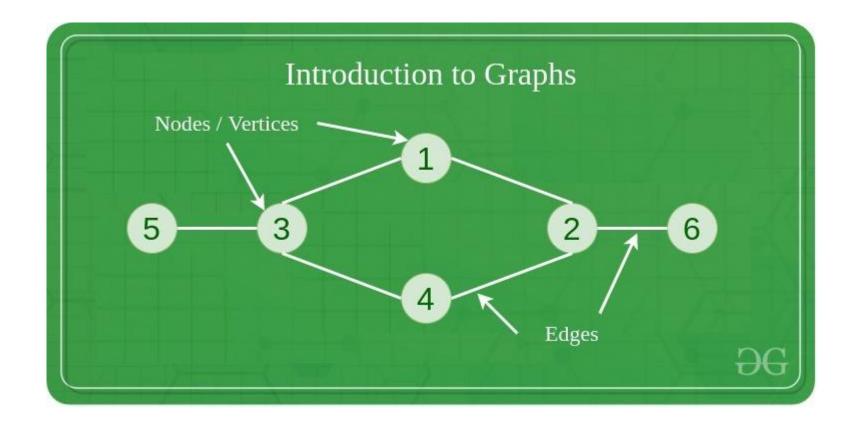
OOP & data struct

14. Graph

BY SOMSIN THONGKRAIRAT

Graph



Graph

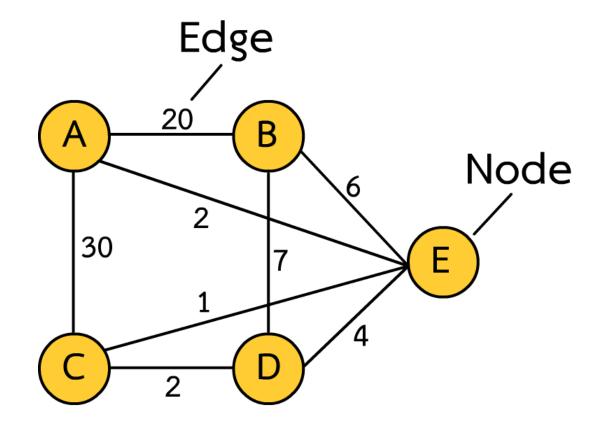
- fundamental non-linear data structure consisting of node(vertex) and path(edge)
- node(vertex) is an item or element in structure
- path(edge) is the information that how each node connected

- เป็น structure พื้นฐานแบบ non-linear ที่ประกอบไปด้วย node(vertex) และ path(edge)
- node(vertex) คือ item หรือ element ใน structure
- path(edge) คือข้อมูลที่บอกว่าแต่ละ node เชื่อต่อกันอย่างไร

Component of graph

- Node contain data object
- Edge contain travel cost between node

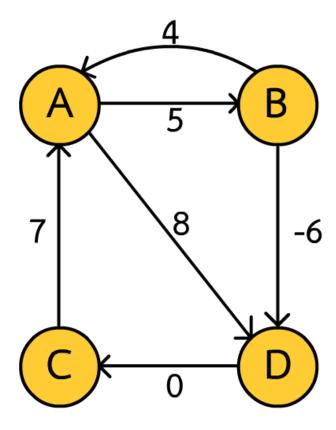
- Node เก็บข้อมูลของแต่ละ object
- Edge เก็บ cost ของการเดินทางระหว่าง node



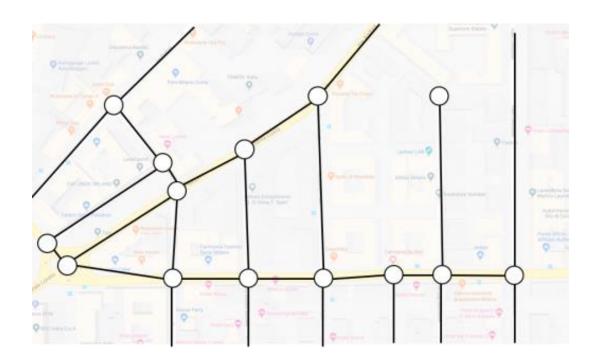
type of graph

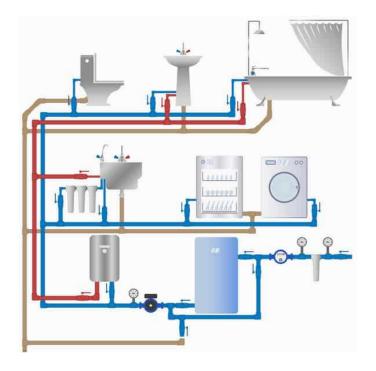
- There are several type of graph classified by rule
- graph มีหลายชนิดจัดประเภทตามกฎ

