### COMS31700 Design Verification

# Hardware Design Languages

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### Hardware Design Languages

- Hardware Design Languages were built with simulation in mind
  - Synthesis and other back-end purposes were added at a later stage
- Most popular languages today (both are IEEE standards)
  - VHDL
  - Verilog
- VHDL:
  - Committee-designed language contracted by U.S. (DoD) (ADA-derived)
  - Functional/logic modeling and simulation language
  - Main differentiator from Verilog is types (e.g. records)
- Verilog:
  - Logic modeling and simulation language
  - Started in EDA industry in the 80's now owned by Cadence
  - Donated to IEEE as a general industry standard
  - SystemVerilog (the next generation of Verilog) is designed to improve abstraction of Verilog
    - Abstraction levels
    - Data types
    - Verification constructs
- Verilog vs. VHDL: personal preferences, EDA tool availability, commercial, business and marketing issues.

### Modeling Levels – Major Dimensions

#### Temporal Dimension:

- continuous (analog)
- gate delay
- clock cycle
- instruction cycle
- events



#### Data Abstraction:

- continuous (analog)
- bit : multiple values
- bit : binary
- abstract value
- composite value ("struct")

# discrete value

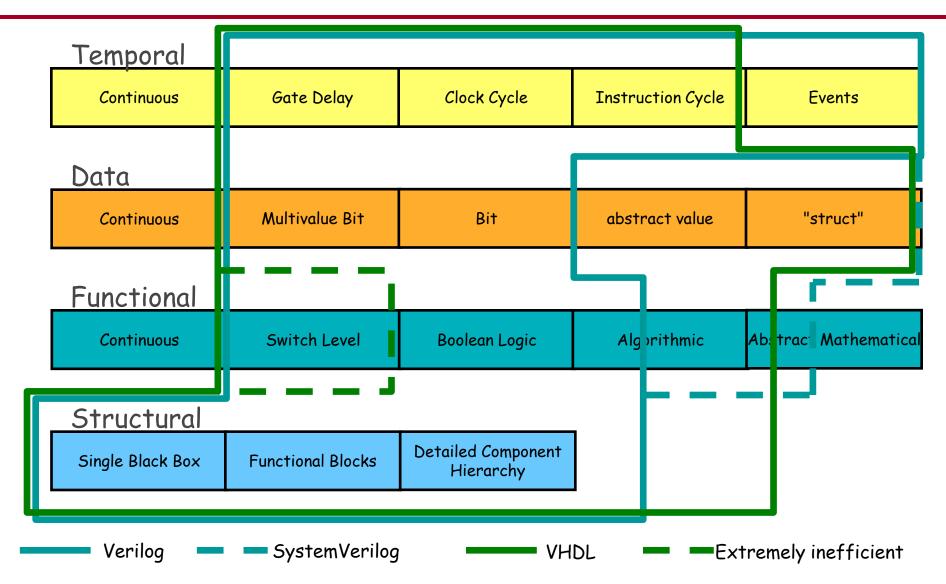
#### Functional Dimension:

- continuous functions (e.g. differential equations)
- Switch-level (transistors as switches)
- Boolean Logic
- Algorithmic (e.g. sort procedure)
- Abstract mathematical formula (e.g. matrix multiplication)

#### Structural Dimension:

- Single black box
- Functional blocks
- Detailed hierarchy with primitive library elements

### Modeling Levels – Major Dimensions



# Verilog for Design Verification

- Assignment calc1 design in Verilog
  - Testbench for calc1 design in Verilog
- Interactive Evita Verilog tutorial (Ch1-4,5-7):
  - Structure of Verilog modules
  - Verilog signal values: 0, 1, x and z (4-valued logic)
  - Verilog signals:
    - nets (used for "connections", no storage capacity)
    - registers (storage capacity, similar to variables in pgr languages)
  - Verilog external signals:
    - ports (input, output or inout, port connecting rules)
  - Coding styles:
    - Structural
    - Dataflow
    - Behavioural (best for verification)

# Continuous Assignment

- Used in Dataflow coding style.
- Keyword assign followed by optional delay declaration
- LHS (target) can be net (scalar or vector) or concatenation of nets
  - NO registers allowed as target for assignment!
- Assignment symbol: =
- RHS is an expression.
  - assign #4 Out = In1 & In2;
- Implicit continuous assignment: wire x = ...;
- Conditional assignment:
  - assign Out = Sel ? In1 : In0;
  - If Sel is 1 then In1 is assigned to Out; if Sel is 0 then Out is In0.
  - If Sel is x/z, evaluate both Inl and Inl, if they are the same then Out is assigned this value, otherwise x/z.

