

below a critical

Hodgkin-Huxley model (neuron's capacitance C set to 1)

$$C\frac{dv}{dt} = -\sum_{k} I_{k} + I$$



 $\sum_{k} I_{k} = g_{Na} m^{3} h(v - E_{Na}) + g_{K} n^{4} (v - E_{K}) + g_{L} (v - E_{L})$ $\begin{aligned} g_{K} &= 36 & \frac{dm}{dt} = \Omega_{m}(v)(1-m) - \beta_{m}(v)m & \Omega_{m} &= (2.5 - 0.1v)/(e^{(2.5 - 0.1v)} - 1) \\ g_{Na} &= 120 & \Omega_{n} &= (0.1 - 0.01v)/(e^{(1-0.1v)} - 1) \\ g_{L} &= 0.3 & \frac{dn}{dt} = \Omega_{n}(v)(1-n) - \beta_{n}(v)n & \Omega_{h} &= 0.07e^{-v/20} \\ E_{L} &= 10.6 & \frac{dh}{dt} = \Omega_{h}(v)(1-h) - \beta_{h}^{i}(v)h & \beta_{h} &= 1/(e^{(3 - 0.1v)} + 1) \end{aligned}$

$$\frac{dn}{dt} = 0.3 \quad \frac{dn}{dt} = \Omega_n(v)(1-n) - \beta_n(v)n \quad \Omega_n = 0.07e^{-v/2}$$

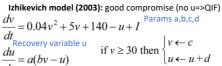
$$E_K = -12$$
 a
 $E_L = 10.6$ $\frac{dh}{dt} = \Omega_h(v)(1 - h) - \beta_h(v)$

$$\frac{(1v)}{(e^{(2.5-0.1v)}-1)}$$
 when the spike and $\frac{(1v)}{(20)}$ if $v \ge 0$

$$\frac{8L - 0.5}{E_K} = \frac{12}{dt} \frac{1}{dt} = \frac{1}{4} \frac{1}$$

$$\beta_m^{i} = 4e^{-\nu/18} \quad \beta_n^{i} = 0.125e^{-\nu/18}$$
$$\beta_h^{i} = 1/(e^{(3-0.1\nu)} + 1)$$

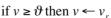
 $\tau \frac{dv}{dt} = a(v_r - v)(v_c - v) + RI$ In the absence of But if it is above v_c Subthreshold profile dendritic current, it increases of v modeled more v decays to the quickly until the accurately in QIF > LIF resting potential neuron fires v., as long as it is



LIF model where v_r is the resting potential, I is the dendritic current, and τ and R are constants. (We'll use τ = 5, R = 1, and v_r = -65mV) $\frac{dv}{r} = v_r - v + RI$ comp. inexpensive but bio inaccurate limited repertoire of sign. behaviours limited repertoire of sign. behaviours

The sub-threshold (before spiking) dynamics of the sodium and potassium currents are approximated by the v_r -v term

The detailed dynamics of the spike itself are ignored. Instead, when the membrane potential reaches a threshold, we record a spike and explicitly reset the neuron

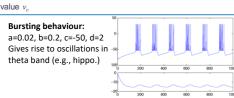


A good value for the threshold 9 is -50r. An instantaneous value to represent the actual spike (a Dirac pulse) can be inserted immediately before the neuron is reset

Adding absolute refractory period alpha: We simply adjust the conditions under which a spike occurs to take account of the time since the last spike

Let t_{snike} be the time of the most recent spike. Then we have

if
$$v \ge \vartheta$$
 and $t - t_{spike} > \alpha$ then
$$\begin{cases} v \leftarrow v_r \\ t_{vnike} \leftarrow t \end{cases}$$



Given y(t), we can approximate the value of $y(t+\delta t)$

Euler: sensitive to dt: small dt better but comp, expensive $y(t + \delta t) = y(t) + \delta t f(y(t))$

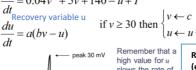
(4th order) Runge-Kutta: not as sensitive + comp. efficient

$$y(t + \delta t) = y(t) + \frac{1}{6}\delta t(k_1 + 2k_2 + 2k_3 + k_4)$$

$$k_1 = f(y(t)) \qquad k_3 = f\left(y(t) + \frac{1}{2}\delta t k_2\right)$$

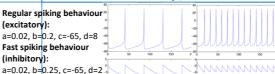
$$k_2 = f\left(y(t) + \frac{1}{2}\delta t k_1\right) k_4 = f(y(t) + \delta t k_3)$$

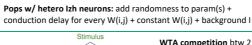
WTA competition to 2 diff. competing stimuli: strongest stimulus wins most of the time. For WTA to work, needs a lot of neurons (800ext, 200inh): else, irregular swapping.

