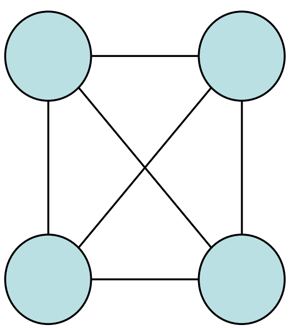
Data Structures
Instructor: Hafiz Tayyeb Javed
Week-15-Lecture-02

22. Minimum Spanning Tree (MST)
Prims Algorithm

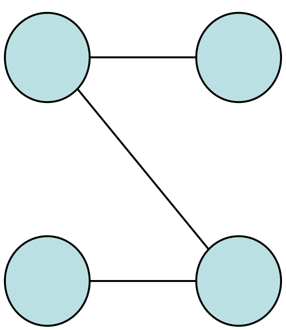
Spanning Trees

- A spanning tree of a graph is a subgraph that contains all the vertices and is a tree

- Formal definition
  - Given a connected graph with $|V| = n$ vertices
  - A spanning tree is defined a collection of $n - 1$ edges which connect all n vertices
  - The n vertices and $n - 1$ edges define a connected sub-graph



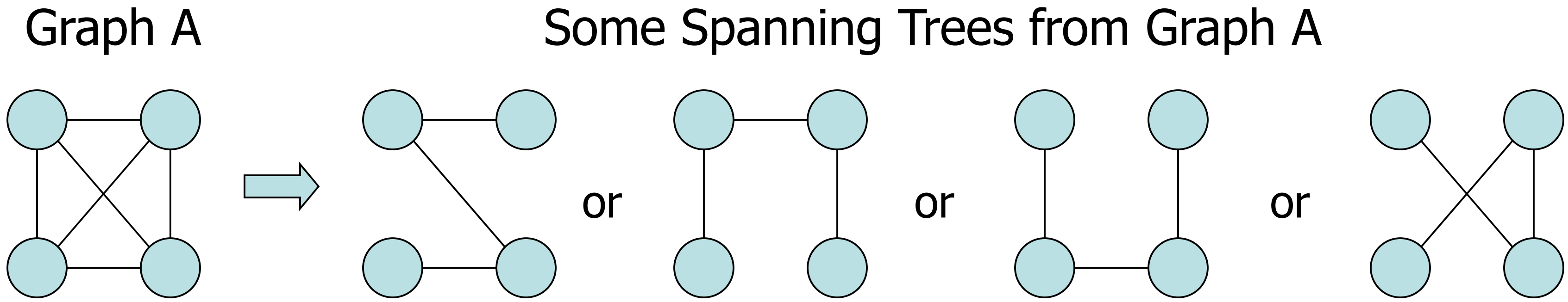
Graph



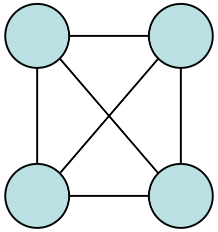
Spanning Tree

Spanning Trees

- A spanning tree is not necessarily unique



Spanning Trees – Example



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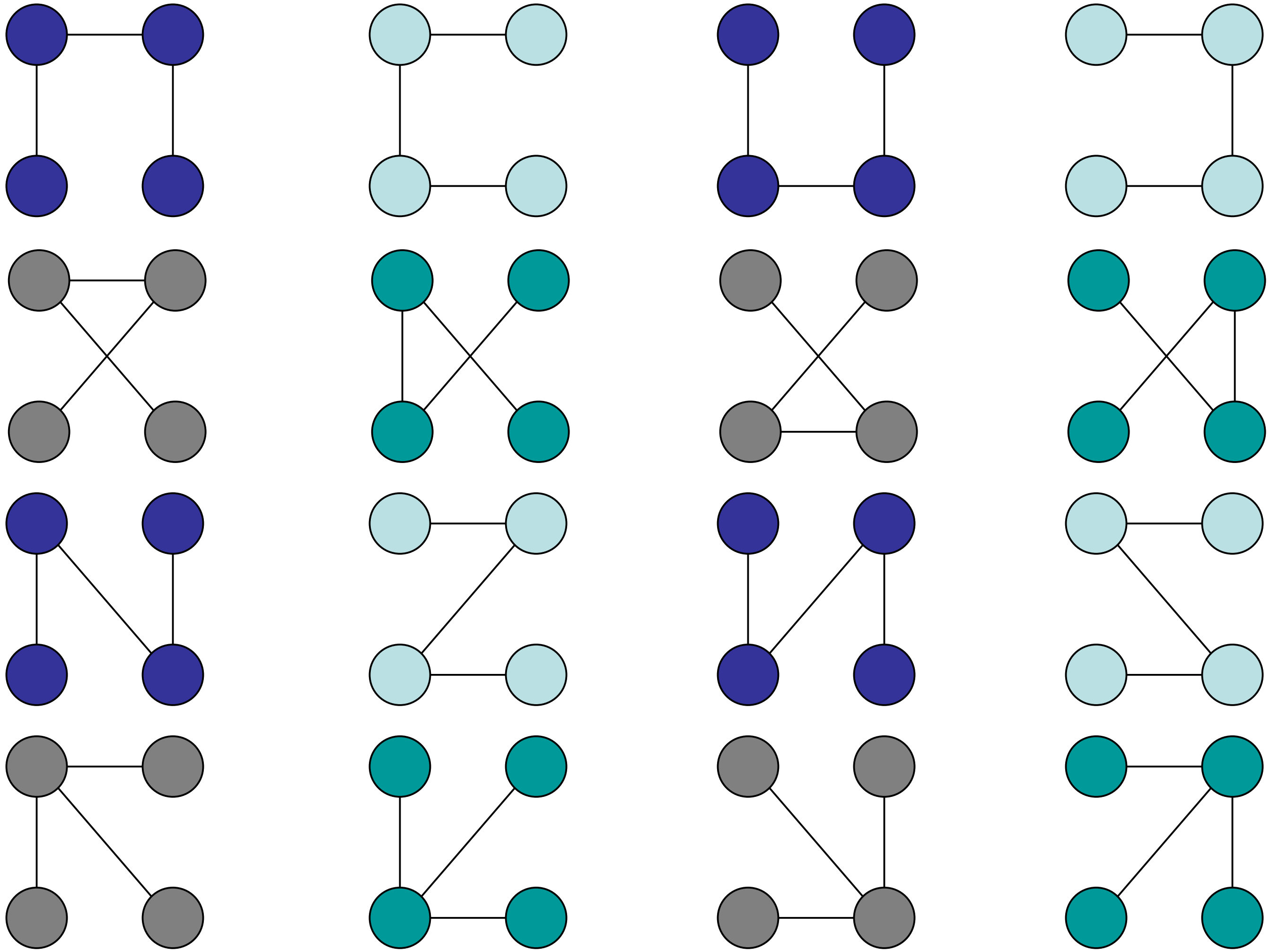
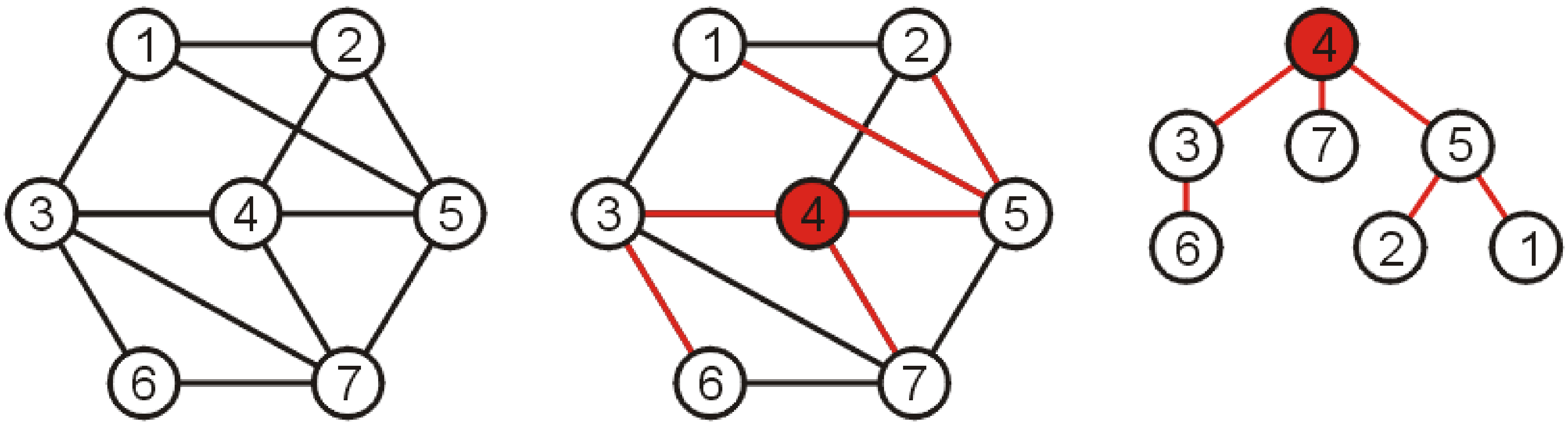


All 16 of its Spanning Trees

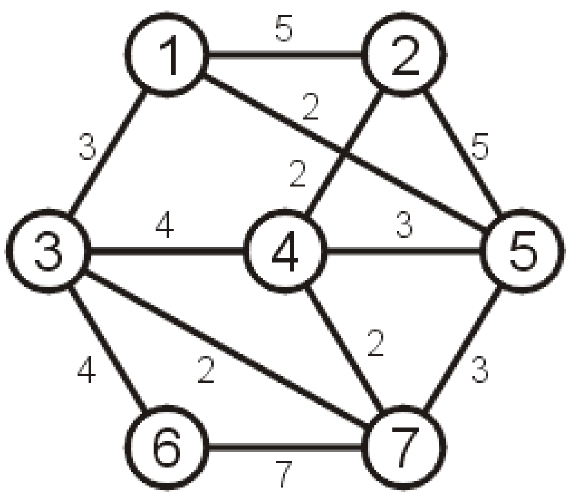
Spanning Trees

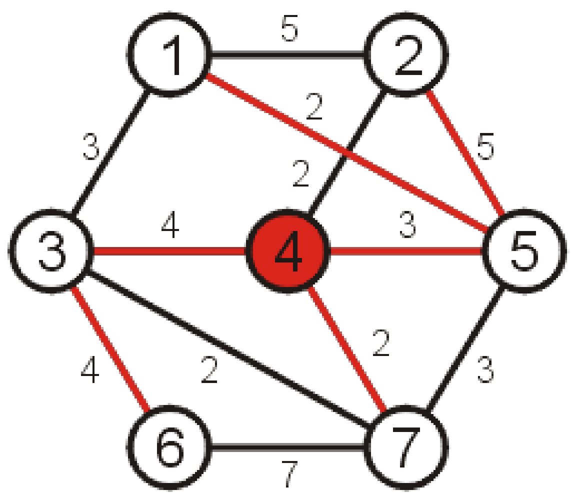
- Why such a collection of `|V|-1` edges is called a tree?
  - If any vertex is taken to be the root, we form a tree by treating the adjacent vertices as children, and so on...