- Direct / undirect
- Non-negative edge / floating point edge
- Cyclic graph / Acyclic graph



Depend on purpose

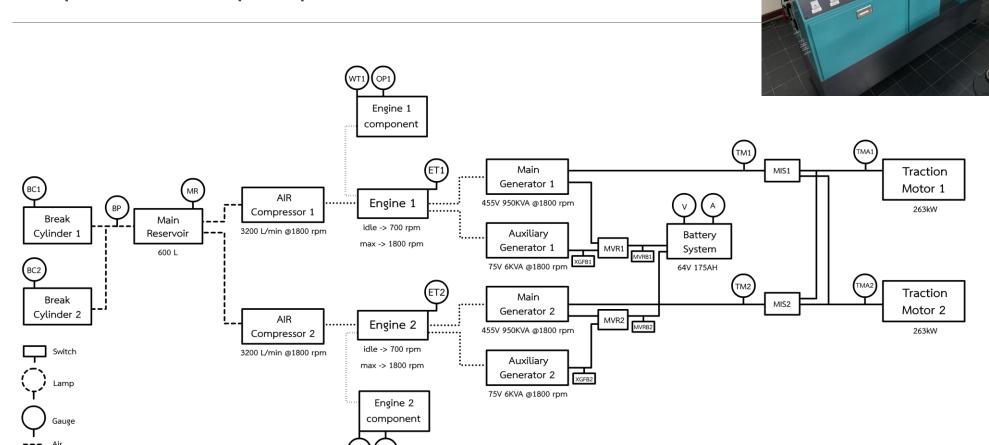




https://dreamcivil.com/systems-of-plumbing/

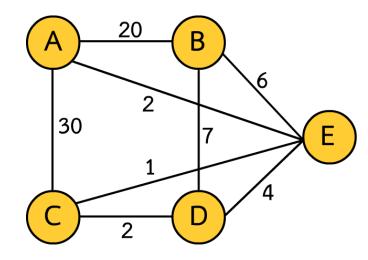
Depend on purpose

Electrical
Other



Our scope

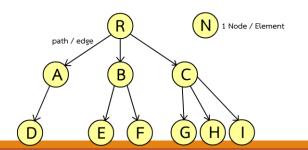
- Non-negative edge undirected cyclicable graph
 - Edge can't be negative cost
 - All edge can use any direction to travel
 - แต่ละ edge จะไม่มี cost ที่ติดลบ
 - ทุก edge ใช้เส้นทางใดก็ได้



