COMSATS UNIVERSITY ISLAMABAD

ATTOCK CAMPUS



Name:

Fahad Saeed

Registration No:

Sp22-BCS-005

Instructor:

Sir Bilal Bukhari

MINI-COMPILER Introduction and Overview

Project Overview This mini compiler is a complete implementation of a programming language compiler written in C#. It demonstrates all major phases of compilation from

compiler written in C#. It demonstrates all major phases of compilation from source code to executable instructions, providing a practical example of compiler construction principles.

Supported Language Features

- Data Types: int, float, bool
- Variables: Declaration and assignment
- Expressions: Arithmetic, logical, and comparison operations
 - Control Flow: if/else statements, while loops Built-in Functions: print statement for output
- Operators:

o Arithmetic: +, -, *, /, % o

Comparison: ==, !=, <, >, <=, >= o

Logical: &&, | |,!

Target Platform

The compiler generates code for a custom stack-based virtual machine (VM), making it platform-independent while maintaining educational clarity.

Educational Value

This implementation serves as a learning tool for understanding:

- Compiler construction phases
- Abstract Syntax Tree (AST) generation
- Symbol table management
- Intermediate code generation

Virtual machine execution

Compiler Architecture and Design

Seven-Phase Architecture

The compiler follows the traditional compilation pipeline with seven distinct phases:

Phase 1: Lexical Analysis (Tokenization)

- Component: Lexer class
- Purpose: Converts source code into tokens
- Output: List of tokens with type, value, and position information
 - Features: Handles keywords, identifiers, literals, operators, and comments

Phase 2: Syntax Analysis (Parsing)

- Component: Parser class
- Purpose: Builds Abstract Syntax Tree from tokens
- Algorithm: Recursive descent parser
 - Grammar: Supports operator precedence and associativity

Phase 3: Semantic Analysis

- Component: SemanticAnalyzer class
- Purpose: Type checking and symbol table management
- Features: Scope management, type compatibility verification, initialization checking

Phase 4: Intermediate Code Generation

- Component: IRGenerator class
- Purpose: Generates platform-independent intermediate representation
- Format: Three-address code style instructions

Phase 5: Optimization

- Component: Optimizer class
- Purpose: Code improvement (basic implementation)
- Note: Currently minimal, designed for future enhancement

Phase 6: Code Generation

Component: CodeGenerator class

Purpose: Translates IR to target machine code

Target: Stack-based VM instructions

Phase 7: Execution

Component: StackVM class

Purpose: Executes generated code

Architecture: Stack-based virtual machine with memory management

Design Patterns Used

• Visitor Pattern: For AST traversal in semantic analysis

* Strategy Pattern: For different optimization techniques

• Interpreter Pattern: For VM instruction execution

Language Specification and Grammar

Lexical Specification Keywords

int, float, bool, void, if, else, while, for, function, return, print, true, false

Operators and Delimiters

Arithmetic: + - * / %

Assignment: =

Comparison: == != < > <= >=

Logical: && | |!

Delimiters:;,(){}.

Literals

Integer: Sequences of digits (e.g., 42, 0, 123)

Float: Numbers with decimal points (e.g., 3.14, 0.5)

Boolean: true or false

```
(BNF-style) bnf program → statement*
statement
              → varDecl | assignment | ifStmt | whileStmt
         | printStmt | block | ";"
varDecl
             → type IDENTIFIER ("=" expression)? ";" assignment
→ IDENTIFIER "=" expression ";" ifStmt → "if" "(" expression ")"
statement ("else" statement)?
whileStmt → "while" "(" expression ")"
statement printStmt → "print" expression ";"
           \rightarrow "{" statement* "}"
block
expression → logicalOr logicalOr → logicalAnd (" | | "
logicalAnd)* logicalAnd → equality ("&&" equality)*
equality \rightarrow comparison (("==" | "!=") comparison)*
comparison \rightarrow term ((">" | ">=" | "<" | "<=") term)* term
\rightarrow factor (("+" | "-") factor)* factor \rightarrow unary (("*" | "/"
| "%") unary)* unary \rightarrow ("!" | "-") unary | primary
              → NUMBER | FLOAT | BOOLEAN | IDENTIFIER
primary
         | "(" expression ")"
type \rightarrow "int" | "float" | "bool"
Operator Precedence (Highest to Lowest)
   1. Unary operators (!, -)
   2. Multiplicative (*, /, %)
   3. Additive (+, -)
```

String: Quoted text with escape sequences Grammar Specification

```
4. Relational (<, >, <=, >=)5. Equality (==, !=)
```

- 6. Logical AND (&&)
- 7. Logical OR (| |)

Type System

- Static Typing: All variables must be declared with explicit types
- Type Compatibility: Automatic promotion from int to float in mixed expressions
- Type Checking: Compile-time verification of type consistency

