

AMPL

A Mathematical Programming Language

The Diet Problem

Item / Nutrient	Protein	Vitamin C	Price
1. Hamburger	28	10	8
2. Sausage	17	0	4
3. Fries	3	15	2
4. Orange juice	1	120	3
Required amount	20	25	

$\min 8x_1 + 4x_2 + 2x_3 + 3x_4$
subject to
 $28x_1 + 17x_2 + 3x_3 + 1x_4 \geq 20$
 $10x_1 + 0x_2 + 15x_3 + 120x_4 \geq 25$
 $x_1, x_2, x_3, x_4 \in \{0,1\}$

```
var x1 binary;  
var x2 binary;  
var x3 binary;  
var x4 binary;  
  
minimize TotalCost: 8 * x1 + 4 * x2 + 2 * x3 + 3 * x4;  
  
subject to Protein: 28 * x1 + 17 * x2 + 3 * x3 + 1 * x4 >= 20;  
subject to VitaminC: 10 * x1 + 0 * x2 + 15 * x3 + 120 * x4 >= 25;
```

The Diet Problem

Item / Nutrient	Protein	Vitamin C	Price
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$$28x_1 + 17x_2 + 3x_3 + 1x_4 \geq 20$$
$$10x_1 + 0x_2 + 15x_3 + 120x_4 \geq 25$$
$$x_1, x_2, x_3, x_4 \in \{0,1\}$$

```
C:\Users\Desktop\Desktop\ampltutorial\ampl_mswin64\ampl.exe
ampl: option solver cplex;
ampl:
ampl: model C:\Users\Desktop\Desktop\ampltutorial\Diet1.mod;
ampl:
ampl: solve;
CPLEX 12.6.3.0: optimal integer solution; objective 9
3 MIP simplex iterations
0 branch-and-bound nodes
No basis.
ampl:
ampl: display x1, x2, x3, x4;
x1 = 0
x2 = 1
x3 = 1
x4 = 1
ampl:
```

The Diet Problem

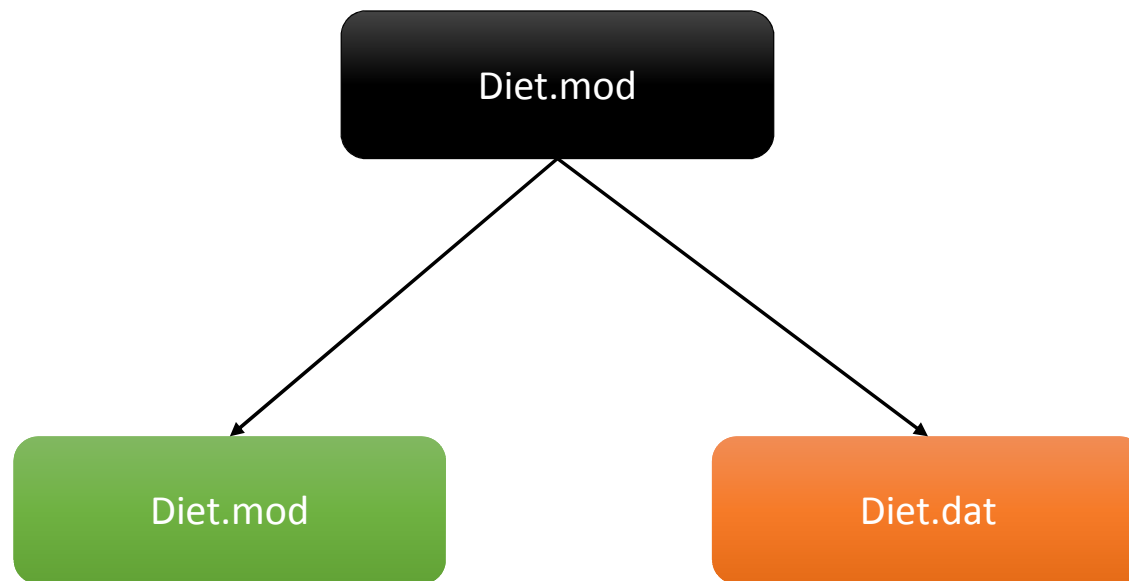
Item / Nutrient	Protein	Vitamin C	Price
1. Hamburger	28	10	8
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 $x_1, x_2, x_3, x_4 \in \{0,1\}$

```
var x1 binary;  
var x2 binary;  
var x3 binary;  
var x4 binary;  
  
minimize TotalCost: 8 * x1 + 4 * x2 + 2 * x3 + 3 * x4;  
  
subject to Protein: 28 * x1 + 17 * x2 + 3 * x3 + 1 * x4 >= 20;  
subject to VitaminC: 10 * x1 + 0 * x2 + 15 * x3 + 120 * x4 >= 25;
```