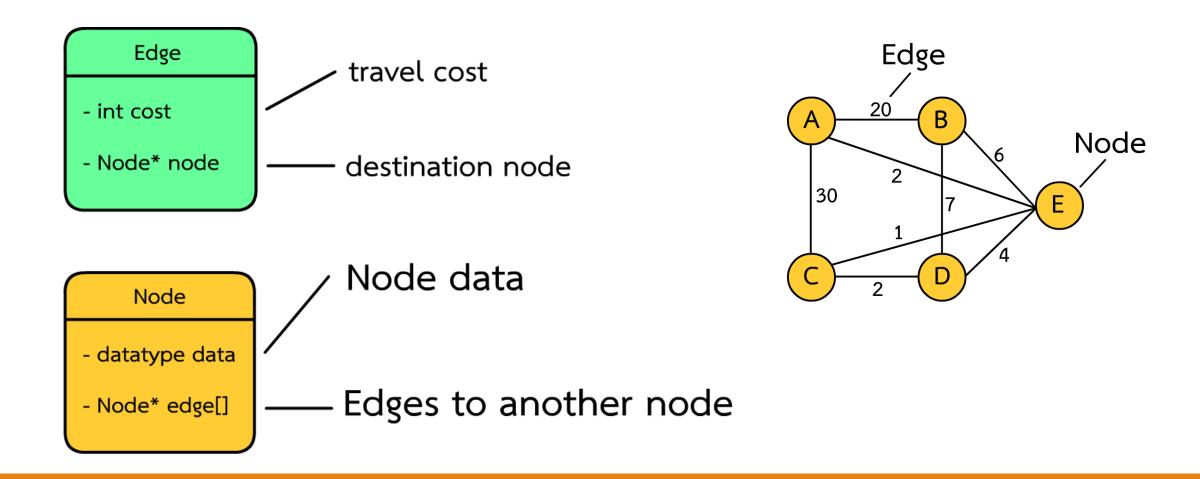
Before we start

- Tree is subset of graphs with restricted rule, So tree always a graph but not all graph will be tree

- Tree เป็น subset ของ graph ที่มีกฎของตัวเอง ดังนั้น tree ทุกตัวจะเป็น graph แต่ไม่ใช่ว่า graph ทุกตัวจะเป็น tree



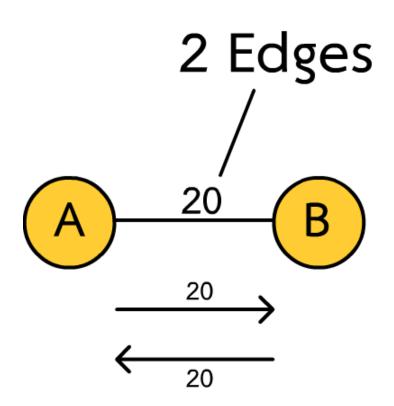
Data structure object base graph (node & edge)

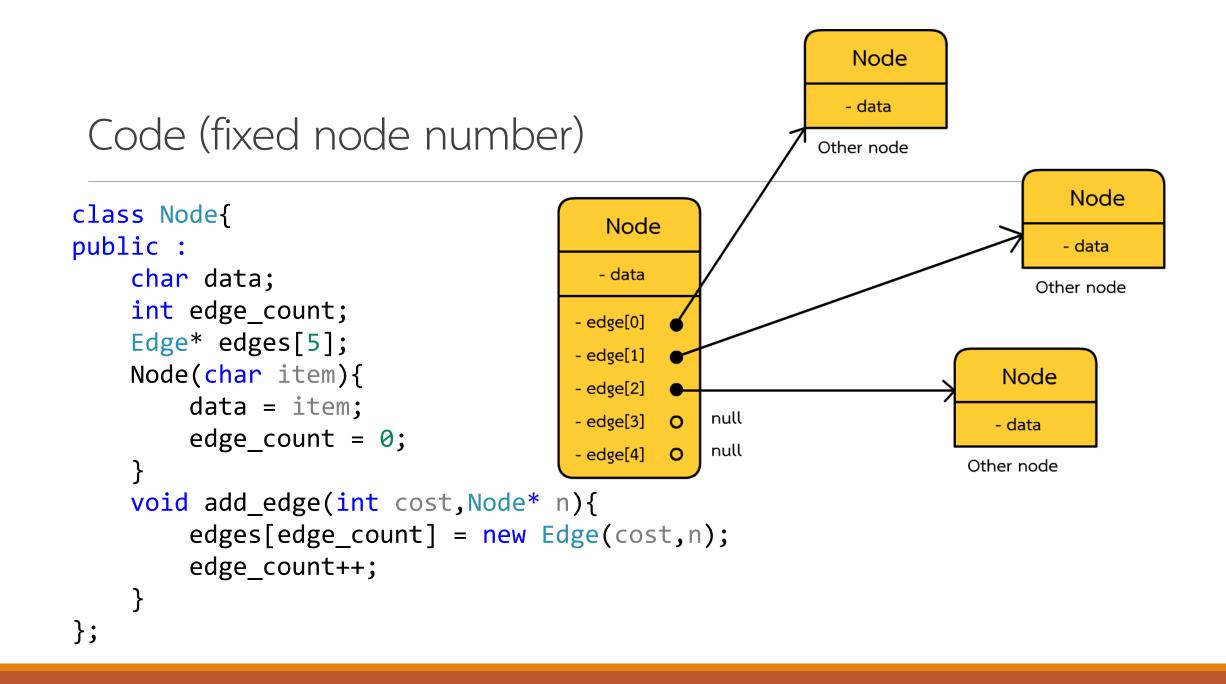


Edge (undirected graph)

