COMS31700 Design Verification:

Block-level Case Study

(with demonstration of ABV and Formal Verification)

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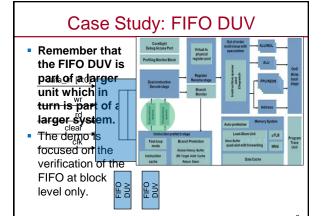
(Acknowledgement: I gratefully acknowledge the support from Jasper Design Automation who provide the licenses for the Formal Verification Tool demonstration.)





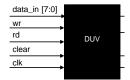
Case Study

Specification
Verification Plan
Directed Testing
(Code Coverage)
Functional Coverage
Assertion-based Verification
Formal Property Checking



Specification

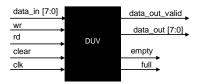
Example DUV Specification - Inputs



- Inputs:
 - wr indicates valid data is driven on the data_in bus
 - data_in is the data to be pushed into the DUV
 - rd pops the next data item from the DUV in the next cycle
 - clear resets the DUV

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Example DUV Specification - Outputs



- Outputs:
 - data_out_valid indicates that valid data is driven on the data_out bus
 - data_out is the data item requested from the DUV
 - empty indicates that the DUV is empty
 - full indicates that the DUV is full

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DUV Specification

- High-Level functional specification of DUV
 - The design is a FIFO.
 - Reading and writing can be done in the same cycle.
 - Data becomes valid for reading one cycle after it is
 - No data is returned for a read when the DUV is empty.
 - Clearing takes one cycle.
 - During clearing read and write are disabled.
 - Inputs arriving during a clear are ignored.
 - The FIFO is 8 entries deep.

Verification Plan

Those who fail to plan plan to fail.

The Verification Plan

- Functions to be verified:
 - Verification Plans are "live" documents For each level in the design verified at that level.

 Verification mains are tive documents.

 They change as our understanding of the DUV increases.
 - In particular, identify corner
- Methods of verification:
 - The Verification Plan is the
- Specification Comr Def
 - for the Verification Process
- Resources required (people) and schedule details:
 - Integrate the verification plan into the overall design plan and estimate the cost of verification.
- Required tools:
 - List the software and hardware necessary to perform verification.

Test Scenarios Matrix - Basic

Test #	Description	Description	
1.1	Check the write functionality		
1.2	Check the read functionality	Develop a Test Scenarios Matrix	
1.3	Check the reset functionality	for our Case	
1.4	Check	Study DUV.	

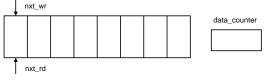
NOTE: "Check that X" should be read as "Create a scenario that allows checking X".

- These generic tests should be broken to more specific tests
 - Test case 1.1.1: Check write when empty
 - Test case 1.1.2: Check write when full - Test case 1.1.3: Check write during reset

White Box View **DUV** Implementation

Example DUV Implementation

- Implementation based on a circular buffer
 - nxt_wr and nxt_rd pointers indicate where the next entry will be written to or read from.
 - data counter indicates the number of valid data items in the FIFO.
 - Complex control logic for pointers and counter.



Functional Coverage

Cross-Product Functional Coverage

[O Lachish, E Marcus, S Ur and A Ziv. Hole Analysis for Functional Coverage Data. In proceedings of the 2002 Design Automation Conference (DAC), June 10-14, 2002, New Orleans, US.]

A cross-product coverage model is composed of the following parts:

- A semantic description of the model (story)
- 2. A list of the attributes mentioned in the story
- 3. A set of all the possible values for each attribute (the attribute value domains)
- 4. A list of restrictions on the legal combinations in the cross-product of attribute values

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FIFO Cross Product Coverage Model

- From a "White Box" verification perspective:
 - The FIFO is implemented using a circular buffer.
 - The circular buffer implementation is based on the following control signals:
 - nxt_rd, nxt_wr, data_counter
 - These signals are used to control the data flow and also the empty and full signals.
 - Verification Plan:
 - Interactions of read and write transactions can create complex and unexpected conditions. All combinations need to be verified to gain confidence in the correctness of the FIFO.
- This is the story.

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FIFO Cross Product Coverage Model

- Attributes relevant to the coverage model:
 - nxt_rd, nxt_wr, data_counter, empty, full signals
- Attri Legal Coverage Tasks:
 - nxt. (0,0,0,1,0)
 - nxt (2,2,8,0,1) - dat (3,7,4,0,0)
 - dat - emu (3,7,4,0,0)
 - full ∈ {0,1}
- Full coverage space:
- 8*8*9*2*2 = **2304** coverage tasks
 - Format: (nxt rd, nxt wr, data counter, empty, full)
 - Find some legal and some illegal examples

. . .

