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	Roll no 1 - 197 Subj + 05.
	Bald 1- CVZ
	EXPERT METUT, (S)
*	Aim - Implement proyour on Produces & consumer with concurrent control.
ě	theosy -
•	- Concessioning conted ensured that multiple processes of thready execute costeely & efficiently when accessing shared regover without propes control, issue like.  Sace condition, deadlock & inconsistently dealer may arise.
	· key concept!
	Place condition:  Occurs when two or more process accers showed  deals simultaneously, leading to un predictable results  Exi- The thready incrementing the same count may  miss same increment.
•	2. Pro Aliania
	A code segment where shared reported are accersed.  Must be protected at to allow any one process at a time.
	3 Solutions 1-
	1) muter Locky - Enguled multial exclusion Conly one.
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	to the second of
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	or companies to Court to the design of south it.
	2) semaphoses 1- Genoralization of muter with counting excelliniting accept to N.
	2) Moniber: High level syndronization constant Clik in jewal
	· Samuphotes 1-
	and the state of t
	Binestey Canting
	Co to 1) ( ~ to o)
	Binary Semaphore are used to check to critical section is occupied as not wherever counting semaphores are used to carring till at empty sloke of
	section is occupied as not whether counting remaphores
	one used for carring full of empty slots &
	oxitical region
	and the state of the military with the state of the
	· Simple analogy 1-
	I measure a sheded pointed (cs) in an office
	isonly one employee can point at a line
	ij Is multiple sty to point they must wait in a Queue
*	Conclusion :-
No. of the last	Thus, we per borned and implemented procpan on commentors
	or concustent execution king should king both Birary &
Control	counting semaphore.
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## **Program (Concurrent Execution):**

```
#include <stdio.h>
#include <pthread.h>
#include <semaphore.h>
#include <unistd.h>
#define BUFFER SIZE 5
#define MAX ITEMS 20 // Total items to produce
int buffer[BUFFER SIZE];
int count = 0;
int in = 0; // Producer index
int out = 0; // Consumer index
sem_t empty, full;
pthread_mutex_t mutex;
void* producer(void* arg) {
  for (int i = 1; i <= MAX ITEMS; i++) {
    sem_wait(&empty); // Wait for empty slot
    pthread mutex lock(&mutex);
    // Produce item
    buffer[in] = i;
    in = (in + 1) % BUFFER SIZE;
    count++;
    printf("Produced: %d\n", i);
    pthread_mutex_unlock(&mutex);
    sem post(&full); // Signal new item available
    sleep(1); // Simulate production time
  return NULL;
}
```

```
void* consumer(void* arg) {
  for (int i = 1; i <= MAX ITEMS; i++) {
    sem wait(&full); // Wait for available item
    pthread mutex lock(&mutex);
    // Consume item
    int item = buffer[out];
    out = (out + 1) % BUFFER SIZE;
    count--;
    printf("Consumed: %d\n", item);
    pthread mutex unlock(&mutex);
    sem post(&empty); // Signal empty slot available
    sleep(2); // Simulate consumption time (slower than
production)
  return NULL;
int main() {
  pthread t prod thread, cons thread;
  // Initialize semaphores and mutex
  sem init(&empty, 0, BUFFER SIZE);
  sem init(&full, 0, 0);
  pthread mutex init(&mutex, NULL);
  // Create threads
  pthread create(&prod thread, NULL, producer, NULL);
  pthread create(&cons thread, NULL, consumer, NULL);
  // Wait for threads to finish
  pthread join(prod thread, NULL);
  pthread join(cons thread, NULL);
```

```
// Cleanup
sem_destroy(&empty);
sem_destroy(&full);
pthread_mutex_destroy(&mutex);

printf("Finished producing and consuming %d items\n",
MAX_ITEMS);
return 0;
}
```

## **Output:**

```
Produced: 1
Consumed: 1
Produced: 2
Produced: 3
Consumed: 2
Produced: 4
Produced: 5
Consumed: 3
Produced: 6
Produced: 7
Consumed: 4
Produced: 8
Produced: 9
Consumed: 5
Produced: 10
Consumed: 6
Produced: 11
Consumed: 7
Produced: 12
Consumed: 8
Produced: 13
```