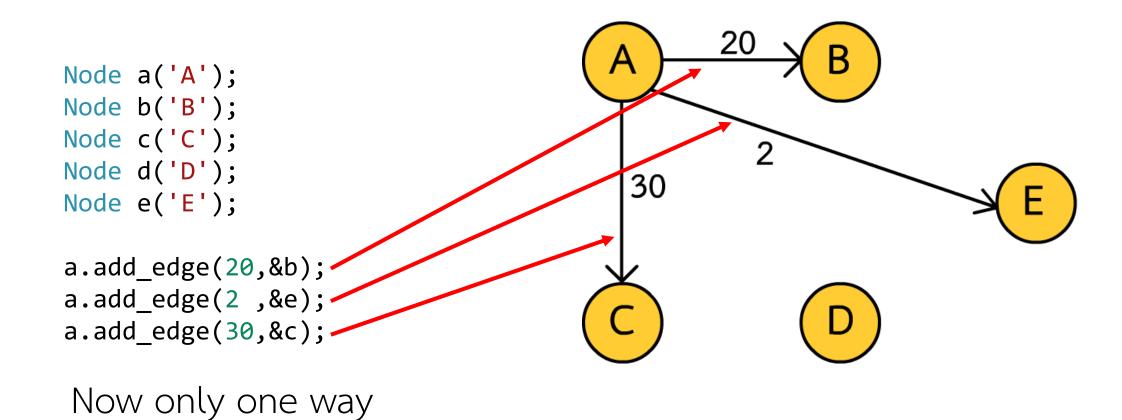
```
class Edge{
    public :
    int cost;
    Node* node;

Edge(int _cost,Node* dest){
        cost = _cost;
        node = dest;
    }
};
```



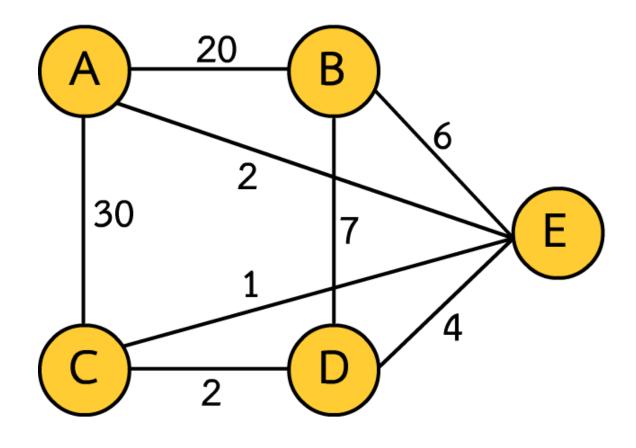


Build graph (each node)



Build graph

```
a.add_edge(20,&b);
a.add_edge(2 ,&e);
a.add_edge(30,&c);
b.add_edge(20,&a);
b.add_edge(7 ,&d);
b.add_edge(6 ,&e);
c.add_edge(30,&a);
c.add_edge(1 ,&e);
c.add_edge(2 ,&d);
d.add_edge(2 ,&c);
d.add_edge(7 ,&b);
d.add_edge(4 ,&e);
e.add_edge(6 ,&b);
e.add_edge(2 ,&a);
e.add_edge(1 ,&c);
e.add_edge(4 ,&d);
```



Print graph (node)

```
void print(){ // method in node
     cout << data << "-> ";
     for(int i=0;i<edge_count;i++){</pre>
           cout << "(" << edges[i]->cost <<","<< (edges[i]->node)->data << ")";</pre>
     cout << endl;</pre>
                                                             Result:
                                                             A \rightarrow (20,B)(2,E)(30,C)
a.print();
b.print();
                                                             B \rightarrow (20,A)(7,D)(6,E)
c.print();
                                                             C \rightarrow (30,A)(1,E)(2,D)
d.print();
                                                             D \rightarrow (2,C)(7,B)(4,E)
e.print();
                                                             E \rightarrow (6,B)(2,A)(1,C)(4,D)
```

Result:

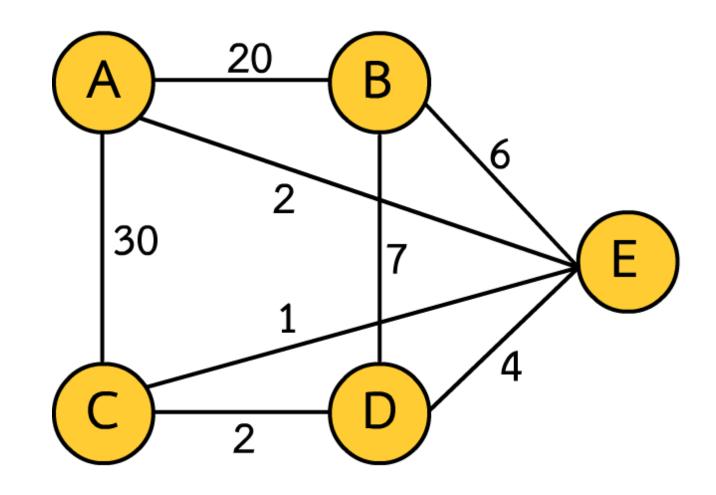
 $A \rightarrow (20,B)(2,E)(30,C)$

 $B \rightarrow (20,A)(7,D)(6,E)$

 $C \rightarrow (30,A)(1,E)(2,D)$

 $D \rightarrow (2,C)(7,B)(4,E)$

 $E \rightarrow (6,B)(2,A)(1,C)(4,D)$



Access

- Can start from any node / สามารถเริ่มใช้จาก node ใดก็ได้

Result:

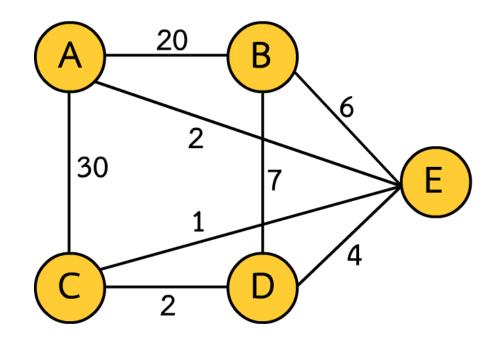
 $A \rightarrow (20,B)(2,E)(30,C)$

 $B \rightarrow (20,A)(7,D)(6,E)$

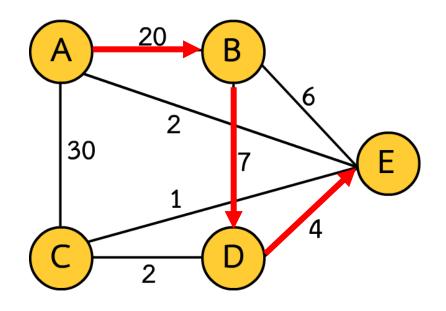
 $C \rightarrow (30,A)(1,E)(2,D)$

 $D \rightarrow (2,C)(7,B)(4,E)$

 $E \rightarrow (6,B)(2,A)(1,C)(4,D)$



example



 $A \rightarrow (20,B)(2,E)(30,C)$

 $B \rightarrow (20,A)(7,D)(6,E)$

 $C \rightarrow (30,A)(1,E)(2,D)$

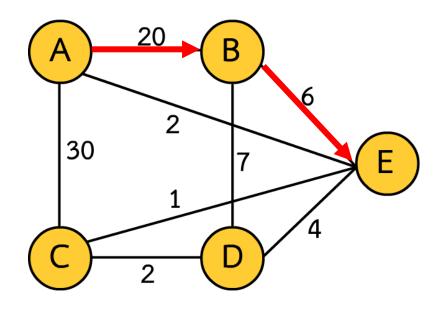
 $D \rightarrow (2,C)(7,B)(4,E)$

 $E \rightarrow (6,B)(2,A)(1,C)(4,D)$

How much total cost for this path from A to E

cout << a.edges[0]->cost + b.edges[1]->cost + d.edges[2]->cost << endl; // 31</pre>

