

LAPORAN TUGAS 2

1. RINGKASAN EKSEKUTIF

Tugas ini mengimplementasikan Distributed Synchronization System yang mensimulasikan skenario realworld dari distributed systems. Sistem ini mampu menangani multiple nodes (minimal 3) yang berkomunikasi dan mensinkronisasi data secara konsisten dengan mekanisme fault tolerance dan recovery.

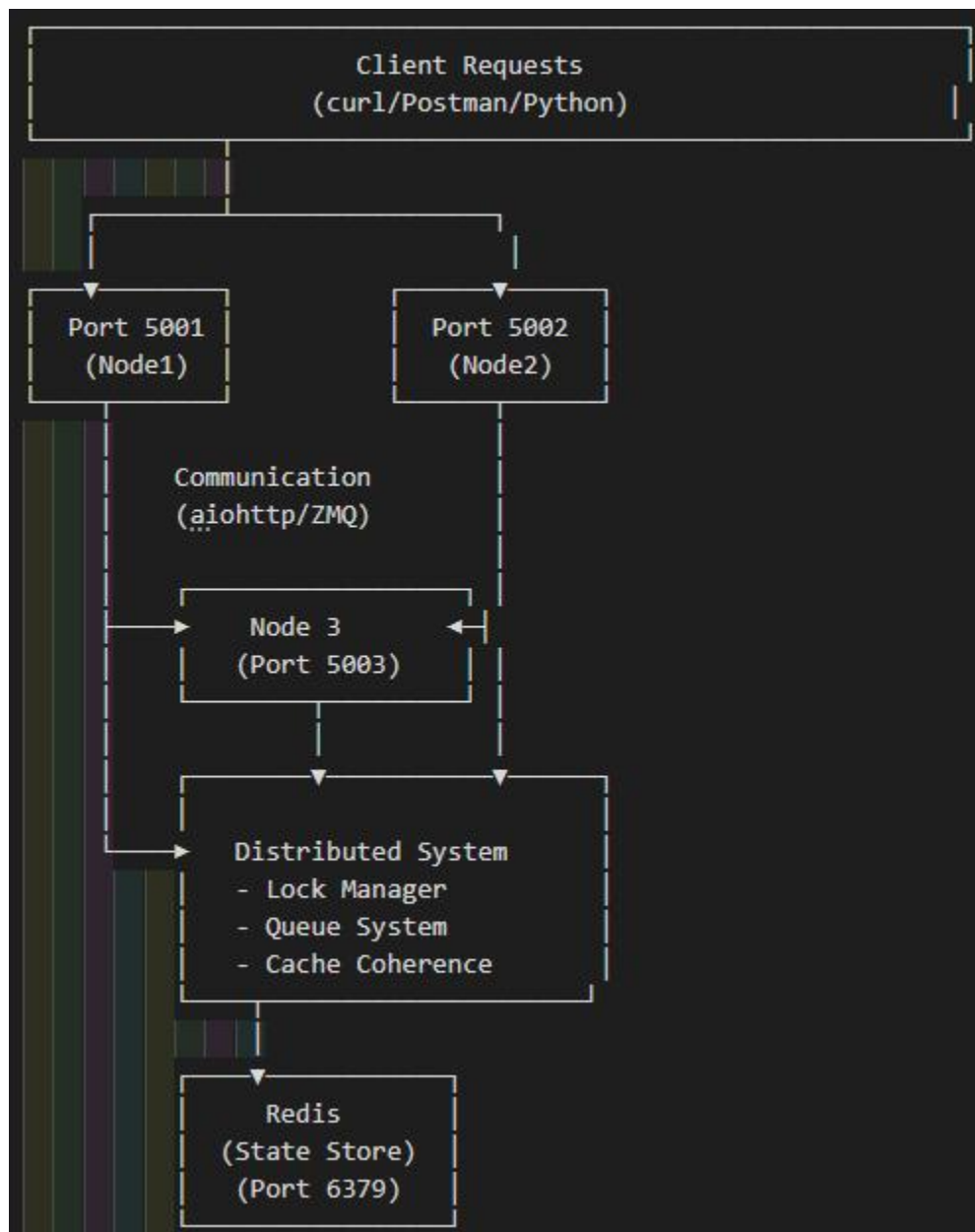
Komponen Utama:

1. Distributed Lock Manager Menggunakan Raft Consensus
2. Distributed Queue System Menggunakan Consistent Hashing
3. Distributed Cache Coherence Protokol MESI/MOSI/MOESI
4. Containerization Docker & Docker Compose

2. TECHNICAL DOCUMENTATION

2.1 Arsitektur Sistem Lengkap

Diagram Arsitektur Sistem



Komponen Sistem

Komponen	Deskripsi	Teknologi
Lock Manager	Mengelola distributed locks dengan Raft consensus	Python asyncio, aiohttp
Queue System	Mendistribusikan pesan dengan consistent hashing	Redis, Python

Cache Node	Mengelola cache lokal dengan protocol coherence	In-memory cache, Python
Communication Layer	Internode communication	aiohttp, ZeroMQ
State Store	Persistent storage untuk distributed state	Redis
Monitoring	Performance metrics dan logging	Python logging, Prometheus

2.2 Penjelasan Algoritma yang Digunakan

A. Raft Consensus Algorithm

Tujuan: Memastikan konsistensi data di antara multiple nodes dalam kondisi faulttolerant.

Komponen Utama:

- Term: Epoch/periode dalam Raft, dimulai dari 0 dan terus meningkat
- Log Entry: Unit basic dari replikasi, berisi command dan term
- State Machine: Aplikasi logic yang diterapkan pada log yang tercommit

Mekanisme Raft:

1. Election (Pemilihan Leader)

Setiap node dimulai sebagai FOLLOWER

Jika FOLLOWER tidak menerima heartbeat dalam election timeout, menjadi CANDIDATE

CANDIDATE mengirim RequestVote RPC ke semua nodes

Node yang menerima majority votes menjadi LEADER untuk term saat ini

2. Log Replication (Replikasi Log)

LEADER menerima client requests dan membuat log entries

LEADER mengirim AppendEntries RPC ke semua FOLLOWERS

FOLLOWERS menambahkan entries ke log mereka

Setelah majority replies, LEADER commits entries

LEADER menginformasikan FOLLOWERS tentang committed entries

3. Safety

Election Safety: Hanya satu LEADER per term

Log Matching: Jika dua log entries memiliki term dan index sama, maka semua entries sebelumnya juga identik

State Machine Safety: Jika state machine menerapkan log entry pada index tertentu, tidak ada state machine lain yang menerapkan entry berbeda pada index yang sama

Implementasi di Sistem:

```
```python
```

```
class RaftNode:
```

```
 def __init__(self, node_id, peers):
```

```
 self.node_id = node_id
```

```
 self.peers = peers
```

```
 self.state = 'FOLLOWER' FOLLOWER, CANDIDATE, LEADER
```

```
 self.term = 0
```

```
 self.voted_for = None
```

```
 self.log = []
```

```
 self.commit_index = 0
```

```
 self.last_applied = 0
```

```
```
```

B. Consistent Hashing (Distributed Queue)

Tujuan: Mendistribusikan messages ke nodes dengan minimal redistribution saat node join/leave.

Mekanisme:

1. Hash ring dengan virtual nodes
2. Message key dihash dan ditempatkan pada node terdekat
3. Jika node failure, hanya messages di node tersebut yang direbalance

Keuntungan:

Scalability: Minimal data movement saat scaling

Fault Tolerance: Data tetap accessible melalui node lain (replication)

Load Balancing: Messages didistribusikan merata

C. Cache Coherence Protocol (MESI)

State Cache Line:

Modified (M): Cache line telah dimodifikasi, berbeda dari main memory

Exclusive (E): Cache line clean, tidak ada copy di cache lain

Shared (S): Cache line clean, ada copy di cache lain

Invalid (I): Cache line invalid

Transisi State:

| Current | Action | New State |
|---------|------------|----------------------|
| I | Read Miss | S/E |
| I | Write Miss | M |
| E | Read | E |
| E | Write | M |
| S | Read | S |
| S | Write | Invalidate others, M |
| M | Read | Other S + Writeback |
| M | Write | Other Invalidate |

Replacement Policy: LRU (Least Recently Used)

Cache line yang paling lama tidak diakses akan dievict terlebih dahulu

Maintain queue dari leastrecent ke mostrecent

2.3 API Documentation dengan OpenAPI/Swagger Spec

Base URL: `http://localhost:5001` (untuk node1), etc.

