常系数非齐次线性递推

求 $f_n = H(n) + \sum_{i=1}^k a_i imes f_{n-i}$,其中deg(H(x)) = m。

仍然考虑齐次递推的方法,构造转移矩阵。

定义初始状态为 $St = \{f_{k-1}, f_{k-2}, \dots, f_0, n^m, n^{m-1}, \dots, n^0\}$ 。

$$A=\left(egin{array}{cc} A_{1k imes k} & 0_{k imes(m+1)} \ A_{2(m+1) imes k} & A_{3(m+1) imes(m+1)} \end{array}
ight)$$
,其中

$$A_{1k imes k} = \left(egin{array}{cccccc} a_1 & 1 & 0 & 0 & \dots & 0 \ a_2 & 0 & 1 & 0 & \dots & 0 \ a_3 & 0 & 0 & 1 & \dots & 0 \ dots & dots & dots & dots & dots & dots \ a_k & 0 & 0 & 0 & \dots & 0 \end{array}
ight)$$

$$A_{2(m+1) imes k} = \left(egin{array}{cccc} h_m & 0 & 0 & \dots & 0 \ h_{m-1} & 0 & 0 & \dots & 0 \ dots & dots & dots & dots & dots \ h_0 & 0 & 0 & \dots & 0 \end{array}
ight)$$

其中 A_3 利用了二项式展开,将 $n^m o (n+1)^m$ 。易知最后要求的是 $(StA^n)_k$ 。

考虑特征多项式 g,由特征多项式的定义,以及分块矩阵的行列式特点,得 $g(x)=|xI-A_1||xI-A_3|$ 。有齐次递推得到的式子,有 $g_1(x)=-x^n+\sum_{i=0}^{k-1}a_{k-i}x^i$,而 $|xI-A_3|$ 只有主对角线上有值,故 $g_2(x)=(x-1)^{m+1}$,二项式系数求一下即可。

 $g(x)=g_1(x)*g_2(x)$,卷一下既得特征多项式的系数,同时也知道了 deg(g)=m+1+k,故还需递推出 $f_k,f_{k+1},\ldots,f_{k+m}$ 的值。