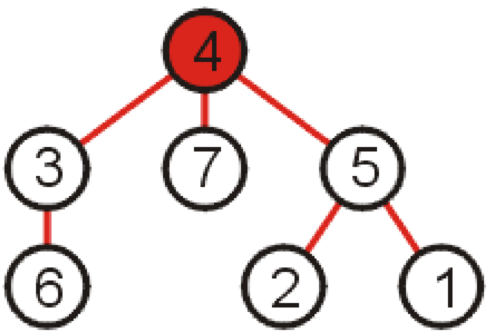


Spanning Tree on Weighted Graphs

- Weight of a spanning tree
  - Sum of the weights on all the edges which comprise the spanning tree



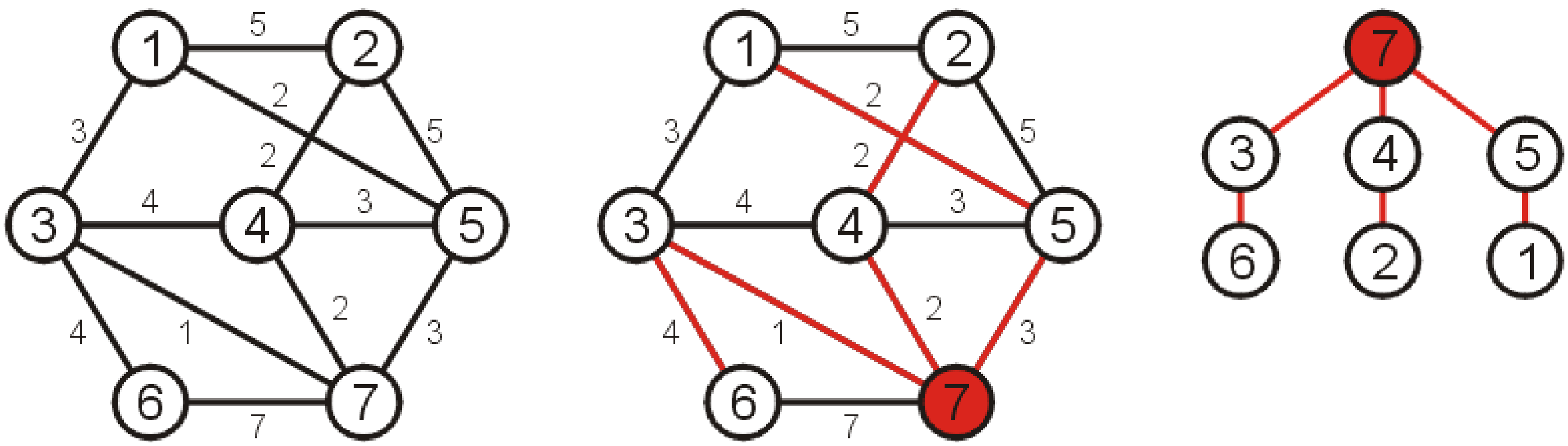




- The weight of this spanning tree is 20

Minimum Spanning Tree (MST)

- Which spanning tree which minimizes the weight?
  - Such a tree is termed a minimum spanning tree



- The weight of this spanning tree is 14

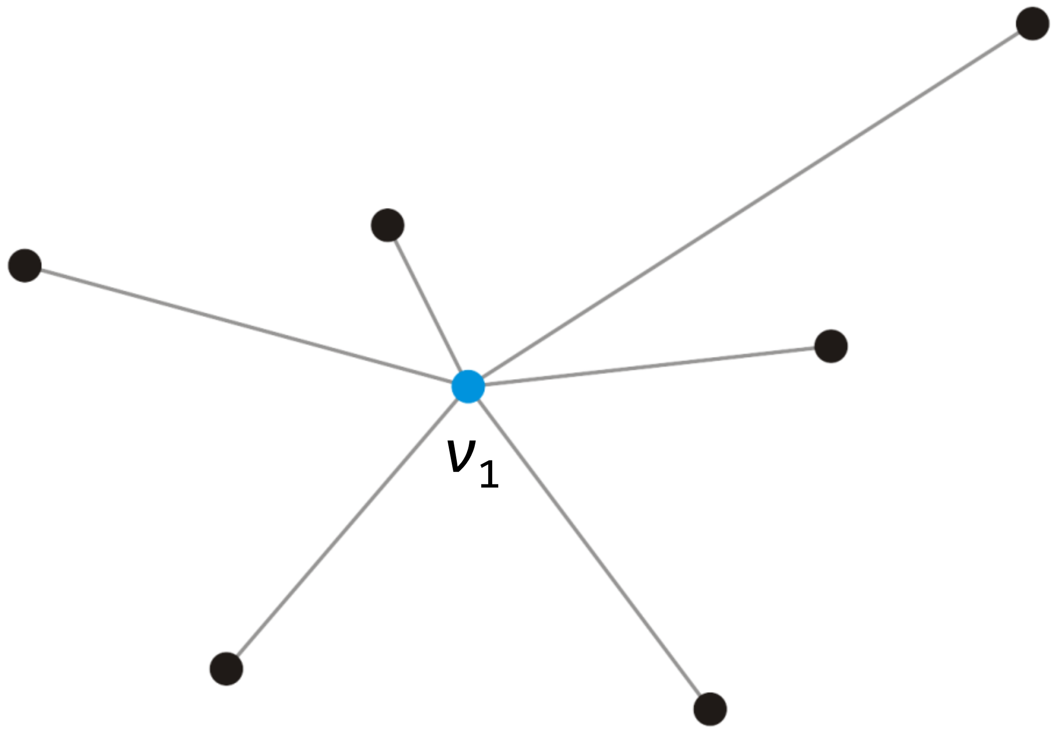
Algorithms For Obtaining MST

- Kruskal's Algorithm
- Prim's Algorithm
- Boruvka's Algorithm

Prim’s Algorithm

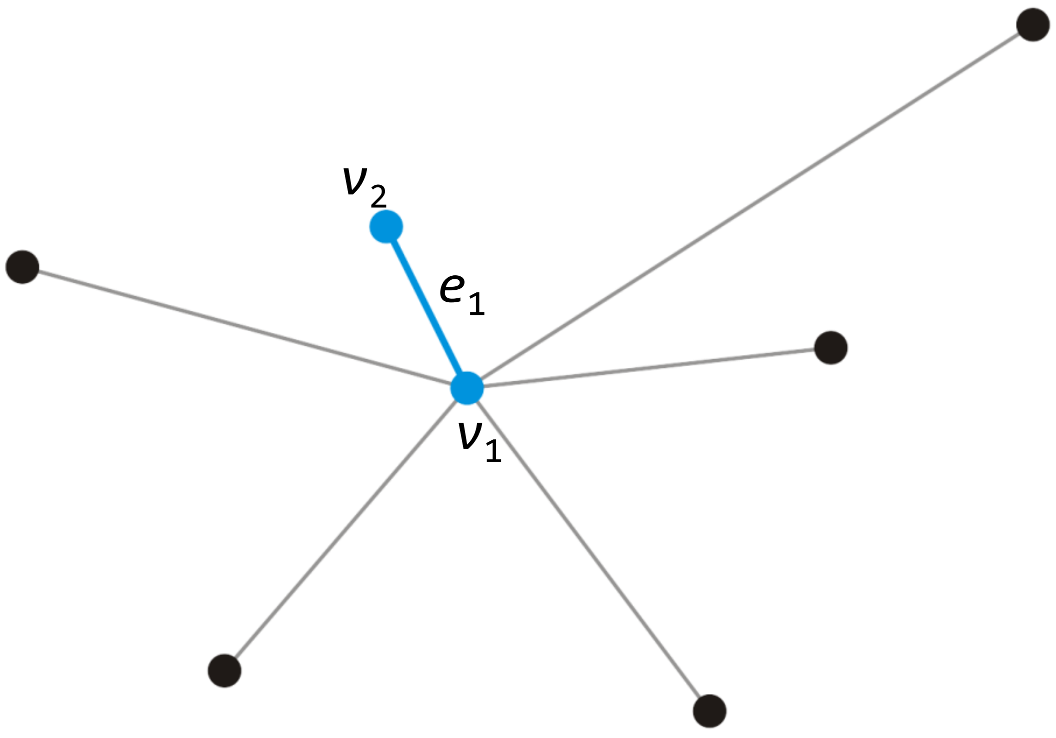
Idea

- Suppose we take a vertex v_1
  - It forms a minimum spanning tree on one vertex



Idea

- Add that adjacent vertex v_2 that has a connecting edge e_1 of minimum weight
  - This forms a minimum spanning tree on our two vertices
  - e_1 must be in any minimum spanning tree containing the vertices v_1 and v_2



Prim's Algorithm

- Start with an arbitrary vertex to form a minimum spanning tree on one vertex
- At each step, add that vertex v not yet in the minimum spanning tree
  - That has an edge with least weight that connects v to the existing minimum spanning sub-tree
- Continue until we have `n - 1` edges and `n` vertices

Prim's Algorithm – Pseudocode

```
MST-Prim(G, w, r) // w is the weight matrix of edges, r is root
    S = Ø
    Q = V[G];  // Insert graph vertices to a Queue
    for each u ∈ Q  // Set distance of all vertices as ∞
        key[u] = ∞;
    key[r] = 0;  // Distance of root is set to 0
    p[r] = NULL; // Parent of root is NULL
    while (Q not empty) {
        u = ExtractMin(Q); // Get the vertex u with min key[u]
        S = S∪{u}
        for each v ∈ Adj[u] {  // Adj is the adjacency list
              if (v ∉ S and w(u,v) < key[v]) {
                p[v] = u;
                key[v] = w(u,v); // weight of an edge (u,v)
            }
        }
  }
```

