SCT211-0510/2021

Chapter 4: Number Theory and Cryptography

4.1 Divisibility and Modular Arithmetic

Divisibility:

- An integer a divides b (denoted a | b) if there exists an integer c such that b = a ×
 c.
- Example: 4 | 20 because 20 = 4 × 5.
- Divisibility has the following properties:
 - o Reflexive: Every number divides itself.
 - \circ Transitive: If $a \mid b$ and $b \mid c$, then $a \mid c$.
 - \circ Multiplicative: If $a \mid b$ and $a \mid c$, then $a \mid (b + c)$ and $a \mid (b c)$.

The Division Algorithm:

- For any integers a and d > 0, there exist unique integers q and r such that:
 - a = dq + r, where $0 \le r < d$.
- This forms the basis for Euclidean division and leads to GCD algorithms.

Modular Arithmetic:

- Two integers a and b are congruent modulo m if $m \mid (a b)$, written as $a \equiv b \pmod{m}$.
- Arithmetic operations in modular systems preserve many useful properties:
 - O Addition: $(a + b) \mod m = [(a \mod m) + (b \mod m)] \mod m$
 - \circ Multiplication: (a × b) mod m = [(a mod m) × (b mod m)] mod m
 - Exponentiation: a^k mod m, with optimizations via modular exponentiation.
- Modular arithmetic forms the basis for many algorithms in computer science and cryptography.

Applications:

- Used in digital clocks, hashing, data encryption, calendars, and cyclic structures.
- Example: Digital clocks use mod 12 or mod 24 arithmetic. For instance, 17:00 + 10 hours = (17 + 10) mod 24 = 3:00.

Example:

- Compute 7^4 mod 10:
 - o 7^4 = 2401, and 2401 mod 10 = 1

Additional Example:

- Determine whether 81 ≡ 4 mod 7:
 - \circ 81 4 = 77, and 77 is divisible by 7, so the congruence holds.

4.2 Integer Representations and Algorithms

Number Representation in Different Bases:

- Binary (base 2): Primary system used in computing. Only uses digits 0 and 1.
- Octal (base 8) and Hexadecimal (base 16): Used to compactly represent binary numbers.
- **Decimal (base 10)** is the standard system for human-readable numbers.

Base Conversion Algorithms:

- From decimal to base *b*: repeatedly divide the number by *b* and record remainders.
- From base *b* to decimal: multiply each digit by b raised to the power of its position.

Example: Convert 43 to binary:

- 43 ÷ 2 = 21 remainder 1
- 21 ÷ 2 = 10 remainder 1
- 10 ÷ 2 = 5 remainder 0
- 5 ÷ 2 = 2 remainder 1
- 2 ÷ 2 = 1 remainder 0
- 1 ÷ 2 = 0 remainder 1 → Binary = 101011

Efficient Arithmetic in Binary:

• Computers perform binary addition, subtraction, multiplication and division using simple logic circuits.

Modular Exponentiation:

- A crucial technique to compute a^b mod n efficiently.
- Used in RSA, Diffie-Hellman key exchange, digital signatures.

Square-and-Multiply Algorithm:

 Converts exponent to binary and uses squaring/multiplication only when needed.

Example: Compute 3^13 mod 17 using binary exponentiation.

4.3 Primes and Greatest Common Divisors

Prime Numbers:

- A number p > 1 is prime if its only positive divisors are 1 and p.
- Infinite in number, as proven by Euclid.
- Importance in encryption: RSA and other cryptographic systems rely on prime numbers.

Fundamental Theorem of Arithmetic:

• Every integer greater than 1 is either prime or can be expressed uniquely as a product of prime numbers (up to the order of the factors).

