of following instruction put in x1
 - Jumps to target address
- Procedure return: jump and link register
`jalr x0, 0(x1)`
 - Like jal, but jumps to $0 + \text{address in } x1$
 - Use x0 as rd ($x0$ cannot be changed)
 - Can also be used for computed jumps
 - e.g., for case/switch statements



Leaf Procedure Example

- C code:

```
long long int leaf_example (
    long long int g, long long int h,
    long long int i, long long int j) {
    long long int f;
    f = (g + h) - (i + j);
    return f;
}
```

- Arguments g, ..., j in x10, ..., x13
- f in x20
- temporaries x5, x6
- Need to save x5, x6, x20 on stack

Leaf Procedure Example

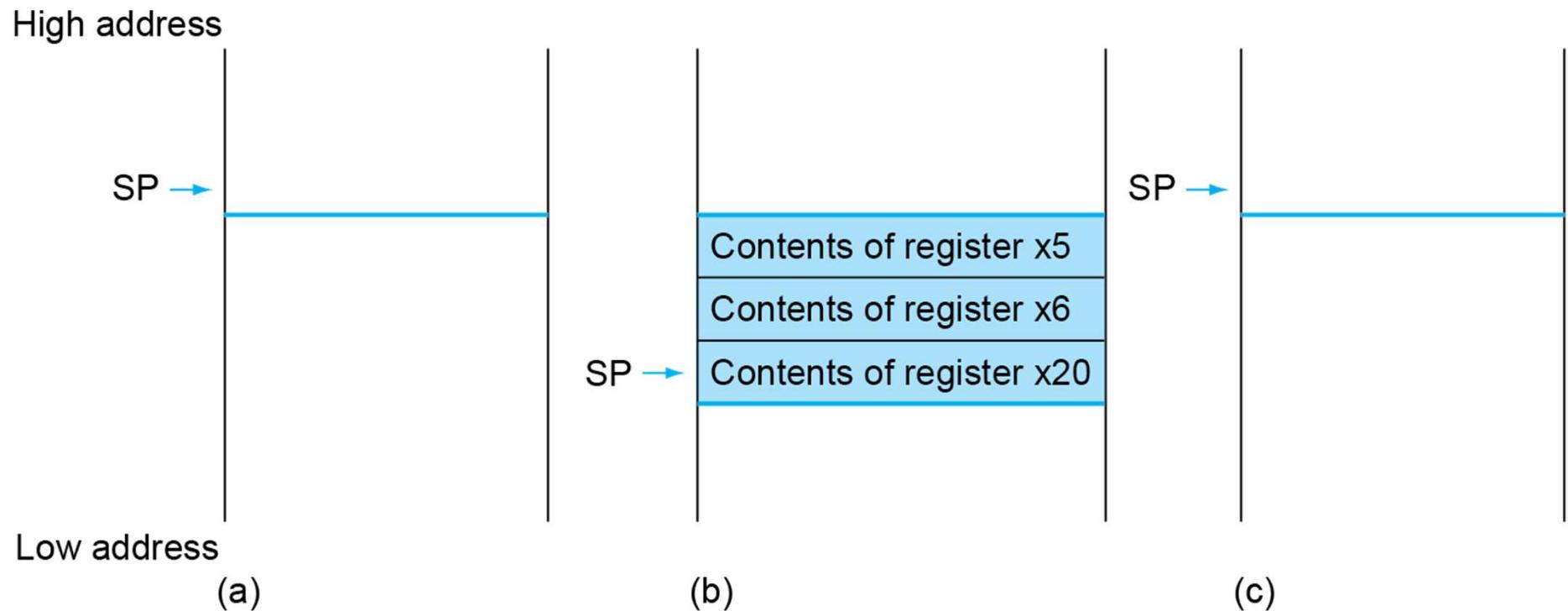
■ RISC-V code:

leaf_example:

addi sp,sp,-24	Save x5, x6, x20 on stack
sd x5,16(sp)	
sd x6,8(sp)	
sd x20,0(sp)	
add x5,x10,x11	$x5 = g + h$
add x6,x12,x1	$x6 = i + j$
sub x20,x5,x6	$f = x5 - x6$
addi x10,x20,0	copy f to return register
ld x20,0(sp)	Resore x5, x6, x20 from stack
ld x6,8(sp)	
ld x5,16(sp)	
addi sp,sp,24	
jalr x0,0(x1)	Return to caller



Local Data on the Stack



Register Usage

- x5 – x7, x28 – x31: temporary registers
 - Not preserved by the callee
- x8 – x9, x18 – x27: saved registers
 - If used, the callee saves and restores them

Non-Leaf Procedures

- Procedures that call other procedures
- For nested call, caller needs to save on the stack:
 - Its return address
 - Any arguments and temporaries needed after the call
- Restore from the stack after the call



Non-Leaf Procedure Example

- C code:

```
long long int fact (long long int n)
{
    if (n < 1) return f;
    else return n * fact(n - 1);
}
```

- Argument n in x10
- Result in x10

Leaf Procedure Example

RISC-V code:

fact:

```
addi sp,sp,-16
sd x1,8(sp)
sd x10,0(sp)
addi x5,x10,-1
bge x5,x0,L1
addi x10,x0,1
addi sp,sp,16
jalr x0,0(x1)

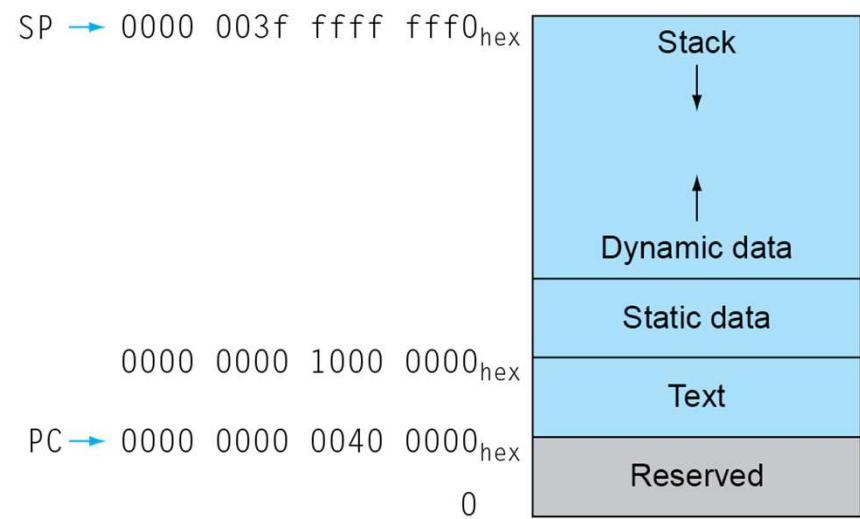
L1: addi x10,x10,-1
jal x1,fact
addi x6,x10,0
ld x10,0(sp)
ld x1,8(sp)
addi sp,sp,16
mul x10,x10,x6
jalr x0,0(x1)
```

Save return address and n on stack
x5 = n - 1
if n >= 1, go to L1
Else, set return value to 1
Pop stack, don't bother restoring values
Return
n = n - 1
call fact(n-1)
move result of fact(n - 1) to x6
Restore caller's n
Restore caller's return address
Pop stack
return n * fact(n-1)
return