Prim's Algorithm – Data Structure

- Associate with each vertex two items of data
  - The minimum distance to the partially constructed tree
    - For a given vertex `v`, `key[v]` represent minimum distance
  - Pointer to the vertex that will form the parent node in resulting tree
    - For a given vertex v, `p[v]` represent parent node

- Initialization
  - Set the distance of all vertices as ∞, e.g., for all $u \in G$, `key[u]=` ∞
  - Set all vertices to being unvisited
    - Add vertices to the Queue
  - Select a root node and set its distance as 0, i.e., `key[r] = 0`
  - Set the parent pointer of root to NULL, i.e., `p[r] = NULL`

Prim's Algorithm – Example

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MST-Prim(G, w, r)
    S = Ø
    Q = V
    for each u ∈ Q
        key[u] = ∞;
    key[r] = 0;
    p[r] = NULL;
    while (Q not empty)
        u = ExtractMin(Q);
        S = S∪{u}
        for each v ∈ Adj[u]
            if (v ∉ S and w(u,v) < key[v])
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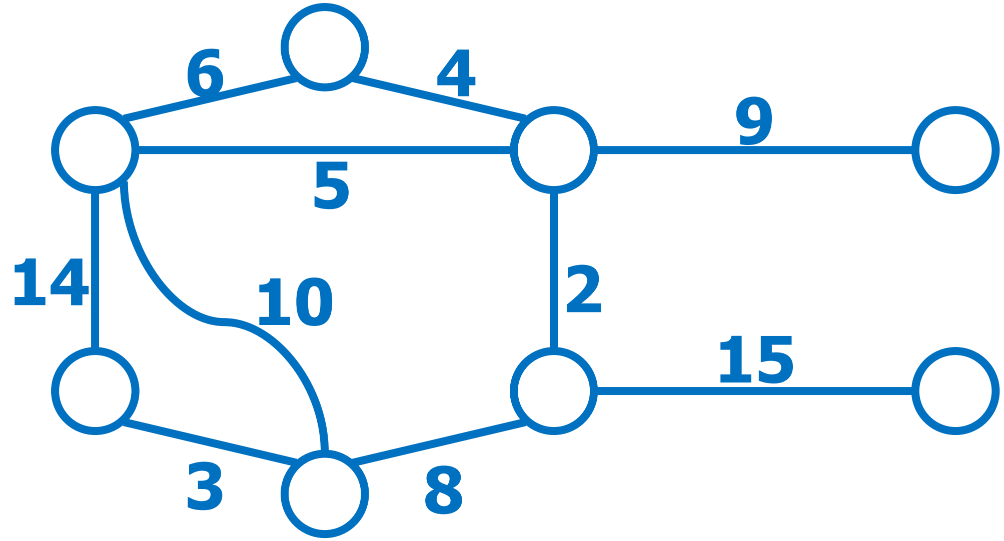

Run on example graph

Prim's Algorithm – Example

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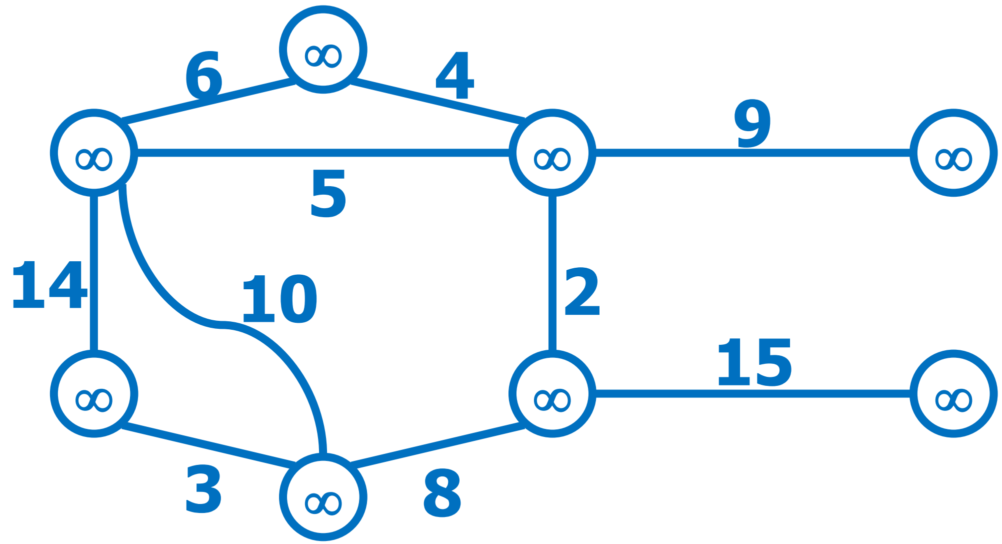

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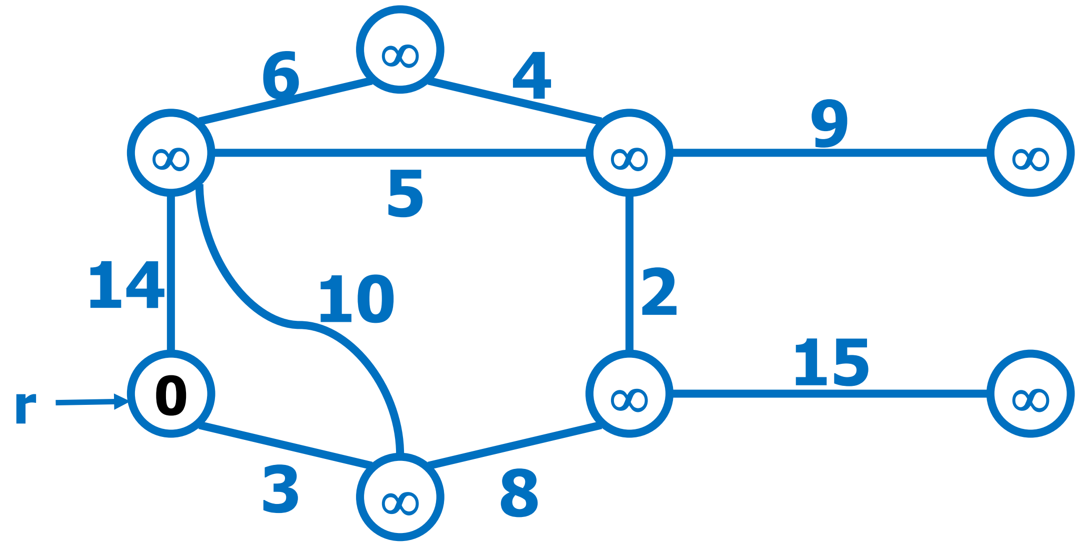

Pick a start vertex r

Prim's Algorithm – Example

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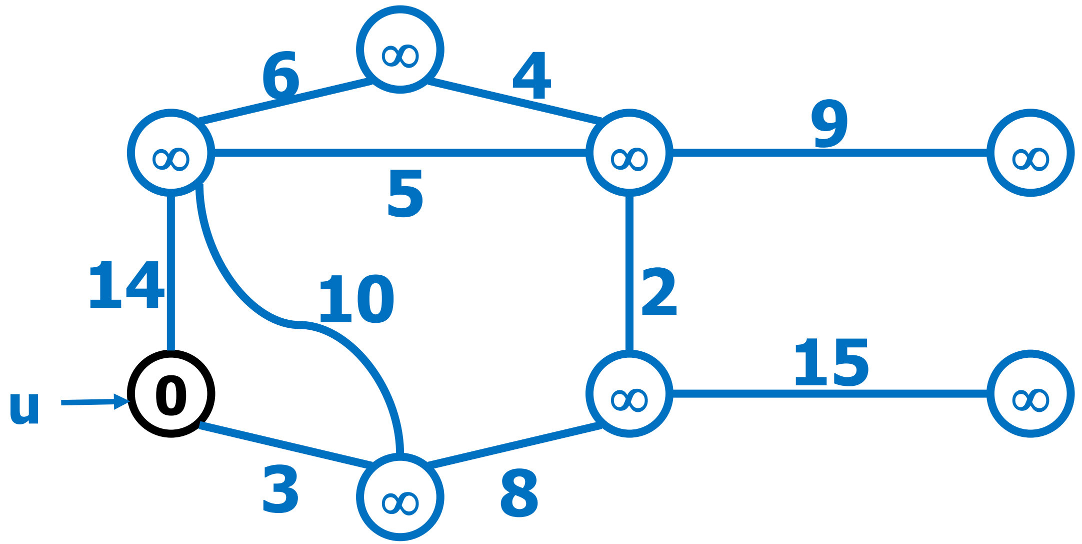

Black vertices have been removed from Q

Prim's Algorithm – Example

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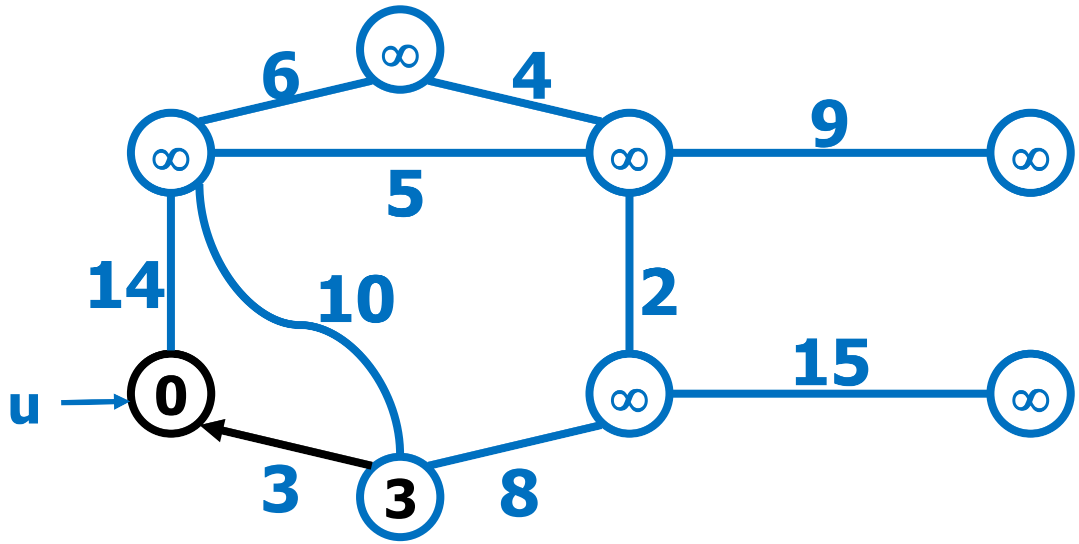

