

# CSE225: Data Structure and Algorithm

## Array, Pointer & Searching

# Array

- Declaration
- Initialization
- Accessing Array Elements
- Inserting and Deleting in a Unsorted Array
- Inserting and Deleting in a sorted Array

# Linear Arrays

- A linear array is a list of a finite number of **n** homogeneous data elements ( that is data elements of the same type) such that.
  - The elements of the arrays are referenced respectively by an index set consisting of **n** consecutive numbers.
  - The elements of the arrays are stored respectively in **successive memory locations.**

# Linear Arrays

- The number **n** of elements is called the length or size of the array.
- The index set consists of the integer **0, 1, 2, ..., n-1**
- **Length** or the number of data elements of the array can be obtained from the index set by

$$\text{Length} = \text{UB} - \text{LB} + 1$$

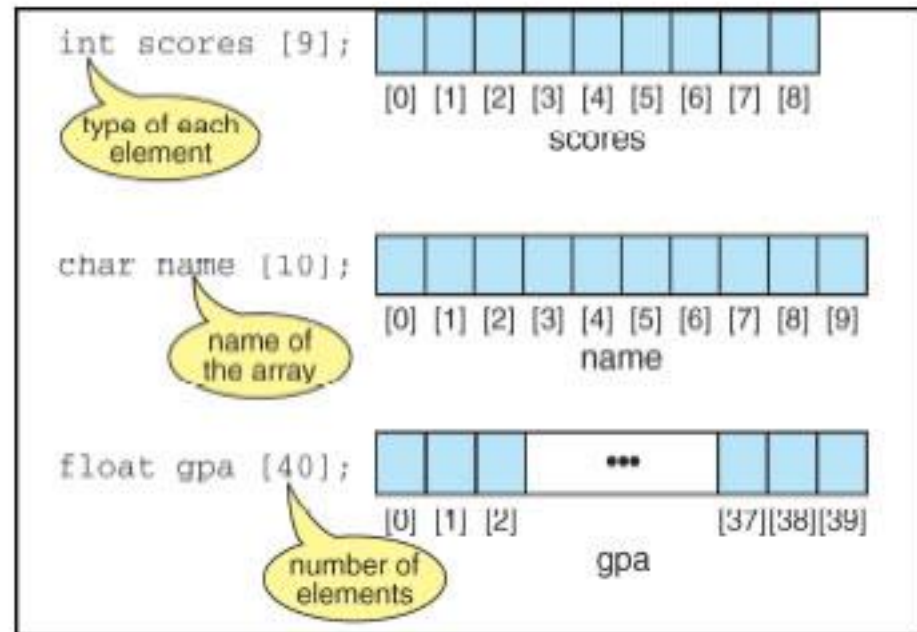
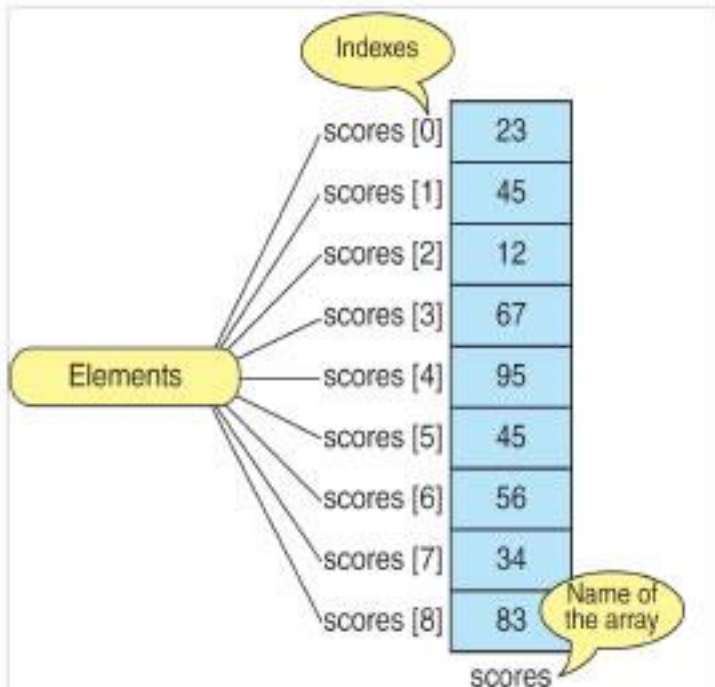
where

**UB** is the **largest index** called the **upper bound** and

**LB** is the **smallest index** called the **lower bound** of the arrays

# Linear Arrays

- Element of an array **A** may be denoted by
  - Subscript notation **A<sub>1</sub>, A<sub>2</sub>, , .... , A<sub>n</sub>**
  - Parenthesis notation **A(1), A(2), .... , A(n)**
  - Bracket notation **A[1], A[2], ..... , A[n]**
- The number **K** in A[K] is called subscript or an index and A[K] is called a **subscripted variable**.