Suppose that you consider tens of nutrients instead of just protein and vitamin C;
and suppose also that the menu includes tens of items instead of just four.

The Diet Problem



Elements of a Mathematical Programming Model

1. Sets
2. Parameters
3. Variables
4. Constraints
5. Objective Function

1. Sets

In AMPL, **set** is a keyword used to define finite collections of elements.

These elements can be numerical or non-numerical, and frequently are used as **indices** (of parameters and variables).

```
set S := 1..5;
```

```
set S := { 1, 2, 3, 10 };
```

```
set S := { 'a', 'b', 'c', 'd' };
```

```
set S;
```

```
param n;
```

```
set S := 1..n;
```

2. Parameters

Parameters are the given constants of your mathematical programming model.
AMPL keyword for parameters is **param**.

```
param budget;  
param capacity := 500;
```

```
# p[j] is the processing time of job j  
set J; #jobs  
param p {J};
```

```
# p[i,j] is the processing time of job j  
# on machine i  
set I; #machines  
set J; #jobs  
param p {I, J};
```


2. Parameters

Data Validation: data types

```
# g[p] is the number of goals scored by  
# player p  
set P; #players  
param g {P} integer;
```

```
# a[i,j] equals 1 if there exists a link  
# between nodes i and j; 0 otherwise  
set N; #nodes  
param a {N,N} binary;
```

2. Parameters

Data Validation: bounds

```
# p[i,j] is the processing time of job j
# on machine i
set I; #machines
set J; #jobs
param p {I, J} >= 0, <= 120;
```

```
# b[i] is the budget of project i.
set I; #projects
param u {I}; #upper bound on project budget
param b {i in I} >= 0, <= u[i];
```

3. Variables

AMPL's keyword for decision variables is **var**.

Declaration of variables is only slightly different than that of the parameters.

3. Variables

Continuous Decision Variables

```
# x is the deviation of a certain measure from a
# constant value
var x;
```



```
# x[j] is the amount of chemical j used in a chemical
# mixture
set J; #chemicals
var x{J} >= 0;
```



```
# we want to ensure that at least 10mg of each
# chemical is used; and no chemical is used more than
# 100mg
set J; #chemicals
var x{J} >= 10, <= 100;
```



```
# or we want to restrict the use each chemical with a
# lower and upper bound specific to that chemical
set J; #chemicals
param lower{J};
param upper{J};
var x{j in J} >= lower[j], <= upper[j];
```

3. Variables

Discrete Decision Variables

```
# x[i] is the number of workers assigned to workstation i
set I; #workstations
var x{I} integer;
```

```
# actually, we know that it is nonnegative
set I; #workstations
var x{I} integer >= 0;
```

```
# y[j,k] equals 1 if job j precedes job k; and 0 otherwise
set J; #jobs
var y{J,J} integer >=0, <=1;
```

```
# or simply
set J; #jobs
var y{J,J} binary;
```

4. Constraints

Constraint declaration in AMPL has the following syntax:

subject to constraintName {forall}: constraintExpression;

Let x_i be the amount of money we invest on project i .