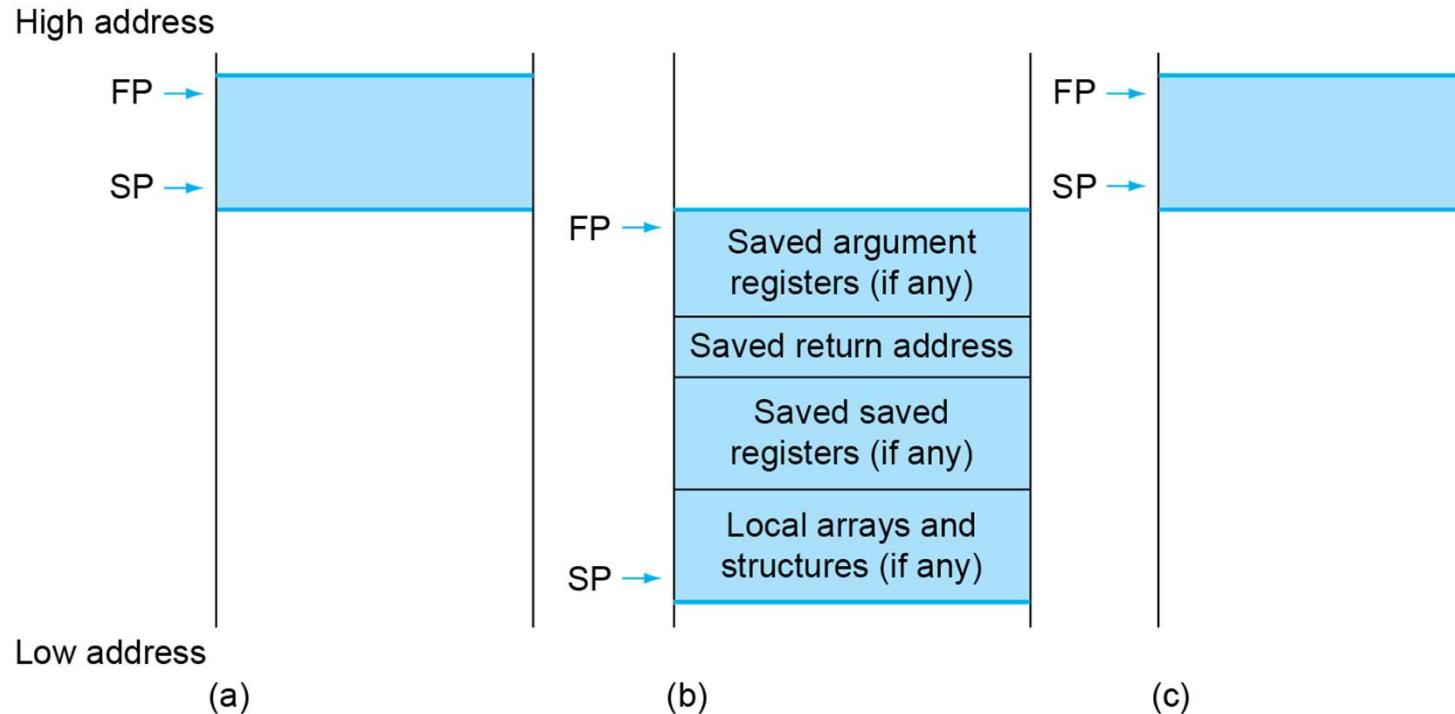


Memory Layout

- Text: program code
- Static data: global variables
 - e.g., static variables in C, constant arrays and strings
 - x3 (global pointer) initialized to address allowing \pm offsets into this segment
- Dynamic data: heap
 - E.g., malloc in C, new in Java
- Stack: automatic storage



Local Data on the Stack



- Local data allocated by callee
 - e.g., C automatic variables
- Procedure frame (activation record)
 - Used by some compilers to manage stack storage

Character Data

- Byte-encoded character sets
 - ASCII: 128 characters
 - 95 graphic, 33 control
 - Latin-1: 256 characters
 - ASCII, +96 more graphic characters
- Unicode: 32-bit character set
 - Used in Java, C++ wide characters, ...
 - Most of the world's alphabets, plus symbols
 - UTF-8, UTF-16: variable-length encodings

Byte/Halfword/Word Operations

- RISC-V byte/halfword/word load/store
 - Load byte/halfword/word: Sign extend to 64 bits in rd
 - `lb rd, offset(rs1)`
 - `lh rd, offset(rs1)`
 - `lw rd, offset(rs1)`
 - Load byte/halfword/word unsigned: Zero extend to 64 bits in rd
 - `lbu rd, offset(rs1)`
 - `lhu rd, offset(rs1)`
 - `lwu rd, offset(rs1)`
 - Store byte/halfword/word: Store rightmost 8/16/32 bits
 - `sb rs2, offset(rs1)`
 - `sh rs2, offset(rs1)`
 - `sw rs2, offset(rs1)`



String Copy Example

- C code:

- Null-terminated string

```
void strcpy (char x[], char y[])
{ size_t i;
  i = 0;
  while ((x[i]=y[i]) != '\0')
    i += 1;
}
```

String Copy Example

RISC-V code:

```
strcpy:  
    addi sp,sp,-8          // adjust stack for 1 doubleword  
    sd    x19,0(sp)        // push x19  
    add   x19,x0,x0        // i=0  
L1:  add   x5,x19,x10    // x5 = addr of y[i]  
    lbu   x6,0(x5)        // x6 = y[i]  
    add   x7,x19,x10    // x7 = addr of x[i]  
    sb    x6,0(x7)        // x[i] = y[i]  
    beq   x6,x0,L2        // if y[i] == 0 then exit  
    addi  x19,x19,1        // i = i + 1  
    jal   x0,L1            // next iteration of loop  
L2:  ld    x19,0(sp)        // restore saved x19  
    addi  sp,sp,8            // pop 1 doubleword from stack  
    jalr x0,0(x1)          // and return
```

32-bit Constants

- Most constants are small
 - 12-bit immediate is sufficient
- For the occasional 32-bit constant
 - **lui rd, constant**
 - Copies 20-bit constant to bits [31:12] of rd
 - Extends bit 31 to bits [63:32]
 - Clears bits [11:0] of rd to 0

lui x19, 976 // 0x003D0

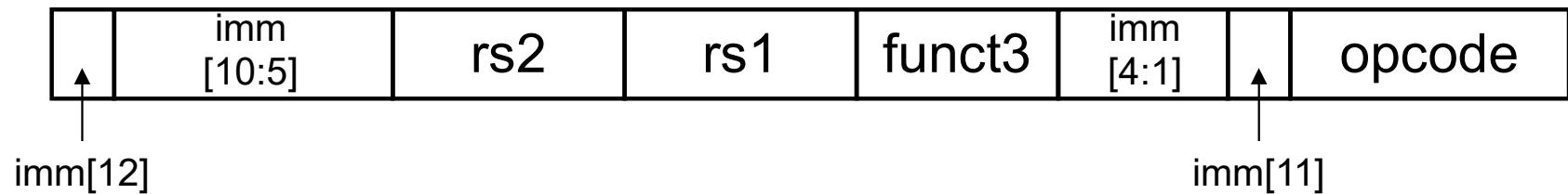
0000 0000 0000 0000	0000 0000 0000 0000	0000 0000 0011 1101 0000	0000 0000 0000
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addi x19, x19, 128 // 0x500