# Declaring and Initializing Arrays

To declare an array:

```
data_type array_name[SIZE];
```

```
int ar_name[10]
```

**data\_type** is a valid data type that must be common to all elements.

**array\_name** is name given to array

**SIZE** is a constant value that defines array maximum capacity.

## Initializing Arrays

Initialization of an array either one by one or using a single statement as follows –

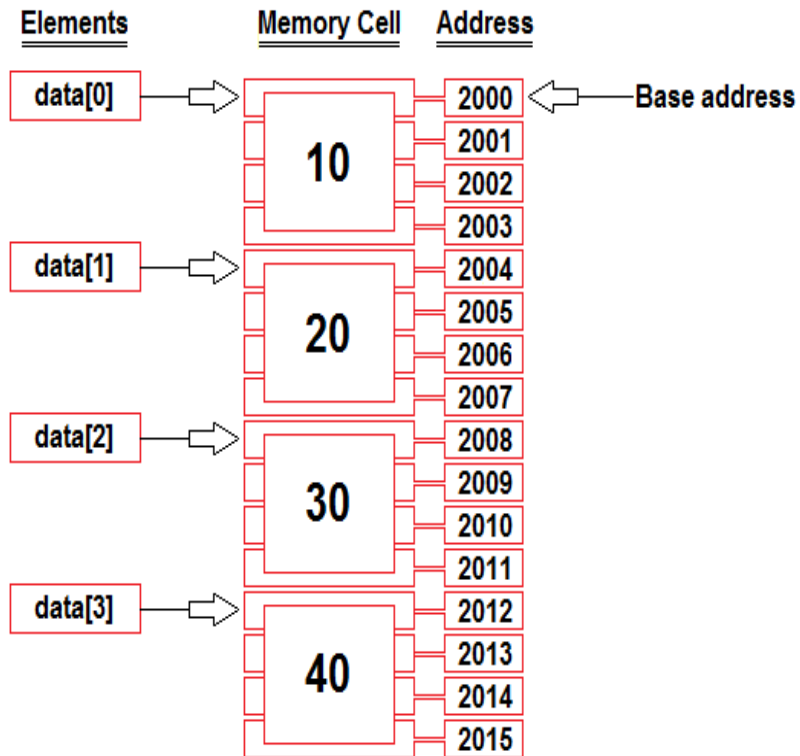
```
int ar[5];
```

```
ar[0]=10; ar[1]=20;
```

```
double balance[5] = { 1000.0, 2.0, 3.4, 7.0, 50.0};
```

# Representation of Linear Array in Memory

# Representation of Linear Array in Memory

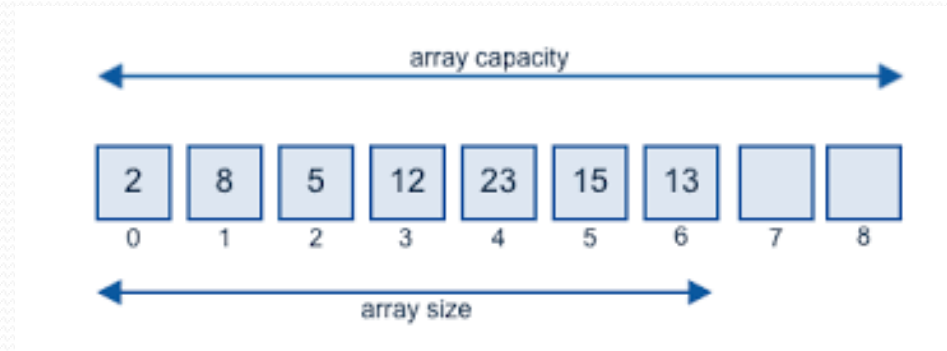
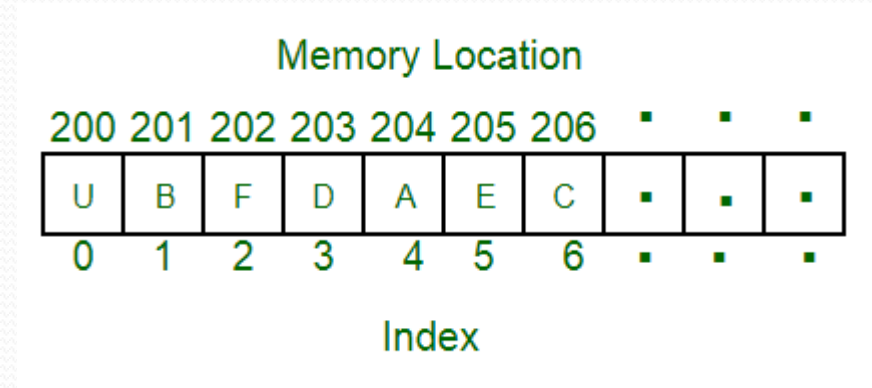


arr[0]	arr[1]	arr[2]	arr[3]	arr[4]
10	20	30	40	50
2000	2004	2008	2012	2016



# Representation of Linear Array in Memory

Address of the memory				indexes of array elements
	1001	20	1	
	1002			
	1003	50	2	
	1004			
	1005	102	3	
	1006			
	1007	600	4	
	1008			
	1009	2	5	
	1010			
	1011	34	6	
	1012			
	1013	500	7	
	1014			
	1015	100	8	
	1016			