Composite Numbers:

- Have more than two positive divisors.
- Example: $12 = 2 \times 2 \times 3$

GCD and LCM:

- The greatest common divisor (GCD) of a and b is the largest number that divides both.
- The least common multiple (LCM) of a and b is the smallest number divisible by both.
- Relationship: gcd(a, b) × lcm(a, b) = a × b

Euclidean Algorithm:

- A classic and efficient algorithm to compute GCD:
 - \circ a = bq + r
 - \circ Repeat with (b, r) until r = 0

Extended Euclidean Algorithm:

- Produces integers x and y such that ax + by = gcd(a, b)
- Critical in finding modular inverses used in RSA.

Example: Find GCD of 99 and 78:

- 99 = 78×1 + 21
- 78 = 21×3 + 15
- 21 = 15×1 + 6
- $15 = 6 \times 2 + 3$
- $6 = 3 \times 2 + 0 \rightarrow GCD = 3$

4.4 Solving Congruences

Linear Congruences:

- General form: $ax \equiv b \pmod{m}$
- Has solution if and only if gcd(a, m) divides b.

Solution Steps:

- 1. Simplify if possible.
- 2. Use the extended Euclidean algorithm to find modular inverse of a.
- 3. Multiply both sides by the inverse modulo m.
- 4. General solution: $x = x_0 + k(m/d)$ for integer k, where d = gcd(a, m)

Chinese Remainder Theorem (CRT):

- If m₁, m₂, ..., m_n are pairwise coprime, and we have:
 - \circ $x \equiv a_1 \pmod{m_1}$
 - o $x \equiv a_2 \pmod{m_2}$
 - o ...
 - \circ $x \equiv a_n \pmod{m_n}$
- Then there exists a unique solution modulo $M = m_1 \times m_2 \times ... \times m_n$
- Applied in solving large systems efficiently.

Fermat's Little Theorem:

- If p is prime and a is not divisible by p, then:
 - o a^{p-1} ≡ 1 mod p
- Used in primality testing, finding inverses in RSA.

Example: Find 3^6 mod 7

• Since 7 is prime, $3^6 \equiv 1 \mod 7$ by Fermat's theorem.

4.5 Applications of Congruences

Hashing:

- Used in data structures (hash tables), authentication systems, and load balancing.
- Example: key mod table size to find storage index.

Check Digits and Error Detection:

- Used in ISBN, credit cards, UPC codes.
- Example: In ISBN-10:
 - o Compute: $(1 \times d_1 + 2 \times d_2 + ... + 10 \times d_{10}) \mod 11 = 0$

Pseudorandom Number Generation:

- Uses linear congruential generators (LCGs):
 - \circ x_{n+1} = (a x_n + c) mod m
- If parameters are chosen properly, can produce long sequences with good randomness properties.

Calendar Algorithms:

- Use modular arithmetic to compute day of week, leap years.
- Example: Zeller's Congruence computes the day of the week for any date.

4.6 Cryptography

Classical Ciphers:

- Caesar Cipher: shift characters using modular arithmetic.
 - \circ E_k(p) = (p + k) mod 26
 - o D_k(c) = (c k) mod 26

Symmetric Key Cryptography:

- Same secret key used by both sender and receiver.
- Fast but key distribution is a challenge.
- Examples: DES, AES.

Public Key Cryptography (RSA):

- Uses two keys: public key for encryption and private key for decryption.
- Based on the difficulty of factoring large semiprimes.

RSA Algorithm:

- 1. Choose large primes p and q
- 2. Compute n = pq
- 3. Compute $\phi(n) = (p-1)(q-1)$
- 4. Choose e such that $gcd(e, \phi(n)) = 1$
- 5. Compute d such that ed = $1 \mod \phi(n)$
- 6. Public key: (e, n), Private key: (d, n)
- 7. Encrypt: $c = m^e \mod n$
- 8. Decrypt: $m = c^d \mod n$

Digital Signatures:

- Verify sender identity and message integrity
- Sign: hash the message and encrypt with private key
- Verify: decrypt signature with public key and compare hashes

Homomorphic Encryption:

- Allows operations on encrypted data
- Example: Add two encrypted salaries without decrypting either