### Continuous Assignment: Execution

- Continuous assignments are always active.
- Concurrency:
  - When any of the operands on RHS changes, assignment is evaluated.
  - Several assignments can be executed concurrently.
  - Race conditions can occur!
    - Two or more assignments, which operate on the same data, read and write the data concurrently.
    - Result, which might be erroneous, depends on which assignment does what when.
- Delays specify time between change of operand on RHS and assignment of resulting value to LHS target.
  - assign #4 Out = In1 & In2;

### Behavioural Coding Style

- Most advanced coding style: flexible and high-level
  - closest to programming languages
  - allows use of conditional statements, case statements, loops, etc.

#### Best for verification, but by no means ideal...

#### Behaviour:

- Actions a circuit is supposed to perform when it is active.
- Algorithmic description: Need "variables" similar to PLs!
  - Abstraction of data storage elements register objects:
    - reg R; one bit register default value x before first assignment
    - time T; can store/manipulate simulation time
    - integer N; by default at least 32 bit stores values signed
    - real R; default value is 0
    - [Other data types, e.g. arrays exist, but are out of the scope of this introduction.]

### Mux421: Behavioural Coding Example

```
module mux421 behavioural (Out, In0, In1, In2, In3, Sel0, Sel1);
  output Out;
  input In0, In1, In2, In3, Sel0, Sel1;
  req Out;
  always @ (Sel1 or Sel0 or In0 or In1 or In2 or In3)
 begin
    case ({Sel1,Sel0})
      2'b00 : Out = In0;
      2'b01 : Out = In1;
      2'b10 : Out = In2;
      2'b11 : Out = In3;
    default : Out = 1/bx;
    endcase
  end
endmodule // mux421 behavioural
```

### **Behavioural Blocks**

- initial and always
  - Can't be nested.
  - Block containing several statements must be grouped using:

```
begin ... end (sequential) orfork ... join (concurrent)
```

- initial block:
  - Used to initialise variables (registers).
  - Executed at (simulation) time 0. Only once!
- always block:
  - Starts executing at time 0.
  - Contents is executed in infinite loop.
    - Means: Executing repeats as long as simulation is running.
  - Multiple blocks are all executed concurrently from time 0.

# Assignment in Behavioural Coding

# Assignment in behavioural coding style is procedural:

- LHS (target) must be a register (reg, integer, real or time) - not a net, a bit or part of a vector of registers.
- NO assign keyword!
- Must be contained within a behavioural (i.e. initial or always) block.
- NOT always active!
  - Target register value is only changed when procedural assignment is executed according to sequence contained in block.
- Delays: indicate time that simulator waits from "finding" the assignment to executing it.

# **Blocking Assignment**

(... as opposed to continuous assignment from dataflow coding style.)

```
reg A;
reg [7:0] Vector;
integer Count;
initial
begin
  A = 1'b0;
  Vector = 8'b0;
  Count = 0;
end
```

Sequential initialisation assignment.

### **Timing Control Evaluation**

- 1. Find procedural assignment
- 2. Wait 5 time units
- 3. Perform A+B
- 4. Wait 10 time units
- 5. Assign result to C
- So, what is the difference between:
  - #10 C = A+B and
  - C = #10 A+B?

### **Events and Wait**

- Events mark changes in nets and registers, e.g. raising/falling edge of clock.
  - @ negedge means from any value to 0
  - @ posedge means from any value to 1
  - @ clk always activates when clock changes
- Wait statement:

```
- wait (condition) stmt;
• wait (EN) #5 C = A + B;
- waits for EN to be 1 before #5 C = A + B;
```

Use wait to block execution by not specifying a statement!

```
- wait (EN); ...
```

# Sensitivity List

```
always @ (posedge Clk or EN) begin ... end
```

- Allows to suspend always blocks.
- Block executes and suspends until signal (one or more) in sensitivity list changes.
- NOTE: or is used to make statement sensitive to multiple signals or events.
- (Don't use sensitivity list to express a logical condition!)
- Common mistake:
  - Forgetting to add relevant signals to sensitivity list!

