COMS31700 Design Verification:

Block-level Case Study

(with demonstration of ABV and Formal Verification)

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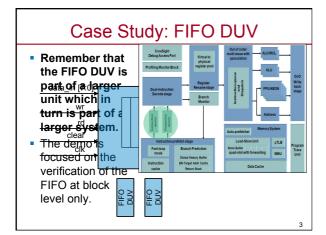
(Acknowledgement: I gratefully acknowledge the support from Jasper Design Automation who provide the licenses for the Formal Verification Tool demonstration.)





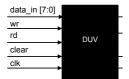
Case Study

Specification
Verification Plan
Directed Testing
(Code Coverage)
Functional Coverage
Assertion-based Verification
Formal Property Checking



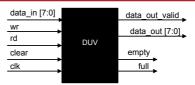
Specification

Example DUV Specification - Inputs



- Inputs:
 - wr indicates valid data is driven on the data_in bus
 - data_in is the data to be pushed into the DUV
 - rd pops the next data item from the DUV in the next cycle
 - clear resets the DUV

Example DUV Specification - Outputs



- Outputs:
 - data_out_valid indicates that valid data is driven on the data_out bus
 - data_out is the data item requested from the DUV
 - empty indicates that the DUV is empty
 - full indicates that the DUV is full

6

DUV Specification

- High-Level functional specification of DUV
 - The design is a FIFO.
 - Reading and writing can be done in the same cycle.
 - Data becomes valid for reading one cycle after it is written.
 - No data is returned for a read when the DUV is empty.
 - Clearing takes one cycle.
 - During clearing read and write are disabled.
 - Inputs arriving during a clear are ignored.
 - The FIFO is 8 entries deep.

Verification Plan

Those who fail to plan, plan to fail.

The Verification Plan

- Functions to be verified:
 - unctions to be verified:

 For each level in the design verified at that level.

 Verification Plans are "live" documents. They change as our understanding of the DUV increases.
 - In particular, identify corner
- Methods of verification: Def

- Comp Def
- The Verification Plan is the Specification
 - for the Verification Process
- In p Resources required (people) and schedule details:
 - Integrate the verification plan into the overall design plan and estimate the cost of verification
- Required tools:
 - List the software and hardware necessary to perform verification.

Test Scenarios Matrix - Basic

| Test # | Description | | |
|--------|-------------------------------|---------------------------------|--|
| 1.1 | Check the write functionality | | |
| 1.2 | Check the read functionality | Develop a Test Scenarios Matrix | |
| 1.3 | Check the reset functionality | for our Case | |
| 1.4 | Check | Study DUV. | |

NOTE: "Check that X" should be read as "Create a scenario that allows checking X"

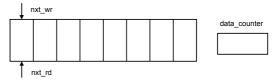
- These generic tests should be broken to more specific tests
 - Test case 1.1.1: Check write when empty
 Test case 1.1.2: Check write when full

 - Test case 1.1.3: Check write during reset

White Box View **DUV** Implementation

Example DUV Implementation

- Implementation based on a circular buffer
 - nxt_wr and nxt_rd pointers indicate where the next entry will be written to or read from.
 - data_counter indicates the number of valid data items in the FIFO.
 - Complex control logic for pointers and counter.



Functional Coverage

Cross-Product Functional Coverage

[O Lachish, E Marcus, S Ur and A Ziv. Hole Analysis for Functional Coverage Data. In proceedings of the 2002 Design Automation Conference (DAC), June 10-14, 2002, New Orleans, US.]

A cross-product coverage model is composed of the following parts:

- 1. A semantic description of the model (story)
- 2. A list of the attributes mentioned in the story
- 3. A set of all the possible values for each attribute (the attribute value domains)
- 4. A list of restrictions on the legal combinations in the cross-product of attribute values

FIFO Cross Product Coverage Model

- From a "White Box" verification perspective:
 - The FIFO is implemented using a circular buffer.
 - The circular buffer implementation is based on the following control signals:
 - nxt rd, nxt wr, data counter
 - These signals are used to control the data flow and also the empty and full signals.
 - Verification Plan:
 - Interactions of read and write transactions can create complex and unexpected conditions. All combinations need to be verified to gain confidence in the correctness of the FIFO.
- This is the story.