example



 $A \rightarrow (20,B)(2,E)(30,C)$

 $B \rightarrow (20,A)(7,D)(6,E)$

 $C \rightarrow (30,A)(1,E)(2,D)$

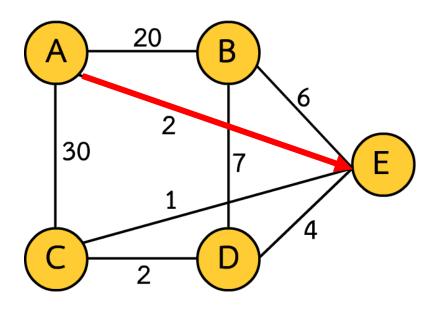
 $D \rightarrow (2,C)(7,B)(4,E)$

 $E \rightarrow (6,B)(2,A)(1,C)(4,D)$

How much total cost for this path from A to E

cout << a.edges[0]->cost + b.edges[2]->cost << endl; // 26</pre>

example



 $A \rightarrow (20,B)(2,E)(30,C)$

 $B \rightarrow (20,A)(7,D)(6,E)$

 $C \rightarrow (30,A)(1,E)(2,D)$

 $D \rightarrow (2,C)(7,B)(4,E)$

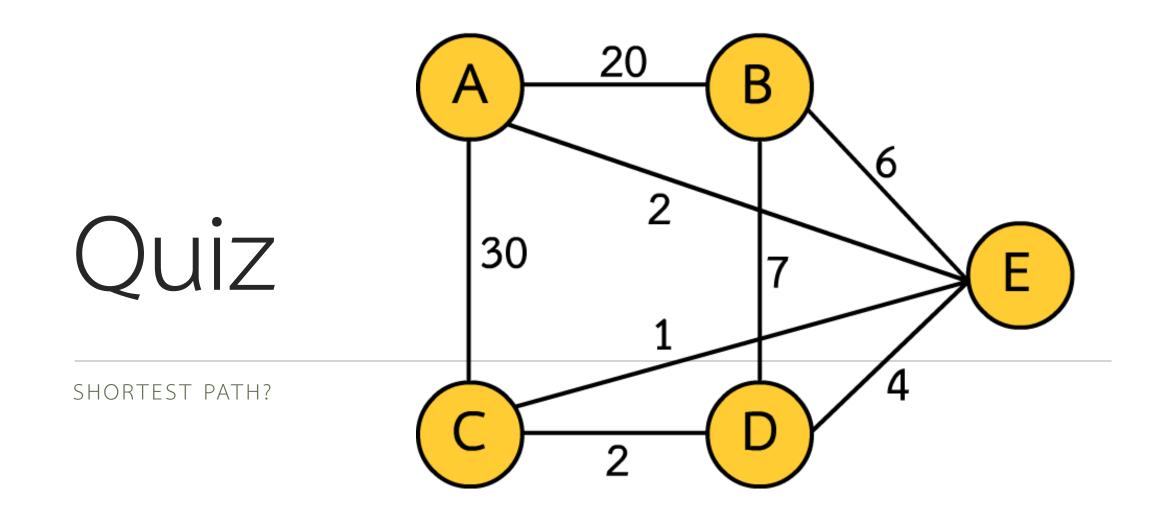
 $E \rightarrow (6,B)(2,A)(1,C)(4,D)$

How much total cost for this path from A to E

cout << a.edges[1]->cost << endl; // 2</pre>

End of graph structure

Нарру?

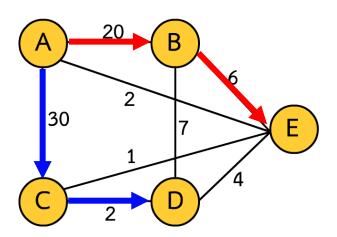


Why graph

- more flexible than tree
- more complex algorithm implementation***
- seem like realistic modelling (such as networking traffic)
- ยืดหยุ่นกว่า tree
- สามารถใช้ algorithm ที่ซับซ้อนได้มากกว่า***
- สามารถ modeling ให้เหมือนกันความเป็นจริงได้ เช่น networking การจราจร

Example program

- second order visited
- List all possibility node that can visit with in 2 move
- แสดง node ทั้งหมดที่เป็นไปได้ ที่สามารถเข้าถึงได้ด้วยการ<u>เดิน</u> 2 ครั้ง

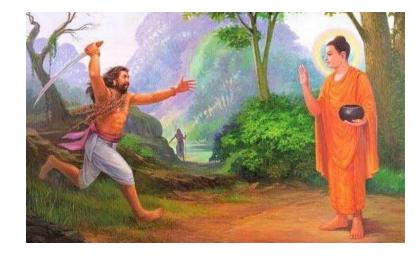


Remark ***

- following algorithm is just a partial algorithm can't use in real life application