# Non-blocking Assignments

- Concurrency can be introduced into sequential statements.
  - Delay is counted down before assignment,
  - BUT control is passed to next statement immediately.
- Non-blocking Assignments allow to model multiple concurrent data transfers after common event.
- A blocking assignment would force sequential execution.

```
A <= #1 1; B <= #2 0; (non-blocking)

A x 1 1 1
B x x 0 0
Time: 0 1 2 3

A = #1 1; B = #2 0; (blocking)

A x 1 1 1
B x x x 0
Time: 0 1 2 3
```

### Approaches to Assignment - I

```
reg [7:0] MyReg;
initial
  fork
    #50    MyReg = 8'hFF;
    #100    MyReg = 8'h01;
    #150    MyReg = 8'h2F;
    #200    MyReg = 8'h00;
    #250    $finish;
    join
```

Concurrent blocking (=)

```
Time: 0 50 100 150 200 250 MyReq[7:0] XX FF 01 2F 00 00
```

# Approaches to Assignment - II

```
reg [7:0] MyReg;
initial
begin
    MyReg <= #50 8'hFF; // pass control, wait, assign
    MyReg <= #50 8'h01;
    MyReg <= #50 8'h2F;
    MyReg <= #50 8'h00;
    #250 $finish;
end</pre>
```

Sequential non-blocking (<=)</li>

```
Time: 0 50 100 150 200 250 MyReg[7:0] XX ??
```

### Approaches to Assignment - III

```
reg [7:0] MyReg;
initial
  begin
    MyReg <= #50    8'hFF; // pass control, wait, assign
    MyReg <= #100   8'h01;
    MyReg <= #150   8'h2F;
    MyReg <= #200   8'h00;
    #250 $finish;
end</pre>
```

Sequential non-blocking (<=)</li>

```
Time: 0 50 100 150 200 250 MyReg[7:0] XX FF 01 2F 00 00
```

# Approaches to Assignment - IV

```
reg [7:0] MyReg;
initial
  begin
    #50 MyReg = 8'hFF; // wait, assign, pass control
  #50 MyReg = 8'h01;
  #50 MyReg = 8'h2F;
  #50 MyReg = 8'h00;
  #250 $finish;
end
```

Sequential blocking (=)

```
Time: 0 50 100 150 200 250 MyReq[7:0] XX FF 01 2F 00 00
```

# HDL vs. Programming Languages

### 3 major new concepts of HDLs compared to PLs:

#### Connectivity:

 Ability to describe a design using simpler blocks and then connecting them together.

#### Time:

- Can specify a delay (in time units of simulator): (WHY?)
  - and #2 (Y3, In3, Sel1, Sel0);
- Concurrency is always assumed! (for structural style this is)
  - No matter in which order primitives/components are specified, a change in value of any input signal activates the component.
  - If 2 or more components are activated concurrently, they perform their actions concurrently.
  - Order of specification does not influence order of activation!
  - (NOTE: Statements inside behavioural blocks may be sequential -more later.)

### **Tasks and Functions**

### Tasks and Functions

- Both are purely behavioural.
  - Can't define nets inside them.
  - Can use logical variables, registers, integers and reals.
- Must be declared within a module.
  - Are local to this module.
  - To share tasks/functions in several modules, specify declaration in separate module and use 'include directive.
- Timing (simulation time)
  - Tasks:
    - No restriction on use of timing; engineer specifies execution.
  - Functions:
    - Execute in ZERO sim time units; no timing/event control allowed.

# Comparing Tasks with Functions

	Tasks	Functions
Timing	can be non-zero sim time	execute in 0 sim time
Calling other tasks or functions	no limit; may enable functions	may not call tasks but may call another function  No recursion!
Arguments	any number; any type; can't return result	at least one input; no output/inout; always results in single return value
Purpose	modularize code	react to some input with single response; only combinatorial code; use as operands in expressions

### **Example Task**

```
task factorial;
  output [31:0] f;
  input [3:0] n;
  integer count; // local variable
  begin
    f = 1;
    for (count=n; count>0; count=count-1)
      f = f * count;
  end
endtask
```

- Invoke task: < task name > (list of arguments);
  - Declaration order determines order of arguments when task is called!

### **Example Function**

```
function ParityCheck;
  input [3:0] Data;
  begin
    ParityCheck = ^Data; // bit-wise xor reduction
  end
endfunction
```

 Result is by default a 1 bit register assigned to implicitly declared local variable that has same name as function.

#### Function calls:

- Are either assigned to a variable, or
- occur in an expression that is assigned to a variable,
- or occur as an argument of another function call.

### System Tasks and Functions

- More than 100 Verilog system tasks/functions.
  - (See Evita Verilog Reference Guide for more information.)
- Can be used in any module without explicit include directive.
- **Syntax:** \$< keyword >
- Most important tasks for verification:
  - \$display, \$monitor
  - \$time, \$stop, \$finish
  - (Also with files: \$fopen, \$fdisplay)

### Summary

- Evita Verilog Tutorial [Ch1-7]
- Verilog HDL IEEE Standard 1364
  - Signals: internal and external (ports)
  - Different coding styles:
    - structural
    - dataflow
    - behavioural
- HDLs: Connectivity, Time and Concurrency
- BOOK: Verilog HDL by Samir Palnikar [in QB Library]
- Next:
  - Introduction to Assignment 1!