FIFO Cross Product Coverage Model

- Attributes relevant to the coverage model:
 - nxt_rd, nxt_wr, data_counter, empty, full signals
- Attribute value domains:
 - nxt_rd ∈ {0,1,2,3,4,5,6,7}
 - nxt_wr ∈ {0,1,2,3,4,5,6,7}
 - data_counter ∈ {0,1,2,3,4,5,6,7,8}
 - empty € {0,1}
 - full ∈ {0,1}
- Full coverage space:
 - 8*8*9*2*2 = 2304 coverage tasks
 - Format: (nxt_rd, nxt_wr, data_counter, empty, full)
 - Find some legal and some illegal examples

FIFO Cross Product Coverage Model

- Attributes relevant to the coverage model:
 - nxt_rd, nxt_wr, data_counter, empty, full signals
- Attri **Legal Coverage Tasks:** (0,0,0,1,0)
 - nxt
 - (2,2,8,0,1)
 - dat – em
- Full coverage space:

full ∈ {0,1}

- 8*8*9*2*2 = 2304 coverage tasks
- Format: (nxt_rd, nxt_wr, data_counter, empty, full)
- Find some legal and some illegal examples

Illegal Coverage Tasks:

(1,1,4,0,0)

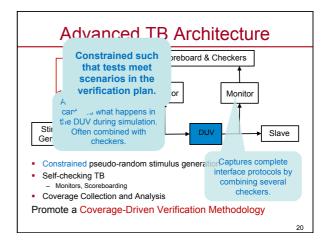
FIFO Cross Product Coverage Model

- - data counter is only important when nxt rd==nxt wr so that one can tell the difference between empty and full

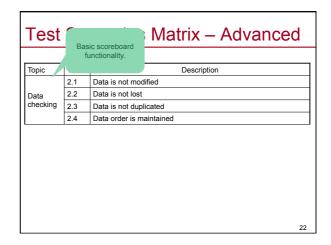
Defining meaningful functional coverage requires design understanding and engineering

- 64 combinations of nxt rd and nxt wr
- 56 cases are for nxt rd!=nxt wr, so er
- 8 cases are for nxt_rd==nxt_wr
- one set of 8 is for data_counter==0, so
- one set of 8 is for data_counter==8, so empty and full are not both asserted at
- Revised format of coverage model:
- (nxt_rd, nxt_wr, empty, full)
- total size of coverage space: 8*8*2*2 = 256 coverage tasks
- Encode assumptions into properties
- Only 56+8+8 = 72 coverage tasks are legal and meaningful.
- It is worth noting the close link of the above coverage model to the properties for formal verification.

Constraint Pseudo Random Test Generation

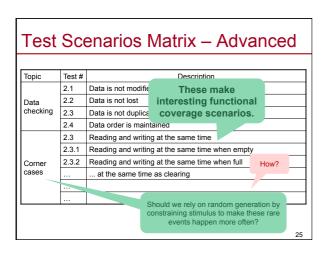


Test Scenarios Matrix — Advanced Topic Test # Description 2.1 Data is not modified 2.2 2.2 2.3 2.4



| rest | Sce | enarios Matrix – Advanced |
|------------------|--------|--------------------------------------|
| | | |
| Topic | Test # | Description |
| Data checking | 2.1 | Data is not modified |
| | 2.2 | Data is not lost |
| | 2.3 | Data is not duplicated |
| | 2.4 | Data order is maintained |
| Corner cases | 2.3 | Reading and writing at the same time |
| | 2.3.1 | |
| | 2.3.2 | |
| | | |
| | | |
| | | |
| | • | |

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| | 2.3.2 | |
| | | |
| | | |
| | | |



Bug Hunting

Given the following bug...