设 F 为 f_i 的 OGF,A 为递推系数 a_i 的 OGF,P 为剩余项的 OGF。于是有 $F=F*A+P\Rightarrow F=\frac{P}{1-A}$ 当 $i\in[k,k+m]$ 时,由递推式子的定义有 $[x^i]P(x)=H(i)$,多点求值即可。而当 $i\in[0,k-1]$,由于 **初始值是题目随机给的**,P 中要减去 F*A 产生的项,即 $[x^i]P(x)=a_i-F*A(i)$,卷积一下求个逆得到 f_0,f_1,\ldots,f_{k+m} 。最后套个齐次线性递推的模板就出来了。

```
//code sourced from kinesis
#include<bits/stdc++.h>
using namespace std;
#define _{REP(i,a,b)} for(int i = (a); i >= (int)(b); --i)
#define REP(i,a,b) for(int i = (a); i \leftarrow (int)(b); ++i)
#define UREP(i,u) for(int i = p[(u)]; i + 1; i = edge[i].next)
//iterator: for(int u:x),x is container
#define x(p) (p).first
#define y(p) (p).second
#define pii pair<int,int>
#define mp(x,y) make_pair((x), (y))
#define sign(x) (fabs(x) < eps ? 0 : ((x) > 0 ? 1 : -1))
#define 11 long long
#define L7 __int128//1<<7 bit
#define ull unsigned long long
const int inf = 0x3f3f3f3f;
const 11 inff = 0x3f3f3f3f3f3f3f3f3f3;
const int mod = 998244353;
const double eps = 1e-9;
#define ri1(x) scanf("%d", &(x))
#define ri2(x,y) scanf("%d%d", &(x), &(y))
#define ri3(x,y,z) scanf("%d%d%d", &(x), &(y), &(z))
#define ri4(a,b,c,d) scanf("%d%d%d%d", &(a), &(b), &(c), &(d))
//#define Debug
#ifdef Debug
#endif // Debug
const int maxn = 131072; //由于倍增的时候有30000*4的多项式存在, 故最长长度开4倍
namespace NTT//优化过的ntt,使用时记得初始化
{
    const int p = 998244353, g = 3;
    int w[maxn<<2],inv[maxn<<2],r[maxn<<2],last;</pre>
    int mod(int x){return x \ge p ? x - p : x;}
    11 qp(11 base,11 n)
    {
        base %= p;
        11 \text{ res} = 1;
        while(n){
            if(n&1) (res *= base) %= p;
            (base *= base) %= p;
            n >>= 1;
        }
        return res;
    }
    void init()
        int lim = maxn << 1;//最长数组的两倍
        inv[1] = 1;
```

```
for(int i=2; i<=\lim_{i\to +} i++) inv[i] = mod(p - 111 * (p / i) * inv[p%i] % p);
        for(int i=1;i<lim;i<<=1)</pre>
        {
             int wn = qp(g,(p-1) / (i << 1));
             for(int j=0,ww=1;j< i;j++,ww=111*ww*wn%p) w[i+j] = ww;
        }
    }
    void ntt(vector<int> &f,int n,int op)
        if(last!=n)
        {
             for(int i=1;i<n;i++) r[i] = (r[i>>1]>>1)|((i&1)?(n>>1):0);
             last=n;
        }
        for(int i=1;i<n;i++) if(i<r[i])swap(f[i],f[r[i]]);</pre>
        for(int i=1;i<n;i<<=1)</pre>
        for(int j=0;j<n;j+=i<<1)</pre>
             for(int k=0;k<i;k++)
                 int x=f[j+k], y=111*f[i+j+k]*w[i+k]%p;
                 f[j+k]=mod(x+y); f[i+j+k]=mod(x-y+p);
             }
        if(op==-1)
        {
             reverse(&f[1],&f[n]);
             for(int i=0;i<n;i++)f[i]=111*f[i]*inv[n]%p;</pre>
        }
    }
}
using NTT::ntt;
11 qp(11 base,11 n)
    base %= mod;
    11 \text{ res} = 1;
    while(n){
        if(n&1) (res *= base) %= mod;
        (base *= base) %= mod;
        n >>= 1;
    return res;
}
void PolyInv(int deg,vector<int> &f,vector<int> &g)//蝴蝶变换的rev[maxn<<2],f.resize(n<<2),g.resi
{
    if(deg==1) \{g[0] = qp(f[0],mod-2); return ;\}
    PolyInv((deg+1)>>1,f,g);
    int \lim = 1;
    while(lim<(2*deg)) lim <<= 1;</pre>
    vector<int> h(lim);
```

```
REP(i,0,deg-1) h[i] = f[i];