- algorithm ต่อไปนี้เป็นเพียงส่วนหนึ่งของ algorithm เท่านั้นไม่สามารถ

นำไปใช้จริงได้



second order visited

```
void second_order_visited(Node* n){
    for(int i = 0;i<n->edge_count;i++){ // like print
        Node* node1 = (n->edges[i])->node;
        int cost1 = (n->edges[i])->cost;
        for(int j=0;j<node1->edge_count;j++){
            Node* node2 = (node1->edges[j])->node;
            int cost2 = cost1 + (node1->edges[j])->cost;
            cout << n->data << "->" << node1->data << "->" << node2->data << " = " << cost2 << end1;
        }
    }
}</pre>
```

second order visited

second_order_visited(&a);

$$A -> B -> A = 40$$

$$A->B->D = 27$$

$$A->B->E = 26$$

$$A -> E -> B = 8$$

$$A -> E -> A = 4$$

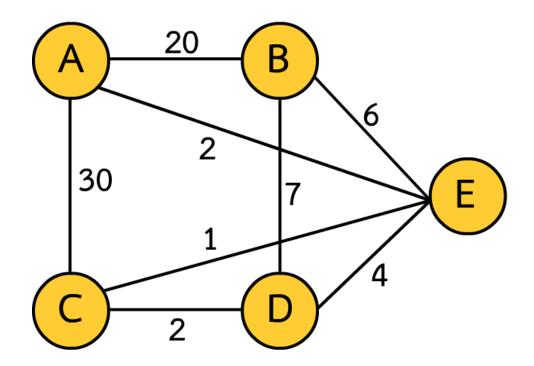
$$A -> E -> C = 3$$

$$A -> E -> D = 6$$

$$A -> C -> A = 60$$

$$A->C->E = 31$$

$$A->C->D = 32$$

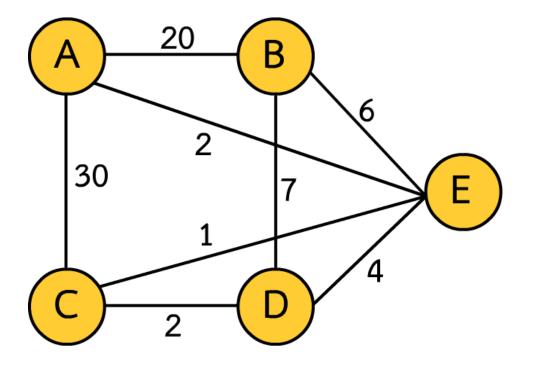


third order visited

```
void third order visited(Node* n){
    for(int i = 0;i<n->edge count;i++){ // like print
        Node* node1 = (n->edges[i])->node;
        int cost1 = (n->edges[i])->cost;
        for(int j=0;j<node1->edge count;j++){
            Node* node2 = (node1->edges[j])->node;
            int cost2 = cost1 + (node1->edges[j])->cost;
            for(int k=0;k< node2->edge_count;k++){
                Node* node3 = (node2->edges[k])->node;
                int cost3 = cost2 + (node2->edges[k])->cost;
                cout << n->data << "->" << node1->data << "->" << node2->data
                     << "->" << node3->data << " = " << cost3 << end1;</pre>
```

third order visited

A->B->A->B=60	A -> E -> A -> E = 6	A->C->E->D = 35
A->B->A->E = 42	A -> E -> A -> C = 34	A->C->D->C = 34
A->B->A->C = 70	A -> E -> C -> A = 33	A->C->D->B=39
A->B->C=29	A -> E -> C -> E = 4	A->C->D->E = 36
A->B->D->B = 34	A->E->C->D = 5	
A->B->D->E = 31	A->E->D->C=8	
A->B->E->B=32	A->E->D->B=13	
A->B->E->A=28	A->E->D->E = 10	
A->B->C=27	A->C->A->B = 80	
A->B->E->D = 30	A->C->A->E = 62	
A -> E -> B -> A = 28	A->C->A->C = 90	
A->E->B->D = 15	A->C->E->B=37	
A -> E -> B -> E = 14	A->C->E->A=33	
A -> E -> A -> B = 24	A->C->E->C=32	



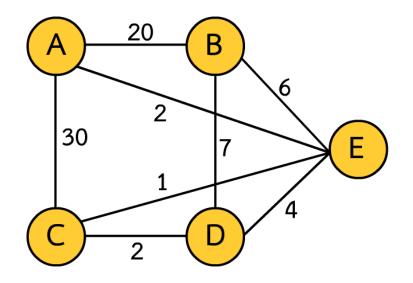
Forth? Fifth? Sixth?

n order visited = graph traversal

- inorder
- postorder
- preorder

nearest neighbor graph

- find nearest node (in first order)
- หา node ที่อยู่ใกล้ที่สุด (ใน first order)



Who is nearest node of 'B'? ใครอยู่ใกล้ B ที่สุด ?