0000 0000 0000 0000	0000 0000 0000 0000	0000 0000 0011 1101 0000	0101 0000 0000
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Branch Addressing

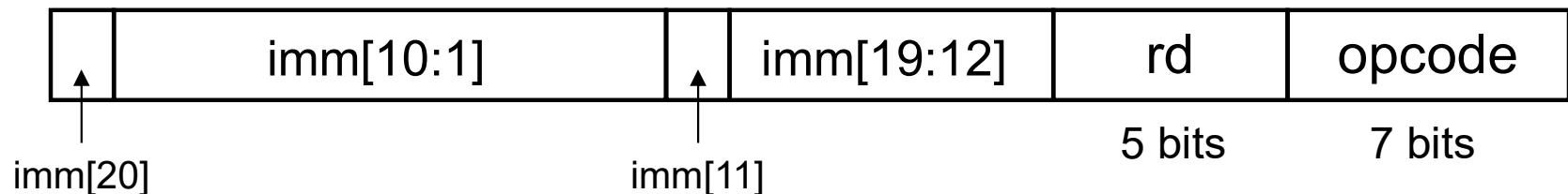
- Branch instructions specify
 - Opcode, two registers, target address
- Most branch targets are near branch
 - Forward or backward
- SB format:



- PC-relative addressing
 - Target address = $PC + \text{immediate} \times 2$

Jump Addressing

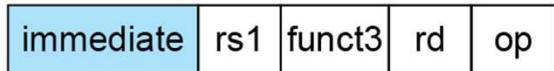
- Jump and link (jal) target uses 20-bit immediate for larger range
- UJ format:



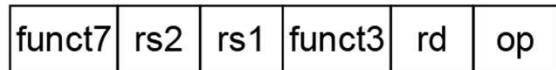
- For long jumps, eg, to 32-bit absolute address
 - lui: load address[31:12] to temp register
 - jalr: add address[11:0] and jump to target

RISC-V Addressing Summary

1. Immediate addressing

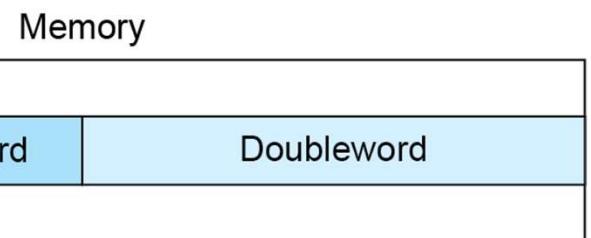
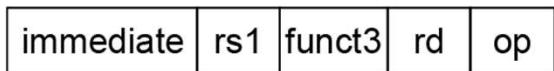


2. Register addressing

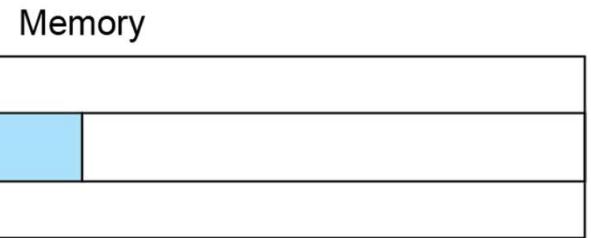
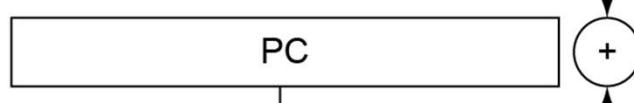
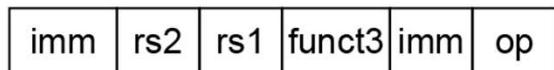


Registers
Register

3. Base addressing



4. PC-relative addressing



RISC-V Encoding Summary

Name (Field Size)	Field						Comments
	7 bits	5 bits	5 bits	3 bits	5 bits	7 bits	
R-type	funct7	rs2	rs1	funct3	rd	opcode	Arithmetic instruction format
I-type	immediate[11:0]		rs1	funct3	rd	opcode	Loads & immediate arithmetic
S-type	immed[11:5]	rs2	rs1	funct3	immed[4:0]	opcode	Stores
SB-type	immed[12,10:5]	rs2	rs1	funct3	immed[4:1,11]	opcode	Conditional branch format
UJ-type	immediate[20,10:1,11,19:12]				rd	opcode	Unconditional jump format
U-type	immediate[31:12]				rd	opcode	Upper immediate format

Synchronization

- Two processors sharing an area of memory
 - P1 writes, then P2 reads
 - Data race if P1 and P2 don't synchronize
 - Result depends of order of accesses
- Hardware support required
 - Atomic read/write memory operation
 - No other access to the location allowed between the read and write
- Could be a single instruction
 - E.g., atomic swap of register \leftrightarrow memory
 - Or an atomic pair of instructions

Synchronization in RISC-V

- Load reserved: `l r . d rd, (rs1)`
 - Load from address in rs1 to rd
 - Place reservation on memory address
- Store conditional: `sc . d rd, (rs1) , rs2`
 - Store from rs2 to address in rs1
 - Succeeds if location not changed since the `l r . d`
 - Returns 0 in rd
 - Fails if location is changed
 - Returns non-zero value in rd

Synchronization in RISC-V

- Example 1: atomic swap (to test/set lock variable)