We have a budget of b ; and hence, sum of our investments cannot exceed b .

$$\sum_{i \in I} x_i \leq b$$

```
param b;  
set I;  
var x{I} >= 0;  
subject to BudgetConstraint: sum{i in I} x[i] <= b;
```

4. Constraints

Consider a network flow problem.

We want to ensure that the difference between the flow leaving a node and the flow entering a node should be equal to the supply of that node.

Let x_{ij} be the amount of flow from node i to node j ; and let b_i be the supply of node i .

$$\sum_{j \in N} x_{ij} - \sum_{k \in N} x_{ki} = b_i \quad , \forall i \in N$$

```
set N;  
param b{N};  
var x{N,N} >= 0;  
subject to FlowBalance {i in N}: sum{j in N} x[i,j] - sum{k in N} x[k,i] = b[i];
```

4. Constraints

Let c_j be the completion time of job j .

And let x_{jk} equal 1 if job j is processed before job k ; and 0 otherwise.

Let parameter M be the big-M; and let parameter p_j be the processing time of job j .

We want to write the big-M constraints to avoid overlapping of tasks in a single machine scheduling problem.

$$M x_{jk} + c_j \geq c_k + p_j \quad ; \forall j \in J, k \in J$$

$$M (1 - x_{jk}) + c_k \geq c_j + p_k \quad ; \forall j \in J, k \in J$$

```
set J;  
param p{J};  
param M;  
var x{J,J} binary;  
var c{J} >= 0;  
subject to Overlap1 {j in J, k in J}: M * x[j,k] + c[j] >= c[k] + p[j];  
subject to Overlap2 {j in J, k in J}: M * (1-x[j,k]) + c[k] >= c[j] + p[k];
```


4. Constraints

Let c_j be the completion time of job j .

And let x_{jk} equal 1 if job j is processed before job k ; and 0 otherwise.

Let parameter M be the big-M; and let parameter p_j be the processing time of job j .

We want to write the big-M constraints to avoid overlapping of tasks in a single machine scheduling problem.

$$\begin{aligned} M x_{jk} + c_j &\geq c_k + p_j && ; \forall j \in J, k \in J \ni k > j \\ M (1 - x_{jk}) + c_k &\geq c_j + p_k && ; \forall j \in J, k \in J \ni k > j \end{aligned}$$

```
set J;  
param p{J};  
param M;  
var x{J,J} binary;  
var c{J} >= 0;  
subject to Overlap1 {j in J, k in J: k > j}: M * x[j,k] + c[j] >= c[k] + p[j];  
subject to Overlap2 {j in J, k in J: k > j}: M * (1-x[j,k]) + c[k] >= c[j] + p[k];
```

4. Constraints

Consider the same scheduling problem; however, this time we want to write the precedence constraints. We are given the parameter $prec_{jk}$ which is equal to 1 if job j has to be processed before job k ; and 0 otherwise.

$$c_k \geq c_j + p_k \quad ; \forall j \in J, k \in J \ni prec_{jk} = 1$$

```
set J;  
param p{J};  
param prec{J,J};  
var c{J} >= 0;  
subject to Precedence {j in J, k in J: prec[j,k]=1}: c[k] >= c[j] + p[k];
```

5. Objective Function

Objective function declaration in AMPL has the following syntax:

minimize/maximize **objectiveName**: **objectiveFunction**;

Let decision variable x_{ij} be the amount of flow from node i to node j .

Let parameter c_{ij} be the cost of unit flow from node i to node j .

$$\min \sum_{i \in N} \sum_{j \in N} c_{ij} x_{ij}$$

```
set N;  
param c{N,N};  
var x{N,N} >= 0;  
minimize TotalCost: sum{i in N, j in N} x[i,j] * c[i,j];
```

Exercise: The Diet Problem

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$$x_1, x_2, x_3, x_4 \in \{0,1\}$$

```
var x1 binary;  
var x2 binary;  
var x3 binary;  
var x4 binary;  
  
minimize TotalCost: 8 * x1 + 4 * x2 + 2 * x3 + 3 * x4;  
  
subject to Protein: 28 * x1 + 17 * x2 + 3 * x3 + 1 * x4 >= 20;  
subject to VitaminC: 10 * x1 + 0 * x2 + 15 * x3 + 120 * x4 >= 25;
```

Define the sets, parameters, variables, constraints and the objective function of the diet problem.