Black arrows indicate parent pointers

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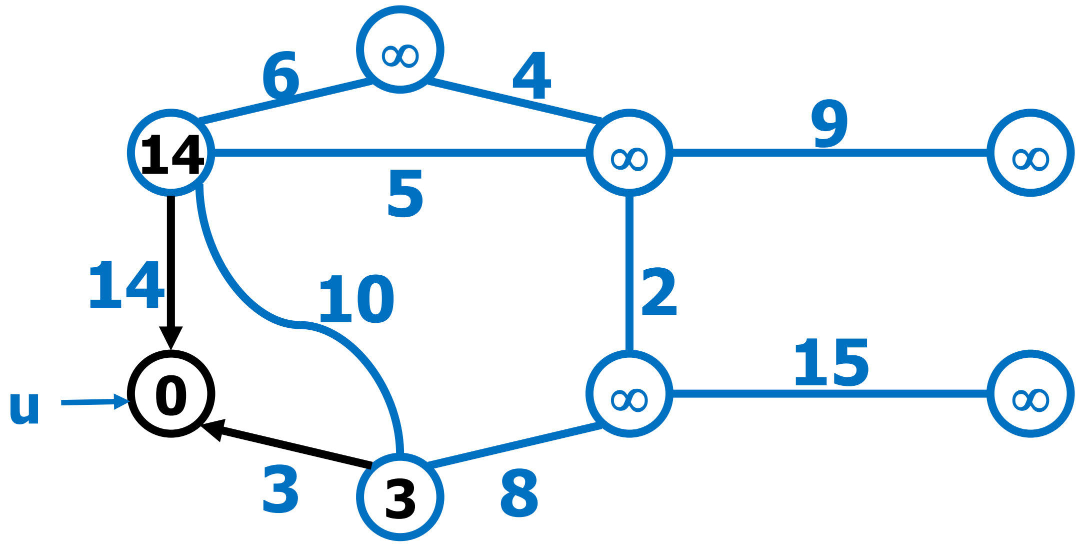

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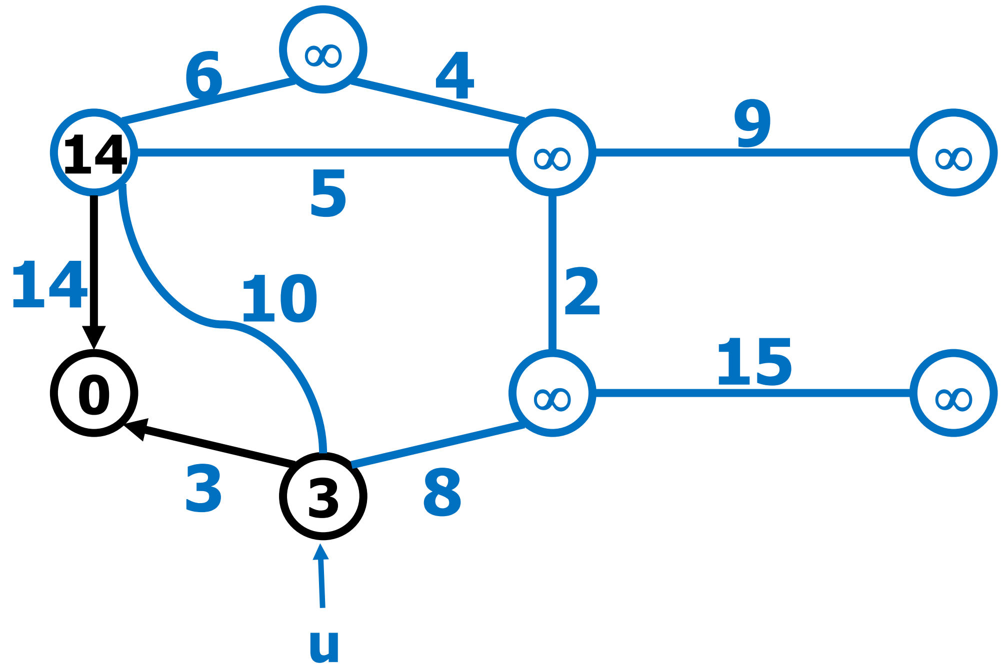

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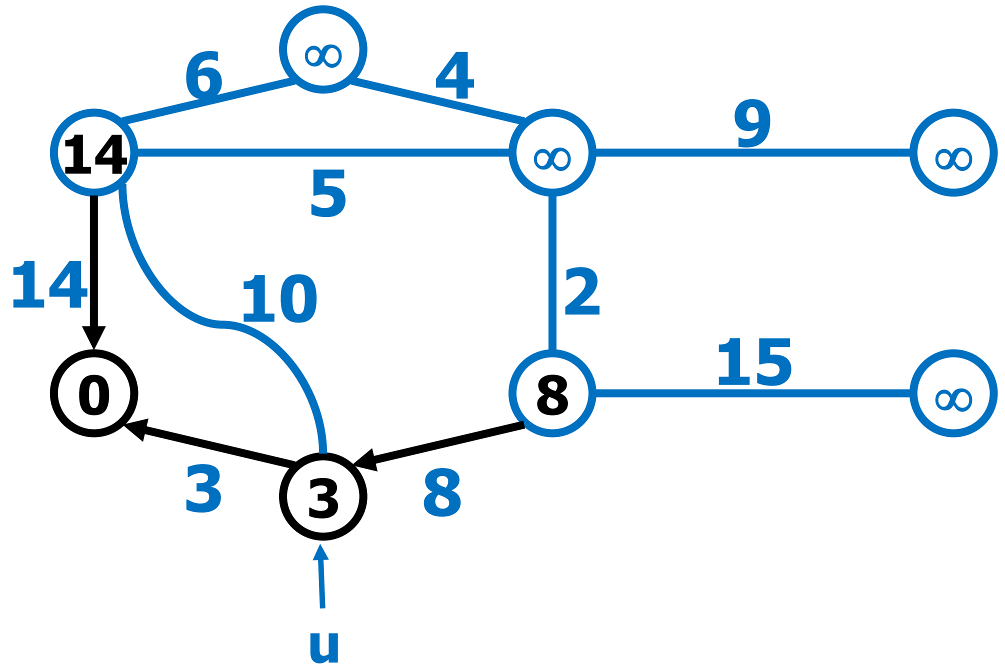

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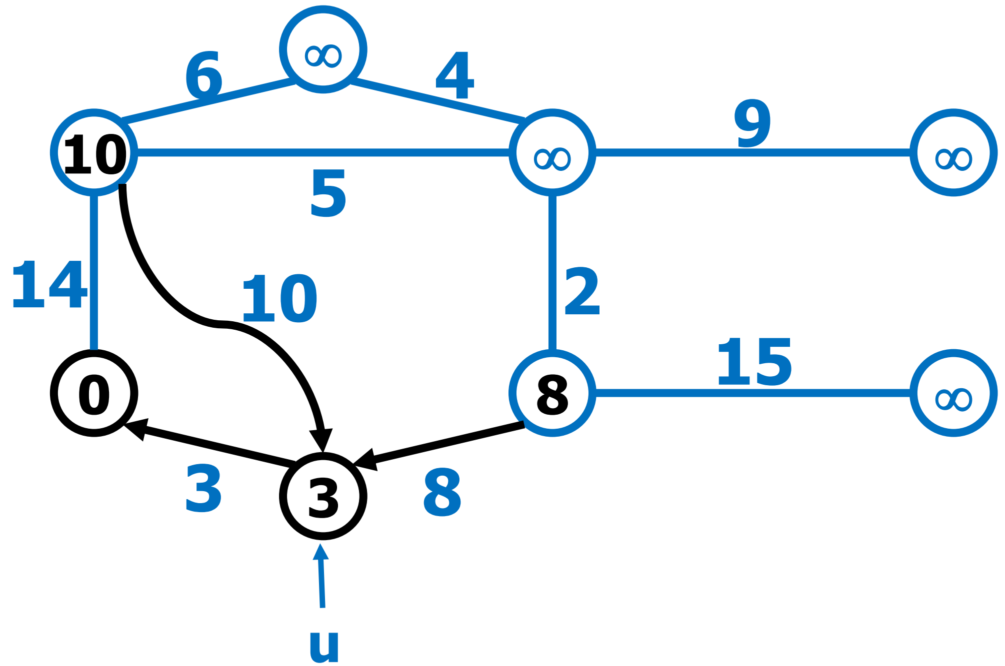

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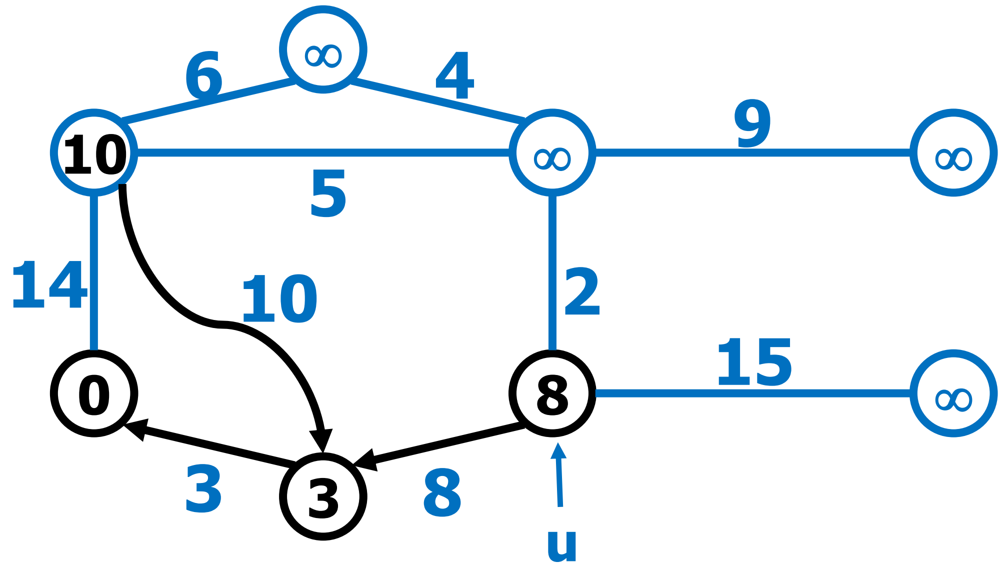