again: `lrd x10,(x20)`

`sc.d x11,(x20),x23 // x11 = status`

`bne x11,x0,again // branch if store failed`

`addi x23,x10,0 // x23 = loaded value`

- Example 2: lock

`addi x12,x0,1 // copy locked value`

again: `lrd x10,(x20) // read lock`

`bne x10,x0,again // check if it is 0 yet`

`sc.d x11,(x20),x12 // attempt to store`

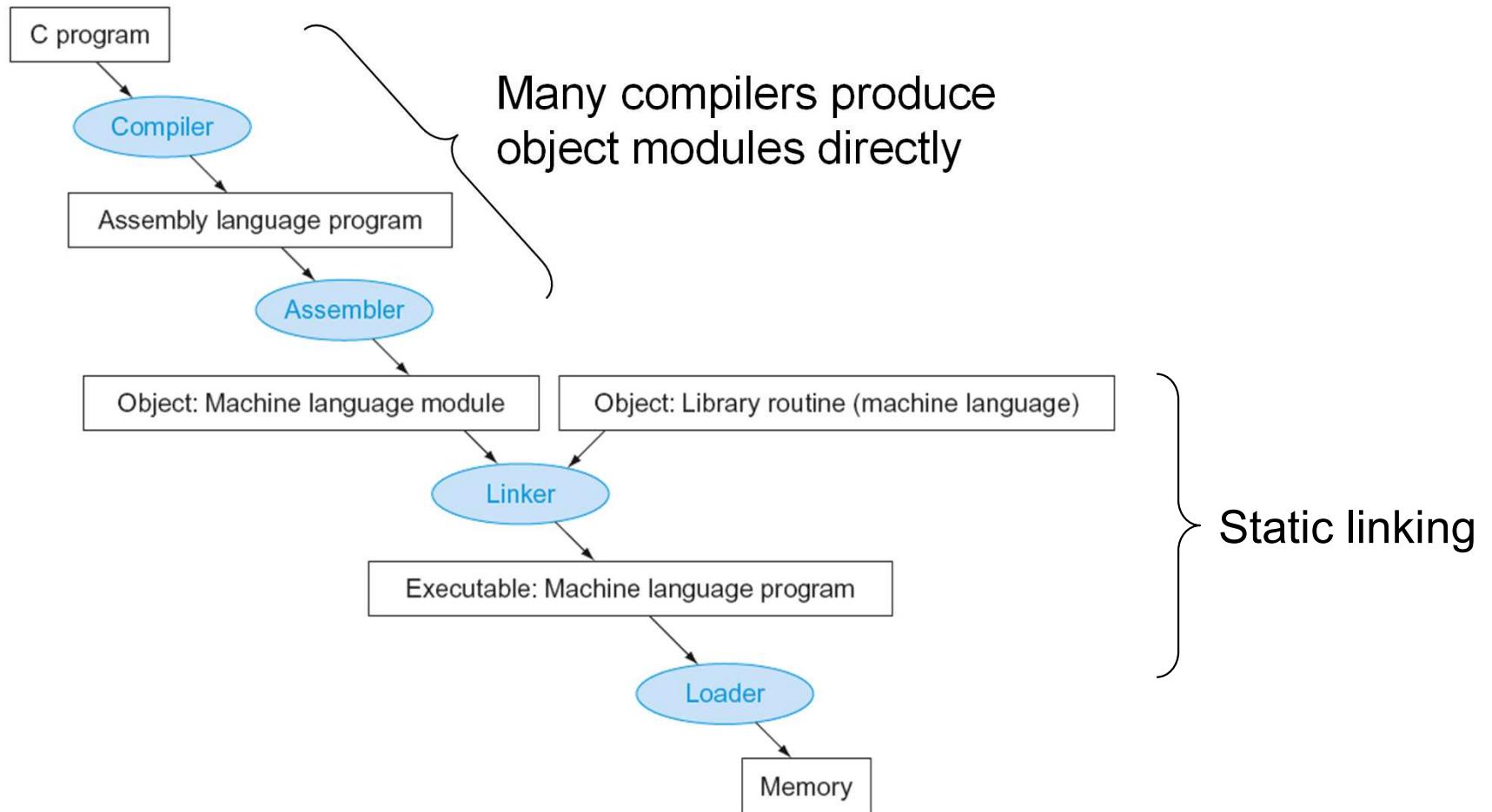
`bne x11,x0,again // branch if fails`

- Unlock:

`sd x0,0(x20) // free lock`



Translation and Startup



Producing an Object Module

- Assembler (or compiler) translates program into machine instructions
- Provides information for building a complete program from the pieces
 - Header: described contents of object module
 - Text segment: translated instructions
 - Static data segment: data allocated for the life of the program
 - Relocation info: for contents that depend on absolute location of loaded program
 - Symbol table: global definitions and external refs
 - Debug info: for associating with source code

Linking Object Modules

- Produces an executable image
 1. Merges segments
 2. Resolve labels (determine their addresses)
 3. Patch location-dependent and external refs
- Could leave location dependencies for fixing by a relocating loader
 - But with virtual memory, no need to do this
 - Program can be loaded into absolute location in virtual memory space

Loading a Program

- Load from image file on disk into memory
 1. Read header to determine segment sizes
 2. Create virtual address space
 3. Copy text and initialized data into memory
 - Or set page table entries so they can be faulted in
 4. Set up arguments on stack
 5. Initialize registers (including sp, fp, gp)
 6. Jump to startup routine
 - Copies arguments to x10, ... and calls main
 - When main returns, do exit syscall

Dynamic Linking

- Only link/load library procedure when it is called
 - Requires procedure code to be relocatable
 - Avoids image bloat caused by static linking of all (transitively) referenced libraries
 - Automatically picks up new library versions