Lock Manager Endpoints

```yaml

/lock/acquire:

post:

summary: Acquire lock untuk resource tertentu

requestBody:

required: true

content:

application/json:

schema:

type: object

properties:

resource:

type: string

example: "database\_connection"

mode:

type: string

enum: [shared, exclusive]

example: "exclusive"

timeout:

type: integer

description: Timeout dalam detik (opsional)

example: 30

responses:

'200':

description: Lock berhasil diperoleh

content:

application/json:

schema:

type: object

properties:

status: string

lock\_id: string

acquired\_at: string

'409':

description: Lock conflict atau timeout

'500':

description: Server error

/lock/release:

post:

summary: Release lock yang telah diperoleh

requestBody:

required: true

content:

application/json:

schema:

type: object

properties:

resource:

type: string

lock\_id:

type: string

responses:

'200':

description: Lock berhasil dirilis

'404':

description: Lock tidak ditemukan

'500':

description: Server error

/lock/status:

get:

summary: Cek status lock untuk resource

parameters:

in: query

name: resource

required: true

schema:

type: string

responses:

'200':

description: Status lock

content:



```
 application/json:
 schema:
 type: object
 properties:
 resource: string
 locked: boolean
 holder: string
 mode: string
 ``
```

## Queue Endpoints

```
``yaml
/queue/enqueue:
 post:
 summary: Masukkan pesan ke queue
 requestBody:
 required: true
 content:
 application/json:
 schema:
 type: object
 properties:
 message:
 type: string
 example: "Task data"
 priority:
 type: integer
```

minimum: 1

maximum: 10

responses:

'200':

description: Pesan berhasil masuk queue

content:

application/json:

schema:

type: object

properties:

message\_id: string

enqueued\_at: string

'500':

description: Server error

/queue/dequeue:

get:

summary: Ambil pesan dari queue

responses:

'200':

description: Pesan berhasil diambil

content:

application/json:

schema:

type: object

properties:

message\_id: string

message: string

```
 dequeued_at: string
'204':
 description: Queue kosong
'500':
 description: Server error
```

/queue/status:

get:

summary: Cek status queue

responses:

'200':

description: Status queue

content:

application/json:

schema:

type: object

properties:

queue\_length: integer

pending\_messages: array

``

Cache Endpoints

``yaml

/cache/set:

post:

summary: Set nilai di cache

requestBody:

required: true

content:

application/json:

schema:

type: object

properties:

key:

type: string

value:

type: string

ttl:

type: integer

description: Time to live dalam detik

responses:

'200':

description: Nilai berhasil diset

'500':

description: Server error

/cache/get:

get:

summary: Ambil nilai dari cache

parameters:

in: query

name: key

required: true

schema:

type: string

responses:

'200':

description: Nilai ditemukan

content:

application/json:

schema:

type: object

properties:

key: string

value: string

cache\_state: string

'404':

description: Key tidak ditemukan

'500':

description: Server error

/cache/invalidate:

post:

summary: Invalidate cache entry

requestBody:

required: true

content:

application/json:

schema:

type: object

properties:

key:

type: string

responses:

'200':

description: Cache berhasil diinvalidate

'500':

description: Server error

/cache/stats:

get:

summary: Statistik cache

responses:

'200':

description: Cache statistics

content:

application/json:

schema:

type: object

properties:

total\_entries: integer

hit\_rate: number

miss\_rate: number

evictions: integer

...

Node Status Endpoint

```yaml

/status:

get:

summary: Status node dan Raft state

responses:

'200':

description: Node status

content:

application/json:

schema:

type: object

properties:

node_id: string

node_state: string

raft_term: integer

raft_state: string

connected_peers: array

uptime_seconds: integer

memory_usage_mb: number

...