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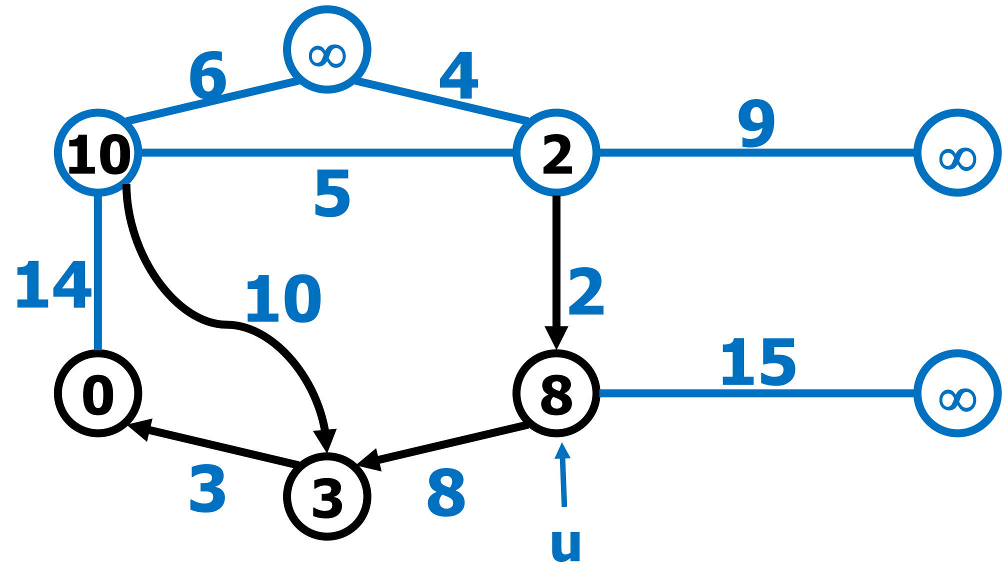

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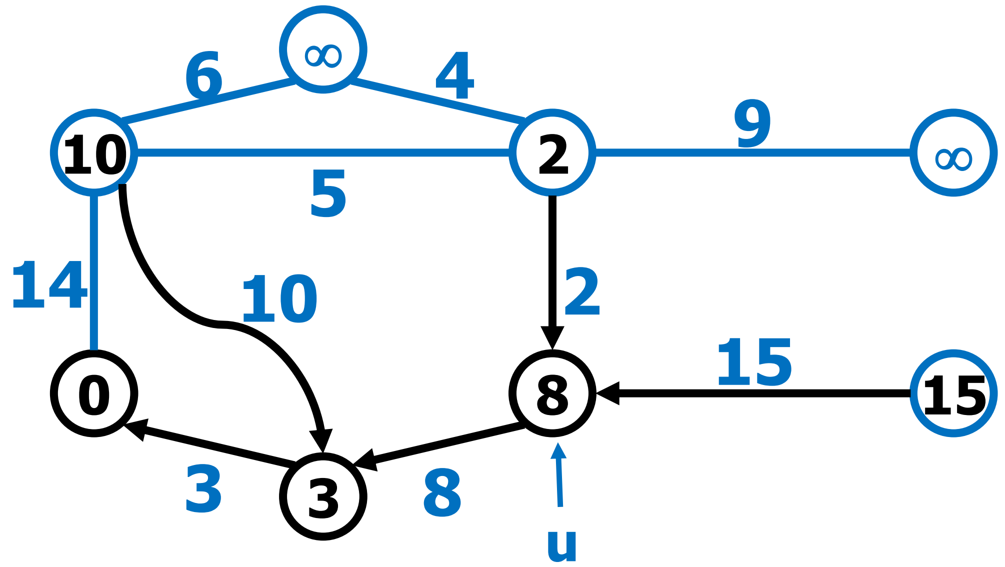

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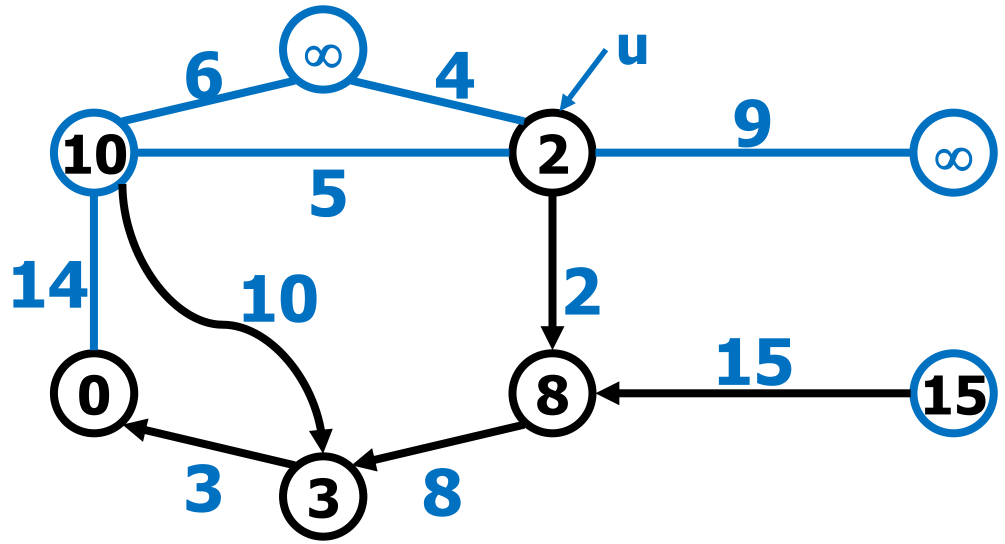

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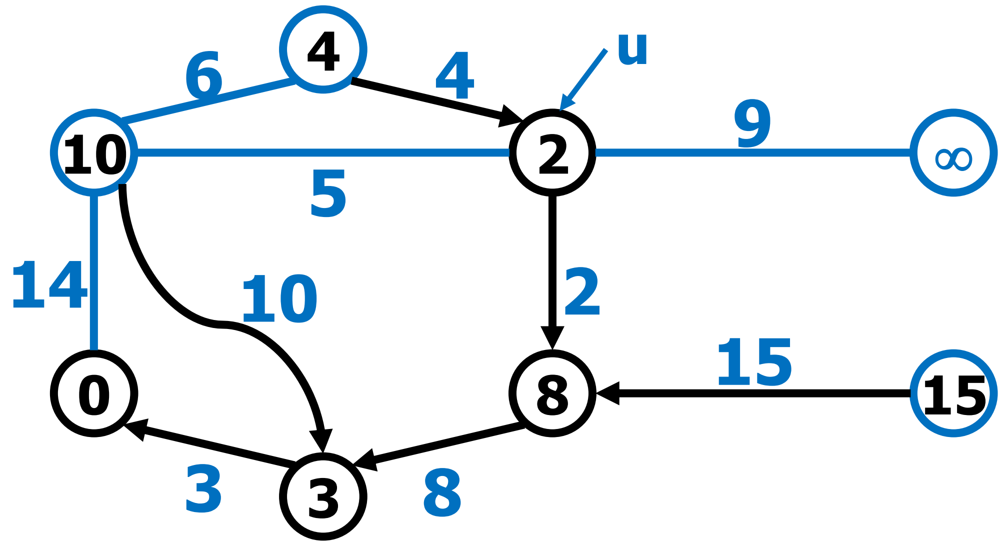

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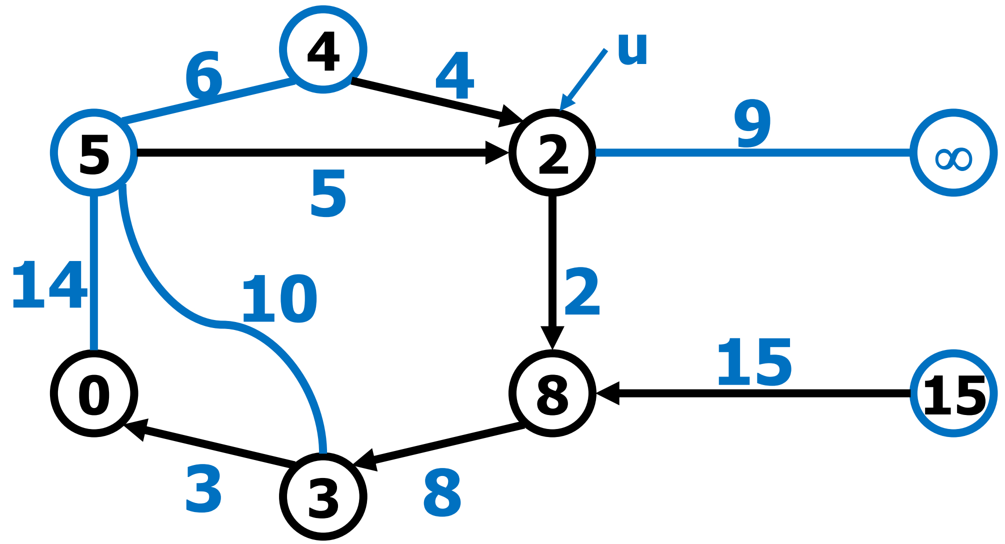