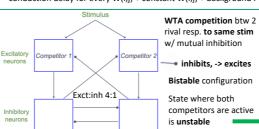


slows the rate of increase of ν , and makes it harder fo the neuron to fire

Gecay with rate a Need to wait until u is down to fire



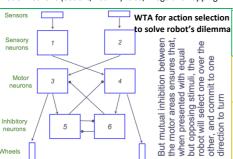




Braitenberg 0000 0000 0000 0000 W/ same stim. 50% of time LHS/RHS dominates. Winner based on tiny rand.diff. in firing pattern at the start of trial period. All acty dies w/o stim

Connectome = description of the brain as a network

(structural or functional). Structural connectome =



For WTA to work, need \mathbf{very} high \mathbf{a} (Izh neurons), i.e., fast dynamics in recovery variable u: else, periodic firing where the 2 competitors are 180deg out of phase w/ each other

A network (or graph) $G = \langle V, E \rangle$ comprises a set V of nodes (or The relation E can also be expressed as a two-dimensional

vertices) and a set $E \subseteq V \times V$ of edges (or arcs or connections) A(i,j) =connectivity matrix A, such that, for all $i, j \in V$

A random (or Erdős-Rényi) network is one in which, for every pair of nodes i and j, dist of randomly

small-world modular network. $\int L(i,j)$ if $(i,j) \in E$ 0 otherwise

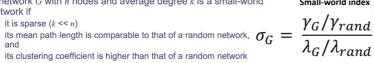
A(j,j)=0

(undirected net.) Degree of node k_i: num of edges it has 2m Network's average k: k = 2n nodes, m edges

 $P(A_{ij} = x) = \begin{cases} p & \text{if } x = 1 \\ 1 - p & \text{if } x = 0 \end{cases}$ assigned

A network G with n nodes and average degree k is a small-world network if

Small-world index



efficiency metrics

Small-world network metrics. where network G = < V. E > path length between nodes i,j: num of edges in shortest path btw nodes i,j **network's mean path length** λ_G : path length avg. over all distinct pairs of nodes clustering coefficient of node i: fraction of the set of all possible edges btw immediate neighbours of i that are actual edges (for leaf node, this is set to 1). **network's clustering coefficient** γ_G : clustering coeff avg. over all nodes

• it is sparse (k << n)

Wheels

• its clustering coefficient is higher than that of a random network

It can be shown that the mean path length λ_{rand} of a random network with n nodes and average degree k is (on average) $\ln(n)/\ln(k)$ and its clustering coefficient γ_{rand} is (on average) k/n

Alternative:

Watts-Strogatz Procedure to build small-world network 1) build ring lattice 2) rewire some connections

A ring lattice with degree k is a set of nodes notionally arranged in a circle, where each node is connected to all its (spatial) neighbours that are less than or equal to k/2 nodes away

Each edge is considered in turn, and with probability pit is rewired

Rewiring an edge (j,i) means deleting (j,i) from E and adding (h,i) for some randomly chosen h



Efficiency between neighbours i and j

Let $\mathit{Eff(i,j)}$ be $1/\lambda$, where λ is the path length in G from node i to node j. This captures the efficiency with which information can be propagated from i to j, the maximum being 1 if i and j are

Global efficiency

Q has a maximum value of 1

The global efficiency of G is then the efficiency averaged over the whole network, defined as

$$Eff_{glob}(G) = \frac{1}{n(n-1)} \sum_{i \neq j} Eff(i,j)$$

Local efficiency of a neighbourhood of a node

of neighbours of l, and $E' \subseteq E$ is the edges that join nodes in l. The efficiency of the neighbourhood of node l is then $Eff_{loo}(E)$. The local efficiency of G is then the neighbourhood efficiency averaged over the whole network, defined as $Eff_{loc}(G) = \frac{1}{n} \sum_{l \in G} Eff_{glob}(G_l)$

$$Eff_{loc}(G) = \frac{1}{n} \sum_{i \in G} Eff_{glob}(G_i)$$

Small-world index peaks when balance between global and local efficiencies

Modular network

One example is modularity. A network is modular if its nodes can be partitioned into subsets (modules, or communities) that

- are highly intra-connected (there are many internal connections within the module)
- but sparsely inter-connected (there are relatively fewer connections between modules)

Modular networks are typically (but not necessarily) also small-

Q = (actual fraction of edges that are within communities) -

n general, suppose the nodes in V are partitioned into belongs. Then we have

$$Q = \frac{1}{2m} \sum_{i,j} \left(A_{ij} - \frac{k_i k_j}{2m} \right) \delta_{c_i c_j}$$

Modularity measure Q

A measure \mathcal{Q} of how modular a given partitioning is for \mathcal{G} is obtained by comparing the number of intra-community edges that actually occur in G with the number that would occur in a comparable network that was randomly connected

(expected fraction of edges that are within communities)

communities, and let c, denote the community to which node

Here's an algorithm for generating a modular network with n nodes, m edges, and C communities. Like the it has two steps First, a set of *C* disconnected communitie is created with n/C nodes Each community has m/C random edges

The second step is a rewiring process, like the Watts-Strogatz procedure again

again

Each existing (intracommunity) edge is
considered, and with
probability p is rewired as
an edge between communities

For rewired edges, the target community and target node within that community are randomly chosen

Note that Q can be negative, if there are fewer connect within a "community" than there are between "community" -Q for high p bc more co. between communities than within community

ne value Q is defined with respect to a given partitioning int

is equal to the maximum obtainable value of Q

Where the 2 fractions are given as (SUM OF):
The actual fraction of intra-community edges is the sum of

2m

In an undirected network, each edge counts twice

for all i and j in the same community, while the expected fraction of intra-community edges is the sum of

$$\frac{k_i}{2m} \times \frac{k_j}{2m}$$

The chances of an edge connecting with *i* and with *j*

for all i and j in the same community, where k_i is the degree of