Code

```
Node* nearest_node(Node* n){
    int nearest_idx = 0;
    for(int i=0;i<n->edge_count;i++){
        if(n->edges[i]->cost < n->edges[nearest_idx]->cost){
            nearest_idx = i;
        }
    }
    return n->edges[nearest_idx]->node;
}
```

*Like find minimum number in set

Code

```
cout << "nearest node of A is " << nearest_node(&a)->data << endl;</pre>
cout << "nearest node of B is " << nearest_node(&b)->data << endl;</pre>
cout << "nearest node of C is " << nearest_node(&c)->data << endl;</pre>
cout << "nearest node of D is " << nearest_node(&d)->data << endl;</pre>
cout << "nearest node of E is " << nearest node(&e)->data << endl;</pre>
Result:
nearest node of A is E
                                                    20
nearest node of B is E
nearest node of C is E
nearest node of D is C
nearest node of E is C
```

question

- what is the shortest path between node x and y with exact 3 move

- เส้นทางใหนที่สั้นที่สุดระหว่าง node x และ y หากกำหนดให้ ต้องเดินทั้งหมด 3 ครั้ง

Brute force algorithm

- travel to the z order visited and collect minimum cost
- ท่อง node โดยใช้วิธี z order visited และเก็บวิธีที่ cost น้อยที่สุด

What if sixth order Really !!?

Easy question (if third order)

- Traversal in third order visited method and collect minimum cost visited node

- ใช้วิธี third order visited และเก็บครั้งที่ใช้ cost น้อยที่สุด

code

```
int third_shortest_path(Node* n, char target){
    int min cost = -1;
    for(int i = 0;i<n->edge_count;i++){ // like print
        Node* node1 = (n->edges[i])->node;
        int cost1 = (n->edges[i])->cost;
        for(int j=0;j<node1->edge count;j++){
            Node* node2 = (node1->edges[j])->node;
            int cost2 = cost1 + (node1->edges[j])->cost;
            for(int k=0;k< node2->edge_count;k++){
                Node* node3 = (node2->edges[k])->node;
                int cost3 = cost2 + (node2->edges[k])->cost;
                if(node3->data == target){
                    if(min cost == -1 || cost3 < min cost){</pre>
                        min_cost = cost3;
    return min_cost;
```

code

```
cout << "Shortest from A to E is cost " << third_shortest_path(&a, 'E') << endl;
cout << "Shortest from A to C is cost " << third_shortest_path(&a, 'C') << endl;
cout << "Shortest from B to C is cost " << third_shortest_path(&b, 'C') << endl;
cout << "Shortest from A to B is cost " << third_shortest_path(&a, 'B') << endl;
Shortest from A to E is cost 4
Shortest from A to C is cost 8
Shortest from B to C is cost 12
Shortest from A to B is cost 13</pre>
```

What is big theta for third shortest path

 $-n^3!!??$

Challenge question

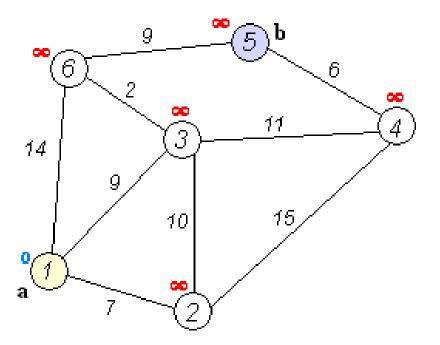
- what is the shortest path between node x and y in z move

- เส้นทางใหนที่สั้นที่สุดระหว่าง node x และ y หากกำหนดให้ เดินได้ z ครั้ง

What is big theta for n path

Dijkstra's algorithm

O(n log n) where n is the number of veritces



Conclude

- Component of graph
- fundamental graph algorithm