Your model file should not include any data.

Exercise: The Diet Problem: Diet2.mod

```
set N;                                #nutrients
set F;                                #foods
param p{F};                           #food price
param r{N};                           #required nutrient amount
param a{F,N};                         #amount of nutrients in foods
var b{F} binary;                      #binary buy/don't buy variable

subject to Nutrient {n in N}: sum{f in F} a[f,n] * b[f] >= r[n];

minimize TotalCost: sum{f in F} p[f] * b[f];
```

The data file: Diet1.dat

```
set F :=  
  Hamburger  
  Sausage  
  Fries  
  "Orange Juice";
```

```
set N :=  
  Protein  
  VitaminC;
```

```
param p :=  
  Hamburger      8  
  Sausage        4  
  Fries          2  
  "Orange Juice" 3;
```

```
param r :=  
  Protein    20  
  VitaminC   25;
```

```
param a:  
      | | | | | | | Protein VitaminC :=  
  Hamburger      28      10  
  Sausage        17       0  
  Fries          3       15  
  "Orange Juice" 1      120;
```

The Diet Solution: Diet1.dat

```
C:\Users\Desktop\Desktop\ampltutorial\ampl_mswin64\ampl.exe
ampl: option solver cplex;
ampl:
ampl: model C:\Users\Desktop\Desktop\ampltutorial\Diet2.mod;
ampl:
ampl: data C:\Users\Desktop\Desktop\ampltutorial\Diet1.dat;
ampl:
ampl: solve;
CPLEX 12.6.3.0: optimal integer solution; objective 9
3 MIP simplex iterations
0 branch-and-bound nodes
No basis.
ampl:
ampl: display b;
b [*] :=
      Fries 1
      Hamburger 0
      'Orange Juice' 1
      Sausage 1
;
ampl: _
```

The Diet Solution: Diet2.dat

```
C:\Users\Desktop\Desktop\ampltutorial\ampl_mswin64\ampl.exe
ampl: option solver cplex;
ampl:
ampl: model C:\Users\Desktop\Desktop\ampltutorial\Diet2.mod;
ampl:
ampl: data C:\Users\Desktop\Desktop\ampltutorial\Diet2.dat;
ampl:
ampl: solve;
CPLEX 12.6.3.0: optimal integer solution; objective 6.66
17 MIP simplex iterations
0 branch-and-bound nodes
No basis.
ampl:
ampl: display b;
b [*] :=
      '1% Lowfat Milk' 1
      'Big Mac' 1
      'Filet-O-Fish' 1
      'Fries, small' 1
      'McGrilled Chicken' 0
      'McLean Deluxe w/ Cheese' 0
      'Orange Juice' 1
      'Quarter Pounder w/ Cheese' 0
      'Sausage McMuffin' 1
;
ampl:
```


Exercise: $1|r_j|\sum w_j C_j$ with time-indexed variables

$$\min \sum_{j \in J} \sum_{t \in T} w_j (t + p_j) x_{jt}$$

subject to

$$\sum_{t \in T} x_{jt} = 1 \quad , \forall j \in J$$

$$\sum_{j \in J} \sum_{\substack{s \in T \ni \\ s \geq \max(0, t+1-p_j) \text{ and } s \leq t}} x_{js} \leq 1 \quad , \forall t \in T \ni t \leq |T| - 1$$