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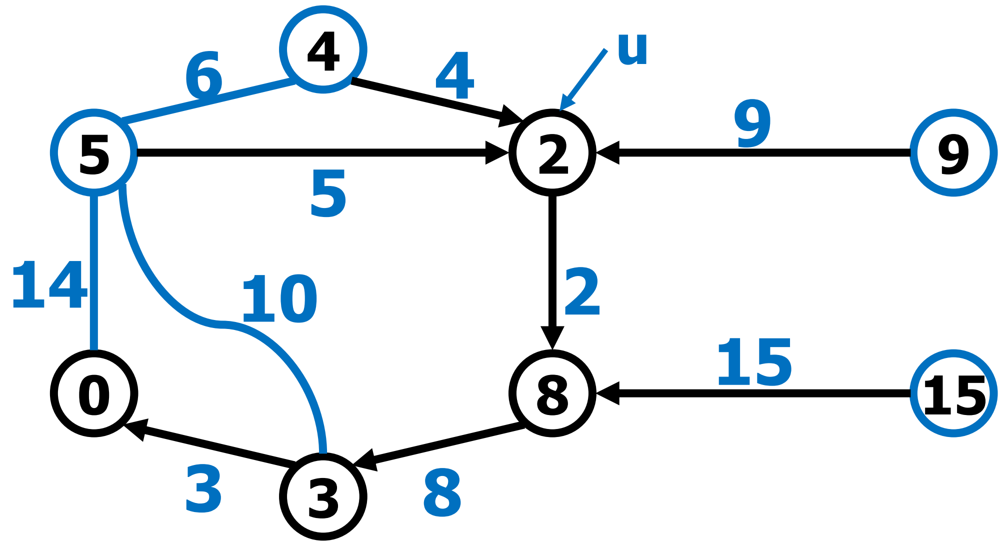

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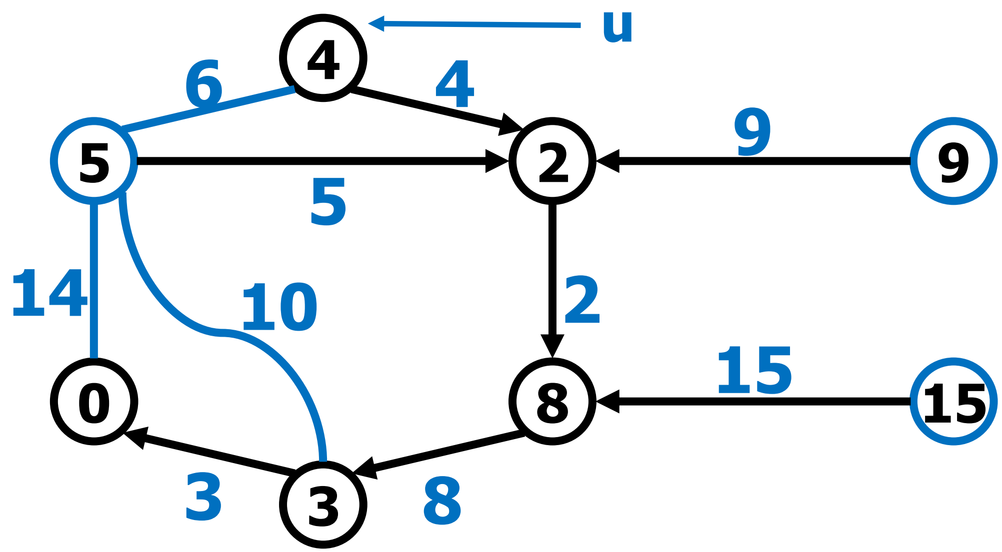

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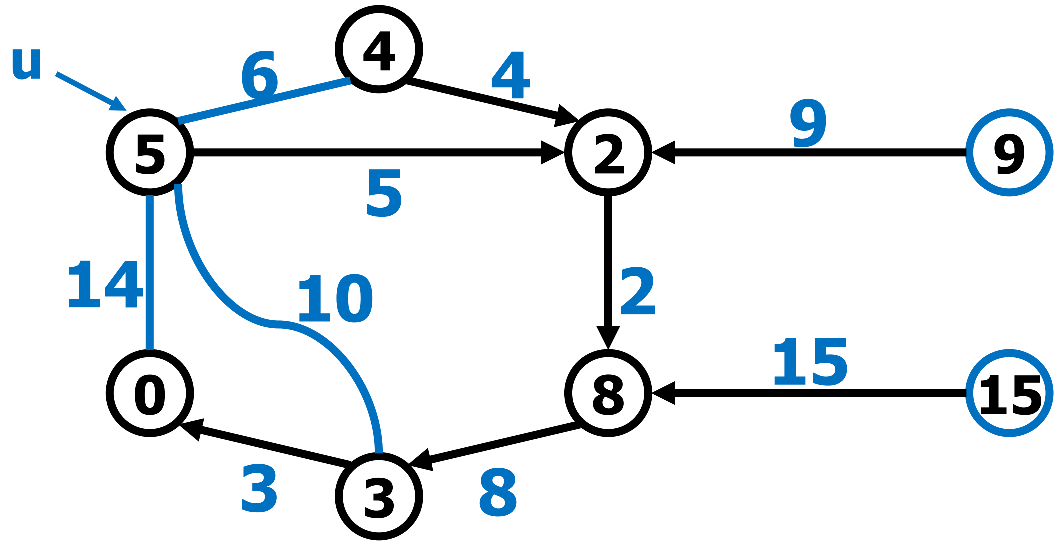

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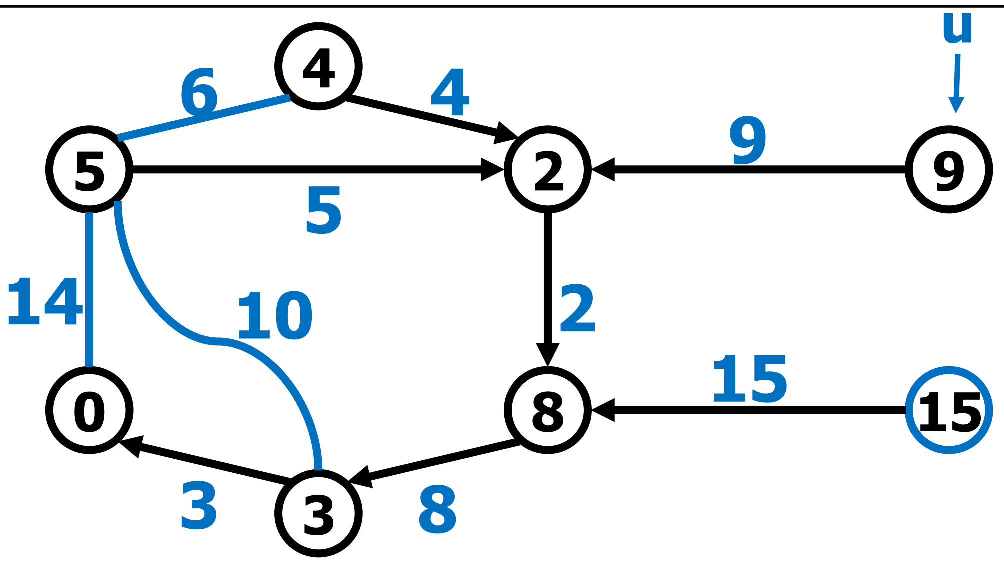

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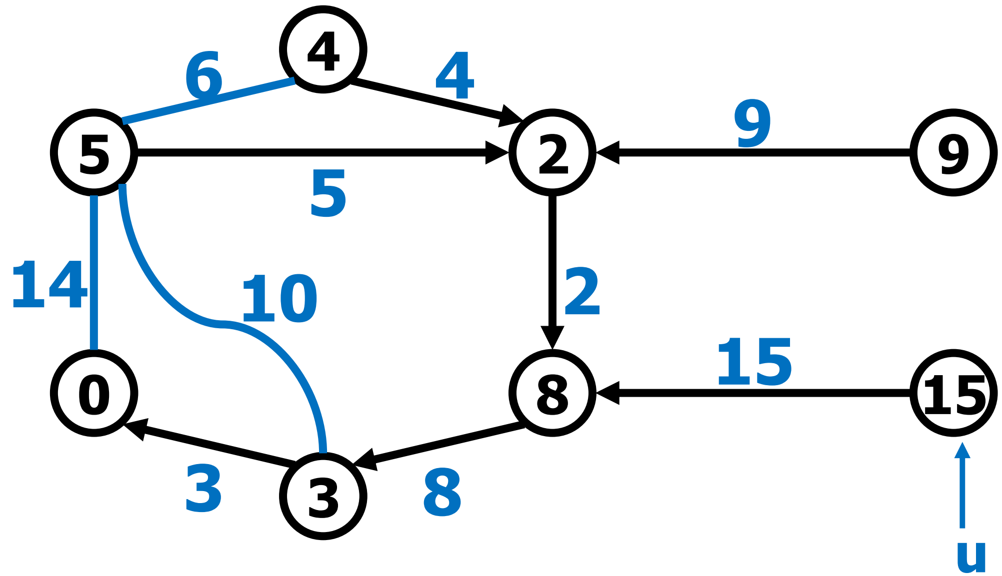

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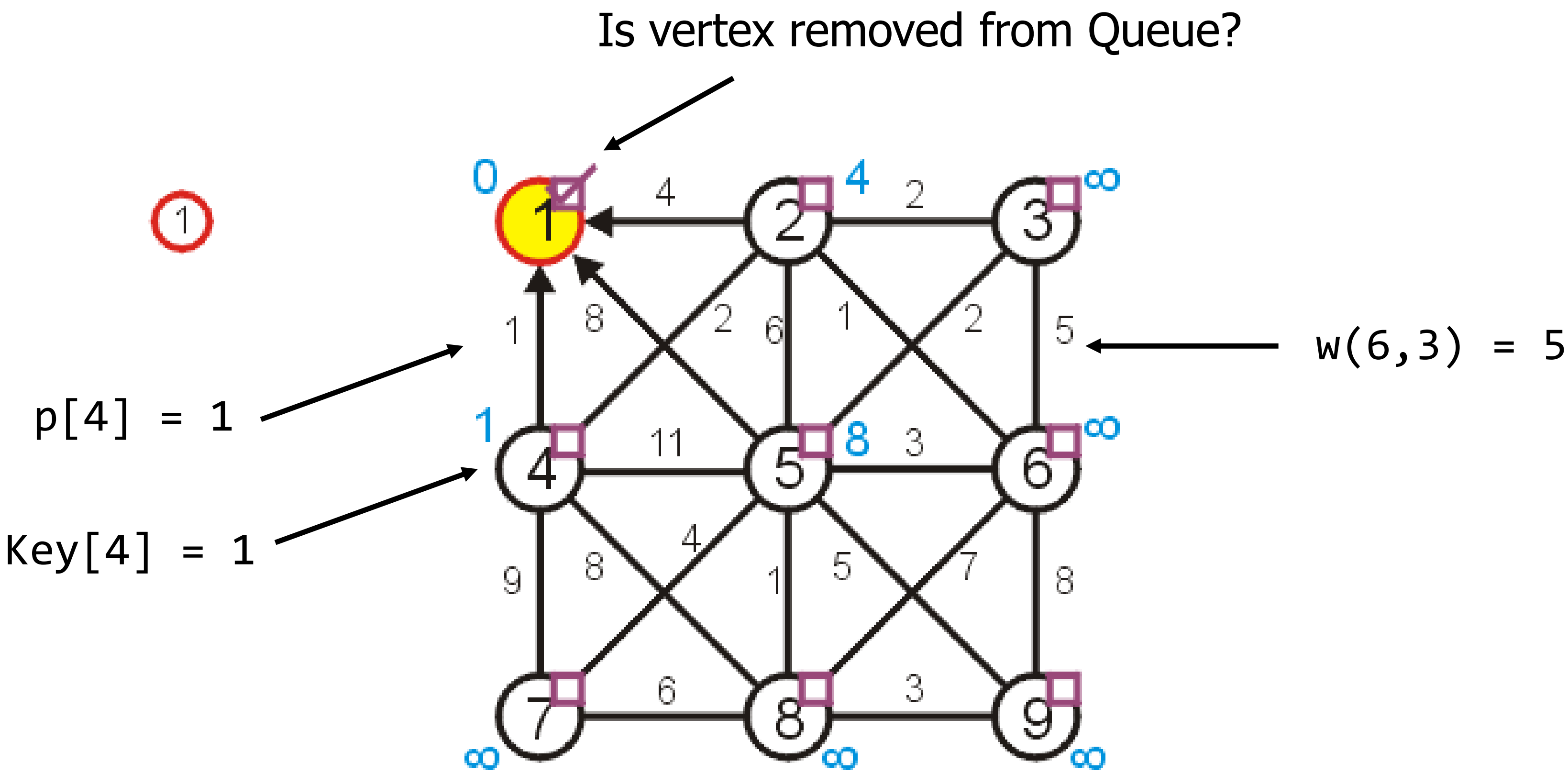