# Representation of Linear Array in Memory

int LA[10]

Let

- **LA** be a linear array in the memory of the computer.
- **LOC(LA[K]) = address of the element LA[K]**
- Computer does not keep track of the address of every element of the array
- Keep track address of the first element of array
- called the **base address** of LA and denoted by **Base(LA)**
- $\text{LOC(LA[K])} = \text{Base(LA)} + w(\text{K} - \text{lower bound})$   
where  $w$  is the number of words per memory cell

# Example 1

Find the address for LA[6] Each element of the array occupy 1 byte

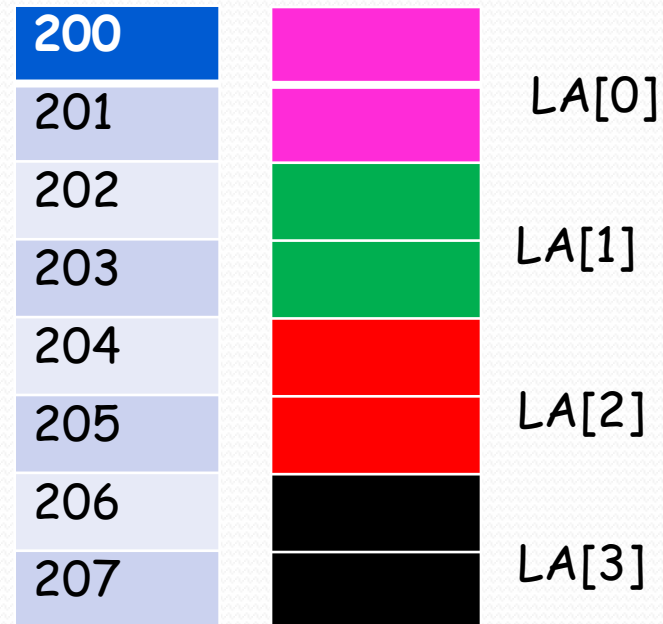
200		LA[1]
201		LA[2]
202		LA[3]
203		LA[4]
204		LA[5]
205		LA[6]
206		LA[7]
207		LA[8]

$$LOC(LA[K]) = \text{Base}(LA) + w(K - \text{lower bound})$$

$$LOC(LA[6]) = 200 + 1(6 - 1) = 205$$

# Example 2

Find the address for LA[16] Each element of the array occupy 2 byte



$$\text{LOC}(\text{LA}[\text{K}]) = \text{Base}(\text{LA}) + w(\text{K} - \text{lower bound})$$

⋮

$$\text{LOC}(\text{LA}[16]) = 200 + 2(16 - 0) = 232$$

# Searching Arrays

- **Linear search:** Compare each element of array with key value
- Search an array for a *key value*
  - Simple
  - Useful for small and unsorted arrays
- Suppose you want to find a number in an unordered sequence
- You have no choice – look through all elements until you have found a match
- This is called linear or sequential search

# Linear Search(LA,N, ITEM)

Linear Search (LA, N, ITEM)

Here LA is a linear array with N elements. This algorithm find an element ITEM into the LA.

1.  $i=0$
2. Repeat steps 3 and 4 while  $i \leq n$  or  $LA[i] == ITEM$
3. IF  $LA[i] == ITEM$  print item found at index  $i$  and exit
4.  $i=i+1$
5. Print item not found and exit.

```
int search(int data[],int n, int
item)
{
for(int i = 0; i<n; i++)
    {
        if(data[i] == item)
            return i;
    }
return -1;
}
```

# Searching Arrays: Binary Search

- Binary search
  - For sorted arrays
  - Compares **middle** element with **key**
    - If equal, match found
    - If **key** < **middle**, looks in first half of array
    - If **key** > **middle**, looks in last half
    - Repeat
  - Very fast; at most  $n$  steps, where  $2^n > \text{number of elements}$ 
    - 30 element array takes at most 5 steps
      - $2^5 > 30$  so at most 5 steps

# Searching Arrays: Binary Search

```
int bsearch(int data[], int n, int value )
{
    int first, middle, last;
    first = 0;
    last = n - 1;
    while (true) {
        middle = (first + last) / 2;
        if (data[middle] == value)
            return middle;
        else if (first >= last)
            return -1;
        else if (value < data[middle])
            last = middle - 1;
        else
            first = middle + 1;
    }
}
```



# Inserting in Unsorted Array

INSERT (LA, N, ITEM) Here LA is a linear array with N elements. This algorithm inserts an element ITEM into the LA.

1. If  $MAX == N$ , print overflow
2. Set  $LA[N] := ITEM$
3. [Reset N.] Set  $N := N + 1$ .
4. Exit.

# Delete in Unsorted Array

DELETE (LA, N, k) Here LA is a linear array with N elements. This algorithm Delete the element ITEM from the LA.