Implementation Details

```
Key Data Structures Token Structure
csharp public class Token { public
TokenType Type { get; set; }
public string Value { get; set; }
public int Line { get; set; } public
int Column { get; set; }
}
```

AST Node Hierarchy

}

- Base Class: ASTNode with type and position information
- Expression Nodes: BinaryOpNode, UnaryOpNode, NumberNode, etc.
- Statement Nodes: VarDeclarationNode, AssignmentNode, IfNode,
 etc. Symbol Table csharp public class Symbol { public string Name { get;
 set; } public DataType Type { get; set; } public int Scope { get; set; }
 public int Address { get; set; } public bool IsInitialized { get; set; }

Error Handling Strategy

Custom Exception: Compiler Exception with line/column information

Error Recovery: Continues parsing after errors when possible

Semantic Errors: Detailed type mismatch and undeclared variable messages

Memory Management

- Stack-based VM: Uses evaluation stack for expression computation
- Variable Storage: Array-based memory with address allocation
- Scope Management: Hierarchical symbol table with scope stack

Intermediate Representation

The IR uses a simple instruction set:

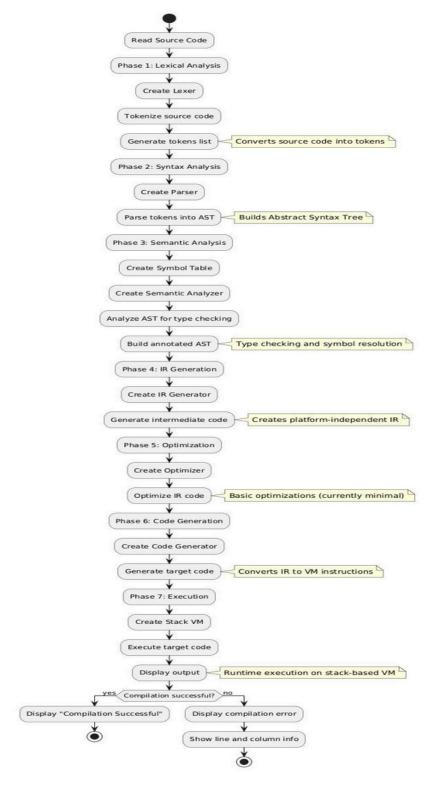
- Data Movement: LOAD_CONST, LOAD_VAR, STORE_VAR
- Arithmetic: ADD, SUB, MUL, DIV, MOD
- Comparison: CMP_EQ, CMP_NEQ, CMP_LT, etc.
- Control Flow: JUMP, JUMP_IF_FALSE, LABEL

I/O: PRINT

Virtual Machine Architecture

- Stack-based Execution: Operands pushed/popped from evaluation stack
- Memory Model: Linear array for variable storage
- Instruction Pointer: Sequential execution with jump capabilities
- Type Handling: Dynamic type conversion during arithmetic operations

Diagrams Activity Diagram:



Sequence Diagram:

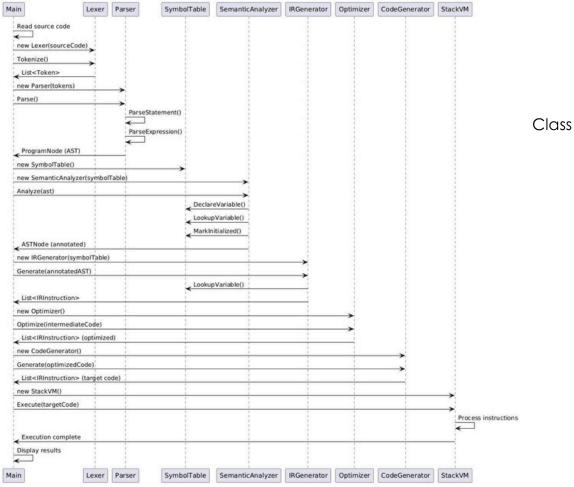
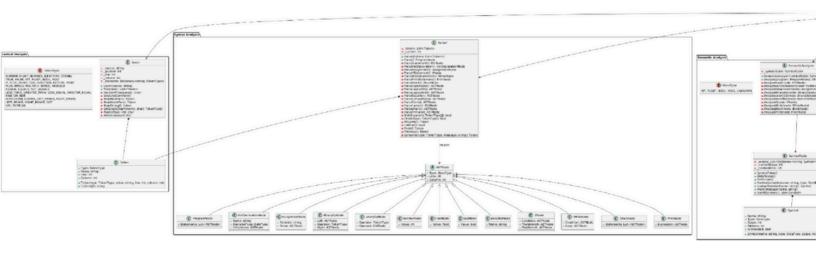


Diagram:

GitHub Repository



https://github.com/UMAR-CUI/Mini-Compiler-SP22.git