

OS实验5-模拟页表

实验目的:

- 1. 创建程序,将以页为基础的虚拟地址翻译成物理地址
- 2. 实现TLB,统计给定地址集上,使用/不使用TLB产生的性能差异
- 3. 实现支持缺页中断,交换的功能

基本设计:

由于我们的程序设计目标比较简单,我们可以认为我们在操作这样一台虚拟机---有64KB的外存,32KB的内存,支持最大64KB的虚拟内存空间,在运行时只加载一个进程,使用"固定分配,局部置换"的方法,将全部内存按256B粒度分割为256个物理页面(frame/block),并为进程初始化一个最大长度为256的页表。

当程序刚开始运行,页表和物理页面被创建,但是页和物理页面间的联系还没有建立起来,物理页面的内容页没有意义。这时在虚拟机里有三个数据结构

1. 页构成的虚拟内存

page	size(byte)
0	256
1	256
255	256

2. 由frame构成的物理内存空间

frame	size(byte)	busy
0	256	false
1	256	false
		false
255	256	false

3. 虚拟机中唯一进程持有的一个页表,现在页表中所有项都不可用 由于虚拟机中只有一个进程,为了方便可以把本来属于页的属性直接绑定到页表上

page	frame	attribute
0	X	not avaliable
	X	not avaliable
255	X	not avaliable

随后,我们运行程序,从address.txt读取到第一个虚拟地址:16916

```
while(fin>>virtualAddr)

pageNum = virtualAddr & 65280;

pageNum = pageNum>>8;

offset = virtualAddr & 255;

visit(pageNum, offset);

}
```

得到pageNum: 66, offset: 20

查找页表,发现对应frame不可用,触发缺页中断,应从外存(BACKING_STORE.bin文件)读取相应frame。

```
frame ReadDisk(int start addr, int size);
newframe = ReadDisk(66*256, 256);
```

接下来检查物理内存空间,我们发现整个物理内存空间都是空闲的,使用最先找到的0号frame。

```
1 memory[0].context = newframe;
2 memory[0].busy = busy;
```

然后我们还需要修改页表

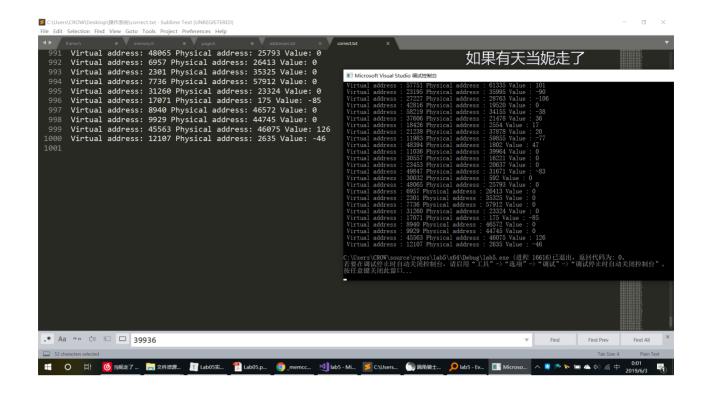
```
pageTable[66].frame = 0;
pageTable[66].attribute = avaliable;
```

然后我们重新进行visit操作,这次在页表中查到66号页表对应0号frame,能获取有效的物理内存,并能读取到数据。

```
1 | physicAddr = 0*256+20 = 20;
```

实验结果

实验目标一(无TLB,物理内存足够,不发生页置换)对应的源代码是"page_adress_translation_1.cpp"文件,其输出与样例"correct.txt"一致。



实验目标二对应"page_adress_translation_2.cpp"。这个版本的代码添加了TLB,并将物理内存缩小到128个 frame。在实验二中,页面的置换使用随机数选择,TLB的更新使用FIFO。在输出时可以发现,当读取超过128个页面时,由于页置换,物理内存地址和correct.txt不再相符,但value仍然相符。

统计结果

TLB hit rate	page fault rate
94.4%	48.6%

```
Virtual address: 50992 Physical address: 6704Value: 0
Virtual address: 10583 Physical address: 27479 Value: 85
Virtual address: 10583 Physical address: 6807Value: 101
Virtual address: 23195 Physical address: 681IValue: -90
Virtual address: 22727 Physical address: 28763 Value: -106
Virtual address: 22727 Physical address: 28763 Value: 0
Virtual address: 42816 Physical address: 19520 Value: 0
Virtual address: 37606 Physical address: 19520 Value: 38
Virtual address: 37606 Physical address: 21478 Value: 36
Virtual address: 18426 Physical address: 2554 Value: 17
Virtual address: 1983 Physical address: 2554 Value: 17
Virtual address: 1983 Physical address: 6868Value: -77
Virtual address: 48394 Physical address: 1802 Value: 47
Virtual address: 30557 Physical address: 1802 Value: 47
Virtual address: 30557 Physical address: 16221 Value: 0
Virtual address: 30557 Physical address: 20637 Value: 0
Virtual address: 49847 Physical address: 2037 Value: 0
Virtual address: 49867 Physical address: 592 Value: 0
Virtual address: 30032 Physical address: 25793 Value: 0
Virtual address: 30032 Physical address: 25793 Value: 0
Virtual address: 301 Physical address: 26413 Value: 0
Virtual address: 301 Physical address: 26793 Value: 0
Virtual address: 301 Physical address: 26909Value: 0
Virtual address: 301 Physical address: 26909Value: 0
Virtual address: 301 Physical address: 3009 Value: 0
Virtual address: 301 Physical address: 3009 Value: 0
Virtual address: 3009 Physical address: 26909Value: 0
Virtual address: 3009 Physical address: 68909Value: 0
Virtual address: 8940 Physical address: 6892Value: 0
Virtual address: 8940 Physical address: 6857Value: 0
Virtual address: 8940 Physical address: 6857Value: 0
Virtual address: 8950 Physical address: 6857Value: 0
Virtual address: 8950 Physical address: 6857Value: 0
Virtual address: 8950 Physical address: 6857Value: 0
Virtual address: 49865 Physical address: 6857Value: 0
Virtual address: 68900 Physical address: 6857Value: 0
Virtual address: 68900 Physical address: 6857Value: 0
Virtual address: 68
```

实验目标3和实验目标源码极度相似,由于实验目标一的代码中已经实现了页的置换(使用随机算法换出),所以只需要更改物理内存大小和添加少量用于统计运算状态的代码。实验目标三对应的是"page_adress_translation_3.cpp"。

统计结果

```
page fault rate
53.2%
```