1. Set  $LA[k] := LA[N-1]$
2. [Reset N.] Set  $N := N - 1$ .
3. Exit.

# INSERT\_SORTL (LA, N, K, ITEM)

Here LA is a sorted array with N elements and K is a positive integer such that  $K < N$ . This algorithm insert an element ITEM from the Kth position in LA.

1.  $j = N$
2. Repeat step while  $LA[j-1] > ITEM$ ;
3. Set  $LA[j] = LA[j-1]$
4.  $j = j - 1$ ;
5. Set  $LA[j] = ITEM$
6. [Reset N.] Set  $N := N + 1$ ;
7. Exit.

0	10
1	20
2	25
3	30
4	40
5	45
6	50
7	60
8	70
9	80

# DELETE\_SORTL (LA, N, K, ITEM)

Here LA is a sorted array with N elements and K is a positive integer such that  $K < N$ . This algorithm Delete an element ITEM from the Kth position in LA.

1. SEARCH(LA, N, K, ITEM)  
Set  $j := k$ .
2. Repeat Steps 3 and 4 while  $j < (N-1)$ .
3. Move  $j^{\text{th}}$  element upward.  
Set  $LA[j] := LA[j+1]$ .
4. [Increase counter.] Set  $j := j+1$ . [End of Step 2 loop.]
5. [Reset N.] Set  $N := N-1$ ;
6. Exit.

0	10
1	20
2	30
3	50
4	60
5	70
6	80
7	90

# Multidimensional Array

- One-Dimensional Array
- Two-Dimensional Array
- Three-Dimensional Array
- Some programming Language allows as many as 7 dimension

# Two-Dimensional Array

- A Two-Dimensional  $m \times n$  array  $A$  is a collection of  $m.n$  data elements
- with property that  $1 \leq J \leq m$  and  $1 \leq K \leq n$

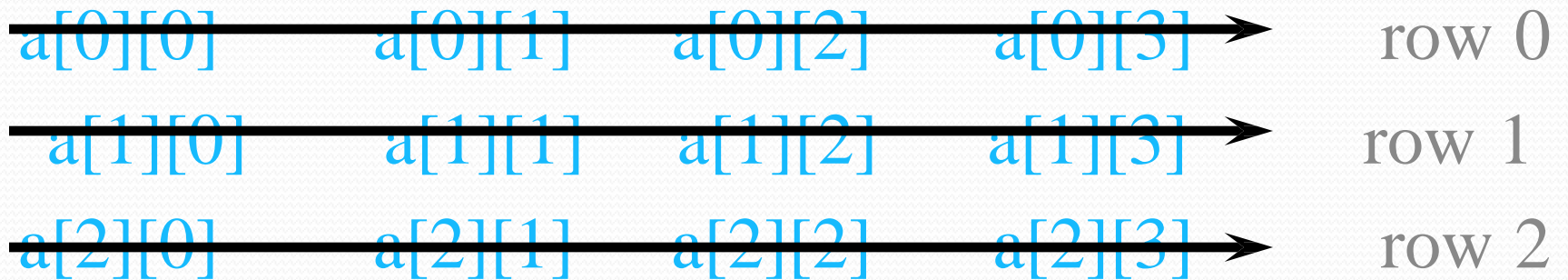
The element of  $A$  with first subscript  $J$  and second subscript  $K$  will be denoted by  $A[J][K]$