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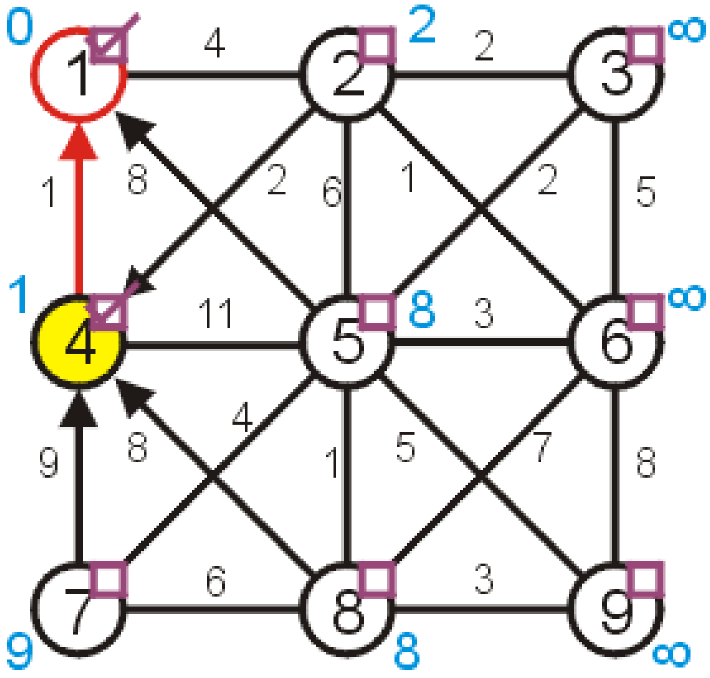

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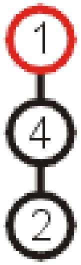
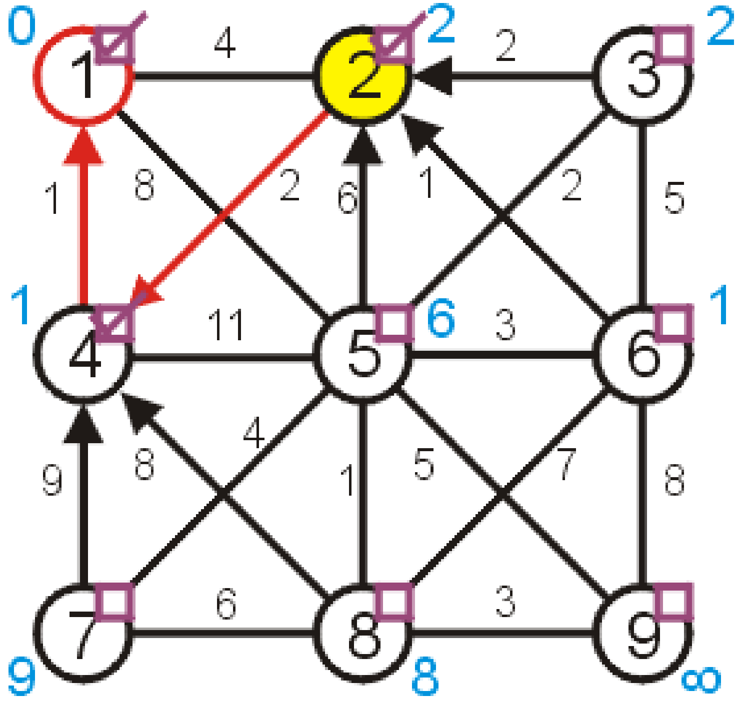
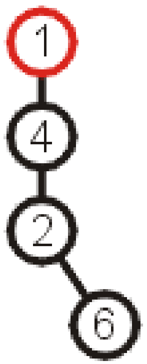
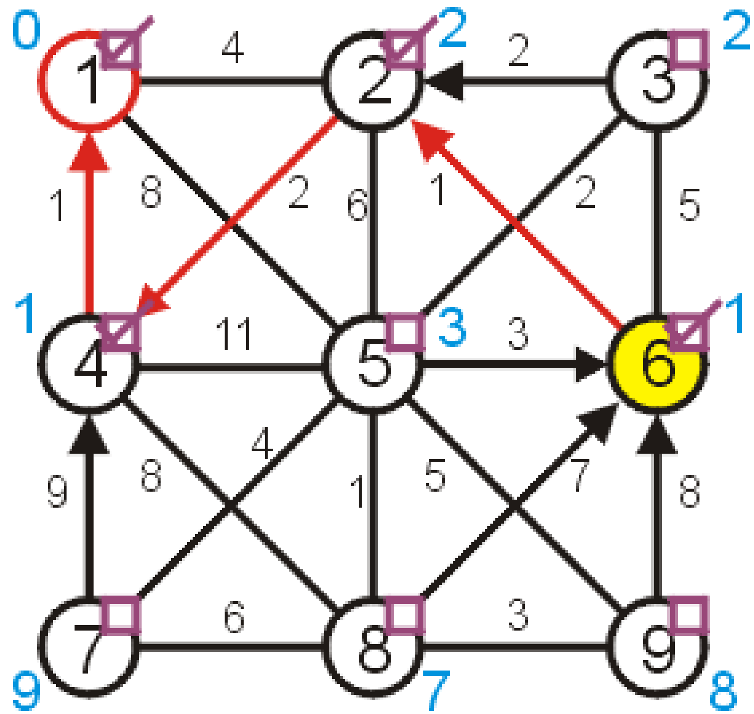