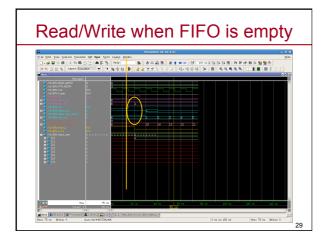
- Concurrent reading from and writing to FIFO
 - Should the data counter change its value?

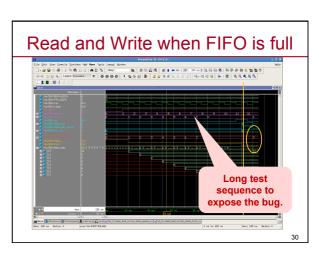
27

Given the following bug...

- Concurrent reading from and writing to FIFO
 - ok, the data counter should not change its value
 - unless reading from and writing to the same data slot
 - this only happens when the FIFO is
 - empty: When the FIFO is empty and there is a write at the same time as a read (from empty), then the read should be ignored.
 - full: When the FIFO is full and there is a read at the same time as a write, then the write (to full) should be ignored.
 - But the logic that controls the value of the data counter does not distinguish these special cases.
 - What do we need to do to find these bugs?

28





ABV

Properties of the DUV

Black box view:

An invariant property.

- Empty and full are never asserted together.
- After clear the FIFO is empty
- After writing 8 data items the FIFO is full.
- Data items are moving through the FIFO unchanged in terms of data content and in terms of data order.
- No data is duplicated.
- No data is lost.
- data_out_valid only for valid data, i.e. no x's in data.

Properties of the DUV

- White box view:
 - The value range of the read and write pointers is between 0 and 7.
 - The data_counter ranges from 0 to 8.
 - The data in the FIFO is not changed during a clear.
 - For each valid read the read pointer is incremented.
 - For each valid write the write pointer is incremented.
 - Data is written only to the slot indicated by nxt wr.
 - Data is read only from the slot indicated by nxt_rd.
 - When reading and writing the data_counter remains unchanged.
 - What about a RW from an empty/full FIFO?

Property Formalization

Formalization of key DUV Assertions

- System Verilog Assertions (SVA) for:
 - Empty and full are never asserted together. property not_empty_and_full;

@ (posed Challenge: After c There are many more simulation. propert interesting @ (posed assertions.

ertions can be nitored during

ertions can also used for formal perty checking.

 On empty after one write the FIFO is no longer empty. property not_empty_after_write_on_empty;
@ (posedge clk) (empty && wr |=> !empty); endproperty

assert property (not empty after write on empty);

Corner case properties

• FIFO empty: When the FIFO is empty and there is a write at the same time as a read (from empty), then the read should be ignored.

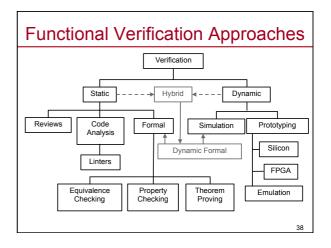
property empty_write_ignore_read;
@ (posedge clk) {empty && wr && rd |=> data_counter == \$past(data_counter)+1}; endproperty

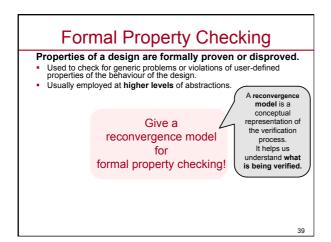
assert property (empty_write_ignore_read); cover property (empty write ignore read);

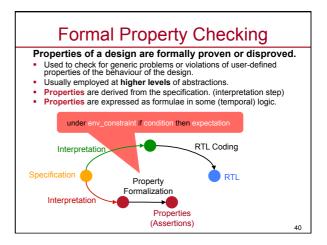
 FIFO full: When the FIFO is full and there is a read at the same time as a write, then the write (to full) should be ignored.