$$x_{jt} = 0 \quad , \forall j \in J, t \in T \ni t < r_j$$

$$x_{jt} \in \{0,1\} \quad , \forall j \in J, t \in T$$

Exercise: $1|r_j|\sum w_j C_j$ with time-indexed variables

```
param maxT;  
set J; #jobs  
set T := 0..maxT; #time intervals  
  
param p{J};  
param r{J};  
param w{J};  
var x{J,T} binary;  
  
minimize WeightedCompTime:  
    sum{j in J, t in T} w[j] * (t + p[j]) * x[j,t];  
  
subject to C1 {j in J}: sum{t in T} x[j,t] = 1;  
  
subject to C2 {t in T: t <= card(T) - 1}:  
    sum{j in J, s in T: s >= max(0, t + 1 - p[j]) && s <= t} x[j,s] <= 1;  
  
subject to C3 {j in J, t in T: t < r[j]}: x[j,t] = 0;
```

Exercise: $1|r_j|\sum w_j C_j$ with time-indexed variables

Solve the following three instances:

- Sch1.dat (7 jobs)
- Sch2.dat (35 jobs)
- Sch3.dat (100 jobs)

with the model:

- Sch.mod.

Exercise: $1|r_j|\sum w_j C_j$ with time-indexed variables (7 jobs)

```
C:\Users\Desktop\Desktop\ampltutorial\ampl_mswin64\ampl.exe
ampl: option solver cplex;
ampl:
ampl: model C:\Users\Desktop\Desktop\ampltutorial\Sch.mod;
ampl:
ampl: data C:\Users\Desktop\Desktop\ampltutorial\Sch1.dat;
ampl:
ampl: solve;
CPLEX 12.6.3.0: optimal integer solution; objective 3771
89 MIP simplex iterations
0 branch-and-bound nodes
No basis.
ampl:
ampl: display {j in J}: sum{t in T} t*x[j,t];
sum{t in T} t*x[1,t] = 150

sum{t in T} t*x[2,t] = 0

sum{t in T} t*x[3,t] = 18

sum{t in T} t*x[4,t] = 73

sum{t in T} t*x[5,t] = 47

sum{t in T} t*x[6,t] = 30

sum{t in T} t*x[7,t] = 57
ampl:
```

Exercise: $1|r_j|\sum w_j C_j$ with time-indexed variables (35 jobs)

```
C:\Users\Desktop\Desktop\ampltutorial\ampl_mswin64\ampl.exe
ampl: option solver cplex;
ampl: model C:\Users\Desktop\Desktop\ampltutorial\Sch.mod;
ampl: data C:\Users\Desktop\Desktop\ampltutorial\Sch2.dat;
ampl:
ampl: option cplex_options 'mipdisplay 2';
ampl:
ampl: solve;
CPLEX 12.6.3.0: mipdisplay 2
Found incumbent of value 71661.960000 after 0.01 sec. (7.35 ticks)
MIP Presolve eliminated 3 rows and 13 columns.
MIP Presolve modified 1 coefficients.
Reduced MIP has 1033 rows, 28011 columns, and 615958 nonzeros.
Reduced MIP has 28011 binaries, 0 generals, 0 SOSs, and 0 indicators.
MIP Presolve eliminated 3 rows and 13 columns.
MIP Presolve modified 1 coefficients.
Reduced MIP has 1030 rows, 27998 columns, and 615700 nonzeros.
Reduced MIP has 27998 binaries, 0 generals, 0 SOSs, and 0 indicators.
Probing time = 0.06 sec. (14.14 ticks)
Clique table members: 1030.
MIP emphasis: balance optimality and feasibility.
MIP search method: dynamic search.
Parallel mode: deterministic, using up to 16 threads.
Root relaxation solution time = 0.63 sec. (691.64 ticks)
```