# 2D Arrays

The elements of a 2-dimensional array `a` is shown as below

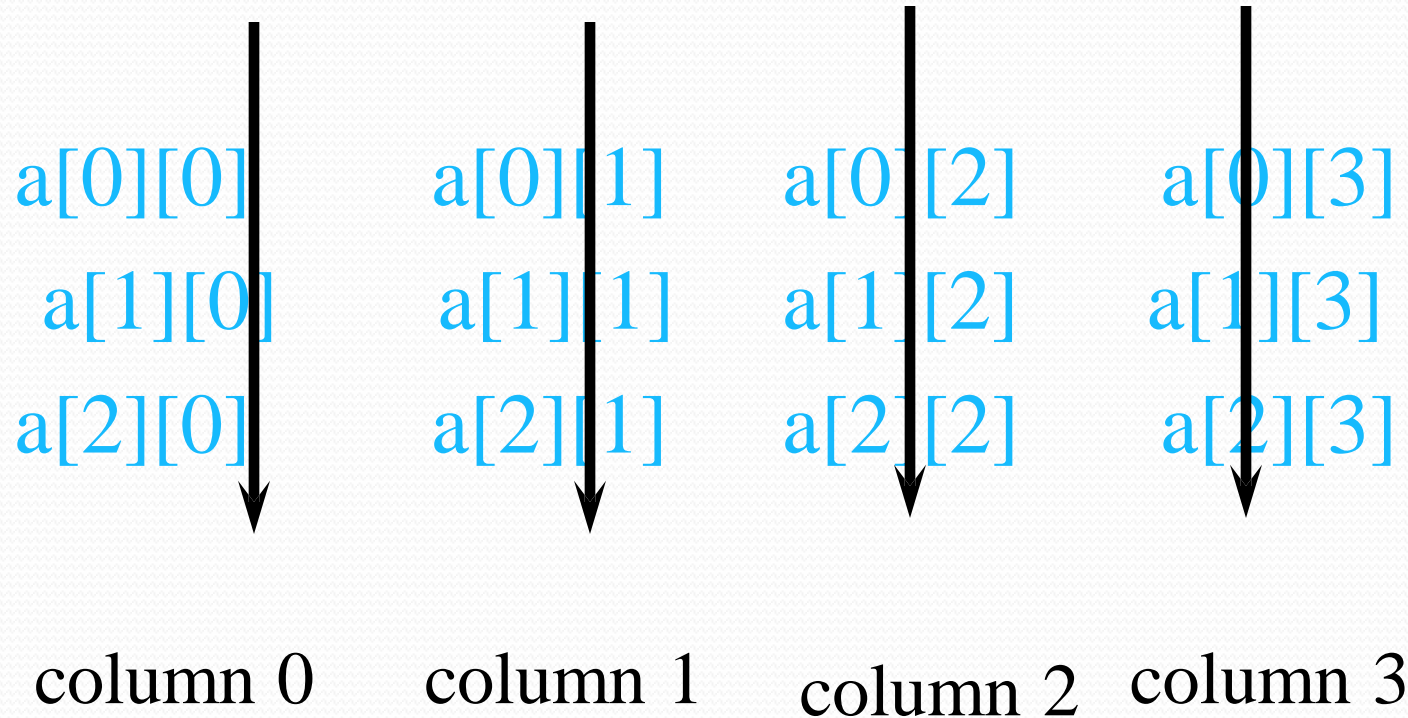
<code>a[0][0]</code>	<code>a[0][1]</code>	<code>a[0][2]</code>	<code>a[0][3]</code>
<code>a[1][0]</code>	<code>a[1][1]</code>	<code>a[1][2]</code>	<code>a[1][3]</code>
<code>a[2][0]</code>	<code>a[2][1]</code>	<code>a[2][2]</code>	<code>a[2][3]</code>

# Rows Of A 2D Array





# Columns Of A 2D Array



# 2D Array

- Let **A** be a two-dimensional array **m x n**
- The array **A** will be represented in the memory by a block of **m x n** sequential memory location
- Programming language will store array **A** either
  - **Column by Column**
  - (Called Column-Major Order) Ex: Fortran, MATLAB
  - **Row by Row**
  - (Called Row-Major Order) Ex: C, C++ , Java

## 2D Array in Memory

A

Subscript

	(1,1)		
	(2,1)	Column	1
	(3,1)		
	(1,2)		
	(2,2)	Column	n 2
	(3,2)		
	(1,3)		
	(2,3)	Colu	mn 3
	(3,3)		
	(1,4)		
	(2,4)	Column	n 4
	(3,4)		

Column-Major Order

A

Subscript

	(1,1)	
	(1,2)	Row 1
	(1,3)	
	(1,4)	
	(2,1)	
	(2,2)	
	(2,3)	Row2
	(2,4)	
	(3,1)	
	(3,2)	Row3
	(3,3)	
	(3,4)	

Row-Major Order

# What is a pointer variable?

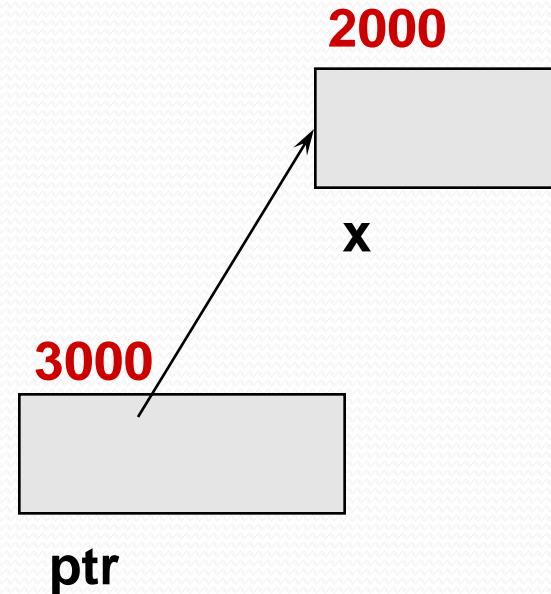
- A pointer variable is a **variable whose value is the address of a location in memory.**
- To declare a pointer variable, you must specify the type of value that the pointer will point to.
- For example,

```
int*    ptr; // ptr will hold the address of an int
char*   q;   // q will hold the address of a char
```

# Using a pointer variable

```
int  x;  
x = 12;
```

```
int* ptr;  
ptr = &x;
```



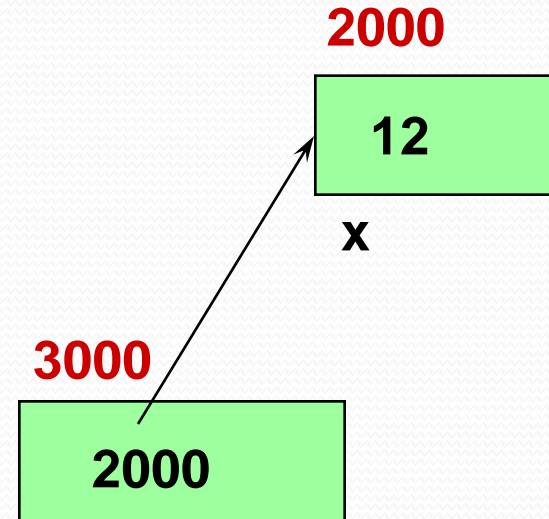
**NOTE:** Because **ptr** holds the address of **x**,  
we say that **ptr** “points to” **x**