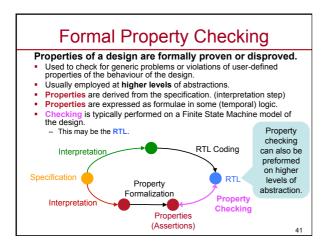
property full_read_ignore_write @ (posedge clk) {full && rd && wr |=> data_counter == \$past(data_counter)-1); endproperty assert property (full_read_ignore_write); cover property (full read ignore write);

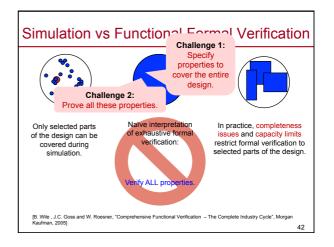
Formal Property Checking





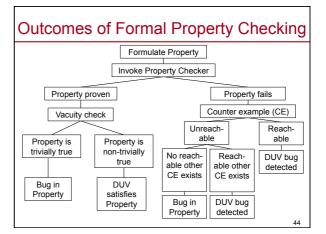


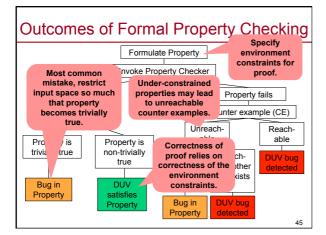




The Role of Formal Property Checking

- Property Checking is the most common form of high-level formal verification used in practice
 - Property checking is fully automatic.
 - Requires the properties to be written.
 - It performs exhaustive verification of the design wrt the specified properties.
 - It provides proofs and can demonstrate the absence of bugs.
 - A counter example is presented for failed properties
 - Used for critical well specified parts of the design
 - Cache coherence protocols, Bus protocols, Interrupt controllers
- Formal Methods can suffer from capacity limits
 - There are tried and trusted techniques to overcome these:
 - Restrict property checking to work over finite small time windows
 - Limit environment behaviour by strengthening constraints
 - Case splits over a set of properties, partitioning and black boxing.





How do you know you've encoded the property right?

- Keep properties and sequences simple.
 - Build complex properties from simple, short properties
 - Peer review properties you write.



- If the property fails, you can investigate the counter example:
 - Is it reachable or not?
- But if the property succeeds, how do you know whether you've encoded the property right?

Formal Property Checking

- Property checking tools can formally verify assertions
 - Basic properties (visualize):
 - Basic functionality
 - Range checks
 - Re-use SV Assertions as properties (check):
 - Empty and full are never asserted together.
 - After clear the FIFO is empty.
 - On empty after one write the FIFO is no longer empty.
 - Understanding counter examples:
 - Debug an assumption
 - Debug a design property

- Note: Closely related to functional coverage.
 - Link from env_constraints to simulation assertions

Formal Property Checking

- Jasper DEMO
 - Formal verification of selected FIFO properties from ABV

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How big is Exhaustive?

- Consider simulating a typical CPU design
 - 500k gates, 20k DFFs, 500 inputs
 70 billion sim cycles,

 - running on 200 linux boxes for a week

 How big: 2³⁶ cycles
- Consider formally verifying this design
- That's a big number!
 - Cycles to simulate the 500k design:Cycles to formally verify a 32-bit adder:

 - Number of stars in universe:
 - Number of atoms in the universe: Possible X combinations in 500k design:
 - Cycles to formally verify the 500k design:

 Input sequences: cycles 2^(inputs-state) = 2²⁰⁵⁰⁰
 What about X's: 2¹⁵⁰⁰⁰ (5,000 X-assignments + 10,000 non-reset DFFs)
 How big: 2²⁰⁵⁰⁰ cycles (2¹⁵⁰⁰⁰ combinations of X is not significant here!) (70 billion) (18 billion billion) (10²¹) (10⁷⁸) (10⁴⁵¹⁵ x 3) (10⁶¹⁷¹) 2⁷⁰ 2²⁶⁰

Summary

- Block-level Case Study
 - Specification

 - Verification PlanDirected Testing
 - (Code Coverage)

 - Functional CoverageAssertion-based Verification
 - Formal Property Checking



- No single method is adequate to cover a whole design in practice.
 - Carefully select the verification methods that maximize ROI.
 - Complement simulation with formal: Integrated approach

Merry Christmas and a Happy New Year



Why red wine is so important for Christmas