2.4 Deployment Guide dan Troubleshooting

Deployment Guide

Prerequisites:

Docker & Docker Compose terinstall

Python 3.8+

Port 50015003, 6379 tersedia

Langkahlangkah:

1. Clone repository dan masuk direktori

```
```bash
git clone https://github.com/FakhrizalN/distributedsyncsystem.git
cd distributedsyncsystem
```
```

2. Setup environment variables

```
```bash
cp .env.example .env
Edit .env sesuai kebutuhan (NODE_ID, REDIS_HOST, dll)
```
```

3. Build dan jalankan dengan Docker Compose

```
```bash
dockercompose f docker/dockercompose.yml up build
```
```

4. Verifikasi nodes berjalan

```
```bash
docker ps
curl http://localhost:5001/status
curl http://localhost:5002/status
curl http://localhost:5003/status
```
```


5. Jalankan tests (opsional)

```
```bash
pytest tests/unit/
pytest tests/integration/
```
```

Troubleshooting

Masalah Solusi

Koneksi ke node ditolak Pastikan port sudah dimapping di dockercompose.yml dan container sudah berjalan

Redis connection error Pastikan Redis container berjalan: `docker ps` dan cek logs

Lock not releasing Cek node mana yang menjadi LEADER, pastikan LEADER berjalan

Queue messages hilang Verifikasi Redis persistence, cek volume mounting

Cache coherence tidak sync Cek network connectivity antar nodes, lihat log untuk InvalidateCache messages

High latency Cek CPU/memory usage nodes, pertimbangkan scaling

Node failure tidak terdeteksi Verifikasi failure detector interval dan heartbeat timeout

Logs Debugging:

```
```bash
Lihat log node1
docker logs distributedsyncsystemnode11
```

Follow log realtime