Illegal Coverage Tasks:

(1,1,4,0,0)

(1,5,0,1,1)

FIFO Cross Product Coverage Model

- Restrictions
 - data_counter is only important when nxt_rd==nxt_wr so that one can tell the difference between empty and full

Defining meaningful functional

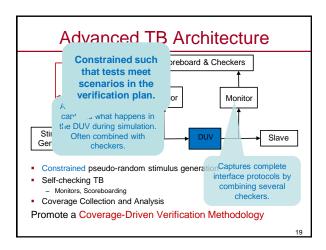
coverage requires design understanding and engineering

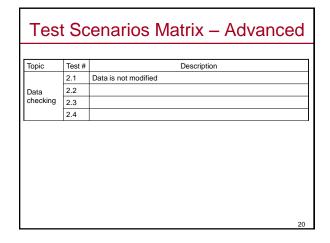
skill

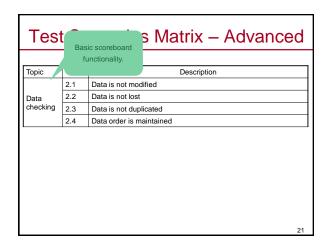
- 64 combinations of nxt_rd and nxt_wr
- 56 cases are for nxt_rd!=nxt_wr. so en
- 8 cases are for nxt_rd==nxt_wr
- one set of 8 is for data_counter==0, so
- one set of 8 is for data_counter==8, so
- empty and full are not both asserted at
- Revised format of coverage model:
 - (nxt_rd, nxt_wr, empty, full)
 - total size of coverage space: 8*8*2*2 = 256 coverage tasks
 - Encode assumptions into properties
- Only 56+8+8 = 72 coverage tasks are legal and meaningful.
- It is worth noting the close link of the above coverage model to the properties for formal verification.

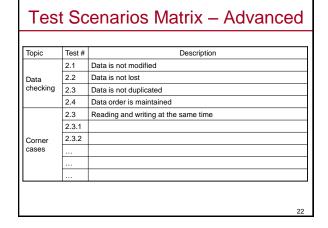
Constraint Pseudo Random
Test Generation

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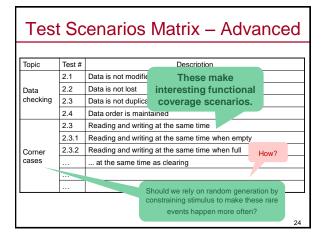








Topic	Test #	Description
Data checking	2.1	Data is not modified
	2.2	Data is not lost
	2.3	Data is not duplicated
	2.4	Data order is maintained
Corner cases	2.3	Reading and writing at the same time
	2.3.1	
	2.3.2	



Bug Hunting

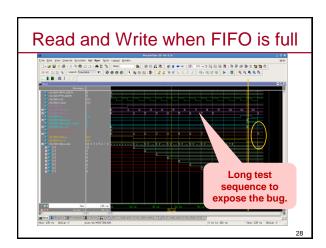
Given the following bug...

Concurrent reading from and writing to FIFO

- ok, the data counter should not change its value
- unless reading from and writing to the same data slot
- this only happens when the FIFO is
 - empty: When the FIFO is empty and there is a write at the same time as a read (from empty), then the read should be ignored.
 - full: When the FIFO is full and there is a read at the same time as a write, then the write (to full) should be ignored.
- But the logic that controls the value of the data counter does not distinguish these special cases.
- What do we need to do to find these bugs?

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Read/Write when FIFO is empty The first two forms instant of the first two forms and the first two fo



ABV

Properties of the DUV Black box view: - Empty and full are never asserted together. - After clear the FIFO is empty. - After writing 8 data items the FIFO is full. - Data items are moving through the FIFO unchanged in terms of data content and in terms of data order. - No data is duplicated. - No data is lost. - data_out_valid only for valid data, i.e. no x's in data.

Properties of the DUV

- White box view:
 - The value range of the read and write pointers is between 0 and 7.
 - The data_counter ranges from 0 to 8.
 - The data in the FIFO is not changed during a clear.
 - For each valid read the read pointer is incremented.
 - For each valid write the write pointer is incremented.
 - Data is written only to the slot indicated by nxt_wr.
 - Data is read only from the slot indicated by nxt_rd.
 - When reading and writing the data_counter remains unchanged.
 - What about a RW from an empty/full FIFO?

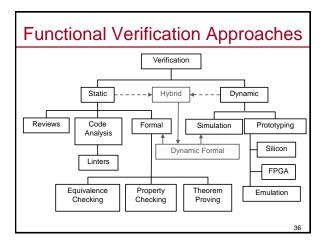
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Property Formalization

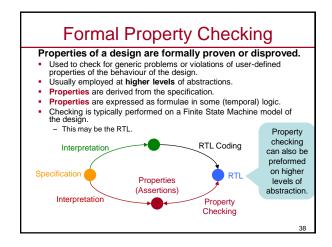
Formalization of key DUV Assertions System Verilog Assertions (SVA) for: Empty and full are never asserted together. property not empty and full; ertions can be @ (posed endprop Challenge: nitored during After c There are many more simulation. propert ertions can also interesting @ (posed used for formal endprop assertions. perty checking. assert property On empty after one write the FIFO is no longer empty. property not empty after write on empty; @ (posedge clk) (empty && wr |=> !empty);