Unary **operator \*** is the deference (indirection) operator

```
int x;  
x = 12;
```

```
int* ptr;  
ptr = &x;
```

```
std::cout << *ptr; ptr
```



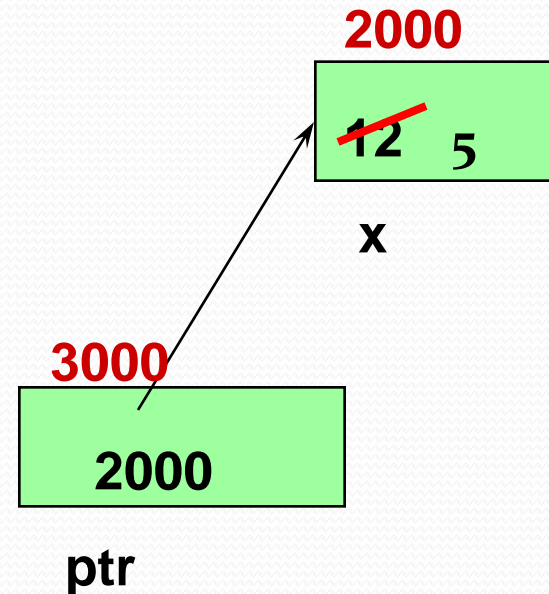
**NOTE:** The value pointed to by `ptr` is denoted by `*ptr`

# Using the dereference operator

```
int  x;  
x = 12;
```

```
int* ptr;  
ptr = &x;
```

```
*ptr = 5;  
// changes the value  
// at address ptr to 5
```



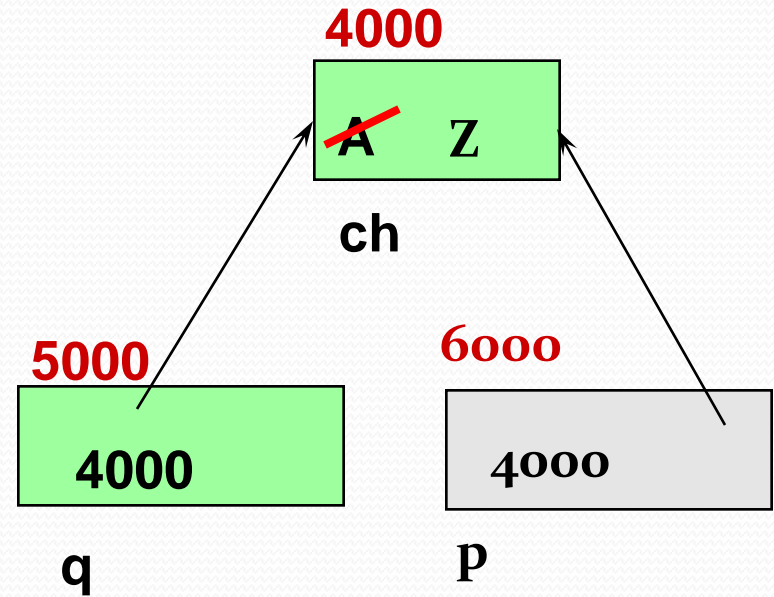
# Another Example

```
char  ch;  
ch = 'A';
```

```
char* q;  
q = &ch;
```

```
*q = 'Z';  
char* p;
```

```
p = q;    // the right side has value 4000  
          // now p and q both point to ch
```





# Dynamically Allocated Data

```
char* ptr;
```

```
ptr = new char;
```

```
*ptr = 'B';
```

```
std::cout << *ptr;
```

2000



ptr

**New** is an operator

# Dynamically Allocated Data

```
char* ptr;
```

```
ptr = new char;
```

```
*ptr = 'B' ;
```

```
std::cout << *ptr;
```

**2000**



**ptr**



**NOTE: Dynamic data has no variable name**

# Dynamically Allocated Data

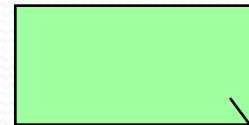
```
char* ptr;
```

```
ptr = new char;
```

```
*ptr = 'B';
```

```
std::cout << *ptr;
```

2000



ptr

'B'

# Dynamically Allocated Data

```
char* ptr;
```

```
ptr = new char;
```

```
*ptr = 'B' ;
```

```
std::cout << *ptr;
```

```
delete ptr;
```

2000

?

ptr

**NOTE:**

**Delete deallocates the memory pointed to by ptr.**



```
// Online C++ compiler to run C++ program online
```

```
#include <iostream>
```

```
using namespace std;
```

```
int main() {
```

```
    int* ptr;
```

```
    ptr = new int;
```

```
    *ptr = 100;
```

```
    cout << *ptr<<endl;
```

```
    delete ptr;
```

```
    cout << *ptr;
```

```
    return 0;
```

```
}
```

## *what does **new** do?*

- takes a pointer variable,
- allocates memory for it to point, and
- leaves the address of the assigned memory in the pointer variable.
- If there is no more memory, the pointer variable is set to `NULL`.