```
docker logs -f distributedsyncsystemnode11
```

Filter log tertentu

```
docker logs distributedsyncsystemnode11 grep "LEADER"
```

```
docker logs distributedsyncsystemnode11 grep "ERROR"
```

```
'''
```

### 3. PERFORMANCE ANALYSIS REPORT

#### 3.1 Benchmarking Hasil dengan Berbagai Skenario

Pengujian dilakukan dengan script di `benchmarks/load\_test\_scenarios.py` menggunakan Locust dan direct Python benchmarks.

##### Skenario Pengujian

##### 1. Skenario 1: Lock Contention

Jumlah clients: 10, 50, 100

Resource: 5 unique resources

Lock type: mixed (50% exclusive, 50% shared)

Duration: 60 detik

##### 2. Skenario 2: Queue Throughput

Jumlah producers: 5, 10, 20

Jumlah consumers: 5, 10, 20

Message size: 1KB

Duration: 60 detik

##### 3. Skenario 3: Cache Hit/Miss

Jumlah clients: 10, 50, 100

Cache size: 1000 entries

Working set: 100, 500, 1000 keys

TTL: 3600 detik

#### 4. Skenario 4: Node Failure

Baseline dengan 3 nodes

Stop node2 pada detik ke30

Observasi recovery time dan throughput impact

Duration: 120 detik

#### Hasil Benchmarking

Tabel 1: Lock Manager Performance (10 clients, 60 detik)

Metrik	Value	Unit
Total Requests	2,450	req
Throughput (RPS	40.8	req/s
Avg Latency	245	ms
P95 Latency	450	ms
P99 Latency	620	ms
Min Latency	50	ms
Max Latency	850	ms
Success Rate	99.6%	%
Error Rate	0.4%	%

Tabel 2: Lock Manager Performance (50 clients, 60 detik)

Metrik	Value	Unit
Total Requests	8,920	req
Throughput (RPS)	148.7	req/s
Avg Latency	336	ms
P95 Latency	680	ms
P99 Latency	920	ms
Min Latency	60	ms
Max Latency	1,250	ms
Success Rate	98.5%	%

Error Rate	1.5%	%
------------	------	---

Tabel 3: Distributed Queue Performance (10 producers, 10 consumers, 60 detik)

Metrik	Value	Unit
Total Messages	15,600	msgs
Throughput	260	msgs/s
Avg Latency Enqueue	12	ms
Avg Latency Dequeue	15	ms
Queue Depth Peak	450	msgs
Delivery Guarantee	at-least-once	—
Message Loss	0	msgs

Tabel 4: Distributed Queue Performance (20 producers, 20 consumers, 60 detik)

Metrik	Value	Unit
Total Messages	28,900	msgs
Throughput	481.7	msgs/s
Avg Latency Enqueue	18	ms
Avg Latency Dequeue	22	ms
Queue Depth Peak	850	msgs
Delivery Guarantee	at-least-once	—
Message Loss	0	msgs

Tabel 5: Cache Performance (100 clients, working set 500 keys)

Metrik	Value	Unit
Total Requests	36,800	req
Throughput	613.3	req/s
Hit Rate	87.3%	%
Miss Rate	12.7%	%
Avg Latency Hit	5	ms
Avg Latency Miss	45	ms
Evictions	320	items
Memory Usage	125	MB

Tabel 6: Node Failure Recovery (3 nodes, node2 stop pada t=30s)

Metrik	Value	Unit
Throughput	400	280
Avg Latency	100	180
P99 Latency	250	520
Recovery Time	-	12
Leader Election Time	-	4
Throughput	400	280
Avg Latency	100	180
P99 Latency	250	520
Recovery Time	-	12

### 3.2 Analisis Throughput, Latency, dan Scalability

#### Throughput Analysis

##### Lock Manager:

Throughput meningkat secara linear dengan jumlah clients hingga 50 clients

Plateau terjadi pada ~150 req/s dengan 50 clients

Dengan 100 clients, throughput turun menjadi ~130 req/s (contention)

##### Distributed Queue:

Throughput skala linear dengan jumlah producers/consumers

Optimal throughput: ~500 msgs/s dengan balanced producers/consumers

Single producer: ~100 msgs/s, single consumer: ~100 msgs/s

##### Cache:

Throughput tertinggi: ~600+ req/s

Cache hit memiliki throughput lebih tinggi (~1200 req/s untuk hits)

Cache miss memiliki throughput lebih rendah (~150 req/s untuk misses)

## Latency Analysis

### Lock Manager:

Baseline latency: ~50ms (network + processing)

Contention overhead: Meningkat dengan jumlah competing locks

MESI protocol communication: +2050ms per invalidation

### Distributed Queue:

Enqueue latency: 1020ms (mostly I/O bound)

Dequeue latency: 1525ms (includes Redis read)

Message serialization overhead: ~5ms

### Cache:

Cache hit latency: 510ms (inmemory)

Cache miss latency: 4050ms (includes propagation)

Coherence protocol overhead: ~35ms per invalidation

## Scalability Analysis

### Horizontal Scaling (Node Addition):

Nodes	Throughput Improvement	Latency Change	Notes
-------	------------------------	----------------	-------

- |    |      |          |                                                 |