	Nodes		Objective	IInf	Best Integer	Cuts/ Best Bound	ItCnt	Gap
	Node	Left						
*	0+	0			71661.9600	0.0000		100.00%
	0	0	54515.0390	191	71661.9600	54515.0390	37	23.93%
*	0+	0			60129.2300	54515.0390		9.34%
	0	0	54587.7261	251	60129.2300	Cuts: 28	524	9.22%
	0	0	54598.5982	206	60129.2300	Cuts: 13	710	9.20%
*	0+	0			55316.8100	54598.5982		1.30%
	0	0	54616.4704	104	55316.8100	Clques: 16	874	1.27%

```
C:\Users\Desktop\Desktop\ampltutorial\ampl_mswin64\ampl.exe
ampl: display {j in J, t in T: x[j,t] = 1} x[j,t];
x[j,t] [+,+]
: 0 4 9 18 24 47 79 95 115 125 141 171 182 217 236 245 266 294 326 :=
2 . . . . . . . . . . . 1 . . . . .
4 . . . . . . . . . . . . . . . 1 . .
6 . . . . . . . . . . . . . . . . .
7 . . . . . . . . . . . . . . . . .
8 . . . 1 . . . . . . . . . . . . .
9 . 1 . . . . . . . . . . . . . . .
11 . . . . . . . . . 1 . . . . . . .
14 1 . . . . . . . . . . . . . . . .
17 . . . . . . . . . . . . . . . 1 .
18 . . . . . . . . . 1 . . . . . . .
20 . . . . . . 1 . . . . . . . . . .
21 . . . . . 1 . . . . . . . . . . .
23 . . 1 . . . . . . . . . . . . . .
25 . . . . 1 . . . . . . . . . . . .
27 . . . . . . . . . . . . . . . . 1
31 . . . . . . . . . . . 1 . . . . .
33 . . . . . . 1 . . . . . . . . . .
34 . . . . . . . . . . . . . . 1 . .
35 . . . . . . . . . . 1 . . . . . .
: 351 380 405 437 443 472 504 526 552 577 609 642 677 705 731 760 :=
1 . . . . . . . . 1 . . . . . . .
3 . . . . . . . . . . . 1 . . . .
5 . . . 1 . . . . . . . . . . . .
```

Exercise: $1|r_j|\sum w_j C_j$ with time-indexed variables (100 jobs)

```
C:\Users\Desktop\Desktop\ampltutorial\ampl_mswin64\ampl.exe
ampl: option solver cplex;
ampl: option cplex_options 'mipdisplay 2';
ampl:
ampl: option show_stats 1;
ampl:
ampl: model C:\Users\Desktop\Desktop\ampltutorial\Sch.mod;
ampl: data C:\Users\Desktop\Desktop\ampltutorial\Sch3.dat;
ampl:
ampl: solve;

Presolve eliminates 64318 constraints and 64318 variables.
Adjusted problem:
185782 variables, all binary
2601 constraints, all linear; 4097778 nonzeros
    100 equality constraints
    2501 inequality constraints
1 linear objective; 185782 nonzeros.
```

```
CPLEX 12.6.3.0: mipdisplay 2
Found incumbent of value 539974.520000 after 0.05 sec. (48.97 ticks)
MIP Presolve eliminated 7 rows and 66 columns.
MIP Presolve modified 8 coefficients.
Reduced MIP has 2594 rows, 185716 columns, and 4096084 nonzeros.
Reduced MIP has 185716 binaries, 0 generals, 0 SOSs, and 0 indicators.
MIP Presolve eliminated 24 rows and 169 columns.
Reduced MIP has 2570 rows, 185547 columns, and 4090174 nonzeros.
Reduced MIP has 185547 binaries, 0 generals, 0 SOSs, and 0 indicators.
Probing time = 0.30 sec. (60.86 ticks)
Clique table members: 2568.
MIP emphasis: balance optimality and feasibility.
MIP search method: dynamic search.
Parallel mode: deterministic, using up to 16 threads.
Root relaxation solution time = 3.59 sec. (4374.31 ticks)
```

	Nodes					Cuts/		
	Node	Left	Objective	IInf	Best Integer	Best Bound	ItCnt	Gap
*	0+	0			539974.5200	0.0000		100.00%
	0	0	420158.4038	518	539974.5200	420158.4038	41	22.19%
*	0+	0			514589.8400	420158.4038		18.35%
	0	0	420294.0867	608	514589.8400	Cuts: 89	1187	18.32%