# The **NULL** Pointer

There is a pointer constant called **NULL** available in `cstddef`.

**NULL** is not a memory address;

it means that the pointer variable points to nothing.

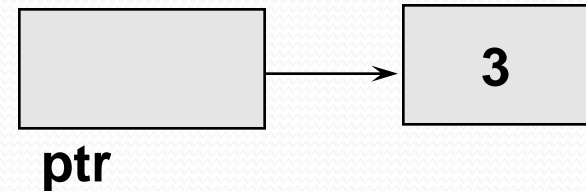
It is an error to dereference a pointer whose value is **NULL**.

It is the programmer's job to check for this.

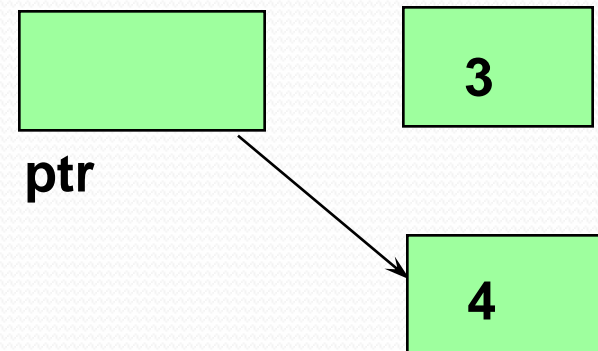
```
while (ptr != NULL)
{
    . . .      // ok to use *ptr here
}
```

# What happens here?

```
int* ptr = new int;  
*ptr = 3;
```



```
ptr = new int;    // changes value of ptr  
*ptr = 4;
```



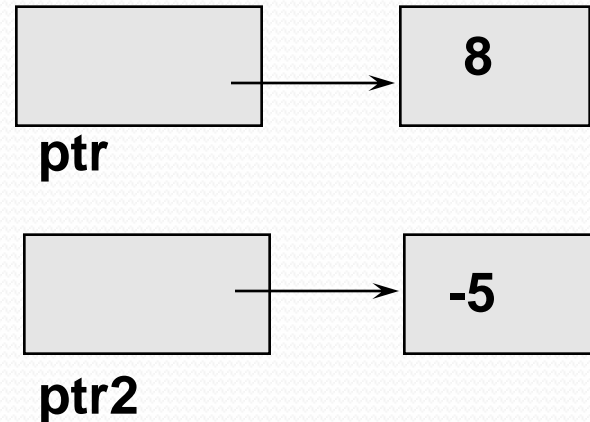


# Memory Leak

A memory leak occurs when dynamic memory (that was created using operator `new`) has been left without a pointer to it by the programmer, and so is inaccessible.

```
int* ptr = new int;  
*ptr = 8;
```

```
int* ptr2 = new int;  
*ptr2 = -5;
```

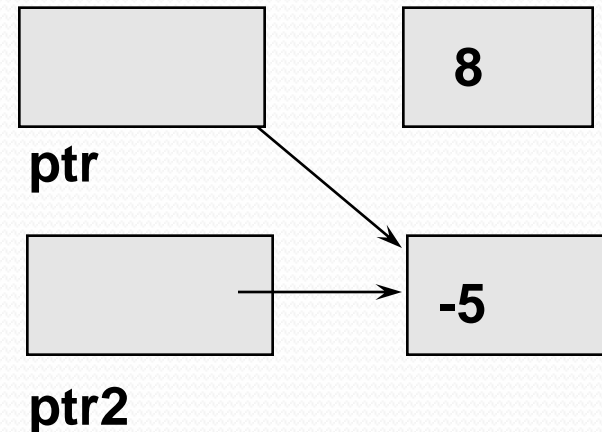
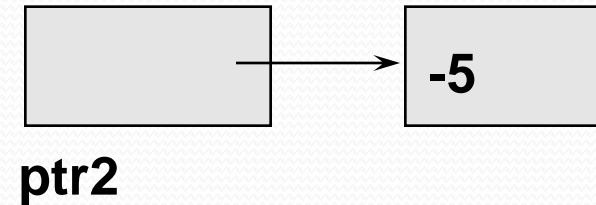
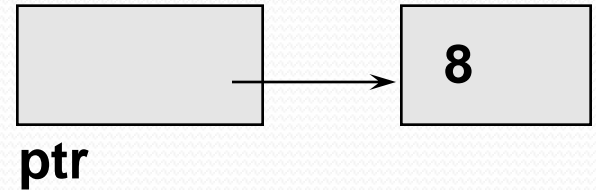


How else can an object become inaccessible?