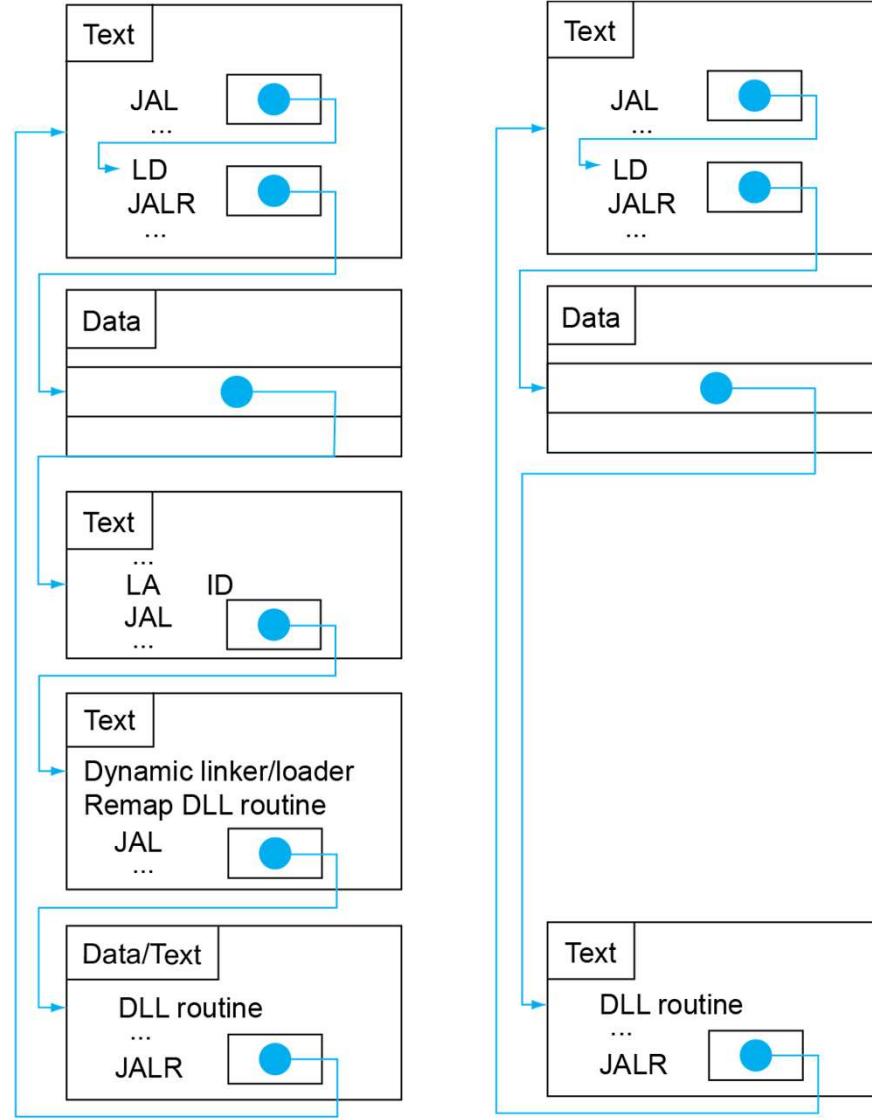
Lazy Linkage

Indirection table

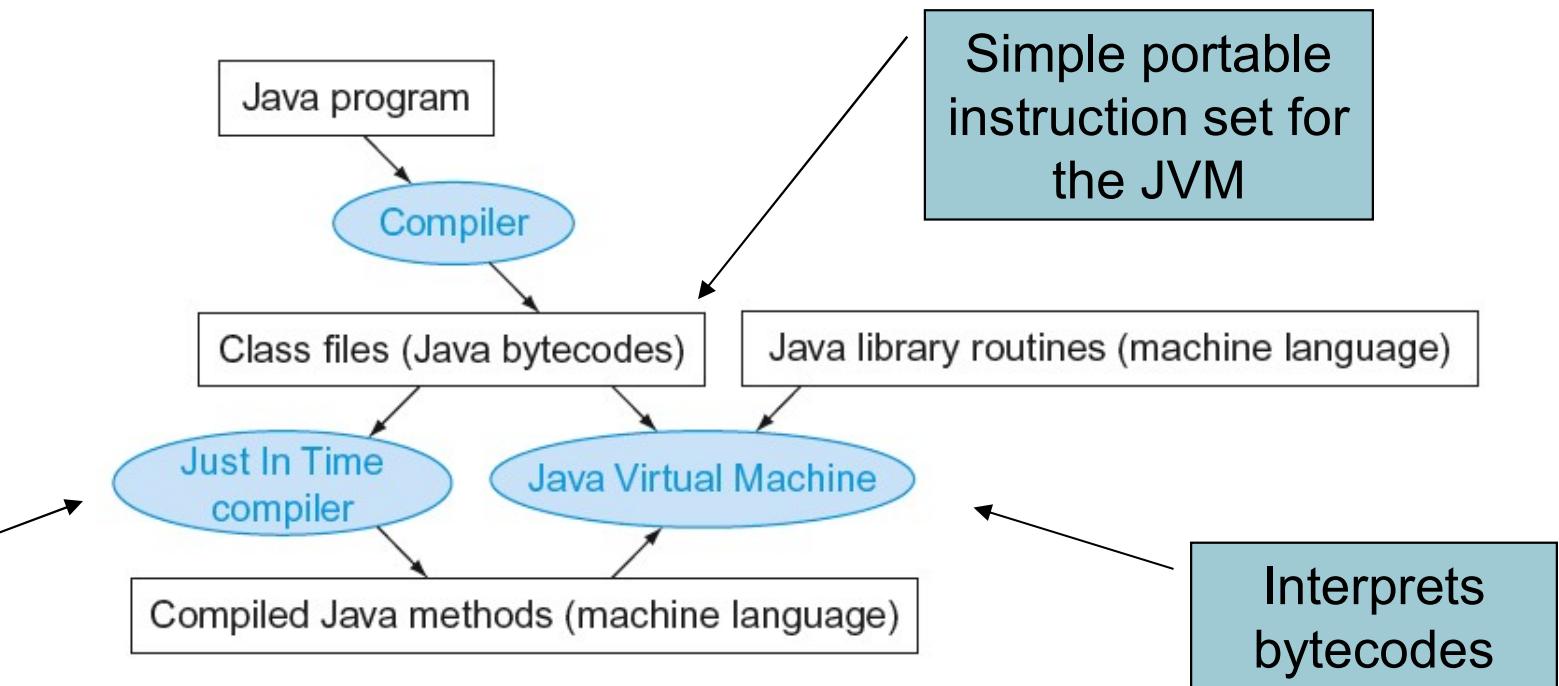
Stub: Loads routine ID,
Jump to linker/loader

Linker/loader code

Dynamically
mapped code



Starting Java Applications



Compiles bytecodes of “hot” methods into native code for host machine

Simple portable instruction set for the JVM

Interprets bytecodes

C Sort Example

- Illustrates use of assembly instructions for a C bubble sort function
- Swap procedure (leaf)

```
void swap(long long int v[],  
          long long int k)  
{  
    long long int temp;  
    temp = v[k];  
    v[k] = v[k+1];  
    v[k+1] = temp;  
}
```

- v in x10, k in x11, temp in x5

The Procedure Swap

swap:

```
slli x6,x11,3      // reg x6 = k * 8
add  x6,x10,x6    // reg x6 = v + (k * 8)
ld   x5,0(x6)      // reg x5 (temp) = v[k]
ld   x7,8(x6)      // reg x7 = v[k + 1]
sd   x7,0(x6)      // v[k] = reg x7
sd   x5,8(x6)      // v[k+1] = reg x5 (temp)
jalr x0,0(x1)      // return to calling routine
```



The Sort Procedure in C

- Non-leaf (calls swap)

```
void sort (long long int v[], size_t n)
{
    size_t i, j;
    for (i = 0; i < n; i += 1) {
        for (j = i - 1;
             j >= 0 && v[j] > v[j + 1];
             j -= 1) {
            swap(v, j);
        }
    }
}
```

- v in x10, n in x11, i in x19, j in x20

The Outer Loop

- Skeleton of outer loop:

- for (*i* = 0; *i* <*n*; *i* += 1) {

- l**i* x19,0 // *i* = 0

- for1tst:**

- bge** x19,x11,exit1 // go to exit1 if x19 ≥ x11 (*i*≥*n*)

- (body of outer for-loop)

- addi** x19,x19,1 // *i* += 1

- j** for1tst // branch to test of outer loop

- exit1:**



The Inner Loop

Skeleton of inner loop:

```
■ for (j = i - 1; j >= 0 && v[j] > v[j + 1]; j -= 1) {
    addi x20,x19,-1      // j = i - 1
for2tst:
    blt x20,x0,exit2    // go to exit2 if x20 < 0 (j < 0)
    slli x5,x20,3        // reg x5 = j * 8
    add x5,x10,x5        // reg x5 = v + (j * 8)
    ld x6,0(x5)          // reg x6 = v[j]
    ld x7,8(x5)          // reg x7 = v[j + 1]
    ble x6,x7,exit2     // go to exit2 if x6 ≤ x7
    mv x21, x10           // copy parameter x10 into x21
    mv x22, x11           // copy parameter x11 into x22
    mv x10, x21           // first swap parameter is v
    mv x11, x20           // second swap parameter is j
    jal x1,swap            // call swap
    addi x20,x20,-1       // j -= 1
    j for2tst             // branch to test of inner loop
exit2:
```



Preserving Registers

- Preserve saved registers:

```
addi sp,sp,-40 // make room on stack for 5 regs  
sd x1,32(sp) // save x1 on stack  
sd x22,24(sp) // save x22 on stack  
sd x21,16(sp) // save x21 on stack  
sd x20,8(sp) // save x20 on stack  
sd x19,0(sp) // save x19 on stack
```

- Restore saved registers:

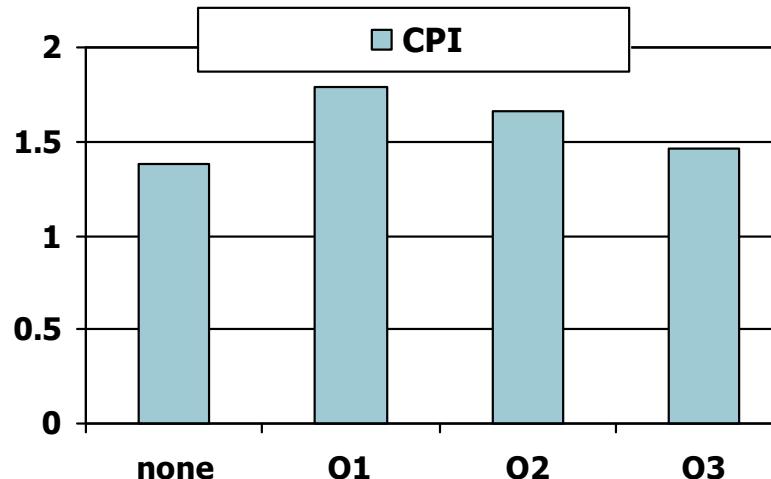
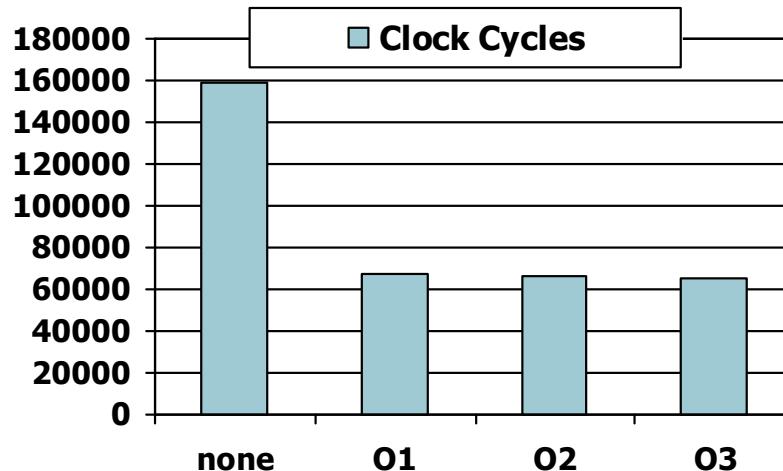
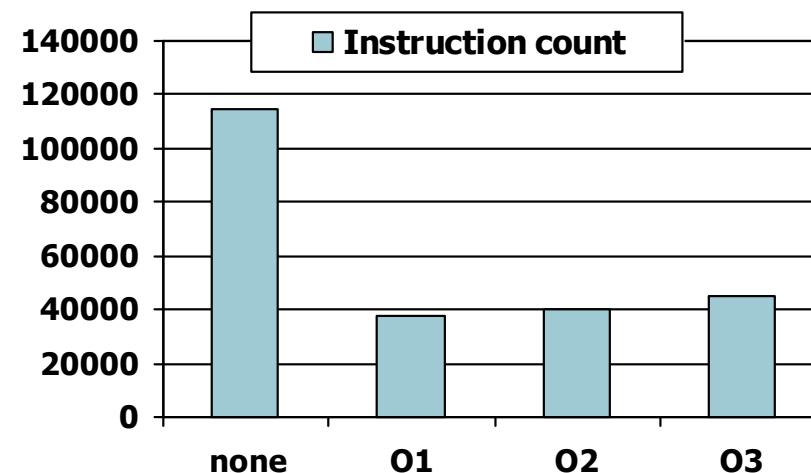
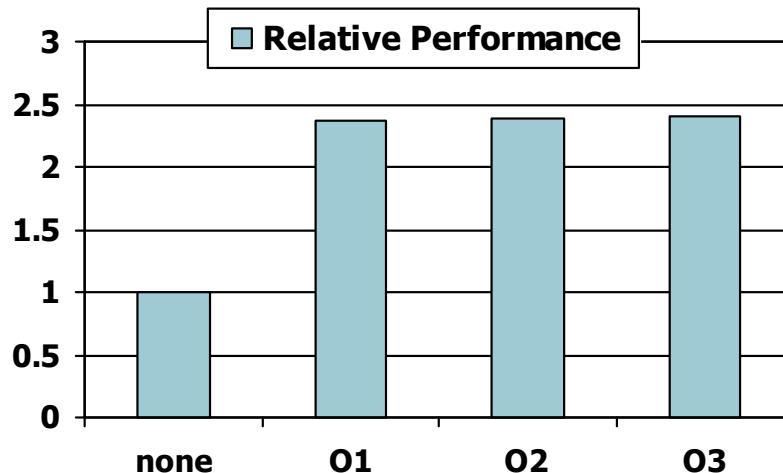
exit1:

```
sd x19,0(sp) // restore x19 from stack  
sd x20,8(sp) // restore x20 from stack  
sd x21,16(sp) // restore x21 from stack  
sd x22,24(sp) // restore x22 from stack  
sd x1,32(sp) // restore x1 from stack  
addi sp,sp, 40 // restore stack pointer  
jalr x0,0(x1)
```