    ntt(g,lim,1),ntt(h,lim,1);
    REP(i,0,lim-1) g[i] = 1LL * (2LL + mod - 1LL * g[i] * h[i] % mod) % mod * g[i] % mod;
    ntt(g, lim, -1);
    REP(i,deg,lim-1) g[i] = 0;
}
namespace Comb
{
    11 Finv[maxn+10], fac[maxn+10], inv[maxn+10];
    11 qp(ll base,ll n) {
        11 \text{ res} = 1;
        base%=mod;
        while(n){
             if(n&1) (res *= base) %= mod;
             (base *= base) %= mod;
            n >>= 1;
        }
        return res;
    }
    void init()
    {
        int n = maxn;
        inv[1]=1;
        for(int i=2;i<=n;++i)inv[i]=((mod-mod/i)*inv[mod%i])%mod;</pre>
        fac[0]=Finv[0]=1;
        for(int i=1;i<=n;++i)fac[i]=fac[i-1]*i%mod,Finv[i]=Finv[i-1]*inv[i]%mod;</pre>
    }
    11 C(11 n,11 m)
    {
        if(m<0||m>n)return 0;
        return fac[n]*Finv[n-m]%mod*Finv[m]%mod;
    }
}
namespace _Rec
{
    namespace mul_Eva
        int x[maxn<<1],y[maxn<<1];</pre>
        vector<int> f(maxn<<1),h(maxn<<1),g[maxn<<1],_g,g_(maxn<<1);</pre>
        void Get_g(int p,int l,int r)
        {
             if(l==r){
                 g[p].resize(2);
                 g[p][0] = 1, g[p][1] = (mod - x[1]) \% mod;
                 return ;
             int mid = l + r \gg 1;
            Get_g(p << 1, 1, mid), Get_g(p << 1 | 1, mid+1, r);
             int \lim = 1, \deg = r-1+2;
```

```
while(lim<deg) lim <<= 1;</pre>
        g[p<<1].resize(lim),g[p<<1|1].resize(lim),g[p].resize(lim);</pre>
        ntt(g[p<<1],lim,1),ntt(g[p<<1|1],lim,1);
        REP(i,0,lim-1) g[p][i] = 1LL * g[p<<1][i] * g[p<<1][i] % mod;
        ntt(g[p],lim,-1);//做完以后除了g[1]每一个g都是点值形式
        g[p].resize(deg);
    }
    vector<int> mulT(vector<int> a,vector<int> b,int n)//除去idft之外的卷积转置,需要保证a和b是
    {
        int lim=a.size();
        for(int i=0;i<lim;i++)a[i]=111*a[i]*b[i]%mod;</pre>
        ntt(a,lim,1);
        return vector<int>(&a[0],&a[n]);//取a[0]-a[n-1]
    }
    void solve(int p,int l,int r,vector<int> h)
    {
        if(l==r){y[1]=h[0];return;}//把y作为存答案的数组
        int mid=(l+r)>>1, lim=1;
        while(lim<(r-l+2))lim<<=1;</pre>
        h.resize(lim);
        ntt(h,lim,-1);
        solve(p << 1, 1, mid, mulT(g[p << 1|1], h, mid-l+1));
        solve(p << 1 | 1, mid+1, r, mulT(g[p << 1], h, r-mid));
    void Get_Y(int m)
        Get_g(1,0,m-1);
        _g = g[1];
        PolyInv(m,_g,g_);
        int \lim = 1;
        while(lim<((m+2)<<1)) lim <<= 1;
        g_.resize(lim),f.resize(lim);
        ntt(g_,lim,1);ntt(f,lim,-1);
        solve(1,0,m-1,mulT(g_,f,m));
    }
namespace Recurrence{
    vector<int> g(maxn<<1),gT(maxn<<1),_gT(maxn<<1);</pre>
    void PolyModulo(int n,int m,vector<int> &f,vector<int> &R)//特别修改的多项式取模
   {
        int \lim = 1, \deg = 2*n - m + 10;
        while(lim<(deg)) lim <<= 1;</pre>
        vector<int> fT(lim),Q(lim);
        REP(i,0,n-1) fT[i] = f[n-1-i];
        ntt(fT,lim,1);
        REP(i,0,lim-1) fT[i] = 1LL * fT[i] * _gT[i] % mod;
        ntt(fT,lim,-1);
        REP(i,0,n-m) Q[i] = fT[i];
        reverse(Q.begin(),Q.begin()+n-m+1);
        \lim = 1, \deg = n + 10;
        while(lim<(deg)) lim <<= 1;</pre>
```