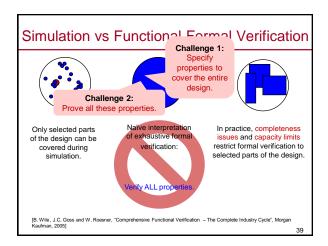
```
Corner case properties
  FIFO empty: When the FIFO is empty and there is a write at the
  same time as a read (from empty), then the read should be ignored.
       property empty_write_ignore_read;
       @(posedge clk) {empty && wr && rd |=>
                           data_counter == $past(data_counter)+1};
       endproperty
       assert property (empty_write_ignore_read);
       cover property (empty_write_ignore_read);
• FIFO full: When the FIFO is full and there is a read at the same time
  as a write, then the write (to full) should be ignored.
      property full_read_ignore_write
      @ (posedge clk) {full && rd && wr |=>
                           data_counter == $past(data_counter)-1};
      endproperty
      assert property (full_read_ignore_write);
      cover property (full_read_ignore_write);
```

Formal Property Checking

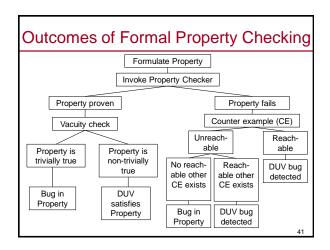


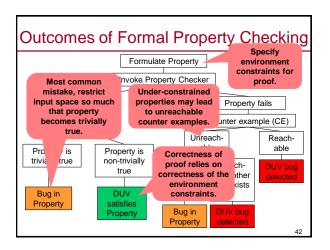
Formal Property Checking Properties of a design are formally proven or disproved. Used to check for generic problems or violations of user-defined properties of the behaviour of the design. Usually employed at higher levels of abstractions. Properties are derived from the specification. Properties are expressed as formulae in some (temporal) logic. Under env_constraint if condition then expectation RTL Coding Properties (Assertions)





The Role of Formal Property Checking Property Checking is the most common form of high-level formal verification used in practice. - Property checking is fully automatic. · Requires the properties to be written. - It performs exhaustive verification of the design wrt the specified properties. - It provides proofs and can demonstrate the absence of bugs. - A counter example is presented for failed properties. Used for critical well specified parts of the design · Cache coherence protocols, Bus protocols, Interrupt controllers Formal Methods can suffer from capacity limits - There are tried and trusted techniques to overcome these: · Restrict property checking to work over finite small time windows. · Limit environment behaviour by strengthening constraints. Case splits over a set of properties, partitioning and black boxing





How do you know you've encoded the property right?

- Keep properties and sequences simple.
 - Build complex properties from simple, short properties
 - Peer review properties you write.



- If the property fails, you can investigate the counter example:
 - Is it reachable or not?
- But if the property succeeds, how do you know whether you've encoded the property right?

Formal Property Checking

- Property checking tools can formally verify assertions.
 - Basic properties (visualize):
 - Basic functionality
 - Range checks
 - Re-use SV Assertions as properties (check):
 - · Empty and full are never asserted together.
 - After clear the FIFO is empty.
 - On empty after one write the FIFO is no longer empty.
 - Understanding counter examples:
 - Debug an assumption
 - Debug a design property

Note: Closely related to functional coverage.

Link from env constraints to simulation assertions.

Formal Property Checking

- Jasper DEMO
 - Formal verification of selected FIFO properties from ABV

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How big is Exhaustive?

- Consider simulating a typical CPU design
- 500k gates, 20k DFFs, 500 inputs
- 70 billion sim cycles, running on 200 linux boxes for a week
- How big: 236 cycles
- Consider formally verifying this design
 - Input sequences: cycles 2(input
 - What about X's: 215000 (5,000 X-assignments + 10,000 non-reset DFFs)
 - How big: 2²⁰⁵⁰⁰ cycles (2¹⁵⁰⁰⁰ combinations of X is not significant here!)
- That's a big number!
 - Cycles to simulate the 500k design:
 - Cycles to formally verify a 32-bit adder: billion)
 - Number of stars in universe:
 - Number of atoms in the universe:
 - Possible X combinations in 500k design:

 - Cycles to formally verify the 500k design:
- (70 billion) 264 (18 billion
- (10^{21}) 2260 (10^{78}) 2¹⁵⁰⁰⁰
- (10⁴⁵¹⁵ x 3) (10^{6171})

Summary

- Block-level Case Study
 - Specification
 - Verification Plan Directed Testing
 - (Code Coverage)
 - Functional Coverage
 - Assertion-based Verification Formal Property Checking
- No single method is adequate to cover a whole design in practice.
 - Carefully select the verification methods that maximize ROI.
 - Complement simulation with formal: Integrated approach

Merry Christmas and a Happy New Year



Why red wine is so important for Christmas