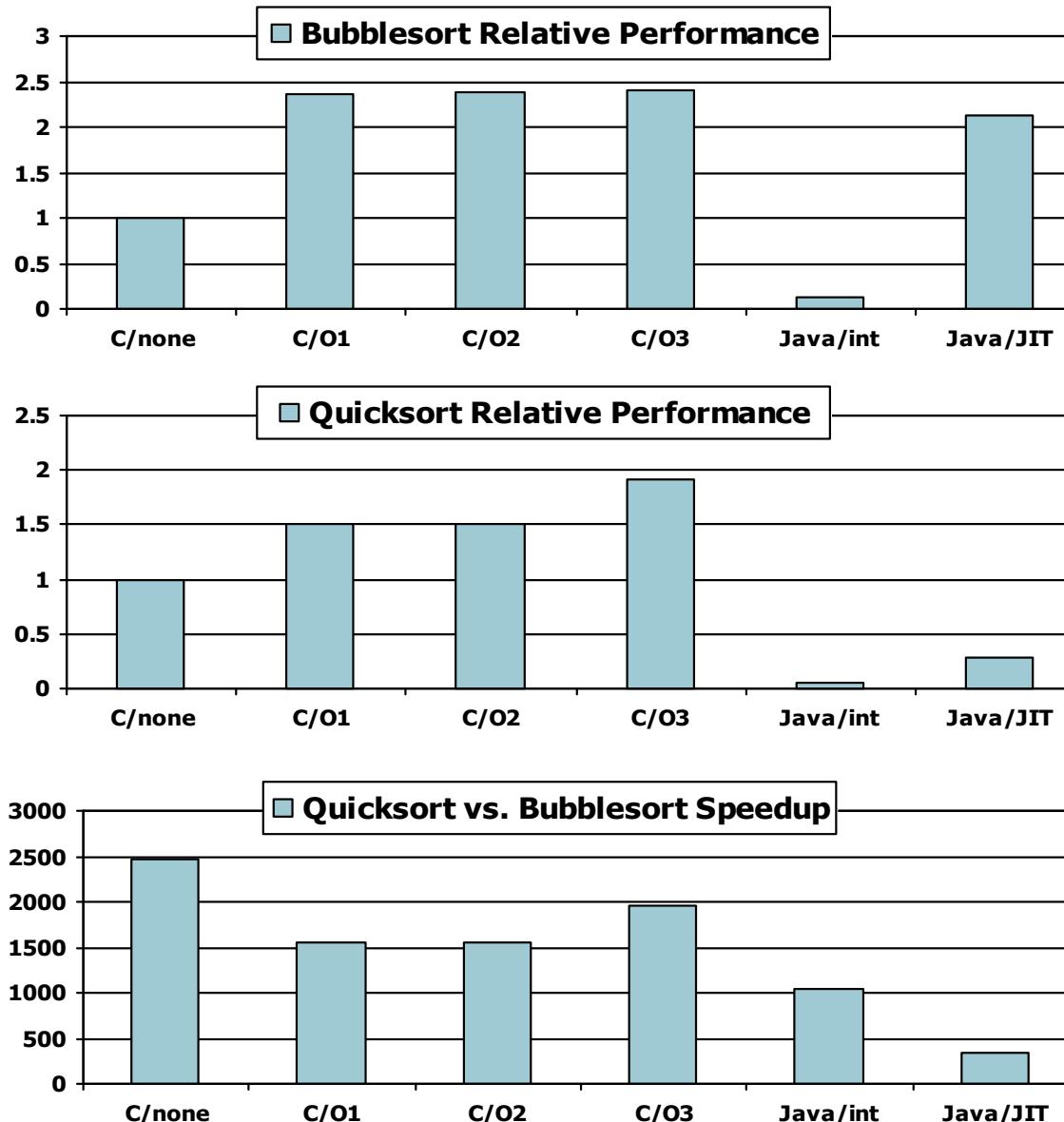


Effect of Compiler Optimization

Compiled with gcc for Pentium 4 under Linux



Effect of Language and Algorithm



Lessons Learnt

- Instruction count and CPI are not good performance indicators in isolation
- Compiler optimizations are sensitive to the algorithm
- Java/JIT compiled code is significantly faster than JVM interpreted
 - Comparable to optimized C in some cases
- Nothing can fix a dumb algorithm!

Arrays vs. Pointers

- Array indexing involves
 - Multiplying index by element size
 - Adding to array base address
- Pointers correspond directly to memory addresses
 - Can avoid indexing complexity

Example: Clearing an Array

```
clear1(int array[], int size) {  
    int i;  
    for (i = 0; i < size; i += 1)  
        array[i] = 0;  
}
```

```
li x5,0          // i = 0  
loop1:  
    slli x6,x5,3   // x6 = i * 8  
    add x7,x10,x6  // x7 = address  
                  // of array[i]  
    sd x0,0(x7)    // array[i] = 0  
    addi x5,x5,1    // i = i + 1  
    blt x5,x11,loop1 // if (i<size)  
                      // go to loop1
```

```
clear2(int *array, int size) {  
    int *p;  
    for (p = &array[0]; p < &array[size];  
         p = p + 1)  
        *p = 0;  
}
```

```
mv x5,x10          // p = address  
                // of array[0]  
    slli x6,x11,3    // x6 = size * 8  
    add x7,x10,x6    // x7 = address  
                  // of array[size]  
loop2:  
    sd x0,0(x5)      // Memory[p] = 0  
    addi x5,x5,8      // p = p + 8  
    bltu x5,x7,loop2 // if (p<&array[size])  
                      // go to loop2
```



Comparison of Array vs. Ptr

- Multiply “strength reduced” to shift
- Array version requires shift to be inside loop
 - Part of index calculation for incremented i
 - c.f. incrementing pointer
- Compiler can achieve same effect as manual use of pointers
 - Induction variable elimination
 - Better to make program clearer and safer



MIPS Instructions

- MIPS: commercial predecessor to RISC-V
- Similar basic set of instructions
 - 32-bit instructions
 - 32 general purpose registers, register 0 is always 0
 - 32 floating-point registers
 - Memory accessed only by load/store instructions
 - Consistent use of addressing modes for all data sizes
- Different conditional branches
 - For <, <=, >, >=
 - RISC-V: blt, bge, bltu, bgeu
 - MIPS: slt, sltu (set less than, result is 0 or 1)
 - Then use beq, bne to complete the branch

Instruction Encoding

Register-register

	31	25 24	20 19	15 14	12 11	7 6	0
RISC-V		funct7(7)	rs2(5)	rs1(5)	funct3(3)	rd(5)	opcode(7)
	31	26 25	21 20	16 15	11 10	6 5	0
MIPS		Op(6)	Rs1(5)	Rs2(5)	Rd(5)	Const(5)	Opx(6)

Load

	31	20 19	15 14	12 11	7 6	0
RISC-V		immediate(12)	rs1(5)	funct3(3)	rd(5)	opcode(7)
	31	26 25	21 20	16 15		0
MIPS		Op(6)	Rs1(5)	Rs2(5)	Const(16)	

Store

	31	25 24	20 19	15 14	12 11	7 6	0
RISC-V		immediate(7)	rs2(5)	rs1(5)	funct3(3)	immediate(5)	opcode(7)
	31	26 25	21 20	16 15			0
MIPS		Op(6)	Rs1(5)	Rs2(5)	Const(16)		

Branch

	31	25 24	20 19	15 14	12 11	7 6	0
RISC-V		immediate(7)	rs2(5)	rs1(5)	funct3(3)	immediate(5)	opcode(7)
	31	26 25	21 20	16 15			0
MIPS		Op(6)	Rs1(5)	Opx/Rs2(5)	Const(16)		

The Intel x86 ISA

- Evolution with backward compatibility
 - 8080 (1974): 8-bit microprocessor
 - Accumulator, plus 3 index-register pairs
 - 8086 (1978): 16-bit extension to 8080
 - Complex instruction set (CISC)
 - 8087 (1980): floating-point coprocessor
 - Adds FP instructions and register stack
 - 80286 (1982): 24-bit addresses, MMU
 - Segmented memory mapping and protection
 - 80386 (1985): 32-bit extension (now IA-32)
 - Additional addressing modes and operations
 - Paged memory mapping as well as segments

The Intel x86 ISA

■ Further evolution...