# Causing a Memory Leak

```
int* ptr = new int;  
*ptr = 8;  
int* ptr2 = new int;  
*ptr2 = -5;  
  
ptr = ptr2;
```

// here the 8 becomes inaccessible



# Using operator delete

The **object or array currently pointed to by the pointer is deallocated**, and the pointer is considered unassigned.

The memory is returned to the free store.

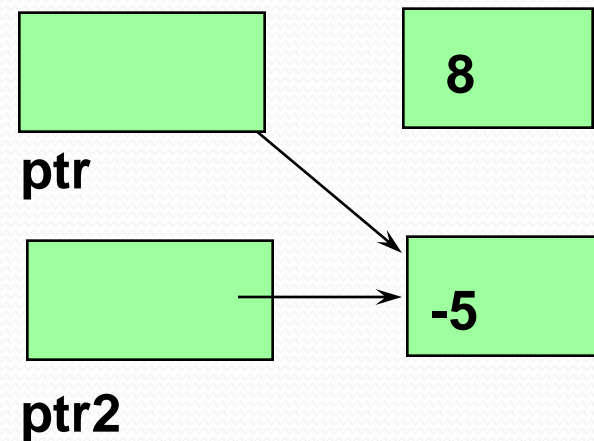
Square brackets are used with delete to deallocate a dynamically allocated array of classes.

# A Dangling Pointer

- occurs when two pointers point to the same object and delete is applied to one of them.

```
int* ptr = new int;  
*ptr = 8;  
int* ptr2 = new int;  
*ptr2 = -5;
```

```
ptr = ptr2;
```



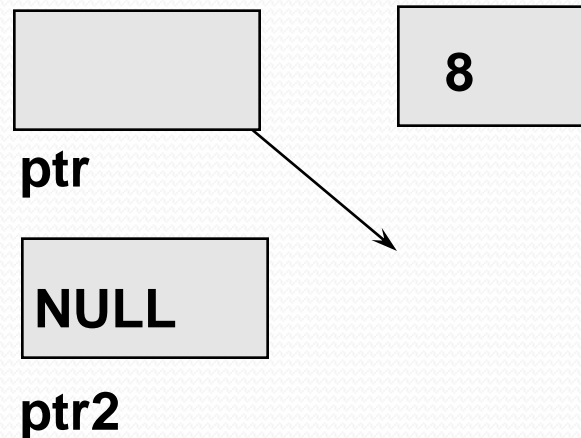
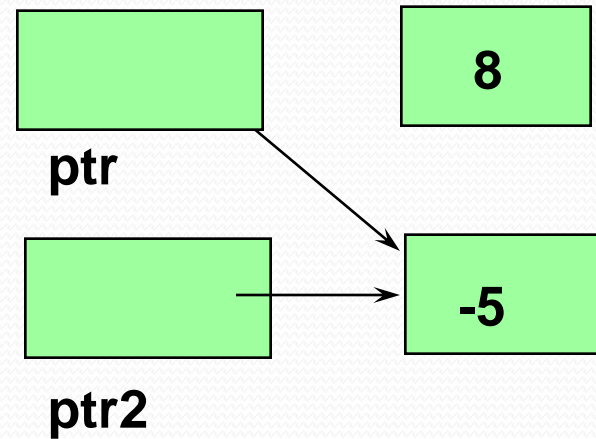
A dangling pointer is a pointer **which points to some non-existing** memory location.

# Leaving a Dangling Pointer

```
int* ptr = new int;  
*ptr = 8;  
int* ptr2 = new int;  
*ptr2 = -5;  
ptr = ptr2;
```

```
delete ptr2;  
ptr2 = NULL;
```

*// ptr is left dangling*



```
#include <iostream>
using namespace std;
```

```
int main() {
    int* ptr = new int;
    *ptr = 8;
    cout<<"ptr = "<<*ptr<<endl;
```

```
    int* ptr2 = new int;
```

```
    *ptr2 = -5;
    ptr = ptr2;
```

```
    cout<<"ptr = "<<*ptr<<endl;
    cout<<"ptr2 = "<<*ptr2<<endl;
```

```
    delete ptr2;
    cout<<"ptr2 = "<<*ptr2<<endl;
```

```
    ptr2 = NULL;
    cout<<"ptr = "<<*ptr<<endl;
```

```
    return 0;
}
```

## Dangling Pointer

# Using Pointer and Reference in the parameter of a Function

## Using Pointer

```
void change(int* i){
    *i = 30;
}

int main(void){
    int a = 40;
    change(&a);
    cout << a;
    return 0;
}
```

The parameter is a pointer;  
Need to pass pointer in the  
argument by using the **&** sign

## Using Reference

```
void change(int& i){
    i = 30;
}

int main(void){
    int a = 40;
    change(a);
    cout << a;
    return 0;
}
```

The parameter is a reference;  
Can use the variable name  
directly

# Remember?

- A list is a homogeneous collection of elements, with a **linear relationship** between elements.
- Each list element (except the first) has a **unique predecessor**, and
- each element (except the last) has a **unique successor**.