}

```
vector<int> f1 = f;
    f1.resize(lim),Q.resize(lim),R.resize(lim);
    ntt(f1,lim,1),ntt(Q,lim,1);
    REP(i,0,lim-1) R[i] = (1LL * f1[i] + mod - 1LL * g[i] * Q[i] % mod) % mod;
    ntt(R,lim,-1);
}
void get_g(int n,int m)
    int lim = 1,deg = 2*n - m + 10;//这里的长度要与mod poly一致
    while(lim<(deg)) lim <<= 1;</pre>
    REP(i,0,m-1) gT[i] = g[m-1-i];
    PolyInv(n-m+1,gT,_gT);
    ntt(_gT,lim,1);
    lim = 1,deg = n + 10;//这里的长度同样要与mod poly一致
    while(lim<(deg)) lim <<= 1;</pre>
    ntt(g,lim,1);
}
vector<int> f(maxn<<1);</pre>
void mul(int k)
{
    int lim = 1;
    while(lim<(k<<1)) lim <<= 1;
    ntt(f,lim,1);
    REP(i,0,lim-1) f[i] = 1LL * f[i] * f[i] % mod;
    ntt(f,lim,-1);
    vector<int> R;
    PolyModulo(2*k+1,k+1,f,R);
    REP(i,0,k-1) f[i] = R[i];
    REP(i,k,lim-1) f[i] = 0;
}
void add(int k)
    _{REP(i,k,1)} f[i] = f[i-1]; f[0] = 0;
}
int st[maxn<<1];//递推系数,初始值,长度均为k
int work(int n,int k)//a_n = sum 1 to k
{
    get_g(2*k+1,k+1);
    f[1] = 1;
    int cur = 1, lg = log2(n), ok = 0;
    _REP(i, lg-1, 0){
        if(k<=(cur<<1)) ok = 1;
        if(!ok) f[cur<<1] = f[cur], f[cur] = 0, cur<<=1;
        else mul(k);
        if(n&(1<<i)){
            if(!ok) f[cur+1] = f[cur], f[cur] = 0, cur+=1;
            else add(k);
        }
    }
```

```
vector<int> R;
        PolyModulo(2*k+1,k+1,f,R);
        REP(i,0,k-1) f[i] = R[i];
        int ans = 0;
        REP(i,0,k-1) (ans += 1LL * f[i] * st[i] % mod) %= mod;
        return (ans + mod) % mod;
    }
}
vector<int> a(maxn<<1),P(maxn<<1);</pre>
int m,k;
void get_P()
{
    int lim = 1;
    while(lim<(k<<1)+2) lim <<= 1;
    vector<int> _A(lim),B(lim);
    REP(i,0,k-1) _A[i] = Recurrence::st[i];
    REP(i,1,k) B[i] = a[i];
    ntt(_A,lim,1),ntt(B,lim,1);
    REP(i, 0, lim-1) _A[i] = 1LL * _A[i] * B[i] % mod;
    ntt( A, lim, -1);
    REP(i,0,k-1) P[i] = (1LL * Recurrence::st[i] - _A[i] + mod) % mod;
    REP(i,0,m) mul_Eva::x[i] = k + i;
    mul_Eva::Get_Y(m+1);
    REP(i,0,m) P[k+i] = mul_Eva::y[i];
}
void get_g()
{
    int op = 1, lim = 1;
    while(lim<(m+k+3)) lim <<= 1;
    vector<int> h1(lim),h2(lim);
    REP(i,0,m+1) \ h1[i] = (1LL * op * Comb::C(m+1,i) + mod) % mod,op = -op;
    REP(i,0,k-1) h2[i] = a[k-i];h2[k] = mod - 1;
    ntt(h1,lim,1),ntt(h2,lim,1);
    REP(i,0,lim-1) Recurrence::g[i] = 1LL * h1[i] * h2[i] % mod;
    ntt(Recurrence::g,lim,-1);
}
void get_st()
    int \lim = 1;
    while(lim<(2*m+2*k+2)) lim <<= 1;
    vector<int> B(lim),_f(lim);
    REP(i,1,k) B[i] = (mod - a[i]) \% mod; B[0] = 1;
    PolyInv(m+k+1,B,_f);
    ntt(_f,lim,1),ntt(P,lim,1);
    REP(i,0,lim-1) P[i] = 1LL * P[i] * _f[i] % mod;
    ntt(P,lim,-1);
    REP(i,0,m+k) Recurrence::st[i] = P[i];
}
int work(int n)
{
    get_P();
```

```
get_g();
        get_st();
        return Recurrence::work(n,m+k+1);
}
int main()
{
   #ifndef ONLINE_JUDGE
   freopen("in.txt","r",stdin);//cf needn't delete this
   #endif // ONLINE_JUDGE
   NTT::init();
   Comb::init();
    int n;
    ri3(n,_Rec::m,_Rec::k);
    REP(i,0,_Rec::k-1) ri1(_Rec::Recurrence::st[i]);
   REP(i,1,_Rec::k) ri1(_Rec::a[i]);
    REP(i,0,_Rec::m) ri1(_Rec::mul_Eva::f[i]);
    printf("%d\n",_Rec::work(n));
    return 0;
}
```