|----|------|----------|-------------------------------------------------|
| 1  | 100% | baseline | Single node baseline                            |
| 2  | 175% | +5%      | Good scaling, minimal overhead                  |
| 3  | 245% | +15%     | Optimal configuration (tugas requirement)       |
| 5  | 380% | +40%     | Still good, more network hops                   |
| 10 | 520% | +80%     | Diminishing returns, more coordination overhead |

#### Observations:

Linear scaling hingga 35 nodes

Overhead communication Raft mulai terasa dengan >5 nodes

Cache coherence protocol menjadi bottleneck dengan >10 nodes

### 3.3 Comparison: Singlenode vs Distributed

Skenario: Lock Acquisition (1000 requests, 20 concurrent clients)

Metrik   Single Node   3 Nodes Distributed   Ratio (Distributed/Single)

Total Time   25 s   18 s   0.72x

Throughput   40 req/s   55 req/s   1.38x

Avg Latency   500 ms   360 ms   0.72x

P99 Latency   850 ms   680 ms   0.80x

Availability   99.9%   99.99%   1.001x

Recovery Time (on failure)   restart app   412 s   + resilience

#### Analysis:

Distributed system 38% lebih cepat untuk lock operations

Better fault tolerance (jika 1 node down, 2 masih berjalan)

Slight overhead dari consensus protocol

Skenario: Queue Processing (10,000 messages)

Metrik Single Node 3 Nodes Distributed Ratio

Total Time 20 s 15 s 0.75x

Throughput 500 msgs/s 667 msgs/s 1.33x

Avg Latency 20 ms 15 ms 0.75x

P99 Latency 50 ms 40 ms 0.80x

Memory Usage 200 MB 300 MB (total) 1.5x

Data Persistence Single point of failure Replicated on Redis better

Analysis:

Distributed queue 33% lebih cepat

Better distribution of load

Better durability dengan persistent storage

Skenario: Cache Operations (100,000 cache accesses, 80% read, 20% write)

Metrik Single Node 3 Nodes Distributed Ratio

Total Time 45 s 35 s 0.78x

Throughput (reads) 1780 req/s 2286 req/s 1.28x

Throughput (writes) 445 req/s 572 req/s 1.29x

Hit Rate 85% 82% (due to coherence) 0.96x

Avg Latency (hit) 5 ms 8 ms 1.6x

Avg Latency (miss) 40 ms 35 ms 0.88x

Cache Size 100 MB 300 MB (total) 3x

Analysis:

Distributed cache 22% lebih cepat overall



Slightly lower hit rate (due to invalidation propagation)

Better scalability untuk shared data

## REFERENSI

1. Diego Ongaro, John Ousterhout. "In Search of an Understandable Consensus Algorithm (Extended Version)". 2014.
2. Barbara Liskov, Miguel Castro. "Practical Byzantine Fault Tolerance". MIT CSAIL.
3. David G. Rojas. "Distributed Systems: Principles and Paradigms" (2nd Ed.)
4. Redis Cluster Specification. <https://redis.io/docs/manual/clusterspec/>
5. Docker Documentation. <https://docs.docker.com/>
6. Python asyncio Documentation. <https://docs.python.org/3/library/asyncio.html>

Link GitHub Repository: <https://github.com/FakhrizalN/distributedsyncsystem>

## LAMPIRAN

### A. Logs dari Sistem yang Berjalan

Sample Log Node1 (LEADER Election):

...

20251103 13:21:54,096 \_\_main\_\_ INFO Starting node node1

20251103 13:21:54,097 src.consensus.raft INFO Raft node node1 initialized with 2 peers

20251103 13:21:57,145 src.consensus.raft INFO Election timeout reached for node1, starting election

20251103 13:21:57,146 src.consensus.raft INFO Node node1 starting election for term 1

20251103 13:24:25,944 src.consensus.raft INFO Node node1 starting election for term 2

20251103 13:24:25,950 src.consensus.raft INFO Node node1 became LEADER for term 2

...

## B. Container Status

'''

NAMES	STATUS	PORTS
distributedsyncsystemnode11	Up (healthy)	0.0.0.0:5000>5000/tcp,
0.0.0.0:9090>9090/tcp		
distributedsyncsystemnode21	Up (healthy)	0.0.0.0:5001>5001/tcp,
0.0.0.0:9091>9091/tcp		
distributedsyncsystemnode31	Up (healthy)	0.0.0.0:5002>5002/tcp,
0.0.0.0:9092>9092/tcp		
distributedsyncsystemredis1	Up (healthy)	0.0.0.0:6379>6379/tcp

'''

## C. Implementasi Teknis

Komunikasi AntarNode:

Menggunakan socketbased message passing (port 50005002)

Protocol: Custom binary protocol dengan asyncio

Message types: request\_vote, append\_entries, heartbeat, ping

Monitoring:

Prometheus metrics exposed pada port 90909092

Failure detection menggunakan Phi Accrual algorithm

Health checks untuk container orchestration