- i486 (1989): pipelined, on-chip caches and FPU
 - Compatible competitors: AMD, Cyrix, ...
- Pentium (1993): superscalar, 64-bit datapath
 - Later versions added MMX (Multi-Media eXtension) instructions
 - The infamous FDIV bug
- Pentium Pro (1995), Pentium II (1997)
 - New microarchitecture (see Colwell, *The Pentium Chronicles*)
- Pentium III (1999)
 - Added SSE (Streaming SIMD Extensions) and associated registers
- Pentium 4 (2001)
 - New microarchitecture
 - Added SSE2 instructions



The Intel x86 ISA

- And further...
 - AMD64 (2003): extended architecture to 64 bits
 - EM64T – Extended Memory 64 Technology (2004)
 - AMD64 adopted by Intel (with refinements)
 - Added SSE3 instructions
 - Intel Core (2006)
 - Added SSE4 instructions, virtual machine support
 - AMD64 (announced 2007): SSE5 instructions
 - Intel declined to follow, instead...
 - Advanced Vector Extension (announced 2008)
 - Longer SSE registers, more instructions
- If Intel didn't extend with compatibility, its competitors would!
 - Technical elegance ≠ market success



Basic x86 Registers

Name	Use
31	0
EAX	GPR 0
ECX	GPR 1
EDX	GPR 2
EBX	GPR 3
ESP	GPR 4
EBP	GPR 5
ESI	GPR 6
EDI	GPR 7
CS	Code segment pointer
SS	Stack segment pointer (top of stack)
DS	Data segment pointer 0
ES	Data segment pointer 1
FS	Data segment pointer 2
GS	Data segment pointer 3
EIP	Instruction pointer (PC)
EFLAGS	Condition codes

Basic x86 Addressing Modes

- Two operands per instruction

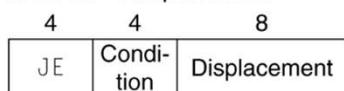
Source/dest operand	Second source operand
Register	Register
Register	Immediate
Register	Memory
Memory	Register
Memory	Immediate

- Memory addressing modes

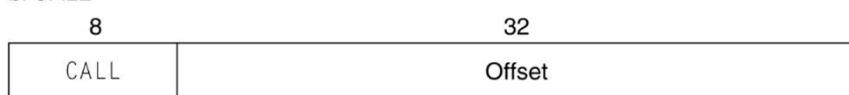
- Address in register
- Address = $R_{base} + \text{displacement}$
- Address = $R_{base} + 2^{\text{scale}} \times R_{index}$ (scale = 0, 1, 2, or 3)
- Address = $R_{base} + 2^{\text{scale}} \times R_{index} + \text{displacement}$

x86 Instruction Encoding

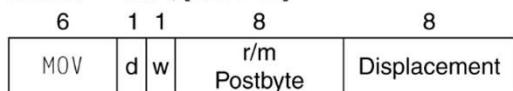
a. JE EIP + displacement



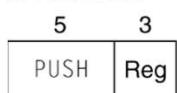
b. CALL



c. MOV EBX, [EDI + 45]



d. PUSH ESI



e. ADD EAX, #6765



f. TEST EDX, #42



Variable length encoding

- Postfix bytes specify addressing mode
- Prefix bytes modify operation
 - Operand length, repetition, locking, ...

Implementing IA-32

- Complex instruction set makes implementation difficult
 - Hardware translates instructions to simpler microoperations
 - Simple instructions: 1–1
 - Complex instructions: 1–many
 - Microengine similar to RISC
 - Market share makes this economically viable
- Comparable performance to RISC
 - Compilers avoid complex instructions

Other RISC-V Instructions

- Base integer instructions (RV64I)
 - Those previously described, plus
 - auipc rd, immed // $rd = (\text{imm} \ll 12) + pc$
 - follow by jalr (adds 12-bit immed) for long jump
 - slt, sltu, slti, sltui: set less than (like MIPS)
 - addw, subw, addiw: 32-bit add/sub
 - sllw, srlw, srlw, slliw, srliw, sraiw: 32-bit shift
- 32-bit variant: RV32I
 - registers are 32-bits wide, 32-bit operations

Instruction Set Extensions

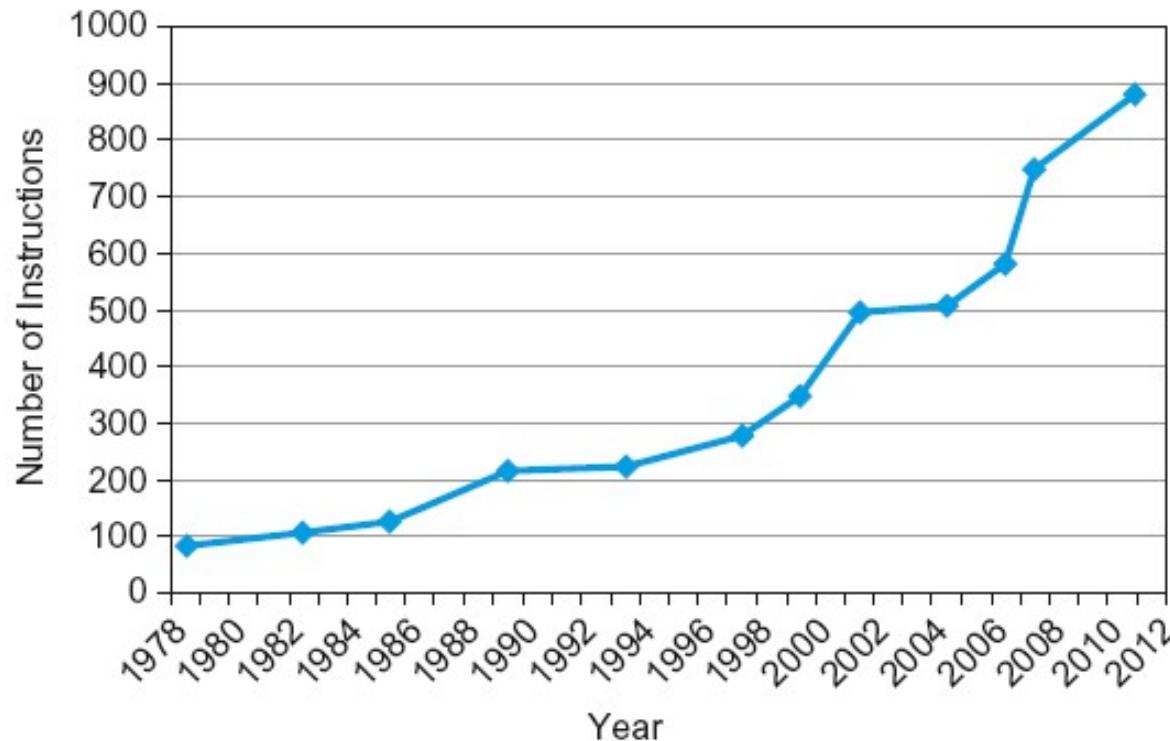
- M: integer multiply, divide, remainder
- A: atomic memory operations
- F: single-precision floating point
- D: double-precision floating point
- C: compressed instructions
 - 16-bit encoding for frequently used instructions

Fallacies

- Powerful instruction \Rightarrow higher performance
 - Fewer instructions required
 - But complex instructions are hard to implement
 - May slow down all instructions, including simple ones
 - Compilers are good at making fast code from simple instructions
- Use assembly code for high performance
 - But modern compilers are better at dealing with modern processors
 - More lines of code \Rightarrow more errors and less productivity

Fallacies

- Backward compatibility \Rightarrow instruction set doesn't change
 - But they do accrete more instructions



x86 instruction set

Pitfalls

- Sequential words are not at sequential addresses
 - Increment by 4, not by 1!
- Keeping a pointer to an automatic variable after procedure returns
 - e.g., passing pointer back via an argument
 - Pointer becomes invalid when stack popped

Concluding Remarks

- Design principles
 - 1. Simplicity favors regularity
 - 2. Smaller is faster
 - 3. Good design demands good compromises
- Make the common case fast
- Layers of software/hardware
 - Compiler, assembler, hardware
- RISC-V: typical of RISC ISAs
 - c.f. x86