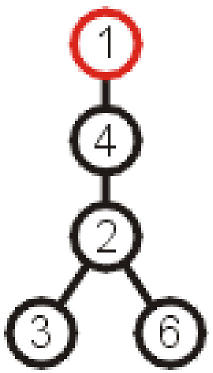


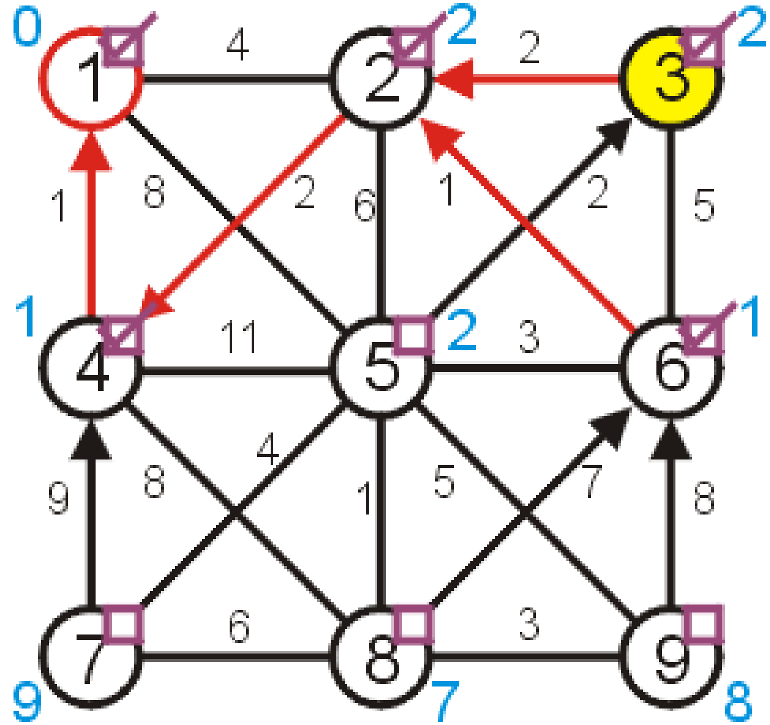
Prim's Algorithm – Example



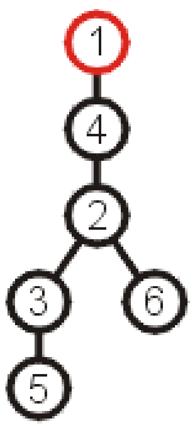


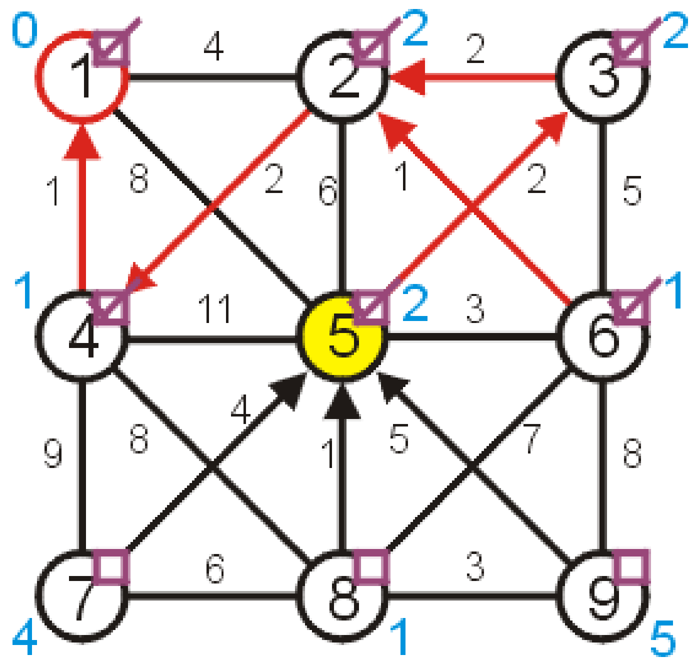
Prim’s Algorithm – Example



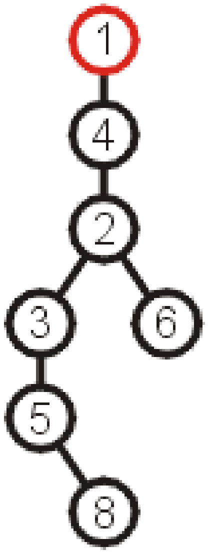


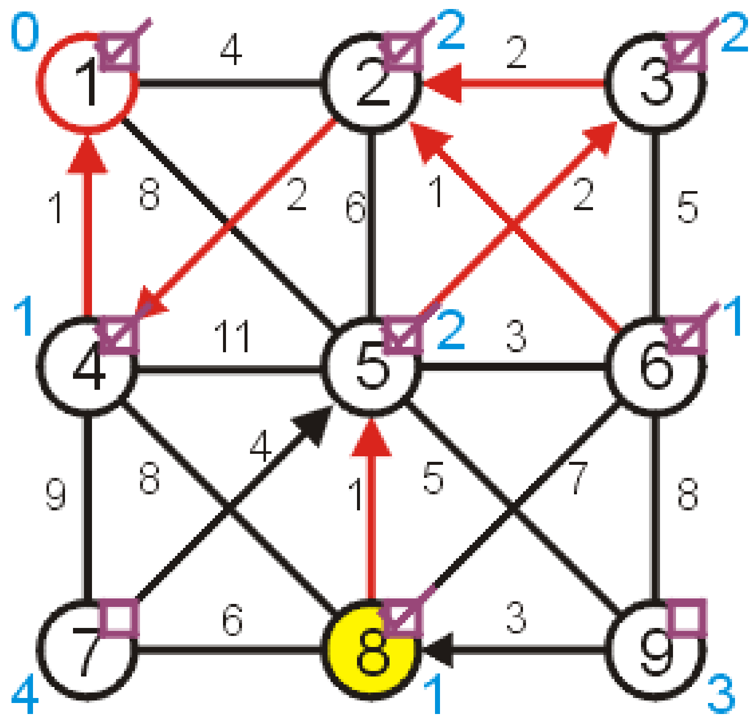
Prim’s Algorithm – Example



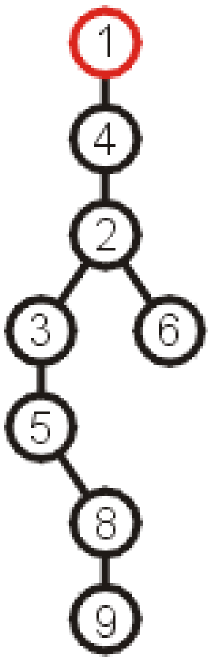


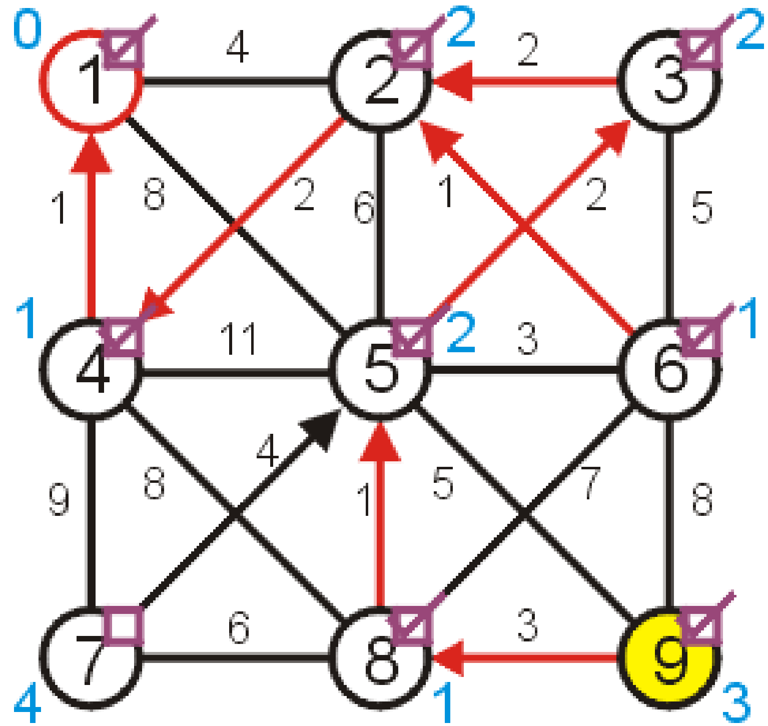
Prim's Algorithm – Example



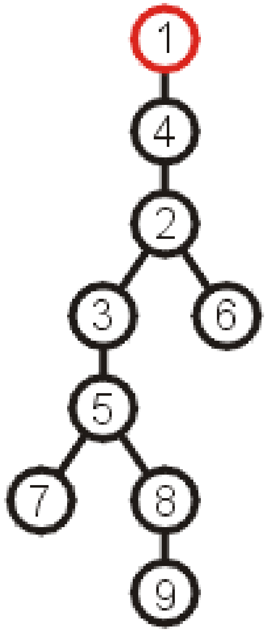


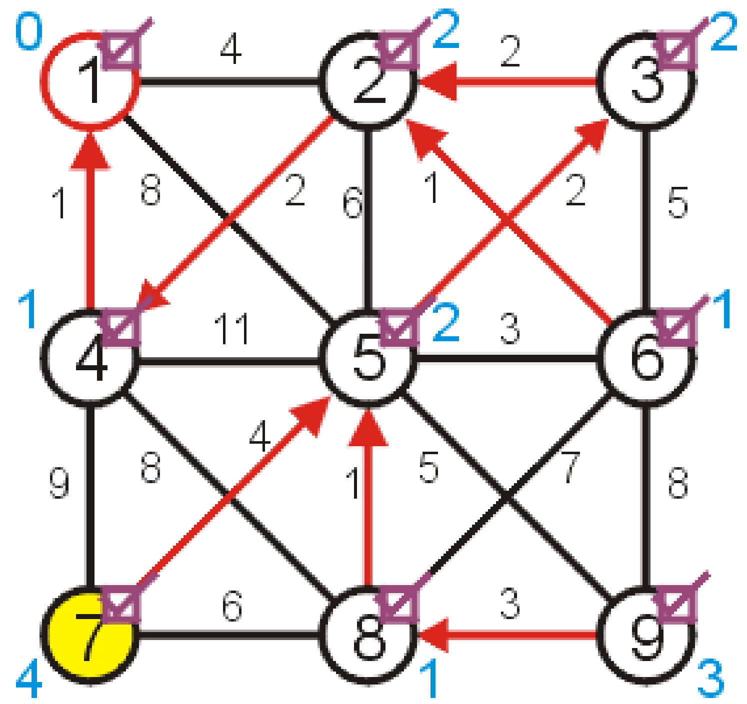
Prim’s Algorithm – Example



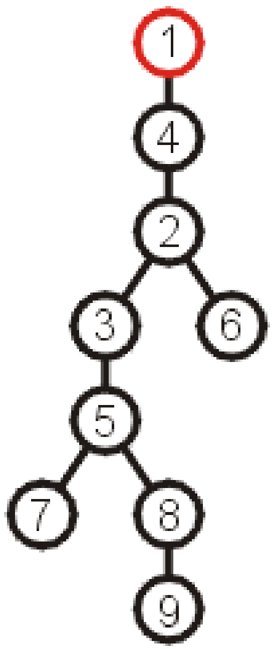


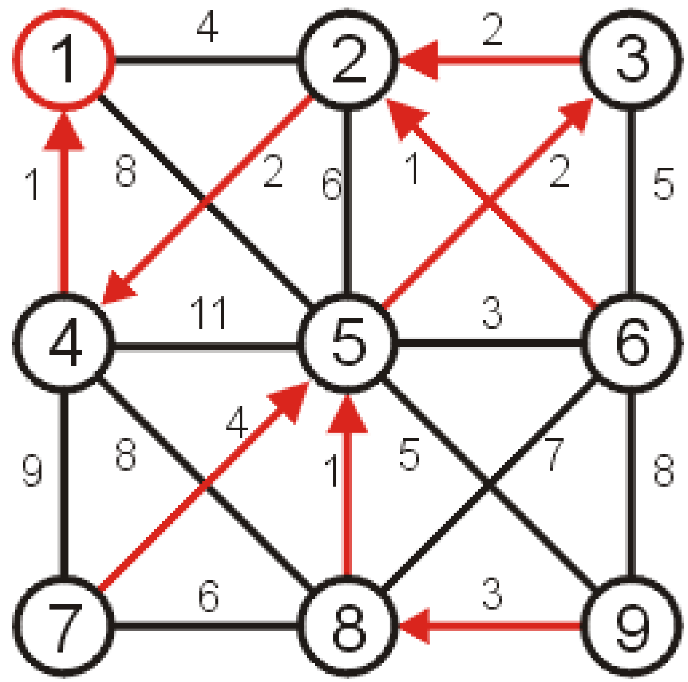
Prim's Algorithm – Example





Prim’s Algorithm – Example





Kruskal’s Algorithm vs Prim’s Algorithm

- In prim’s algorithm, graph must be connected
- Kruskal’s algorithm can function on disconnected graphs too
- Prim's algorithm is significantly faster for dense graphs with more number of edges than vertices
  - No pre-processing required
- Kruskal's algorithm runs faster in the case of sparse graphs
  - N^2 cost for pre-processing (sorting)

Any Question So Far?